

The Network Layer

lec 20-3 The Distance-Vector Routing

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The Distance-Vector (DV) Routing Algorithm

- Distributed
- each node receives some information from one or more of its *directly attached* neighbors, performs a calculation, and then distributes the results of its calculation back to its neighbors
- Iterative
 - this process continues on until no more information is exchanged between neighbors
- Asynchronous
 - it does not require all of the nodes to operate in lockstep with each other.

Bellman-Ford equation

$d_x(y)$ be the cost of the least-cost path from node x to node y .

$$d_x(y) = \min_v \{c(x, v) + d_v(y)\}$$

where the \min_v in the equation is taken over all of x 's neighbors

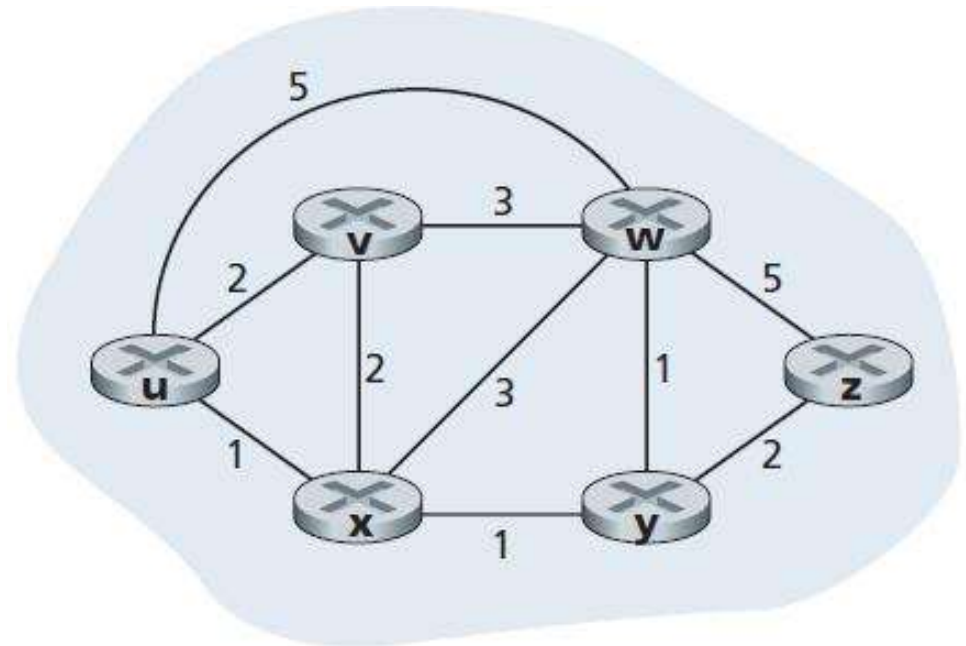
$d_u(z)$: $(u, x) (x, y) (y, z) = 4$

Neighbor of u : v, w and x

1. $c(u, v) + d_v(z)$

2. $c(u, w) + d_w(z)$

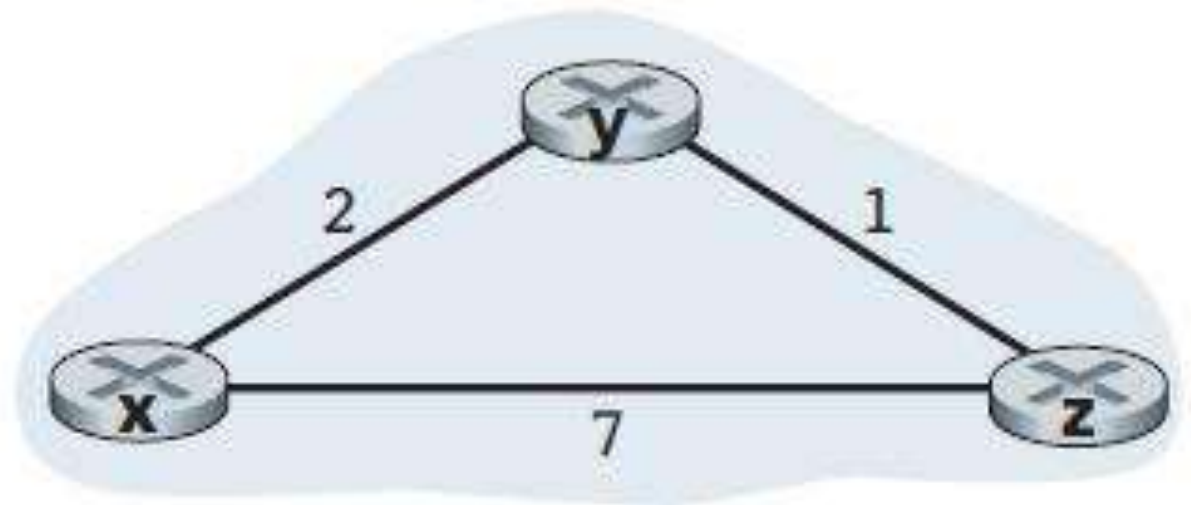
3. $c(u, x) + d_x(z)$



The DV Algorithm

- Each node x begins with $D_x(y)$, an **estimate** of the cost of the least-cost path from x to node y , for all nodes in N .
- $D_x = [D_x(y): y \text{ in } N]$ be node x 's distance vector, which is the vector of cost estimates from x to all other nodes, y , in N .

$$\begin{aligned} D_z &= [D_z(x), D_z(y), D_z(z)] \\ &= [7, 1, 0] \end{aligned}$$



The DV Algorithm

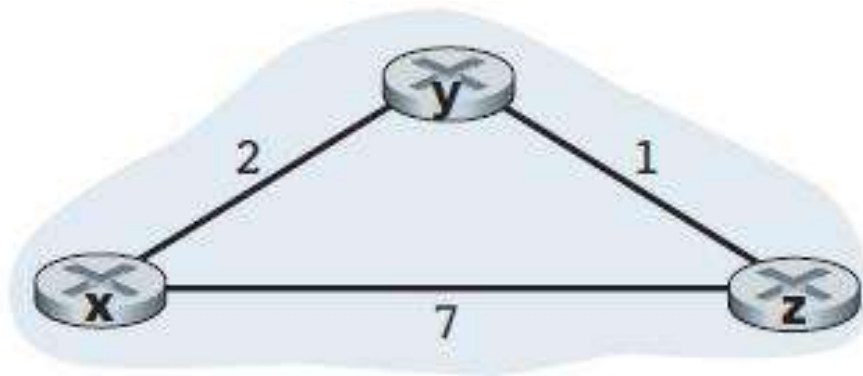
- With the DV algorithm, each node x maintains the following routing information:
 - For each neighbor v , the cost $c(x,v)$ from x to directly attached neighbor, v
 - Node x 's distance vector, that is, $\mathbf{D}_x = [D_x(y): y \text{ in } N]$, containing x 's estimate of its cost to all destinations, y , in N
 - The distance vectors of each of its neighbors, that is, $\mathbf{D}_v = [D_v(y): y \text{ in } N]$ for each neighbor v of x

The DV Algorithm run by node x

```
1  Initialization:
2    for all destinations y in N:
3       $D_x(y) = c(x,y)$  /* if y is not a neighbor then  $c(x,y) = \infty$  */
4    for each neighbor w
5       $D_w(y) = ?$  for all destinations y in N
6    for each neighbor w
7      send distance vector  $D_x = [D_x(y): y \text{ in } N]$  to w
8
9  loop
10    wait (until I see a link cost change to some neighbor w or
11          until I receive a distance vector from some neighbor w)
12
13    for each y in N:
14       $D_x(y) = \min_v \{c(x,v) + D_v(y)\}$ 
15
16    if  $D_x(y)$  changed for any destination y
17      send distance vector  $D_x = [D_x(y): y \text{ in } N]$  to all neighbors
18
19  forever
```


The DV algorithm is *decentralized* and does not use such global information.

Indeed, the only information a node will have is the costs of the links to its directly attached neighbors and information it receives from these neighbors.



Node x table

		cost to		
		x	y	z
from	x	0	2	7
	y	∞	∞	∞
	z	∞	∞	∞

		cost to		
		x	y	z
from	x	0	2	3
	y	2	0	1
	z	7	1	0

		cost to		
		x	y	z
from	x	0	2	3
	y	2	0	1
	z	3	1	0

Node y table

		cost to		
		x	y	z
from	x	∞	∞	∞
	y	2	0	1
	z	∞	∞	∞

		cost to		
		x	y	z
from	x	0	2	7
	y	2	0	1
	z	7	1	0

		cost to		
		x	y	z
from	x	0	2	3
	y	2	0	1
	z	3	1	0

Node z table

		cost to		
		x	y	z
from	x	∞	∞	∞
	y	∞	∞	∞
	z	7	1	0

		cost to		
		x	y	z
from	x	0	2	7
	y	2	0	1
	z	3	1	0

		cost to		
		x	y	z
from	x	0	2	3
	y	2	0	1
	z	3	1	0

Time

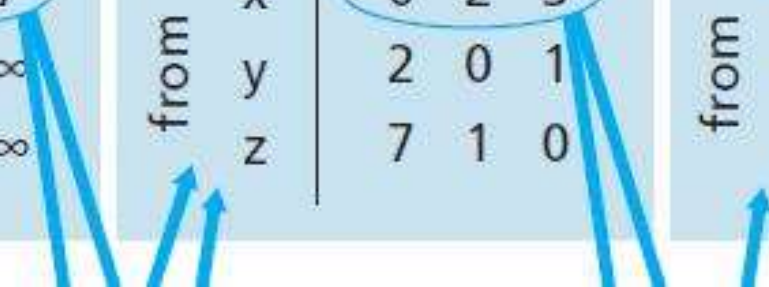
At node x, after receiving updates

Node x table

		cost to		
		x	y	z
from	x	0	2	7
	y	∞	∞	∞
	z	∞	∞	∞

		cost to		
		x	y	z
from	x	0	2	3
	y	2	0	1
	z	7	1	0

		cost to		
		x	y	z
from	x	0	2	3
	y	2	0	1
	z	3	1	0



$$D_x(x) = 0$$

$$D_x(y) = \min\{c(x,y) + D_y(y), c(x,z) + D_z(y)\} = \min\{2 + 0, 7 + 1\} = 2$$

$$D_x(z) = \min\{c(x,y) + D_y(z), c(x,z) + D_z(z)\} = \min\{2 + 1, 7 + 0\} = 3$$

