

# The Transfer Layer

UDP Service

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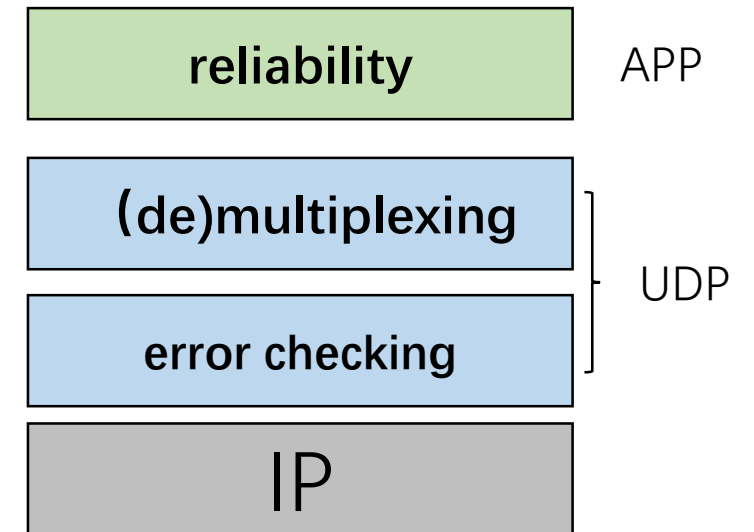
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# Connectionless Transport: UDP

a no-frills, bare-bones transport protocol, with just a small enhance of IP service

- Connectionless
  - no handshaking between sending and receiving transport-layer entities before sending a segment.
- No reliability?
  - Can be built into the application itself.



# Why UDP?

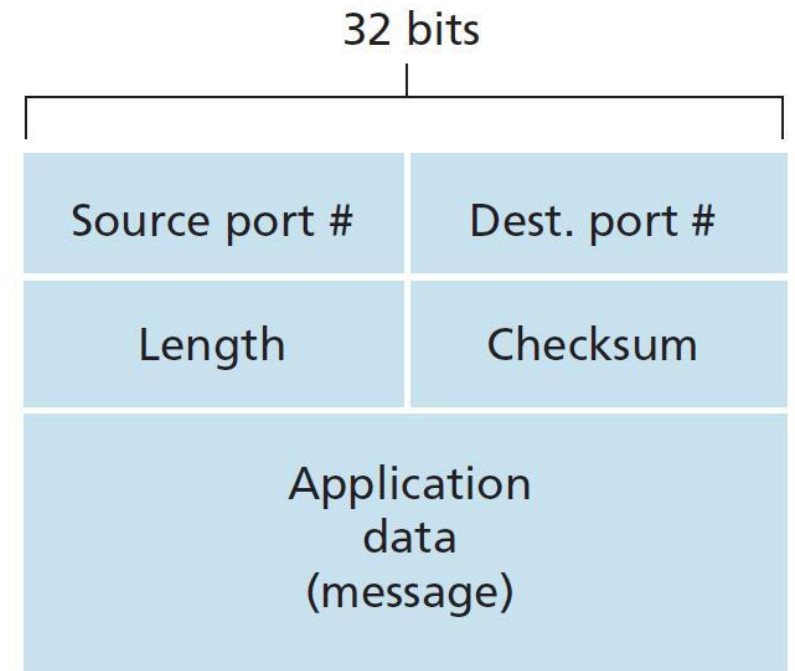
- Finer application-level control over what data is sent, and when.
  - Consider real-time applications requiring a minimum sending rate.
- No connection establishment
  - Comparing DNS and Web
- No connection state
  - support many more active clients (than TCP)
- Small packet header overhead
  - UDP: 8 bytes, TCP: 20 bytes

| Application            | Application-Layer Protocol | Underlying Transport Protocol |
|------------------------|----------------------------|-------------------------------|
| Electronic mail        | SMTP                       | TCP                           |
| Remote terminal access | Telnet                     | TCP                           |
| Web                    | HTTP                       | TCP                           |
| File transfer          | FTP                        | TCP                           |
| Remote file server     | NFS                        | Typically UDP                 |
| Streaming multimedia   | typically proprietary      | UDP or TCP                    |
| Internet telephony     | typically proprietary      | UDP or TCP                    |
| Network management     | SNMP                       | Typically UDP                 |
| Routing protocol       | RIP                        | Typically UDP                 |
| Name translation       | DNS                        | Typically UDP                 |

**Figure 3.6** ♦ Popular Internet applications and their underlying transport protocols

# UDP Segment Structure

- **Port** number fields for (de)multiplexing
- **Length** field specifies the number of bytes in the UDP segment (header plus data).
- **Checksum** is used by the receiving host to check whether errors have been introduced into the segment.
  - i.e. error detection.
- Only error detection, no error correction



# UDP Checksum

- UDP at the sender side performs the 1s complement of the sum of all the 16-bit words in the segment, with any overflow encountered during the sum being **wrapped around**.

- RFC 1071

Sender

0110011001100000  
0101010101010101  
1000111100001100

Receiver

0110011001100000  
0101010101010101  
1011101110110101  
1011101110110101  
1000111100001100  
0100101011000010

0100101011000010  
↓ 1s complement  
1011010100111101

0100101011000010  
1011010100111101  
1111111111111111

?

# Why error detection at UDP? The **end-end principle**

- Bit errors
  - During transmitting
  - During buffering.
- Many link layer protocols (including the popular Ethernet protocol) also provide error checking.
- But **No Guarantee** that all the links between source and destination provide error checking;
- Bit errors could be introduced when a segment is stored in a router's memory