Massachusetts Institute of Technology

Department of Electrical Engineering & Computer Science

6.041/6.431: Probabilistic Systems Analysis (Fall 2010)

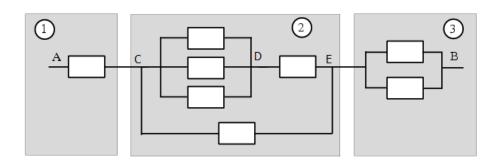
Tutorial 1 Solutions September 16/17, 2010

- 1. If $A \subset B$, then $\mathbf{P}(B \cap A) = \mathbf{P}(A)$ But we know that in order for A and B to be independent, $\mathbf{P}(B \cap A) = \mathbf{P}(A)\mathbf{P}(B)$. Therefore, A and B are independent if and only if $\mathbf{P}(B) = 1$ or $\mathbf{P}(A) = 0$. This could happen, for example, if B is the universe or if A is empty.
- 2. This problem is similar in nature to Example 1.24, page 40. In order to compute the success probability of individual sub-systems, we make use of the following two properties, derived in that example:
 - If a serial sub-system contains m components with success probabilities $p_1, p_2...p_m$, then the probability of success of the entire sub-system is given by

$$\mathbf{P}(\text{whole system succeeds}) = p_1 p_2 p_3 ... p_m$$

• If a parallel sub-system contains m components with success probabilities $p_1, p_2...p_m$, then the probability of success of the entire sub-system is given by

P(whole system succeeds) =
$$1 - (1 - p_1)(1 - p_2)(1 - p_3)...(1 - p_m)$$



Let $\mathbf{P}(X \to Y)$ denote the probability of a successful connection between node X and Y. Then,

$$\mathbf{P}(A \to B) = \mathbf{P}(A \to C)\mathbf{P}(C \to E)\mathbf{P}(E \to B)$$
 (since they are in series)
 $\mathbf{P}(A \to C) = p$
 $\mathbf{P}(C \to E) = 1 - (1 - p)(1 - \mathbf{P}(C \to D)\mathbf{P}(D \to E))$
 $\mathbf{P}(E \to B) = 1 - (1 - p)^2$

The probabilities $\mathbf{P}(C \to D)$, $\mathbf{P}(D \to E)$ can be similarly computed as

$$\mathbf{P}(C \to D) = 1 - (1 - p)^3$$

$$\mathbf{P}(D \to E) = p$$

The probability of success of the entire system can be obtained by substituting the subsystem success probabilities:

$$\mathbf{P}(A \to B) = p \left(1 - (1 - p)(1 - (1 - (1 - p)^3)p\right) \left(1 - (1 - p)^2\right).$$

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3. The Chess Problem.

= 0.7992

- (a) i. $P(2\text{nd Rnd Req}) = (0.6)^2 + (0.4)^2 = 0.52$
 - ii. **P**(Bo Wins 1st Rnd) = $(0.6)^2 = 0.36$
 - iii. $\mathbf{P}(\text{Al Champ}) = 1 \mathbf{P}(\text{Bo Champ}) \mathbf{P}(\text{Ci Champ})$ = $1 - (0.6)^2 * (0.5)^2 - (0.4)^2 * (0.3)^2 = 0.8956$
- (b) i. $\mathbf{P}(\text{Bo Challenger}|2\text{nd Rnd Req}) = \frac{(0.6)^2}{0.52} = \frac{0.36}{0.52} = 0.6923$ ii. $\mathbf{P}(\text{Al Champ}|2\text{nd Rnd Req}) = \int (\text{Al } \cap \text{Bo}|2\text{nd}) + \int (\text{Al } \cap \text{Ci}|2\text{nd}) + \int (\text{Al } \cap |\text{Ind}|2\text{nd})$ $= \mathbf{P}(\text{Al Champ}|\text{Bo Challenger}, 2\text{nd Rnd Req}) \times \mathbf{P}(\text{Bo Challenger}|2\text{nd Rnd Req})$ $+ \mathbf{P}(\text{Al Champ}|\text{Ci Challenger}, 2\text{nd Rnd Req}) \times \mathbf{P}(\text{Ci Challenger}|2\text{nd Rnd Req})$ $= (1 (0.5)^2) \times 0.6923 + (1 (0.3)^2) \times 0.3077$ in sub-samplestane

(c)
$$\mathbf{P}((\text{Bo Challenger})|\{(2\text{nd Rnd Req}) \cap (\text{One Game})\}) = \frac{(0.6)^2*(0.5)}{(0.6)^2*(0.5)+(0.4)^2*(0.7)}$$

$$= \frac{(0.6)^2(0.5)}{0.2920} = 0.6164$$

$$= \frac{p(\beta_0 \mid 2\text{nd} \mid 0\text{ne})}{p(2\text{nd} \mid 0\text{ne})} = \frac{p(2\text{nd} \mid 0\text{ne} \mid \beta_0) \cdot p(\beta_0)}{p(2\text{nd} \mid 0\text{ne} \mid \beta_0) \cdot p(\beta_0) + p(2\text{nd} \mid 0\text{ne} \mid \beta_0) \cdot p(\beta_0)}$$

$$= \frac{0.5 \cdot (0.6)^2}{0.5 \cdot (0.6)^2} + 0.7 \cdot (0.4)^2$$

Bo
$$\subseteq$$
 2nol, $C_i \subseteq$ 2nol
Bo \cup $C_i = 2nol$
One \subseteq 2nol
Inol = 2nol²
Inol + 2nol = |nol + Bo + C_i = Ω
P(Bo | 2nol \cap One) = P(Bo | One)

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