

CHAPTER 2

Exercise 2.1

(a) $Y = \bar{A}\bar{B} + A\bar{B} + AB$

(b) $Y = \bar{A}\bar{B}\bar{C} + ABC$

(c) $Y = \bar{A}\bar{B}\bar{C} + \bar{A}B\bar{C} + A\bar{B}\bar{C} + A\bar{B}C + ABC$

(d)

$$Y = \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}\bar{C}D + \bar{A}\bar{B}C\bar{D} + \bar{A}B\bar{C}\bar{D} + \bar{A}B\bar{C}D + \bar{A}BC\bar{D} + \bar{A}BCD$$

(e)

$$Y = \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}C\bar{D} + \bar{A}B\bar{C}\bar{D} + \bar{A}B\bar{C}D + \bar{A}BC\bar{D} + \bar{A}BCD + ABC\bar{D} + ABCD$$

Exercise 2.3

(a) $Y = (A + \bar{B})$

(b)

$$Y = (A + B + \bar{C})(A + \bar{B} + C)(A + \bar{B} + \bar{C})(\bar{A} + B + C)(\bar{A} + B + \bar{C})(\bar{A} + \bar{B} + C)$$

(c) $Y = (A + B + \bar{C})(A + \bar{B} + \bar{C})(\bar{A} + \bar{B} + C)$

(d)

$$Y = (A + \bar{B} + C + \bar{D})(A + \bar{B} + C + \bar{D})(A + \bar{B} + \bar{C} + D)(A + \bar{B} + \bar{C} + \bar{D})(\bar{A} + B + C + \bar{D})(\bar{A} + B + \bar{C} + D)(\bar{A} + \bar{B} + C + D)(\bar{A} + \bar{B} + \bar{C} + \bar{D})$$

(e)

$$Y = (A + B + C + \bar{D})(A + B + \bar{C} + D)(A + \bar{B} + C + D)(A + \bar{B} + \bar{C} + \bar{D})(\bar{A} + B + C + D)(\bar{A} + B + \bar{C} + \bar{D})(\bar{A} + \bar{B} + C + D)(\bar{A} + \bar{B} + \bar{C} + \bar{D})$$

Exercise 2.5

(a) $Y = A + \bar{B}$

(b) $Y = \bar{A}\bar{B}\bar{C} + ABC$

(c) $Y = \bar{A}\bar{C} + A\bar{B} + AC$

(d) $Y = \bar{A}\bar{B} + \bar{B}\bar{D} + ACD$

(e)

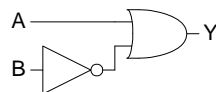
$$Y = \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}CD + \bar{A}B\bar{C}\bar{D} + \bar{A}BC\bar{D} + A\bar{B}\bar{C}\bar{D} + A\bar{B}C\bar{D} + AB\bar{C}\bar{D} + ABCD$$

This can also be expressed as:

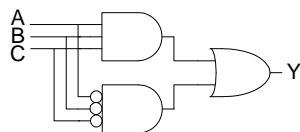
$$Y = (\bar{A} \oplus B)(\bar{C} \oplus D) + (A \oplus B)(C \oplus D)$$

Exercise 2.7

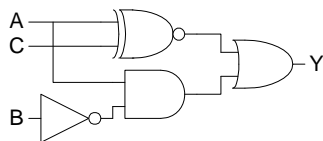
(a)



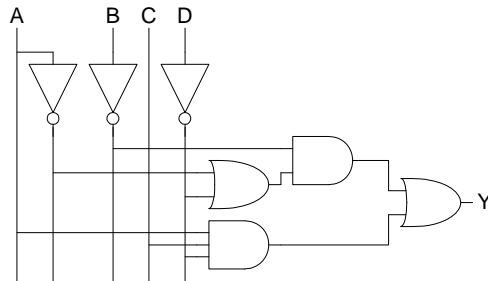
(b)



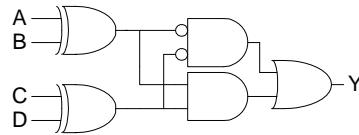
(c)



(d)



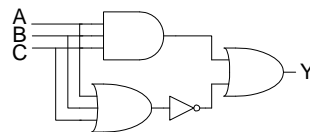
(e)



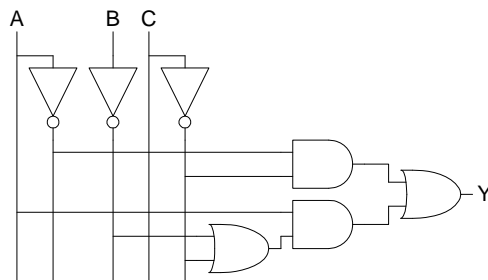
Exercise 2.9

(a) Same as 2.7(a)

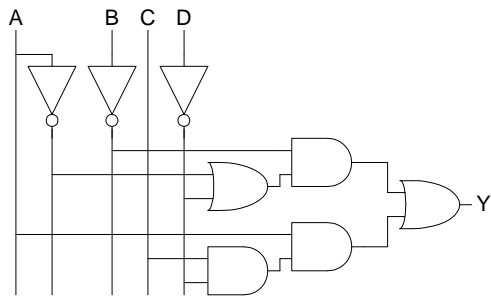
(b)



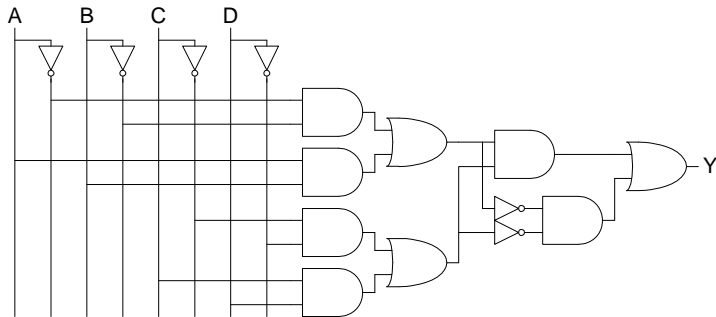
(c)



(d)

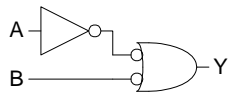


(e)

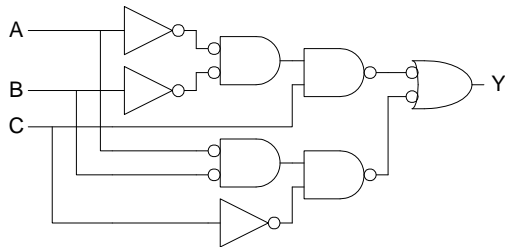


Exercise 2.11

(a)



(b)

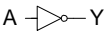


The diagram shows a logic circuit for a 3-input majority gate. It has three inputs: A, B, and C. The circuit consists of the following components and connections:

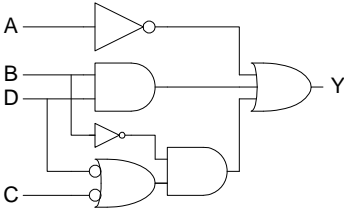
- Two 3-input AND gates at the top. The first 3-input AND gate has inputs A, B, and C. The second 3-input AND gate has inputs A, B, and the output of a NOT gate.
- A NOT gate with input B and output connected to the third input of the second 3-input AND gate.
- A 2-input AND gate with inputs A and C. Its output is connected to the first input of a 3-input OR gate.
- The output of the first 3-input AND gate is connected to the second input of the 3-input OR gate.
- The output of the 2-input AND gate (A and C) is also connected to the third input of the 3-input OR gate.
- The output of the 3-input OR gate is labeled Y.

(a) $Y = AC + \overline{B}C$
 (b) $Y = \overline{A}$
 (c) $Y = \overline{A} + \overline{B} \overline{C} + \overline{B} \overline{D} + BD$

(b)



(c)

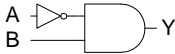


Exercise 2.17

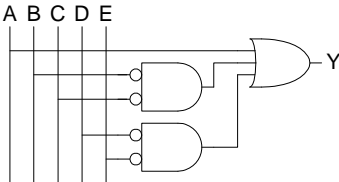
(a) $Y = B + \bar{A}\bar{C}$



(b) $Y = \bar{A}B$



(c) $Y = A + \bar{B}\bar{C} + \bar{D}\bar{E}$

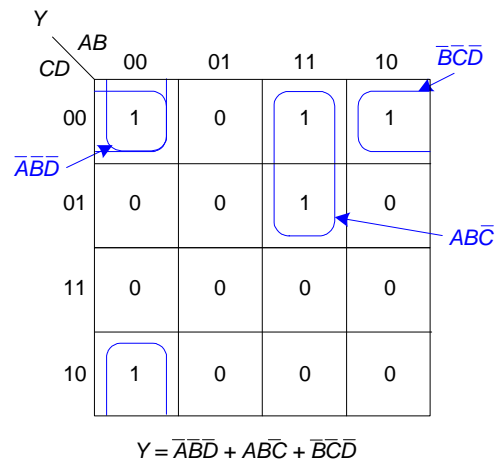
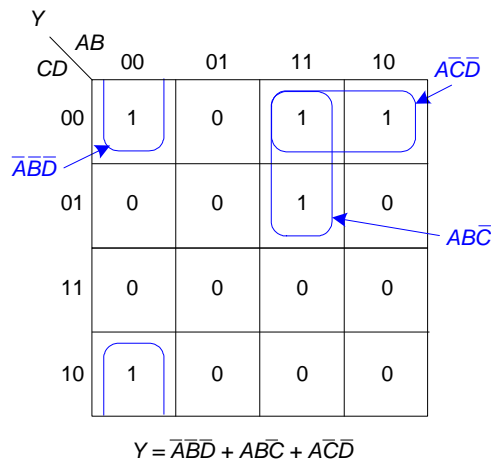


Exercise 2.19

4 gigarows = 4×2^{30} rows = 2^{32} rows, so the truth table has 32 inputs.

Exercise 2.21

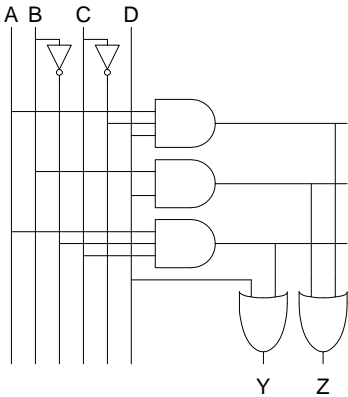
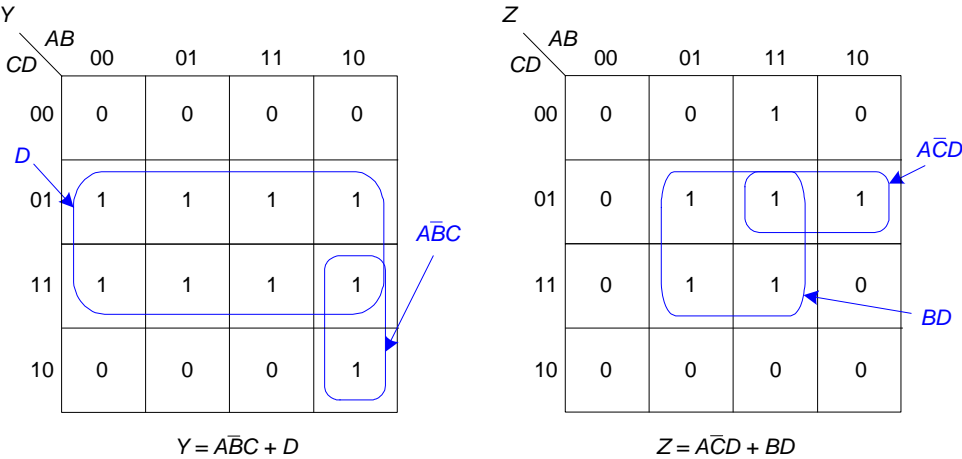
Ben is correct. For example, the following function, shown as a K-map, has two possible minimal sum-of-products expressions. Thus, although $A\bar{C}\bar{D}$ and $\bar{B}\bar{C}\bar{D}$ are both prime implicants, the minimal sum-of-products expression does not have both of them.



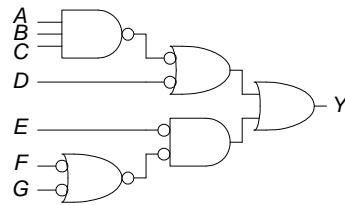
Exercise 2.23

B_2	B_1	B_0	$\overline{B_2 \bullet B_1 \bullet B_0}$	$\overline{B_2} + \overline{B_1} + \overline{B_0}$
0	0	0	1	1
0	0	1	1	1
0	1	0	1	1
0	1	1	1	1
1	0	0	1	1
1	0	1	1	1
1	1	0	1	1
1	1	1	0	0

Exercise 2.25



Exercise 2.27

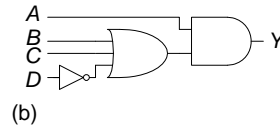
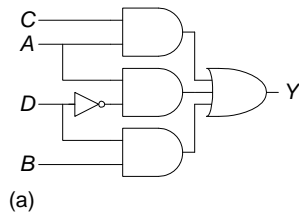


$$Y = ABC + \bar{D} + (\bar{F} + \bar{G})\bar{E}$$

$$= ABC + \bar{D} + \bar{E}\bar{F} + \bar{E}\bar{G}$$

Exercise 2.29

Two possible options are shown below:



Exercise 2.31

$$Y = \bar{A}D + A\bar{B}\bar{C}\bar{D} + BD + CD = A\bar{B}\bar{C}\bar{D} + D(\bar{A} + B + C)$$

Exercise 2.33

The equation can be written directly from the description:

$$E = \bar{S}A + AL + H$$

Exercise 2.35

Decimal Value	A_3	A_2	A_1	A_0	D	P
0	0	0	0	0	0	0
1	0	0	0	1	0	0
2	0	0	1	0	0	1
3	0	0	1	1	1	1
4	0	1	0	0	0	0
5	0	1	0	1	0	1
6	0	1	1	0	1	0
7	0	1	1	1	0	1
8	1	0	0	0	0	0
9	1	0	0	1	1	0
10	1	0	1	0	0	0
11	1	0	1	1	0	1
12	1	1	0	0	1	0
13	1	1	0	1	0	1
14	1	1	1	0	0	0
15	1	1	1	1	1	0

P has two possible minimal solutions:

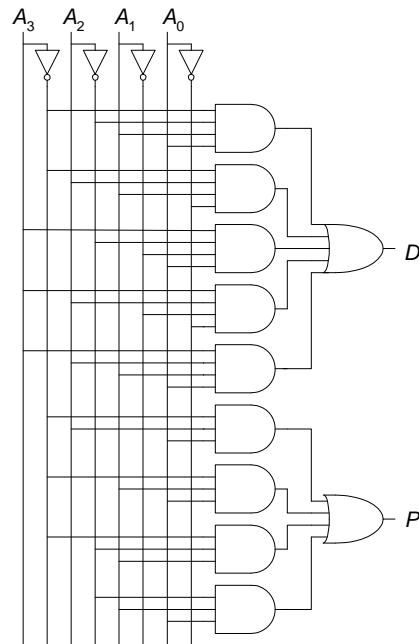
D $A_{1:0}$	$A_{3:2}$			
	00	01	11	10
00	0	0	1	0
01	0	0	0	1
11	1	0	1	0
10	0	1	0	0

$$D = \overline{A}_3 \overline{A}_2 A_1 A_0 + \overline{A}_3 A_2 A_1 \overline{A}_0 + A_3 \overline{A}_2 \overline{A}_1 A_0 + A_3 A_2 \overline{A}_1 \overline{A}_0 + A_3 A_2 A_1 A_0$$

P $A_{1:0}$	$A_{3:2}$			
	00	01	11	10
00	0	0	0	0
01	0	1	1	0
11	1	1	0	1
10	1	0	0	0

$$P = \overline{A}_3 \overline{A}_2 A_0 + \overline{A}_3 A_1 A_0 + \overline{A}_3 \overline{A}_2 A_1 + A_2 A_1 A_0$$
$$P = \overline{A}_3 A_1 A_0 + \overline{A}_3 \overline{A}_2 A_1 + \overline{A}_2 A_1 A_0 + A_2 \overline{A}_1 A_0$$

Hardware implementations are below (implementing the first minimal equation given for P).



Exercise 2.37

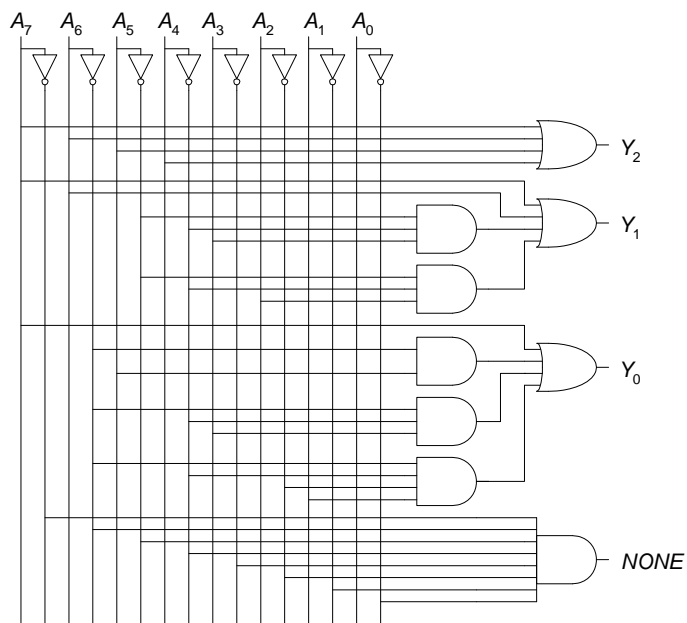
The equations and circuit for $Y_{2:0}$ is the same as in Exercise 2.25, repeated here for convenience.

A_7	A_6	A_5	A_4	A_3	A_2	A_1	A_0	Y_2	Y_1	Y_0
0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	X	0	0	1
0	0	0	0	0	1	X	X	0	1	0
0	0	0	0	1	X	X	X	0	1	1
0	0	0	1	X	X	X	X	1	0	0
0	0	1	X	X	X	X	X	1	0	1
0	1	X	X	X	X	X	X	1	1	0
1	X	X	X	X	X	X	X	1	1	1

$$Y_2 = A_7 + A_6 + A_5 + A_4$$

$$Y_1 = A_7 + A_6 + \overline{A_5}\overline{A_4}A_3 + \overline{A_5}\overline{A_4}A_2$$

$$Y_0 = A_7 + \overline{A_6}A_5 + \overline{A_6}\overline{A_4}A_3 + \overline{A_6}\overline{A_4}\overline{A_2}A_1$$



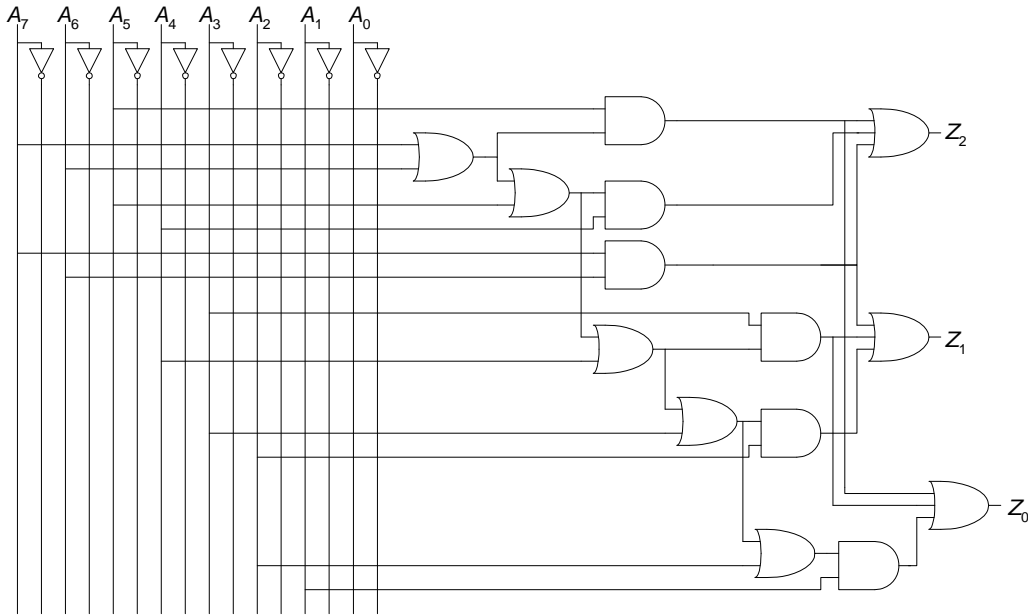
The truth table, equations, and circuit for $Z_{2:0}$ are as follows.

A_7	A_6	A_5	A_4	A_3	A_2	A_1	A_0	Z_2	Z_1	Z_0
0	0	0	0	0	0	1	1	0	0	0
0	0	0	0	0	1	0	1	0	0	0
0	0	0	0	1	0	0	1	0	0	0
0	0	0	1	0	0	0	1	0	0	0
0	0	1	0	0	0	0	1	0	0	0
0	1	0	0	0	0	0	1	0	0	0
1	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	1	1	X	0	0	1
0	0	0	0	1	0	1	X	0	0	1
0	0	0	1	0	0	1	X	0	0	1
0	0	1	0	0	0	1	X	0	0	1
0	1	0	0	0	0	1	X	0	0	1
1	0	0	0	0	0	1	X	0	0	1
0	0	0	0	1	1	X	X	0	1	0
0	0	0	1	0	1	X	X	0	1	0
0	0	1	0	0	1	X	X	0	1	0
0	1	0	0	0	1	X	X	0	1	0
1	0	0	0	0	1	X	X	0	1	0
0	0	0	1	1	X	X	X	0	1	1
0	0	1	0	1	X	X	X	0	1	1
0	1	0	0	1	X	X	X	0	1	1
1	0	0	0	1	X	X	X	0	1	1
0	0	1	1	X	X	X	X	1	0	0
0	1	0	1	X	X	X	X	1	0	0
1	0	0	1	X	X	X	X	1	0	0
0	1	1	X	X	X	X	X	1	0	1
1	0	1	X	X	X	X	X	1	0	1
1	1	X	X	X	X	X	X	1	1	0

$$Z_2 = A_4(A_5 + A_6 + A_7) + A_5(A_6 + A_7) + A_6A_7$$

$$Z_1 = A_2(A_3 + A_4 + A_5 + A_6 + A_7) + A_3(A_4 + A_5 + A_6 + A_7) + A_6A_7$$

$$Z_0 = A_1(A_2 + A_3 + A_4 + A_5 + A_6 + A_7) + A_3(A_4 + A_5 + A_6 + A_7) + A_5(A_6 + A_7)$$



Exercise 2.39

$$Y = A + \overline{C \oplus D} = A + CD + \overline{C}\overline{D}$$

Exercise 2.41

A	B	C	Y
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

(a)

(b)

A	Y
0	\overline{BC}
1	BC

(c)

Exercise 2.43

$$t_{pd} = 3t_{pd_NAND2} = \mathbf{60\ ps}$$

$$t_{cd} = t_{cd_NAND2} = \mathbf{15\ ps}$$

Exercise 2.45

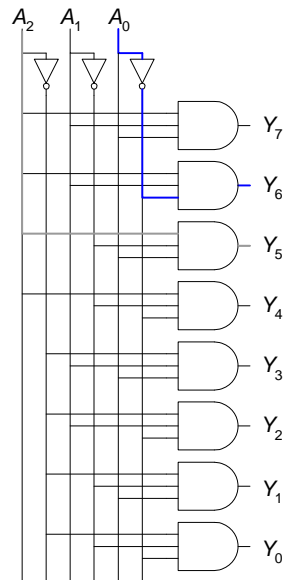
$$t_{pd} = t_{pd_NOT} + t_{pd_AND3}$$

$$= 15\ ps + 40\ ps$$

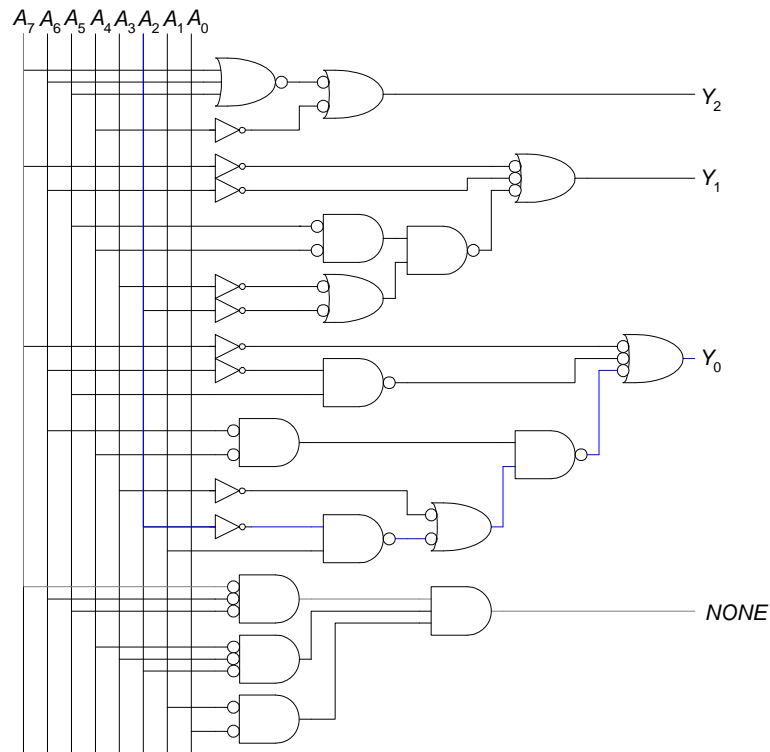
$$= \mathbf{55\ ps}$$

$$t_{cd} = t_{cd_AND3}$$

$$= \mathbf{30\ ps}$$



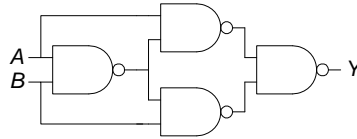
Exercise 2.47



$$\begin{aligned}
 t_{pd} &= t_{pd_INV} + 3t_{pd_NAND2} + t_{pd_NAND3} \\
 &= [15 + 3(20) + 30] \text{ ps} \\
 &= \mathbf{105 \text{ ps}}
 \end{aligned}$$

$$\begin{aligned}
 t_{cd} &= t_{cd_NOT} + t_{cd_NAND2} \\
 &= [10 + 15] \text{ ps} \\
 &= \mathbf{25 \text{ ps}}
 \end{aligned}$$

Question 2.1



Question 2.3

A tristate buffer has two inputs and three possible outputs: 0, 1, and Z. One of the inputs is the data input and the other input is a control input, often called the *enable* input. When the enable input is 1, the tristate buffer transfers the data input to the output; otherwise, the output is high impedance, Z. Tristate buffers are used when multiple sources drive a single output at different times. One and only one tristate buffer is enabled at any given time.

Question 2.5

A circuit's contamination delay might be less than its propagation delay because the circuit may operate over a range of temperatures and supply voltages, for example, 3-3.6 V for LVC MOS (low voltage CMOS) chips. As temperature increases and voltage decreases, circuit delay increases. Also, the circuit may have different paths (critical and short paths) from the input to the output. A gate itself may have varying delays between different inputs and the output, affecting the gate's critical and short paths. For example, for a two-input NAND gate, a HIGH to LOW transition requires two nMOS transistor delays, whereas a LOW to HIGH transition requires a single pMOS transistor delay.

