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Blind Remote Exploitation

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Abstract



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whoami

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Objective

- How to write a generic exploit
- How to build a ROP chain in time of exploitation
- How to hack proprietary closed-binary services
- How to bypass anti-exploitation techniques



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Disclaimer

- Security really exists?
- How do you expect to get security without understanding the hacker's mind?
- For whom is this talk?
- Penetration testers nowadays



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Requirements

- What must we know?
 - Linux syscalls
 - x86 architecture
 - C programming
 - Sockets and network programming
 - Basics of software exploitation
 - Return Oriented Programming
- The same idea can be ported to another Operational Systems!



Introduction



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Scenarios

1. Hacking an open-source server for which the binary is unknown
2. Hacking a known vulnerability in a library thought to be used in a proprietary closed-binary software
3. Hacking a proprietary closed-binary services



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Scenario example

1. **Hacking an open-source server for which the binary is unknown**
 - Proftpd Telnet IAC buffer overflow (CVE-2010-4221) vulnerability as example
2. Hacking a known vulnerability in a library thought to be used in a proprietary closed-binary software
3. Hacking a proprietary closed-binary services

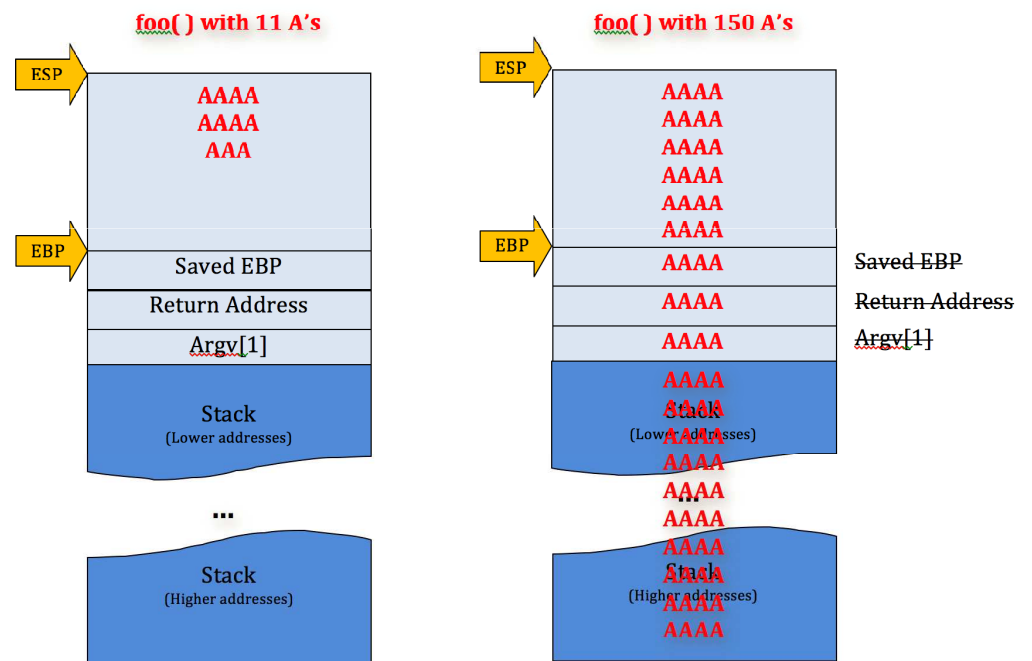


Stack Overflow overview

- Linux uses `cdecl` as function call convention

```
void foo(int argc, char *argv[]) {
    int variable;
    char buffer[24];

    ...
    strcpy(buffer, argv[1]);
    ...
}
```



Anti-exploitation techniques

- NX – non-executable stack
- AAAS – ASCII Armored Address Space
- ASLR – Address Space Layout Randomization
- Stack Canary (or stack cookie)



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Bypass Anti-Exploitation



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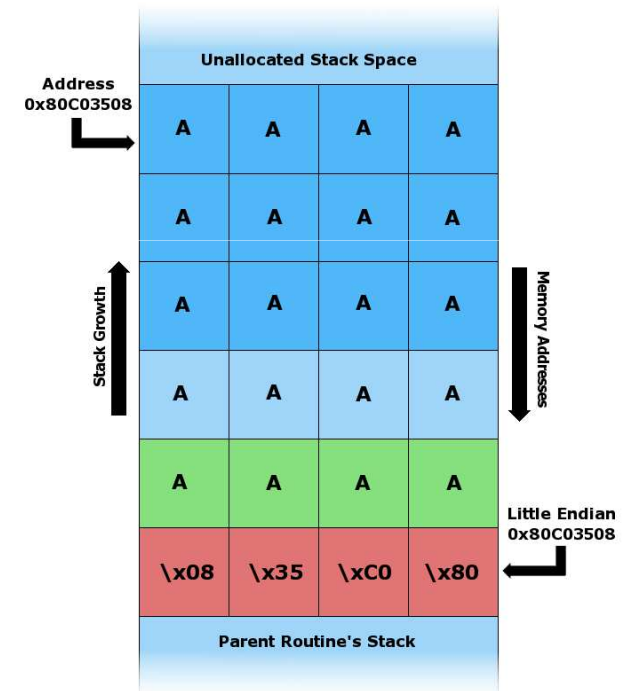
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Bypassing stack canary

- Process that `fork()` each connection
- Brute force byte-by-byte
- 1020 possible combinations to discover the canary on 32-bits system
- `fork()` followed by `execve()` will not work!



Return Oriented Programming

- Chain gadgets to execute malicious code
- A gadget is a suite of instructions which end by the branch instruction **ret** (Intel) or the equivalent on ARM.

- Intel examples:

- **pop** **eax** ; **ret**
- **xor** **ebx, ebx** ; **ret**

- ARM examples:

- **pop** {**r4, pc**}
- **str** **r1, [r0]** ; **bx lr**

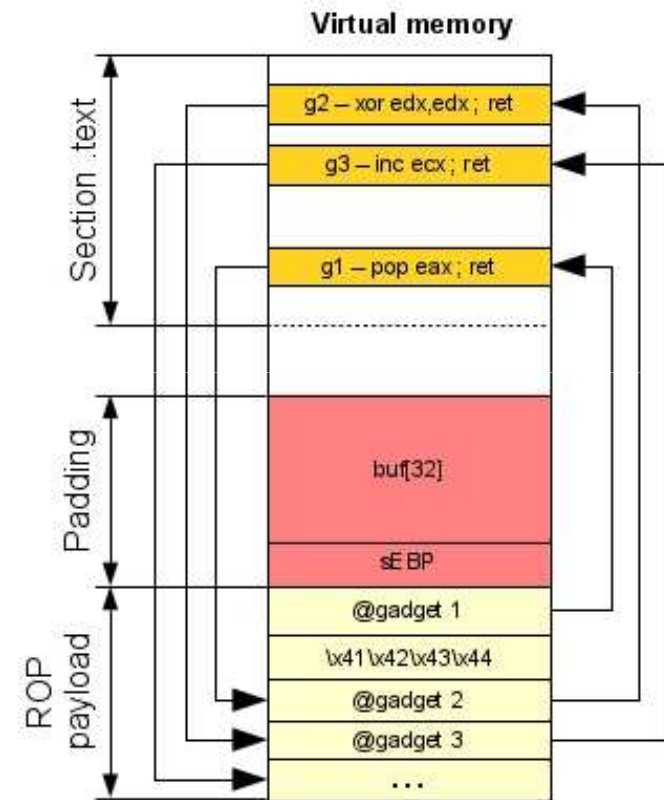
- Objective: Use gadgets instead of classical shellcode



Attack using ROP

- **Gadget1** is executed and returns
- **Gadget2** is executed and returns
- **Gadget3** is executed and returns
- And so on until all instructions that you want are executed
- So, the real execution is:

```
pop eax
xor edx, edx
inc ecx
```



Blind Return Oriented Programming

- Perform ROP in time of exploitation
- The target server must restart after crash
- Bruteforce byte-by-byte to find **write()** PLT address
- Perform ret2plt with **write()** function to leak the server's binary
- Find more gadgets to build our ROP chain to exploit

Blind Remote Exploitation



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CVE-2010-4221

- Stack overflow on function `pr_netio_telnet_gets()` from the file "`src/netio.c`"
- The server `fork()` each connection
- The server restore after crash



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BROP on CVE-2010-4221

- Bruteforce byte-by-byte `write()` PLT address
- `ret2plt` with `write(1, 0x080532d8, 0xffffffff);`
- Leak server's binary
- Search `mmap64()` and `memcpy()` PLT entry
- Search gadgets:
 - `pop; pop; pop; ret;`
 - `add esp, 20h; pop; pop; ret;`
- Search byte offsets from our shellcode copy routine



Exploiting CVE-2010-4221

- Execute `mmap64()` to map a read, write and executable memory
- Use `memcpy()` to copy our shellcode copy routine to the memory mapped
- The shellcode copy routine will copy our shellcode from stack to a executable memory and will pass the flow to the shellcode.



Exploiting CVE-2010-4221

```
void *mmap(void *addr, size_t length, int prot, int flags, int fd, off_t offset)
```

ROP chain:

```
mmap() plt address
add esp 20h; pop; pop; ret;
\x00\x00\x00\x10 ----> start address
\x00\x10\x00\x00 ----> length 1000 bytes
\x07\x00\x00\x00 ----> 07 = set read, write and executable permission
\x32\x00\x00\x00 ----> 32 bit flag
\xff\xff\xff\xff ----> file descriptor (0xffffffff probably will not be in use)
\x00\x00\x00\x00 --|
\x00\x00\x00\x00   |
\x00\x00\x00\x00   | -> offset
\x00\x00\x00\x00   |
\x00\x00\x00\x00   |
\x00\x00\x00\x00 --|
```



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Exploiting CVE-2010-4221

```
void *memcpy(void *dest, const void *src, size_t n)
```

First ROP chain:

```
memcpy() plt address
pop; pop; pop; ret;
\x00\x00\x00\x10 ----> destiny
      xxxx          ----> found byte addr (from shellcode copy routine)
\x02\x00\x00\x00 ----> number of bytes to copy
```

Second ROP chain:

```
memcpy() plt address
pop; pop; pop; ret;
\x02\x00\x00\x10 ----> destiny
      yyyy          ----> found byte addr (from shellcode copy routine)
\x01\x00\x00\x00 ----> number of bytes to copy
```



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DEMO

Talk is cheap, show me the code!!!



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Pros

- We will be able to inject any shellcode that we want with the same technique
- We will be able to exploit any server that runs the vulnerable application, even if it runs on a different distro with different compiling parameters
- We will be able to bypass all principals anti-exploitation methods
- The same technique will works on our another scenarios



Cons

- The server can restart after crash
- In other words: the server must be `fork()` for each connection
- The server cannot use `fork()` followed by `execve()` on creating a child connection
- This type of exploitation may make some noise!
- This attack may take a while to get success



Conclusion



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Hacker's vision

- We can build generic ways to exploit applications
- We can attack even with old vulnerabilities
- We can attack a library vulnerability, even if we can't replicate the server's scenario
- We can defeat anti-exploitation protections
- We can attack a application that we don't have the binary



Security really exists?

- Computer vulnerabilities are conceptual
- Always will exist vulnerabilities
- Always will have a way to exploit them

Security is an illusion!!!



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What we can do now?

- Security is a **continuous** process
- Security is: **hardware** + **software** + **people**
- The attacker's vision shows our weakness

We should always be steps ahead of the attackers!!!



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