

RISC-V External Debug Support

Version 0.13

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# Preface

**Warning! This draft specification will change before being accepted as standard, so implementations made to this draft specification will likely not conform to the future standard.**

## Acknowledgments

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# Chapter 1

## Introduction

When a design progresses from simulation to hardware implementation, a user’s control and understanding of the system’s current state drops dramatically. To help bring up and debug low level software and hardware, it is critical to have good debugging support built into the hardware. When a robust OS is running on a core, software can handle many debugging tasks. However, in many scenarios, hardware support is essential.

This document outlines a standard architecture for external debug support on RISC-V platforms. This architecture allows a variety of implementations and tradeoffs, which is complementary to the wide range of RISC-V implementations. At the same time, this specification defines common interfaces to allow debugging tools and components to target a variety of platforms based on the RISC-V ISA.

System designers may choose to add additional hardware debug support, but this specification defines a standard interface for common functionality.

### 1.1 Terminology

A *platform* is a single integrated circuit consisting of one or more *components*. Some components may be RISC-V cores, while others may have a different function. Typically they will all be connected to a single system bus. A single RISC-V core contains one or more hardware threads, called *harts*.

### 1.2 About This Document

#### 1.2.1 Structure

This document contains two parts. The main part of the document is the specification, which is given in the numbered sections. The second part of the document is a set of appendices. The information in the appendix is intended to clarify and provide examples, but is not part of the

actual specification.

### 1.2.2 Register Definition Format

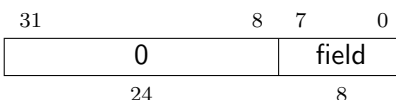
All register definitions in this document follow the format shown below. A simple graphic shows which fields are in the register. The upper and lower bit indices are shown to the top left and top right of each field. The total number of bits in the field are shown below it.

After the graphic follows a table which for each field lists its name, description, allowed accesses, and reset value. The allowed accesses are listed in Table 1.1.

Table 1.1: Register Access Abbreviations

R	Read-only.
R/W	Read/Write.
R/W0	Read/Write. Only writing 0 has an effect.
R/W1	Read/Write. Only writing 1 has an effect.
R/W1C	Read/Write. For each bit in the field, writing 1 clears that bit. Writing 0 has no effect.
W	Write-only. When read this field returns 0.
W1	Write-only. Only writing 1 has an effect.

#### 1.2.2.1 Long Name (shortname, at 0x123)



Field	Description	Access	Reset
field	Description of what this field is used for.	R/W	15

## 1.3 Background

There are several use cases for dedicated debugging hardware, both internal to a CPU core and with an external connection. This specification addresses the use cases listed below. Implementations can choose not to implement every feature, which means some use cases might not be supported.

- Debugging low-level software in the absence of an OS or other software.
- Debugging issues in the OS itself.
- Bootstrapping a system to test, configure, and program components before there is any executable code path in the system.



- Accessing hardware on a system without a working CPU.

In addition, even without a hardware debugging interface, architectural support in a RISC-V CPU can aid software debugging and performance analysis by allowing hardware triggers and breakpoints. This specification aims to define common resources which can be used for different cases.

When debugging software, this specification distinguishes between two forms of external debugging. The first is *halt mode* debugging, where an external debugger halts some or all components of a platform and inspects their state while they are in stasis. The debugger can read and/or modify state, then direct the hardware to execute a single instruction, or continue to run freely.

The second is *run mode* debugging. In this mode a software debug agent runs on a component (eg. triggered by a timer interrupt or breakpoint on a RISC-V core) which transfers data to or from the debugger without halting the component, only briefly interrupting its program flow. This functionality is essential if the component is controlling some real-time system (like a hard drive) where long timing delays could lead to physical damage. This requires additional software support (both on the system as well as on the debugger), and efficient communication channels between the component and the debugger.

## 1.4 Supported Features

The debug interface described in this specification supports the following features:

1. RV32, RV64, and future RV128 are all supported.
2. Any hart in the platform can be independently debugged.
3. A debugger can discover almost<sup>1</sup> everything it needs to know itself, without user configuration.
4. Each hart can be debugged from the very first instruction executed.
5. A RISC-V hart can be halted when a software breakpoint instruction is executed.
6. Hardware single-step can execute one instruction at a time.
7. Debug functionality is independent of the debug transport used.
8. The debugger does not need to know anything about the microarchitecture of the cores it is debugging.
9. Arbitrary subsets of harts can be halted and resumed simultaneously. (Optional)
10. Arbitrary instructions can be executed on a halted hart. That means no new debug functionality is needed when a core has additional or custom instructions or state, as long as there exist programs that can move that state into GPRs. (Optional)
11. Registers can be accessed without halting. (Optional)

---

<sup>1</sup>Notable exceptions include information about the memory map and peripherals.

12. A running hart can be directed to execute a short sequence of instructions, with little overhead. (Optional)
13. A system bus master allows memory access without involving any hart. (Optional)
14. A RISC-V hart can be halted when a trigger matches the PC, read/write address/data, or an instruction opcode. (Optional)

While both the mechanism to execute arbitrary instructions and the system bus master are optional, at least one of them must be implemented. Otherwise there is no mechanism to access memory.

This document does not suggest a strategy or implementation for hardware test, debugging or error detection techniques. Scan, BIST, etc. are out of scope of this specification, but this specification does not intend to limit their use in RISC-V systems.

## Chapter 2

# System Overview

Figure 2.1 shows the main components of External Debug Support. Blocks shown in dotted lines are optional.

The user interacts with the Debug Host (eg. laptop), which is running a debugger (eg. gdb). The debugger communicates with a Debug Translator (eg. OpenOCD, which may include a hardware driver) to communicate with Debug Transport Hardware (eg. Olimex USB-JTAG adapter). The Debug Transport Hardware connects the Debug Host to the Platform’s Debug Transport Module (DTM). The DTM provides access to the Debug Module (DM) using the Debug Module Interface (DMI).

The DM allows the debugger to halt any hart in the platform. Abstract commands provide access to GPRs. Additional registers are accessible through abstract commands or by writing programs to the optional Program Buffer.

The Program Buffer allows the debugger to execute arbitrary instructions on a hart. This mechanism can be used to access memory. An optional system bus access block allows memory accesses without using a RISC-V hart to perform the access.

Each RISC-V hart may implement a Trigger Module. When trigger conditions are met, harts will halt and inform the debug module that they have halted.



Figure 2.1: RISC-V Debug System Overview

## Chapter 3

# Debug Module (DM)

The Debug Module implements a translation interface between abstract debug operations and their specific implementation. It might support the following operations:

1. Give the debugger necessary information about the implementation. (Required)
2. Allow any individual hart to be halted and resumed. (Required)
3. Provide status on which harts are halted. (Required)
4. Provide read and write access to a halted hart's GPRs. (Required)
5. Provide access to a reset signal that allows debugging from the very first instruction after reset. (Required)
6. Provide access to other hart registers. (Optional)
7. Provide a Program Buffer to force the hart to execute arbitrary instructions. (Optional)
8. Allow multiple harts to be halted, resumed, and/or reset at the same time. (Optional)
9. Allow direct System Bus Access. (Optional)

In order to implement memory access, a target must implement either the Program Buffer or System Bus Access.

A single DM can debug up to 1024 harts.

### 3.1 Debug Module Interface (DMI)

The Debug Module is a slave to a bus called the Debug Module Interface (DMI). The master of the bus is the Debug Transport Module(s). The Debug Module Interface can be a trivial bus with one master and one slave, or use a more full-featured bus like TileLink or the AMBA Advanced Peripheral Bus. The details are left to the system designer.

The DMI uses between 7 and 32 address bits. It supports read and write operations. The bottom of the address space is used for the DM. Extra space can be used for custom debug devices, other cores, additional DMs, etc.

The Debug Module is controlled via register accesses to its DMI address space.

Table 3.1: Debug Module Interface Address Space

0x00 – 0x3f	Registers described in Section 3.11.
0x40 – 0x5f	This is called the ‘halt region’. These 32 addresses for 32-bit words provide access to the halt bit for up to 1024 harts. If the hart is halted, the bit is 1. Otherwise the bit is 0. The bit for hart 0 is the LSB in the 32-bit word at 0x40. The bit for hart 1023 is the MSB in the 32-bit word at 0x5f.

## 3.2 Reset Control

The Debug Module controls a global reset signal, `ndmreset` (non-debug module reset), which can reset, or hold in reset, every component in the platform, except for the Debug Module and Debug Transport Modules. Exactly what is affected by this reset is implementation dependent, as long as it is possible to debug programs from the first instruction executed. The Debug Module’s own state and registers should only be reset at power-up and while `dmactive` in `dmcontrol` is 0. The halt state of harts should be maintained across system reset provided that `dmactive` is 1, although trigger CSRs may be cleared.

Due to clock and power domain crossing issues, it may not be possible to perform arbitrary DMI accesses across system reset. While `ndmreset` or any external reset is asserted, the only supported DM operation is accessing `dmcontrol`. The behavior of other accesses is undefined.

## 3.3 Selecting Harts

Up to 1024 harts can be connected to a single DM. The debugger selects a hart, and then subsequent halt, resume, reset, and debugging commands are specific to that hart.

A debugger can enumerate all the harts attached to the DM by selecting each hart starting from 0 until nonexistent in `dmstatus` is 1.

The debugger can discover the mapping between hart indices and `mhartid` by using the interface to read `mhartid`, or by reading the system’s Device Tree.

### 3.3.1 Selecting a Single Hart

All debug modules must support selecting a single hart. The debugger can select a hart by writing its index to `hartsel`. Hart indexes start at 0 and are contiguous until the final index.

### 3.3.2 Selecting Multiple Harts

Debug Modules may optionally implement a Hart Array Mask register to allow selecting multiple harts at once. The debugger can set bits in the hart array mask register using `hawindowssel` and `hawindow`, then apply actions to all selected harts by setting `hasel`. If this feature is supported, multiple harts can be halted, resumed, and reset simultaneously.

Only the actions initiated by `dmcontrol` can apply to multiple harts at once, Abstract Commands apply only to the hart selected by `hartsel`.

## 3.4 Run Control

The Debug Module can halt selected harts and allow them to run again using the `dmcontrol` register. After selecting harts, it sets `haltreq`, then waits for `allhalted` to indicate the harts are halted before clearing `haltreq` to 0. Setting `haltreq` has no effect on harts which are already halted.

To resume, the debugger sets `resumereq`. This action sets each selected hart's `resumeack` bit to 0. Once a hart has resumed, it sets its `resumeack` bit to 1. Thus, the debugger should wait for `allresumeack` to indicate all harts have resumed before clearing `resumereq` to 0. If a debugger reads `anyresumeack` and `anyhalted` in the same operation, then some selected harts must have resumed and then halted again.

---

*When waiting for a hart to resume, a debugger should examine `allresumeack`, not `allrunning` or `allhalted`, because the hart may immediately halt again due to trigger or step conditions.*

Setting `resumereq` on hart which is running clears `resumeack` but may cause a hart which halts in the future to immediately resume. Debuggers should not depend on this behavior and should not set `resumereq` for running harts.

To halt or resume multiple harts at the same time, the debugger first sets the hart's bits in the hart array mask register, then follows the same procedure but with `hasel` set to 1. Depending on the desired operation, the debugger might consider the `any*` versions of the status instead of `all*`.

When halt or resume is requested, a hart must respond in less than one second, unless it is unavailable. (How this is implemented is not further specified. A few clock cycles will be a more typical latency).

## 3.5 Abstract Commands

The DM supports a set of abstract commands, most of which are optional. Depending on the implementation, the debugger may be able to perform some abstract commands even when the selected hart is not halted. Debuggers can only determine which abstract commands are supported by a given hart in a given state by attempting them and then looking at `cmderr` in `abstractcs` to see if they were successful.

Debuggers execute abstract commands by writing them to `command`. Debuggers can determine whether an abstract command is complete by reading `busy` in `abstractcs`. If the command takes arguments, the debugger must write them to the `data` registers before writing to `command`. If a command returns results, the Debug Module must ensure they are placed in the `data` registers before `busy` is cleared. Which `data` registers are used for the arguments is described in Table 3.2. In all cases the least-significant word is placed in the lowest-numbered `data` register.

Table 3.2: Use of Data Registers

XLEN	arg0/return value	arg1	arg2
32	<code>data0</code>	<code>data1</code>	<code>data2</code>
64	<code>data0</code> , <code>data1</code>	<code>data2</code> , <code>data3</code>	<code>data4</code> , <code>data5</code>
128	<code>data0</code> – <code>data3</code>	<code>data4</code> – <code>data7</code>	<code>data8</code> – <code>data11</code>

### 3.5.1 Abstract Command Listing

This section describes each of the different abstract commands and how their fields should be interpreted when they are written to `command`.

Each abstract command is a 32-bit value. The top 8 bits contain `cmdtype` which determines the kind of command. Table 3.3 lists all commands.

Table 3.3: Meaning of `cmdtype`

<code>cmdtype</code>	Command	Page
0	Access Register Command	<a href="#">10</a>
1	Quick Access	<a href="#">11</a>

#### 3.5.1.1 Access Register

This command gives the debugger access to CPU registers and program buffer. It performs the following sequence of operations:

1. Copy data from the register specified by `regno` into the `arg0` region of `data`, if `write` is clear and `transfer` is set.
2. Copy data from the `arg0` region of `data` into the register specified by `regno`, if `write` is set and `transfer` is set.
3. Execute the Program Buffer, if `postexec` is set.

If any of these operations fail, `cmderr` is set and none of the remaining steps are executed. An implementation may detect an upcoming failure early, and fail the overall command before it reaches the step that would cause failure.

Debug Modules must implement this command and must support read and write access to all GPRs when the selected hart is halted. Debug Modules may optionally support accessing other registers,



or accessing registers when the hart is running. If this command is supported for a register while the hart is running, it must also be supported for a register while the hart is halted. Each individual register (aside from GPRs) may be supported differently across read, write, and halt status.

---

*The encoding of size was chosen to match sbaccess in sbcs.*

31	24	23	22	20	19	18	17	16	15	0
cmdtype	0	size	0	postexec	transfer	write	regno			
8	1	3	1	1	1	1	16			

Field	Description
cmdtype	This is 0 to indicate Access Register Command.
size	2: Access the lowest 32 bits of the register. 3: Access the lowest 64 bits of the register. 4: Access the lowest 128 bits of the register. If size specifies a size larger than the register's actual size, then the access must fail. If a register is accessible, then reads of size less than or equal to the register's actual size must be supported.
postexec	When 1, execute the program in the Program Buffer exactly once after performing the transfer, if any.
transfer	0: Don't do the operation specified by write. 1: Do the operation specified by write. This bit can be used to just execute the Program Buffer without having to worry about placing valid values into size or regno.
write	When transfer is set: 0: Copy data from the specified register into arg0 portion of data. 1: Copy data from arg0 portion of data into the specified register.
regno	Number of the register to access, as described in Table 3.4. dpc may be used as an alias for PC if this command is supported on a non-halted hart.

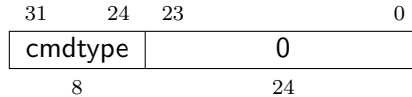
### 3.5.1.2 Quick Access

Perform the following sequence of operations:

1. If the hart is halted, the command sets cmderr to halt/resume and does not continue.
2. Halt the hart. If the hart halts for some other reason (e.g. breakpoint), the command sets cmderr to halt/resume and does not continue.
3. Execute the Program Buffer. If an exception occurs, cmderr is set to exception and the program buffer execution ends, but the quick access command continues.

4. Resume the hart.

Implementing this command is optional.



Field	Description
cmdtype	This is 1 to indicate Quick Access command.

Table 3.4: Abstract Register Numbers

0x0000 – 0x0fff	CSRs. The “PC” can be accessed here through <a href="#">dpc</a> .
0x1000 – 0x101f	GPRs
0x1020 – 0x103f	Floating point registers
0xc000 – 0xffff	Reserved for non-standard extensions and internal use.

## 3.6 Program Buffer

To support executing arbitrary instructions on a halted hart, a Debug Module can include a Program Buffer that a debugger can write small programs to. Systems that support all necessary functionality using abstract commands only may choose to omit the Program Buffer.

A debugger can write a small program to the Program Buffer, and then execute it exactly once with the Access Register Abstract Command, setting the `postexec` bit in [command](#). The debugger can write whatever program it likes (including jumps out of the Program Buffer), but the program must end with `ebreak` or `ebreak.c`. To save hardware, an implementation may support an implied `ebreak` that is executed when a hart runs off the end of the Program Buffer. This is indicated in `impebreak`. With this feature, a Program Buffer of just 2 32-bit words can offer efficient debugging.

If `progbufsize` is 1, the Program Buffer may only hold a single instruction, and `impebreak` must be 1. This instruction can be a 32-bit instruction, or a compressed instruction in the lower 16 bits accompanied by a compressed `nop` in the upper 16 bits.

If the debugger executes a program that does not terminate with an `ebreak` instruction, the hart will remain in Debug Mode until it is reset.

While these programs are executed, the hart does not leave Debug Mode (see Section 4.1). If an exception is encountered during execution of the Program Buffer, no more instructions are executed, the hart remains in Debug Mode, and `cmderr` is set to 3 (**exception error**). If the debugger executes a program that doesn’t terminate, then it loses control of the hart.

Executing the Program Buffer may clobber [dpc](#). If that is the case, it must be possible to read/write [dpc](#) using an abstract command with `postexec` not set. The debugger must attempt to save [dpc](#) between halting and executing a Program Buffer, and then restore [dpc](#) before leaving Debug Mode.

---

*Allowing Program Buffer execution to clobber `dpc` allows for direct implementations that don't have a separate PC register, and do need to use the PC when executing the Program Buffer.*

The Program Buffer may be implemented as RAM which is accessible to the hart as RAM memory. A debugger can determine if this is the case by executing small programs that attempt to write and read back relative to `pc` while executing from the Program Buffer. If so, the debugger has more flexibility in what it can do with the program buffer.

## 3.7 Overview of States

Figure 3.1 shows a conceptual view of the states passed through by a hart during run/halt debugging as influenced by the different fields of `dmcontrol`, `abstractcs`, `abstractauto`, and `command`.

## 3.8 System Bus Access

When a Program Buffer is present, a debugger can access the system bus by having a RISC-V hart perform the accesses it requires. A Debug Module may also include a System Bus Access block to provide memory access without involving a hart, regardless of whether Program Buffer is implemented. The System Bus Access block uses physical addresses.

---

*Implementing a System Bus Access block has several benefits even when a Debug Module also implements a Program Buffer. First, it is possible to access memory in a running system with minimal impact. Second, it may improve performance when accessing memory. Third, it may provide access to devices that a hart does not have access to.*

## 3.9 Quick Access

Depending on the task it is performing, some harts can only be halted very briefly. There are several mechanisms that allow accessing resources in such a running system with a minimal impact on the running hart.

First, an implementation may allow some abstract commands to execute without halting the hart.

Second, the Quick Access abstract command can be used to halt a hart, quickly execute the contents of the Program Buffer, and let the hart run again. Combined with instructions that allow Program Buffer code to access the `data` registers, as described in 3.11.3, this can be used to quickly perform a memory or register access. For some systems this will be too intrusive, but many systems that can't be halted can bear an occasional hiccup of a hundred or less cycles.

Third, if the System Bus Access block is implemented, it can be used while a hart is running to access system memory.



Figure 3.1: Run/Halt Debug State Machine. As only a small amount of state is visible to the debugger, the states and transitions are conceptual.

### 3.10 Security

To protect intellectual property it may be desirable to lock access to the Debug Module. To allow access during a manufacturing process and not afterwards, a reasonable solution could be to add a fuse bit to the Debug Module that can be used to be permanently disable it. Since this is technology specific, it is not further addressed in this spec.

Another option is to allow the DM to be unlocked only by users who have an access key. A few bits in [dmstatus](#) and [authdata](#) can support an arbitrarily complex authentication mechanism. When authenticated is clear, the DM must not interact with the rest of the platform in any way.

### 3.11 Debug Module DMI Registers

When read, unimplemented Debug Module DMI Registers return 0. Writing them has no effect.

Table 3.5: Debug Module Debug Bus Registers

Address	Name	Page
0x04	Abstract Data 0	<a href="#">24</a>
0x0f	Abstract Data 11	
0x10	Debug Module Control	<a href="#">17</a>
0x11	Debug Module Status	<a href="#">16</a>
0x12	Hart Info	<a href="#">19</a>
0x13	Halt Summary	<a href="#">20</a>
0x14	Hart Array Window Select	<a href="#">21</a>
0x15	Hart Array Window	<a href="#">21</a>
0x16	Abstract Control and Status	<a href="#">21</a>
0x17	Abstract Command	<a href="#">23</a>
0x18	Abstract Command Autoexec	<a href="#">23</a>
0x19	Device Tree Addr 0	<a href="#">24</a>
0x1a	Device Tree Addr 1	
0x1b	Device Tree Addr 2	
0x1c	Device Tree Addr 3	
0x20	Program Buffer 0	<a href="#">24</a>
0x2f	Program Buffer 15	
0x30	Authentication Data	<a href="#">25</a>
0x38	System Bus Access Control and Status	<a href="#">25</a>
0x39	System Bus Address 31:0	<a href="#">26</a>
0x3a	System Bus Address 63:32	<a href="#">27</a>
0x3b	System Bus Address 95:64	<a href="#">27</a>
0x3c	System Bus Data 31:0	<a href="#">27</a>
0x3d	System Bus Data 63:32	<a href="#">28</a>
0x3e	System Bus Data 95:64	<a href="#">28</a>
0x3f	System Bus Data 127:96	<a href="#">29</a>

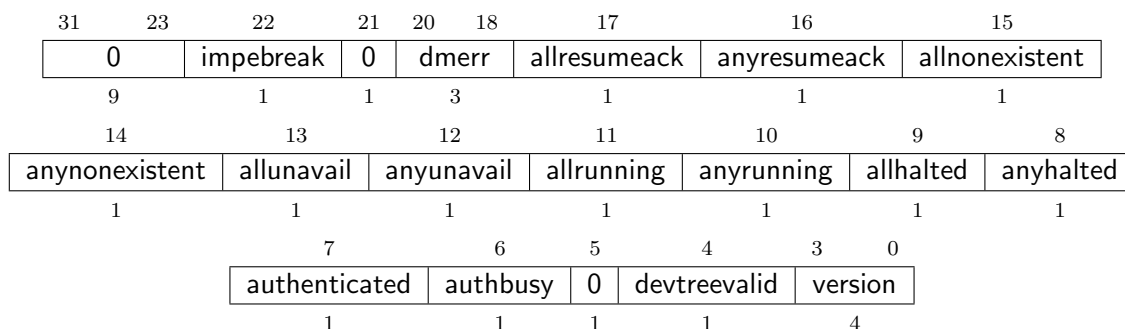
### 3.11.1 Debug Module Status (dmstatus, at 0x11)

The address of this register will not change in the future, because it contains **version**. It has changed from version 0.11 of this spec.

This register reports status for the overall debug module as well as the currently selected harts, as defined in **hasel**.

Harts are nonexistent if they will never be part of this system, no matter how long a user waits. Eg. in a simple single-hart system only one hart exists, and all others are nonexistent.

Harts are unavailable if they might exist/become available at a later time. Eg. in a multi-hart system some might temporarily be powered down, or a system might support hot-swapping harts.



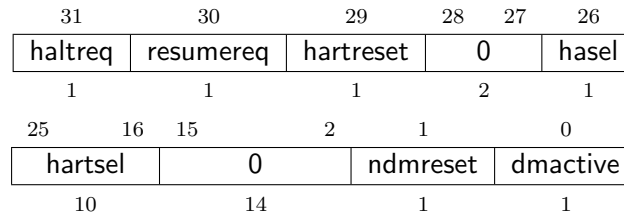
Field	Description	Access	Reset
impebreak	If 1, then there is an implicit <b>ebreak</b> instruction at the non-existent word immediately after the Program Buffer. This saves the debugger from having to write the <b>ebreak</b> itself, and allows the Program Buffer to be one word smaller. This must be 1 when <b>progbuFSIZE</b> is 1.	R	Preset
dmerr	Gets set if the Debug Module was accessed incorrectly. 0 (none): No error. 1 (badaddr): There was an access to an unimplemented Debug Module address. 7 (other): An access failed for another reason.	R/W1C	0

Continued on next page

allresumeack	This field is 1 when all currently selected harts have acknowledged the previous <code>resumereq</code> .	R	-
anyresumeack	This field is 1 when any currently selected hart has acknowledged the previous <code>resumereq</code> .	R	-
allnonexistent	This field is 1 when all currently selected harts do not exist in this system.	R	-
anynonexistent	This field is 1 when any currently selected hart does not exist in this system.	R	-
allunavail	This field is 1 when all currently selected harts are unavailable.	R	-
anyunavail	This field is 1 when any currently selected hart is unavailable.	R	-
allrunning	This field is 1 when all currently selected harts are running.	R	-
anyrunning	This field is 1 when any currently selected hart is running.	R	-
allhalted	This field is 1 when all currently selected harts are halted.	R	-
anyhalted	This field is 1 when any currently selected hart is halted.	R	-
authenticated	0 when authentication is required before using the DM. 1 when the authentication check has passed. On components that don't implement authentication, this bit must be preset as 1.	R	Preset
authbusy	0: The authentication module is ready to process the next read/write to <code>authdata</code> . 1: The authentication module is busy. Accessing <code>authdata</code> results in unspecified behavior. <code>authbusy</code> only becomes set in immediate response to an access to <code>authdata</code> .	R	0
devtreevalid	0: <code>devtreeaddr0</code> – <code>devtreeaddr3</code> hold information which is not relevant to the Device Tree. 1: <code>devtreeaddr0</code> – <code>devtreeaddr3</code> registers hold the address of the Device Tree.	R	Preset
version	0: There is no Debug Module present. 1: There is a Debug Module and it conforms to version 0.11 of this specification. 2: There is a Debug Module and it conforms to version 0.13 of this specification. 15: There is a Debug Module but it does not conform to any available version of this spec.	R	2

### 3.11.2 Debug Module Control (`dmcontrol`, at `0x10`)

This register controls the overall debug module as well as the currently selected harts, as defined in `hasel`.



Field	Description	Access	Reset
haltreq	Halt request signal for all currently selected harts. When set to 1, each selected hart will halt if it is not currently halted. Writing 1 or 0 has no effect on a hart which is already halted, but the bit should be cleared to 0 before the hart is resumed. Writes apply to the new value of <code>hartsel</code> and <code>hasel</code> .	R/W	0
resumereq	Resume request signal for all currently selected harts. When set to 1, each selected hart will resume if it is currently halted. This bit is ignored while <code>haltreq</code> is set. Writes apply to the new value of <code>hartsel</code> and <code>hasel</code> .	R/W	0
hartreset	This optional bit controls reset to all the currently selected harts. To perform a reset the debugger writes 1, and then writes 0 to deassert the reset signal. If this feature is not implemented, the bit always stays 0, so after writing 1 the debugger can read the register back to see if the feature is supported. Writes apply to the new value of <code>hartsel</code> and <code>hasel</code> .	R/W	0
hasel	Selects the definition of currently selected harts. 0: There is a single currently selected hart, that selected by <code>hartsel</code> . 1: There may be multiple currently selected harts – that selected by <code>hartsel</code> , plus those selected by the hart array mask register. An implementation which does not implement the hart array mask register should tie this field to 0. A debugger which wishes to use the hart array mask register feature should set this bit and read back to see if the functionality is supported.	R/W	0
Continued on next page			



<b>hartsel</b>	The DM-specific index of the hart to select. This hart is always part of the currently selected harts.	R/W	0
<b>ndmreset</b>	This bit controls the reset signal from the DM to the rest of the system. The signal should reset every part of the system, including every hart, except for the DM and any logic required to access the DM. To perform a system reset the debugger writes 1, and then writes 0 to deassert the reset.	R/W	0
<b>dmactive</b>	<p>This bit serves as a reset signal for the Debug Module itself.</p> <p>0: The module's state, including authentication mechanism, takes its reset values (the <b>dmactive</b> bit is the only bit which can be written to something other than its reset value).</p> <p>1: The module functions normally.</p> <p>No other mechanism should exist that may result in resetting the Debug Module after power up, including the platform's system reset or Debug Transport reset signals.</p> <p>A debugger may pulse this bit low to get the debug module into a known state.</p> <p>Implementations may use this bit to aid debugging, for example by preventing the Debug Module from being power gated while debugging is active.</p>	R/W	0

### 3.11.3 Hart Info (hartinfo, at 0x12)

This register gives information about the hart currently selected by **hartsel**.

This register is optional. If it is not present it should read all-zero.

If this register is included, the debugger can do more with the Program Buffer by writing programs which explicitly access the **data** and/or **dscratch** registers.

This entire register is read-only.

31	24	23	20	19	17	16	15	12	11	0
0	nscratch	0	dataaccess	datasize	dataaddr					
8	4	3	1	4	12					

Field	Description	Access	Reset
Continued on next page			

<b>nscratch</b>	Number of <b>dscratch</b> registers available for the debugger to use during program buffer execution, starting from <b>dscratch0</b> . The debugger can make no assumptions about the contents of these registers between commands.	R	Preset
<b>dataaccess</b>	0: The <b>data</b> registers are shadowed in the hart by CSR registers. Each CSR register is XLEN bits in size, and corresponds to a single argument, per Table 3.2. 1: The <b>data</b> registers are shadowed in the hart's memory map. Each register takes up 4 bytes in the memory map.	R	Preset
<b>datasize</b>	If <b>dataaccess</b> is 0: Number of CSR registers dedicated to shadowing the <b>data</b> registers. If <b>dataaccess</b> is 1: Number of 32-bit words in the memory map dedicated to shadowing the <b>data</b> registers.	R	Preset
<b>dataaddr</b>	If <b>dataaccess</b> is 0: The number of the first CSR dedicated to shadowing the <b>data</b> registers. If <b>dataaccess</b> is 1: Signed address of RAM where the <b>data</b> registers are shadowed, to be used to access relative to <b>zero</b> .	R	Preset

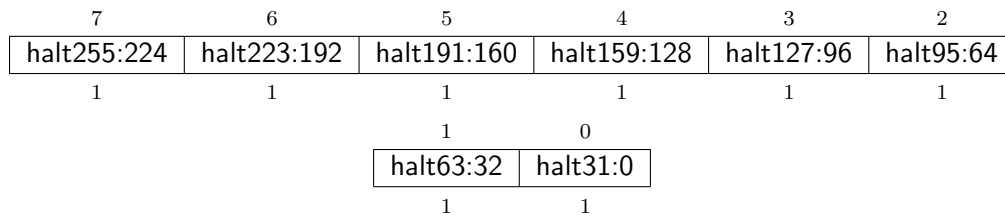
### 3.11.4 Halt Summary (haltsum, at 0x13)

This register contains a summary of which harts are halted.

Each bit contains the logical OR of 32 halt bits. When there are a large number of harts in the system, the debugger can first read this register, and then read from the halt region (0x40–0x5f) to determine which hart is the one that is halted.

This entire register is read-only.

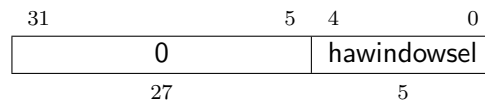
31	30	29	28	27	26
halt1023:992	halt991:960	halt959:928	halt927:896	halt895:864	halt863:832
1	1	1	1	1	1
25	24	23	22	21	20
halt831:800	halt799:768	halt767:736	halt735:704	halt703:672	halt671:640
1	1	1	1	1	1
19	18	17	16	15	14
halt639:608	halt607:576	halt575:544	halt543:512	halt511:480	halt479:448
1	1	1	1	1	1
13	12	11	10	9	8
halt447:416	halt415:384	halt383:352	halt351:320	halt319:288	halt287:256
1	1	1	1	1	1



### 3.11.5 Hart Array Window Select (hawindowssel, at 0x14)

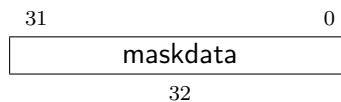
This register selects which of the 32-bit portion of the hart array mask register is accessible in [hawindow](#).

The hart array mask register provides a mask of all harts controlled by the debug module. A hart is part of the currently selected harts if the corresponding bit is set in the hart array mask register and hasel in [dmcontrol](#) is 1, or if the hart is selected by [hartsel](#).

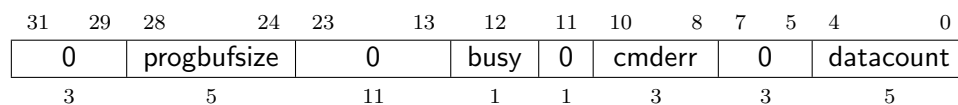


### 3.11.6 Hart Array Window (hawindow, at 0x15)

This register provides R/W access to a 32-bit portion of the hart array mask register. The position of the window is determined by [hawindowssel](#). I.e. bit 0 refers to hart  $\text{hawindowssel} * 32$ , while bit 31 refers to hart  $\text{hawindowssel} * 32 + 31$ .



### 3.11.7 Abstract Control and Status (abstractcs, at 0x16)



Field	Description	Access	Reset
progbuysize	Size of the Program Buffer, in 32-bit words. Valid sizes are 0 - 16.	R	Preset

Continued on next page

busy	1: An abstract command is currently being executed. This bit is set as soon as <code>command</code> is written, and is not cleared until that command has completed.	R	0
cmderr	Gets set if an abstract command fails. The bits in this field remain set until they are cleared by writing 1 to them. No abstract command is started until the value is reset to 0. 0 (none): No error. 1 (busy): An abstract command was executing while <code>command</code> , <code>abstractcs</code> , <code>abstractauto</code> was written, or when one of the <code>data</code> or <code>progbuf</code> registers was read or written. 2 (not supported): The requested command is not supported. A command that is not supported while the hart is running may be supported when it is halted. 3 (exception): An exception occurred while executing the command (eg. while executing the Program Buffer). 4 (halt/resume): An abstract command couldn't execute because the hart wasn't in the expected state (running/halted). 7 (other): The command failed for another reason.	R/W1C	0
Continued on next page			

datacount	Number of <b>data</b> registers that are implemented as part of the abstract command interface. Valid sizes are 0 - 12.	R	Preset
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### 3.11.8 Abstract Command (command, at 0x17)

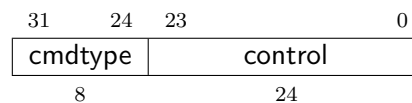
Writes to this register cause the corresponding abstract command to be executed.

Writing while an abstract command is executing causes **cmderr** to be set.

If **cmderr** is non-zero, writes to this register are ignored.

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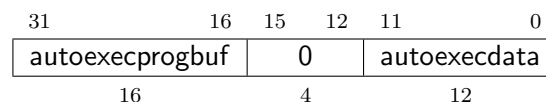
*cmderr inhibits starting a new command to accommodate debuggers that, for performance reasons, send several commands to be executed in a row without checking **cmderr** in between. They can safely do so and check **cmderr** at the end without worrying that one command failed but then a later command (which might have depended on the previous one succeeding) passed.*



Field	Description	Access	Reset
cmdtype	The type determines the overall functionality of this abstract command.	W	0
control	This field is interpreted in a command-specific manner, described for each abstract command.	W	0

### 3.11.9 Abstract Command Autoexec (abstractauto, at 0x18)

This register is optional. Including it allows more efficient burst accesses. Debugger can attempt to set bits and read them back to determine if the functionality is supported.



Field	Description	Access	Reset
autoexecprogbuf	When a bit in this field is 1, read or write accesses the corresponding <b>progbuf</b> word cause the command in <a href="#">command</a> to be executed again.	R/W	0
Continued on next page			

autoexecdata	When a bit in this field is 1, read or write accesses the corresponding <b>data</b> word cause the command in <b>command</b> to be executed again.	R/W	0
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### 3.11.10 Device Tree Addr 0 (devtreeaddr0, at 0x19)

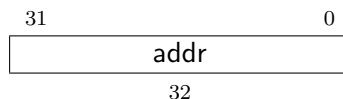
When **devtreevalid** is set, reading this register returns bits 31:0 of the Device Tree address. Reading the other **devtreeaddr** registers returns the upper bits of the address.

When system bus mastering is implemented, this must be an address that can be used with the System Bus Access module. Otherwise, this must be an address that can be used to access the Device Tree from the hart with ID 0.

If **devtreevalid** is 0, then the **devtreeaddr** registers hold identifier information which is not further specified in this document.

The Device Tree itself is described in the RISC-V Privileged Specification.

This entire register is read-only.



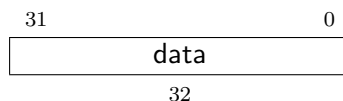
### 3.11.11 Abstract Data 0 (data0, at 0x04)

Basic read/write registers that may be read or changed by abstract commands.

Accessing them while an abstract command is executing causes **cmderr** to be set.

Attempts to write them while **busy** is set does not change their value.

The values in these registers may not be preserved after an abstract command is executed. The only guarantees on their contents are the ones offered by the command in question. If the command fails, no assumptions can be made about the contents of these registers.

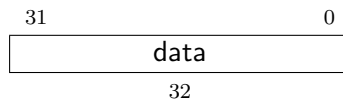


### 3.11.12 Program Buffer 0 (progbuf0, at 0x20)

The **progbuf** registers provide read/write access to the optional program buffer.

Accessing them while an abstract command is executing causes **cmderr** to be set.

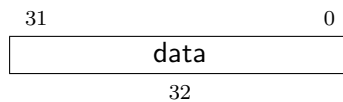
Attempts to write them while `busy` is set does not change their value.



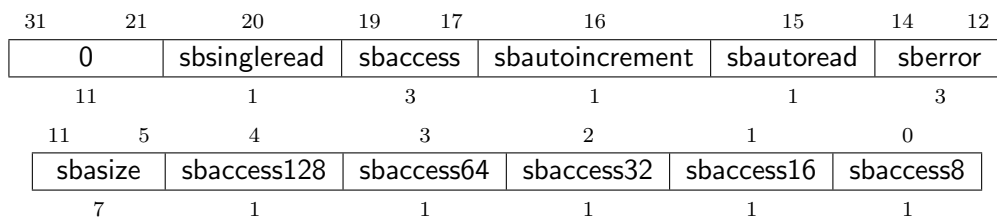
### 3.11.13 Authentication Data (`authdata`, at `0x30`)

This register serves as a 32-bit serial port to the authentication module.

When `authbusy` is clear, the debugger can communicate with the authentication module by reading or writing this register. There is no separate mechanism to signal overflow/underflow.



### 3.11.14 System Bus Access Control and Status (`sbc`s, at `0x38`)



Field	Description	Access	Reset
<code>sbingleread</code>	When a 1 is written here, triggers a read at the address in <code>sbaddress</code> using the access size set by <code>sbaccess</code> .	W1	0
<code>sbaccess</code>	Select the access size to use for system bus accesses triggered by writes to the <code>sbaddress</code> registers or <code>sbdta0</code> . 0: 8-bit 1: 16-bit 2: 32-bit 3: 64-bit 4: 128-bit If an unsupported system bus access size is written here, the DM does not perform the access and <code>sberror</code> is set to 3.	R/W	2

Continued on next page

<b>sbautoincrement</b>	When 1, <b>sbaddress</b> is incremented by the access size (in bytes) selected in <b>sbaccess</b> after every system bus access.	R/W	0
<b>sbautoread</b>	When 1, every read from <b>sbdatab0</b> automatically triggers a system bus read at the (possibly auto-incremented) address.	R/W	0
<b>sberror</b>	When the debug module's system bus master causes a bus error, this field gets set. The bits in this field remain set until they are cleared by writing 1 to them. While this field is non-zero, no more system bus accesses can be initiated by the debug module. 0: There was no bus error. 1: There was a timeout. 2: A bad address was accessed. 3: There was some other error (eg. alignment). 4: The system bus master was busy when one of the <b>sbaddress</b> or <b>sbdatab0</b> registers was written, or <b>sbdatab0</b> was read when it had stale data.	R/W1C	0
<b>sbsize</b>	Width of system bus addresses in bits. (0 indicates there is no bus access support.)	R	Preset
<b>sbaccess128</b>	1 when 128-bit system bus accesses are supported.	R	Preset
<b>sbaccess64</b>	1 when 64-bit system bus accesses are supported.	R	Preset
<b>sbaccess32</b>	1 when 32-bit system bus accesses are supported.	R	Preset
<b>sbaccess16</b>	1 when 16-bit system bus accesses are supported.	R	Preset
<b>sbaccess8</b>	1 when 8-bit system bus accesses are supported.	R	Preset

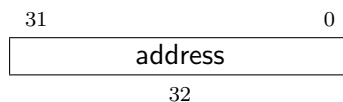
### 3.11.15 System Bus Address 31:0 (**sbaddress0**, at 0x39)

If **sbsize** is 0, then this register is not present.

When the system bus master is busy, writes to this register will set **sberror**.

If **sberror** is 0 and **sbautoread** is set then the system bus master will start to read after updating the address from **address**. The access size is controlled by **sbaccess** in **sbus**.

If **sbsingleread** is set, the bit is cleared.

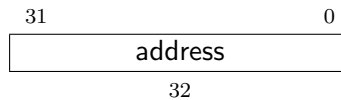


Field	Description	Access	Reset
Continued on next page			



address	Accesses bits 31:0 of the physical address in <b>sbaddress</b> .	R/W	0
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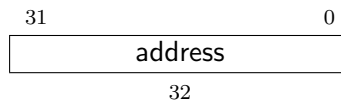
### 3.11.16 System Bus Address 63:32 (sbaddress1, at 0x3a)



Field	Description	Access	Reset
address	Accesses bits 63:32 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.11.17 System Bus Address 95:64 (sbaddress2, at 0x3b)

If **sbasize** is less than 65, then this register is not present.



Field	Description	Access	Reset
address	Accesses bits 95:64 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.11.18 System Bus Data 31:0 (sbddata0, at 0x3c)

If all of the **sbaccess** bits in **sbcs** are 0, then this register is not present.

Any successful system bus read updates the data in this register, and marks it no longer stale.

If **sberror** isn't 0 then accesses do nothing.

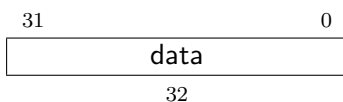
Writes to this register:

1. If the bus master is busy then accesses set **sberror**, and don't do anything else.
2. Start a bus write of **sbddata** to **sbaddress**.
3. If **sbautoincrement** is set, increment **sbaddress**.

Reads from this register:

1. If the register is marked stale, then set **sberror** and don't do anything else.
2. "Return" the data.
3. Mark the register stale.
4. If **sbautoincrement** is set, increment **sbaddress**.
5. If **sbautoread** is set, start another system bus read.

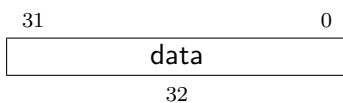
Only **sbdata0** has this behavior. The other **sbdata** registers have no side effects. On systems that have buses wider than 32 bits, a debugger should access **sbdata0** after accessing the other **sbdata** registers.



Field	Description	Access	Reset
data	Accesses bits 31:0 of <b>sbdata</b> .	R/W	0

### 3.11.19 System Bus Data 63:32 (**sbdata1**, at 0x3d)

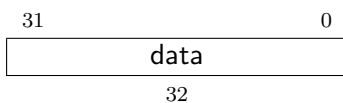
If **sbaccess64** and **sbaccess128** are 0, then this register is not present.



Field	Description	Access	Reset
data	Accesses bits 63:32 of <b>sbdata</b> (if the system bus is that wide).	R/W	0

### 3.11.20 System Bus Data 95:64 (**sbdata2**, at 0x3e)

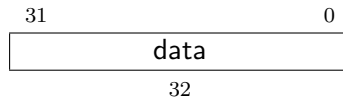
This register only exists if **sbaccess128** is 1.



Field	Description	Access	Reset
<b>data</b>	Accesses bits 95:64 of <b>sbdata</b> (if the system bus is that wide).	R/W	0

### 3.11.21 System Bus Data 127:96 (sbdata3, at 0x3f)

This register only exists if **sbaccess128** is 1.



Field	Description	Access	Reset
<b>data</b>	Accesses bits 127:96 of <b>sbdata</b> (if the system bus is that wide).	R/W	0



## Chapter 4

# RISC-V Debug

Modifications to the RISC-V core to support debug are kept to a minimum. There is a special execution mode (Debug Mode) and a few extra CSRs. The DM takes care of the rest.

### 4.1 Debug Mode

Debug Mode is a special processor mode used only when the core is halted for external debugging. How Debug Mode is implemented is not specified here.

When executing code from the Program Buffer, the processor stays in Debug Mode and the following apply:

1. All operations are executed at machine mode privilege level, except that `mprv` in `mstatus` is ignored.
2. All interrupts are masked.
3. Exceptions don't update any registers. That includes `cause`, `epc`, `badaddr`, `dpc`, and `mstatus`. They do end execution of the Program Buffer.
4. No action is taken if a trigger matches.
5. Trace is disabled.
6. Counters may be stopped, depending on `stopcount` in `dcsr`.
7. Timers may be stopped, depending on `stoptime` in `dcsr`.
8. The `wfi` instruction acts as a `nop`.
9. Almost all instructions that change the privilege level have undefined behavior. This includes `ecall`, `mret`, `hret`, `sret`, and `uret`. (To change the privilege level, the debugger can write `prv` in `dcsr`). The only exception is `ebreak`. When that is executed in Debug Mode, it halts the processor again but without updating `dpc` or `dcsr`.

## 4.2 Load-Reserved/Store-Conditional Instructions

The reservation registered by an `lr` instruction on a memory address may be lost when entering Debug Mode or while in Debug Mode. This means that there may be no forward progress if Debug Mode is entered between `lr` and `sc` pairs.

## 4.3 Reset

If the halt signal is asserted when a core comes out of reset, the core must enter Debug Mode before executing any instructions, but after performing any initialization that would usually happen before the first instruction is executed.

### 4.3.1 `dret` Instruction

To return from Debug Mode, a new instruction is defined: `dret`. It has an encoding of 0x7b200073. On harts which support this instruction, executing `dret` in Debug Mode changes `pc` to the value stored in `dpc`. The current privilege level is changed to that specified by `prv` in `dcsr`. The hart is no longer in debug mode.

It is not necessary for the debugger to know whether an implementation supports `dret`, as the Debug Module will ensure that it is executed if necessary. It is defined in this specification only to reserve the opcode and allow for reusable Debug Module implementations.

## 4.4 Core Debug Registers

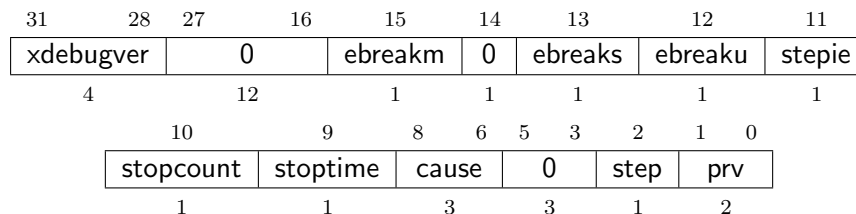
The supported Core Debug Registers must be implemented for each hart that can be debugged.

These registers are only accessible from Debug Mode.

Table 4.1: Core Debug Registers

Address	Name	Page
0x7b0	Debug Control and Status	<a href="#">33</a>
0x7b1	Debug PC	<a href="#">34</a>
0x7b2	Debug Scratch Register 0	
0x7b3	Debug Scratch Register 1	

## 4.4.1 Debug Control and Status (dcsr, at 0x7b0)



Field	Description	Access	Reset
xdebugver	0: There is no external debug support. 4: External debug support exists as it is described in this document. 15: There is external debug support, but it does not conform to any available version of this spec.	R	Preset
ebreakm	When 1, <b>ebreak</b> instructions in Machine Mode enter Debug Mode.	R/W	0
ebreaks	When 1, <b>ebreak</b> instructions in Supervisor Mode enter Debug Mode.	R/W	0
ebreaku	When 1, <b>ebreak</b> instructions in User/Application Mode enter Debug Mode.	R/W	0
stepie	0: Interrupts are disabled during single stepping. 1: Interrupts are enabled during single stepping. Implementations may hard wire this bit to 0. The debugger must read back the value it writes to check whether the feature is supported. If not supported, interrupt behavior can be emulated by the debugger.	R/W	0
stopcount	0: Increment counters as usual. 1: Don't increment any counters while in Debug Mode or on <b>ebreak</b> instructions that cause entry into Debug Mode. These counters include the <b>cycle</b> and <b>instret</b> CSRs. This is preferred for most debugging scenarios. An implementation may choose not to support writing to this bit. The debugger must read back the value it writes to check whether the feature is supported.	R/W	Preset

Continued on next page

stoptime	<p>0: Increment timers as usual.</p> <p>1: Don't increment any hart-local timers while in Debug Mode.</p> <p>An implementation may choose not to support writing to this bit. The debugger must read back the value it writes to check whether the feature is supported.</p>	R/W	Preset
cause	<p>Explains why Debug Mode was entered.</p> <p>When there are multiple reasons to enter Debug Mode in a single cycle, the cause with the highest priority is the one written.</p> <p>1: An <b>ebreak</b> instruction was executed. (priority 3)</p> <p>2: The Trigger Module caused a breakpoint exception. (priority 4)</p> <p>3: The debugger requested entry to Debug Mode. (priority 2)</p> <p>4: The hart single stepped because <b>step</b> was set. (priority 1)</p> <p>Other values are reserved for future use.</p>	R	0
step	<p>When set and not in Debug Mode, the hart will only execute a single instruction and then enter Debug Mode. If the instruction does not complete due to an exception, the hart will immediately enter Debug Mode before executing the trap handler, with appropriate exception registers set.</p>	R/W	0
prv	<p>Contains the privilege level the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.4. A debugger can change this value to change the hart's privilege level when exiting Debug Mode.</p> <p>Not all privilege levels are supported on all harts. If the encoding written is not supported or the debugger is not allowed to change to it, the hart may change to any supported privilege level.</p>	R/W	0

#### 4.4.2 Debug PC (dpc, at 0x7b1)

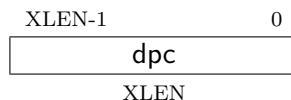
Upon entry to debug mode, [dpc](#) is updated with the virtual address of the next instruction to be executed. The behavior is described in more detail in Table 4.2.



Table 4.2: Virtual address in DPC upon Debug Mode Entry

Cause	Virtual Address in DPC
<b>ebreak</b>	Address of the <b>ebreak</b> instruction
single step	Address of the instruction that would be executed next if no debugging was going on. Ie. <code>pc + 4</code> for 32-bit instructions that don't change program flow, the destination PC on taken jumps/branches, etc.
trigger module	If <code>timing</code> is 0, the address of the instruction which caused the trigger to fire. If <code>timing</code> is 1, the address of the next instruction to be executed at the time that debug mode was entered.
halt request	Address of the next instruction to be executed at the time that debug mode was entered

When resuming, the hart's PC is updated to the virtual address stored in `dpc`. A debugger may write `dpc` to change where the hart resumes.



#### 4.4.3 Debug Scratch Register 0 (dscratch0, at 0x7b2)

Optional scratch register that can be used by implementations that need it. A debugger must not write to this register unless `hartinfo` explicitly mentions it (the Debug Module may use this register internally).

#### 4.4.4 Debug Scratch Register 1 (dscratch1, at 0x7b3)

Optional scratch register that can be used by implementations that need it. A debugger must not write to this register unless `hartinfo` explicitly mentions it (the Debug Module may use this register internally).

### 4.5 Virtual Debug Registers

Virtual debug registers are a requirement on the debugger SW/interface, not on the Core designer.

Users of the debugger shouldn't need to know about the core debug registers, but may want to change things affected by them. A virtual register is one that doesn't exist directly in the hardware, but that the debugger exposes as if it does.

Table 4.3: Virtual Core Debug Registers

Address	Name	Page
virtual	Privilege Level	36

Table 4.4: Privilege Level Encoding

Encoding	Privilege Level
0	User/Application
1	Supervisor
3	Machine

#### 4.5.1 Privilege Level (priv, at virtual)

User can read this register to inspect the privilege level that the hart was running in when the hart halted. User can write this register to change the privilege level that the hart will run in when it resumes.

This register contains `priv` from `dcsr`, but in a place that the user is expected to access. The user should not access `dcsr` directly, because doing so might interfere with the debugger.



Field	Description	Access	Reset
<b>priv</b>	Contains the privilege level the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.4, and matches the privilege level encoding from the RISC-V Privileged ISA Specification. A user can write this value to change the hart's privilege level when exiting Debug Mode.	R/W	0

## Chapter 5

# Trigger Module

Triggers can cause a breakpoint exception, entry into Debug Mode, or a trace action without having to execute a special instruction. This makes them invaluable when debugging code from ROM. They can trigger on execution of instructions at a given memory address, or on the address/data in loads/stores. These are all features that can be useful without having the Debug Module present, so the Trigger Module is broken out as a separate piece that can be implemented separately.

Each trigger may support a variety of features. A debugger can build a list of all triggers and their features as follows:

1. Write 0 to `tselect`.
2. Read back `tselect` to confirm this trigger exists. If not, exit.
3. Read `tdata1`, and possible `tdata2` and `tdata3` depending on the trigger type.
4. If type is 0, this trigger doesn't exist. Exit the loop.
5. Repeat, incrementing the value in `tselect`.

---

*There are two ways to check whether a given trigger is the last one to support these implementations:*

1. *When no hardware triggers are implemented at all, all related registers return 0. The algorithm above terminates when checking type.*
2. *When 2 triggers are implemented, `tselect` is just a single bit that selects one of the two. When the debugger writes 2, it reads back as 0 which terminates the enumeration.*

### 5.1 Trigger Registers

The trigger registers are only accessible in machine and Debug Mode to prevent untrusted user code from causing entry into Debug Mode without the OS's permission.

Table 5.1: Trigger Registers

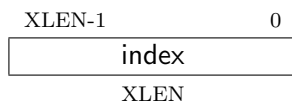
Address	Name	Page
0x7a0	Trigger Select	<a href="#">38</a>
0x7a1	Trigger Data 1	<a href="#">38</a>
0x7a1	Match Control	<a href="#">39</a>
0x7a1	Instruction Count	<a href="#">43</a>
0x7a2	Trigger Data 2	<a href="#">39</a>
0x7a3	Trigger Data 3	<a href="#">39</a>

### 5.1.1 Trigger Select (tselect, at 0x7a0)

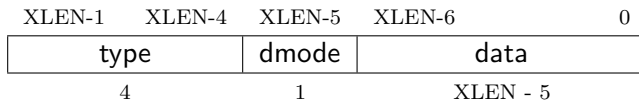
This register determines which trigger is accessible through the other trigger registers. The set of accessible triggers must start at 0, and be contiguous.

Writes of values greater than or equal to the number of supported triggers may result in a different value in this register than what was written. Debuggers should read back the value to confirm that what they wrote was a valid index.

Since triggers can be used both by Debug Mode and M Mode, the debugger must restore this register if it modifies it.



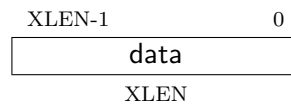
### 5.1.2 Trigger Data 1 (tdata1, at 0x7a1)



Field	Description	Access	Reset
<b>type</b>	0: There is no trigger at this <a href="#">tselect</a> . 1: The trigger is a legacy SiFive address match trigger. These should not be implemented and aren't further documented here. 2: The trigger is an address/data match trigger. The remaining bits in this register act as described in <a href="#">mcontrol</a> . 3: The trigger is an instruction count trigger. The remaining bits in this register act as described in <a href="#">icount</a> . 15: This trigger exists (so enumeration shouldn't terminate), but is not currently available. Other values are reserved for future use.	R	Preset
<b>dmode</b>	0: Both Debug and M Mode can write the <b>tdata</b> registers at the selected <a href="#">tselect</a> . 1: Only Debug Mode can write the <b>tdata</b> registers at the selected <a href="#">tselect</a> . Writes from other modes are ignored. This bit is only writable from Debug Mode.	R/W	0
<b>data</b>	Trigger-specific data.	R/W	Preset

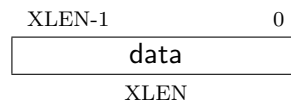
### 5.1.3 Trigger Data 2 (tdata2, at 0x7a2)

Trigger-specific data.



### 5.1.4 Trigger Data 3 (tdata3, at 0x7a3)

Trigger-specific data.



### 5.1.5 Match Control (mcontrol, at 0x7a1)

This register is accessible as [tdata1](#) when **type** is 2.

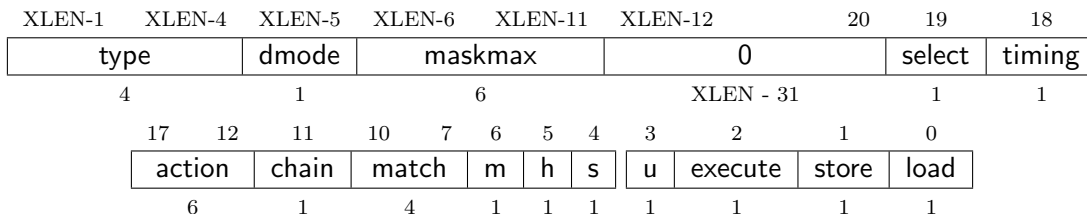
Writing unsupported values to any field in this register results in the reset value being written

instead. When a debugger wants to use a feature, it must write the appropriate value and then read back the register to determine whether it is supported.

Address and data trigger implementation are heavily dependent on how the processor core is implemented. To accommodate various implementations, execute, load, and store address/data triggers may fire at whatever point in time is most convenient for the implementation. The debugger may request specific timings as described in timing. Table 5.2 suggests timings for the best user experience.

Table 5.2: Suggested Breakpoint Timings

Match Type	Suggested Trigger Timing
Execute Address	Before
Execute Instruction	Before
Execute Address+Instruction	Before
Load Address	Before
Load Data	After
Load Address+Data	After
Store Address	Before
Store Data	Before
Store Address+Data	Before



Field	Description	Access	Reset
maskmax	Specifies the largest naturally aligned powers-of-two (NAPOT) range supported by the hardware. The value is the logarithm base 2 of the number of bytes in that range. A value of 0 indicates that only exact value matches are supported (one byte range). A value of 63 corresponds to the maximum NAPOT range, which is $2^{63}$ bytes in size.	R	Preset

Continued on next page

select	0: Perform a match on the virtual address. 1: Perform a match on the data value loaded/stored, or the instruction executed.	R/W	0
timing	0: The action for this trigger will be taken just before the instruction that triggered it is executed, but after all preceding instructions are committed. 1: The action for this trigger will be taken after the instruction that triggered it is executed. It should be taken before the next instruction is executed, but it is better to implement triggers and not implement that suggestion than to not implement them at all. Most hardware will only implement one timing or the other, possibly dependent on <b>select</b> , <b>execute</b> , <b>load</b> , and <b>store</b> . This bit primarily exists for the hardware to communicate to the debugger what will happen. Hardware may implement the bit fully writable, in which case the debugger has a little more control. Data load triggers with <b>timing</b> of 0 will result in the same load happening again when the debugger lets the core run. For data load triggers, debuggers must first attempt to set the breakpoint with <b>timing</b> of 1. A chain of triggers that don't all have the same <b>timing</b> value will never fire (unless consecutive instructions match the appropriate triggers).	R/W	0
Continued on next page			

action	<p>Determines what happens when this trigger matches.</p> <p>0: Raise a breakpoint exception. (Used when software wants to use the trigger module without an external debugger attached.)</p> <p>1: Enter Debug Mode. (Only supported when <code>dmode</code> is 1.)</p> <p>2: Start tracing.</p> <p>3: Stop tracing.</p> <p>4: Emit trace data for this match. If it is a data access match, emit appropriate Load/Store Address/Data. If it is an instruction execution, emit its PC.</p> <p>Other values are reserved for future use.</p>	R/W	0
chain	<p>0: When this trigger matches, the configured action is taken.</p> <p>1: While this trigger does not match, it prevents the trigger with the next index from matching.</p>	R/W	0
match	<p>0: Matches when the value equals <code>tdata2</code>.</p> <p>1: Matches when the top M bits of the value match the top M bits of <code>tdata2</code>. M is XLEN-1 minus the index of the least-significant bit containing 0 in <code>tdata2</code>.</p> <p>2: Matches when the value is greater than (unsigned) or equal to <code>tdata2</code>.</p> <p>3: Matches when the value is less than (unsigned) <code>tdata2</code>.</p> <p>4: Matches when the lower half of the value equals the lower half of <code>tdata2</code> after the lower half of the value is ANDed with the upper half of <code>tdata2</code>.</p> <p>5: Matches when the upper half of the value equals the lower half of <code>tdata2</code> after the upper half of the value is ANDed with the upper half of <code>tdata2</code>.</p> <p>Other values are reserved for future use.</p>	R/W	0

Continued on next page



m	When set, enable this trigger in M mode.	R/W	0
h	When set, enable this trigger in H mode.	R/W	0
s	When set, enable this trigger in S mode.	R/W	0
u	When set, enable this trigger in U mode.	R/W	0
execute	When set, the trigger fires on the virtual address or opcode of an instruction that is executed.	R/W	0
store	When set, the trigger fires on the virtual address or data of a store.	R/W	0
load	When set, the trigger fires on the virtual address or data of a load.	R/W	0

### 5.1.6 Instruction Count (icount, at 0x7a1)

This register is accessible as `tdata1` when type is 3.

**Warning! `icount` is just a proposal. So far nobody has commented on it, so it could very easily be removed or changed in the future.**

Writing unsupported values to any field in this register results in the reset value being written instead. When a debugger wants to use a feature, it must write the appropriate value and then read back the register to determine whether it is supported.

---

*This trigger type is intended to be used as a single step that's useful both for external debuggers and for software monitor programs. For that case it is not necessary to support count greater than 1. The only two combinations of the mode bits that are useful in those scenarios are u by itself, or m, h, s, and u all set.*

*If the hardware limits count to 1, and changes mode bits instead of decrementing count, this register can be implemented with just 2 bits. One for u, and one for m, h, and s tied together. If only the external debugger or only a software monitor needs to be supported, a single bit is enough.*

XLEN-1	XLEN-4	XLEN-5	XLEN-6	24	23	10	9	8	7	6	5	0
type	dmode	0	count	m	h	s	u	action				
4	1	XLEN - 29	14	1	1	1	1	6				

Field	Description	Access	Reset
count	When count is decremented to 0, the trigger fires. Instead of changing count from 1 to 0, it is also acceptable for hardware to clear m, h, s, and u. This allows count to be hard-wired to 1 if this register just exists for single step.	R/W	1
Continued on next page			

m	When set, every instruction completed or exception taken in M mode decrements <code>count</code> by 1.	R/W	0
h	When set, every instruction completed or exception taken in in H mode decrements <code>count</code> by 1.	R/W	0
s	When set, every instruction completed or exception taken in S mode decrements <code>count</code> by 1.	R/W	0
u	When set, every instruction completed or exception taken in U mode decrements <code>count</code> by 1.	R/W	0
action	<p>Determines what happens when this trigger matches.</p> <p>0: Raise a breakpoint exception. (Used when software wants to use the trigger module without an external debugger attached.)</p> <p>1: Enter Debug Mode. (Only supported when <code>dmode</code> is 1.)</p> <p>2: Start tracing.</p> <p>3: Stop tracing.</p> <p>4: Emit trace data for this match. If it is a data access match, emit appropriate Load/Store Address/Data. If it is an instruction execution, emit its PC.</p> <p>Other values are reserved for future use.</p>	R/W	0

## Chapter 6

# Debug Transport Module (DTM)

Debug Transport Modules provide access to the DM over one or more transports (eg. JTAG or USB).

There may be multiple DTMs in a single platform. Ideally every component that communicates with the outside world includes a DTM, allowing a platform to be debugged through every transport it supports. For instance a USB component could include a DTM. This would trivially allow any platform to be debugged over USB. All that is required is that the USB module already in use also has access to the Debug Module Interface.

Using multiple DTMs at the same time is not supported. It is left to the user to ensure this does not happen.

This specification defines a JTAG DTM in Section 6.1. Additional DTMs may be added in future versions of this specification.

## 6.1 JTAG Debug Transport Module

This Debug Transport Module is based around a normal JTAG Test Access Port (TAP). The JTAG TAP allows access to arbitrary JTAG registers by first selecting one using the JTAG instruction register (IR), and then accessing it through the JTAG data register (DR).

### 6.1.1 JTAG Background

JTAG refers to IEEE Std 1149.1-2013. It is a standard that defines test logic that can be included in an integrated circuit to test the interconnections between integrated circuits, test the integrated circuit itself, and observe or modify circuit activity during the components normal operation. This specification uses the latter functionality. The JTAG standard defines a Test Access Port (TAP) that can be used to read and write a few custom registers, which can be used to communicate with debug hardware in a component.

### 6.1.2 JTAG DTM Registers

JTAG TAPs used as a DTM must have an IR of at least 5 bits. When the TAP is reset, IR must default to 00001, selecting the IDCODE instruction. A full list of JTAG registers along with their encoding is in Table 6.1. If the IR actually has more than 5 bits, then the encodings in Table 6.1 should be extended with 0's in their most significant bits. The only regular JTAG registers a debugger might use are BYPASS and IDCODE, but this specification leaves IR space for many other standard JTAG instructions. Unimplemented instructions must select the BYPASS register.

Table 6.1: JTAG DTM TAP Registers

Address	Name	Description	Page
0x00	BYPASS	JTAG recommends this encoding	47 48
0x01	IDCODE	JTAG recommends this encoding	
0x10	DTM Control and Status	For Debugging	
0x11	Debug Module Interface Access	For Debugging	
0x12	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x13	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x14	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x15	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x16	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x17	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x1f	BYPASS	JTAG requires this encoding	

### 6.1.3 IDCODE (at 0x01)

This register is selected (in IR) when the TAP state machine is reset. Its definition is exactly as defined in IEEE Std 1149.1-2013.

This entire register is read-only.

31	28	27	12	11	1	0
Version		PartNumber		Manufld		1
4		16		11		1

Field	Description	Access	Reset
Version	Identifies the release version of this part.	R	Preset
PartNumber	Identifies the designer's part number of this part.	R	Preset
Manufld	Identifies the designer/manufacturer of this part. Bits 6:0 must be bits 6:0 of the designer/manufacturer's Identification Code as assigned by JEDEC Standard JEP106. Bits 10:7 contain the modulo-16 count of the number of continuation characters (0x7f) in that same Identification Code.	R	Preset

### 6.1.4 DTM Control and Status (dtmcs, at 0x10)

The size of this register will remain constant in future versions so that a debugger can always determine the version of the DTM.

31	18	17	16	15	14	12	11	10	9	4	3	0
0	dmihardreset	dmireset	0	idle	dmistat	abits	version					
14	1	1	1	3	2	6	4					

Field	Description	Access	Reset
dmihardreset	Writing 1 to this bit does a hard reset of the DTM, causing the DTM to forget about any outstanding DMI transactions. In general this should only be used when the Debugger has reason to expect that the outstanding DMI transaction will never complete (e.g. a reset condition caused an inflight DMI transaction to be cancelled).	W1	0
dmireset	Writing 1 to this bit clears the sticky error state and allows the DTM to retry or complete the previous transaction.	W1	0
idle	This is a hint to the debugger of the minimum number of cycles a debugger should spend in Run-Test/Idle after every DMI scan to avoid a ‘busy’ return code (dmistat of 3). A debugger must still check dmistat when necessary. 0: It is not necessary to enter Run-Test/Idle at all. 1: Enter Run-Test/Idle and leave it immediately. 2: Enter Run-Test/Idle and stay there for 1 cycle before leaving. And so on.	R	Preset

Continued on next page

<b>dmistat</b>	0: No error. 1: Reserved. Interpret the same as 2. 2: An operation failed (resulted in <b>op</b> of 2). 3: An operation was attempted while a DMI access was still in progress (resulted in <b>op</b> of 3).	R	0
<b>abits</b>	The size of <b>address</b> in <a href="#">dmi</a> .	R	Preset
<b>version</b>	0: Version described in spec version 0.11. 1: Version described in spec version 0.13 (and later?), which reduces the DMI data width to 32 bits. 15: Version not described in any available version of this spec.	R	1

### 6.1.5 Debug Module Interface Access (**dmi**, at 0x11)

This register allows access to the Debug Module Interface (DMI).

In Update-DR, the DTM starts the operation specified in **op** unless the current status reported in **op** is sticky.

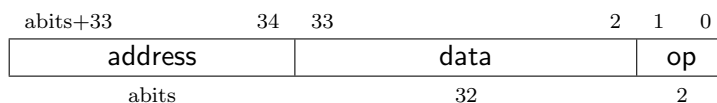
In Capture-DR, the DTM updates **data** with the result from that operation, updating **op** if the current **op** isn't sticky.

See Section [B.1](#) and Table [B.1](#) for examples of how this is used.

---

*The still-in-progress status is sticky to accommodate debuggers that batch together a number of scans, which must all be executed or stop as soon as there's a problem.*

*For instance a series of scans may write a Debug Program and execute it. If one of the writes fails but the execution continues, then the Debug Program may hang or have other unexpected side effects.*



Field	Description	Access	Reset
<b>address</b>	Address used for DMI access. In Update-DR this value is used to access the DM over the DMI.	R/W	0
<b>data</b>	The data to send to the DM over the DMI during Update-DR, and the data returned from the DM as a result of the previous operation.	R/W	0
Continued on next page			

op	<p>When the debugger writes this field, it has the following meaning:</p> <p>0: Ignore <b>data</b> and <b>address</b>. (nop)</p> <p>Don't send anything over the DMI during Update-DR. This operation should never result in a busy or error response. The address and data reported in the following Capture-DR are undefined.</p> <p>1: Read from <b>address</b>. (read)</p> <p>2: Write <b>data</b> to <b>address</b>. (write)</p> <p>3: Reserved.</p> <p>When the debugger reads this field, it means the following:</p> <p>0: The previous operation completed successfully.</p> <p>1: Reserved.</p> <p>2: A previous operation failed. The data scanned into <b>dmi</b> in this access will be ignored. This status is sticky and can be cleared by writing <b>dmireset</b> in <b>dtmcs</b>.</p> <p>This indicates that the DM itself responded with an error. Note: there are no specified cases in which the DM would respond with an error, and DMI is not required to support returning errors.</p> <p>3: An operation was attempted while a DMI request is still in progress. The data scanned into <b>dmi</b> in this access will be ignored. This status is sticky and can be cleared by writing <b>dmireset</b> in <b>dtmcs</b>. If a debugger sees this status, it needs to give the target more TCK edges between Update-DR and Capture-DR. The simplest way to do that is to add extra transitions in Run-Test/Idle. (The DTM, DM, and/or component may be in different clock domains, so synchronization may be required. Some relatively fixed number of TCK ticks may be needed for the request to reach the DM, complete, and for the response to be synchronized back into the TCK domain.)</p>	R/W	2
----	---	-----	---

### 6.1.6 BYPASS (at 0x1f)

1-bit register that has no effect. It is used when a debugger does not want to communicate with this TAP.

This entire register is read-only.



### 6.1.7 Recommended JTAG Connector

To make it easy to acquire debug hardware, this spec recommends a connector that is compatible with the Atmel AVR JTAG Connector, as described below.

The connector is a .05”-spaced, gold-plated male header with .016” thick hardened copper or beryllium bronze square posts (SAMTEC FTSH-105 or equivalent). Female connectors are compatible 20 $\mu$ m gold connectors.

Viewing the male header from above (the pins pointing at your eye), a target’s connector looks as it does in Table 6.2. The function of each pin is described in Table 6.3.

Table 6.2: JTAG Connector Diagram

TCK	1	2	GND
TDO	3	4	VCC
TMS	5	6	(SRSTn)
(NC)	7	8	(TRSTn)
TDI	9	10	GND

Table 6.3: JTAG Connector Pinout

1	TCK	JTAG TCK signal, driven by the debug adapter. This pin must be clearly marked in both male and female headers.
5	TMS	JTAG TMS signal, driven by debug adapter.
9	TDI	JTAG TDI signal, driven by the debug adapter.
3	TDO	JTAG TDO signal, driven by the target.
8	TRSTn	Test Reset (optional, only used by some devices. Used to reset the JTAG TAP Controller).
4	VCC	Reference voltage for logic high. A debug adapter may attempt to draw up to 20mA from this pin to power itself, but a target is not obligated to provide that power.
2, 10	GND	Target ground.
6	SRSTn	Active-low reset signal, driven by the debug adapter. Asserting reset should reset any RISC-V cores as well as any other peripherals on the PCB. It should not reset the debug logic. Although connecting this pin is optional, it is recommended as it allows the debugger to hold the target device in a reset state, which may be essential to debug some scenarios. If not implemented in a target, this pin must not be connected.



Target connectors may be shrouded. In that case the key slot should be next to pin 5. Female headers should have a matching key.

Debug adapters should be tagged or marked with their isolation voltage threshold (i.e. unisolated, 250V, etc.).

All debug adapter pins other than GND should be current-limited to 20mA.



# Appendix A

## Hardware Implementations

Below are two possible implementations. A designer could choose one, mix and match, or come up with their own design.

### A.1 Abstract Command Based

Halting happens by stalling the processor execution pipeline.

Muxes on the register file(s) allow for accessing GPRs and CSRs using the Access Register abstract command.

System Bus Access allows main memory access.

### A.2 Execution Based

This implementation only implements the Access Register abstract command for GPRs on a halted hart, and relies on the Program Buffer for all other operations.

This method uses the processor’s existing pipeline and ability to execute from arbitrary memory locations to avoid modifications to a processor’s datapath. When `haltreq` is set, the Debug Module raises a special interrupt to the selected hart(s). This interrupt causes each hart to enter Debug Mode and jump to a defined memory region that is serviced by the DM. When taking this exception, `pc` is saved to `dpc` and `cause` is updated in `dcsr`.

The code in the Debug Module causes the hart to execute a “park loop”. In the park loop the hart writes its `mhartid` to a memory location within the Debug Module to indicate that it is halted. To allow the DM to individually control one out of several halted harts, each hart polls for flags in a DM-controlled memory location to determine whether the debugger wants it to execute the Program Buffer or perform a resume.

To execute an abstract command, the DM first populates some internal words of program buffer

according to `command`. When `transfer` is set, the debugger populates these words with `lw <gpr>, 0x400(zero)` or `sw 0x400(zero), <gpr>`. 64- and 128-bit accesses use `ld/sd` and `lq/sq` respectively. If `transfer` is not set, these instructions are populated as `nops`. If `execute` is set, execution continues to the debugger-controlled Program Buffer, otherwise the debug module causes a `ebreak` to execute immediately.

When `ebreak` is executed (indicating the end of the Program Buffer code) the hart returns to its park loop. If an exception is encountered, the hart jumps to a defined debug exception address within the Debug Module. The code at that address causes the hart to write to an address in the Debug Module which indicates exception. Then the hart jumps back to the park loop. The DM infers from the write that there was an exception, and sets `cmderr` appropriately.

To resume execution, the debug module sets a flag which causes the core to execute a `dret`. When `dret` is executed, `pc` is restored from `dpc` and normal execution resumes at the privilege set by `prv`.

`data0` etc. are mapped into regular memory at an address relative to `zero` with only a 12-bit `imm`. The exact address is an implementation detail that a debugger must not rely on. For example, the `data` registers might be mapped to `0x400`.

For additional flexibility, `progbuf0`, etc. are mapped into regular memory immediately preceding `data0`, in order to form a contiguous region of memory which can be used for either program execution or data transfer.

## Appendix B

# Debugger Implementation

This section details how an external debugger might use the described debug interface to perform some common operations on RISC-V cores using the JTAG DTM described in Appendix ???. All these examples assume a 32-bit core but it should be easy to adapt the examples to 64- or 128-bit cores.

To keep the examples readable, they all assume that everything succeeds, and that they complete faster than the debugger can perform the next access. This will be the case in a typical JTAG setup. However, the debugger must always check the sticky error status bits after performing a sequence of actions. If it sees any that are set, then it should attempt the same actions again, possibly while adding in some delay, or explicit checks for status bits.

### B.1 Debug Module Interface Access

To read an arbitrary Debug Module register, select `dmi`, and scan in a value with `op` set to 1, and `address` set to the desired register address. In Update-DR the operation will start, and in Capture-DR its results will be captured into `data`. If the operation didn't complete in time, `op` will be 3 and the value in `data` must be ignored. The busy condition must be cleared by writing `dmireset` in `dtmcs`, and then the second scan must be performed again. This process must be repeated until `op` returns 0. In later operations the debugger should allow for more time between Capture-DR and Update-DR.

To write an arbitrary Debug Bus register, select `dmi`, and scan in a value with `op` set to 2, and `address` and `data` set to the desired register address and data respectively. From then on everything happens exactly as with a read, except that a write is performed instead of the read.

It should almost never be necessary to scan IR, avoiding a big part of the inefficiency in typical JTAG use.

## B.2 Main Loop

A debugger continuously monitors `haltsum` to see if any harts have spontaneously halted.

## B.3 Halting

To halt a hart, the debugger sets `hartsel` and `haltreq`. Then it waits for `allhalted` to become 1.

## B.4 Running

First, the debugger should restore any registers that it has clobbered. Once that's done, it can let the core run by setting `resumereq`.

## B.5 Single Step

Using the hardware single step feature is almost the same as regular running. The debugger just sets `step` in `dcsr` before letting the core run. The core behaves exactly as in the running case, except that interrupts may be disabled (depending on `stepie`) and it only fetches and executes a single instruction before re-entering Debug Mode.

## B.6 Accessing Registers

### B.6.1 Using Abstract Command

Read `s0` using abstract command:

Op	Address	Value	Comment
Write	<code>command</code>	size = 2, transfer, 0x1008	Read <code>s0</code>
Read	<code>data0</code>	-	Returns value that was in <code>s0</code>

Write `mstatus` using abstract command:

Op	Address	Value	Comment
Write	<code>data0</code>	new value	
Write	<code>command</code>	size = 2, transfer, write, 0x300	Write <code>mstatus</code>

### B.6.2 Using Program Buffer

Abstract commands are used to exchange data with GPRs. Using this mechanism, other registers can be accessed by moving their value into/out of GPRs.

Write `mstatus` using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>csrw s0, MSTATUS</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	new value	
Write	<code>command</code>	size = 2, postexec, transfer, write, 0x1008	Write <code>s0</code> , then execute program buffer

Read `f1` using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>fmv.x.s s0, f1</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>command</code>	<code>postexec</code>	Execute program buffer
Write	<code>command</code>	transfer 0x1008	read <code>s0</code>
Read	<code>data0</code>	-	Returns the value that was in <code>f1</code>

## B.7 Reading Memory

### B.7.1 Using System Bus Access

Read a word from memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbaddress0</code>	address	
Write	<code>sbc</code>	<code>sbaccess = 2, sbsingleread</code>	Perform a read
Read	<code>sbdatab0</code>	-	Value read from memory

Read block of memory using system bus access:

Op	Address	Value	Comment
Write	<a href="#">sbaddress0</a>	address	
Write	<a href="#">sbcs</a>	sbaccess = 2, sbsingleread, sbautoread, sbautoincrement	Turn on autoread and autoincrement, and perform a read
Read	<a href="#">sbddata0</a>	-	Value read from memory
Read	<a href="#">sbddata0</a>	-	Next value read from memory
...	...	...	...
Write	<a href="#">sbcs</a>	0	Clear sbautoread
Read	<a href="#">sbddata0</a>	-	Get last value read from memory.

### B.7.2 Using Program Buffer

Read a word from memory using program buffer:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	lw s0, 0(s0)	
Write	<a href="#">progbuf1</a>	ebreak	
Write	<a href="#">data0</a>	address	
Write	<a href="#">command</a>	write, postexec, 0x1008	Write s0, then execute program buffer
Write	<a href="#">command</a>	0x1008	Read s0
Read	<a href="#">data0</a>	-	Value read from memory

Read block of memory using program buffer:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	lw s1, 0(s0)	
Write	<a href="#">progbuf1</a>	addi s0, s0, 4	
Write	<a href="#">progbuf2</a>	ebreak	
Write	<a href="#">data0</a>	address	
Write	<a href="#">command</a>	write, postexec, 0x1008	Write s0, then execute program buffer
Write	<a href="#">command</a>	postexec, 0x1009	Read s1, then execute program buffer
Write	<a href="#">abstractauto</a>	autoexecdata [0]	Set autoexecdata [0]
Read	<a href="#">data0</a>	-	Get value read from memory, then execute program buffer
Read	<a href="#">data0</a>	-	Get next value read from memory, then execute program buffer
...	...	...	...
Write	<a href="#">abstractauto</a>	0	Clear autoexecdata [0]
Read	<a href="#">data0</a>	-	Get last value read from memory.

TODO: Table [B.1](#) shows the scans involved in reading a single word using this method.



Table B.1: Memory Read Timeline

	JTAG State	Activity
TODO	TODO	TODO

## B.8 Writing Memory

### B.8.1 Using System Bus Access

Write a word to memory using system bus access:

Op	Address	Value	Comment
Write	<a href="#">sbaddress0</a>	address	
Write	<a href="#">sbdata0</a>	value	

Write block of memory using system bus access:

Op	Address	Value	Comment
Write	<a href="#">sbaddress0</a>	address	
Write	<a href="#">sbcs</a>	sbaccess = 2, sbautoincrement	Turn on autoincrement
Write	<a href="#">sbdata0</a>	value0	
Write	<a href="#">sbdata0</a>	value1	
...	...	...	...
Write	<a href="#">sbdata0</a>	valueN	

### B.8.2 Using Program Buffer

Write a word to memory using program buffer:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	sw s1, 0(s0)	
Write	<a href="#">progbuf1</a>	ebreak	
Write	<a href="#">data0</a>	value	
Write	<a href="#">command</a>	write, 0x1008	Write s0
Write	<a href="#">data0</a>	address	
Write	<a href="#">command</a>	write, postexec, 0x1009	Write s1, then execute program buffer

Write block of memory using program buffer:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	<code>sw s1, 0(s0)</code>	
Write	<a href="#">progbuf1</a>	<code>addi s0, s0, 4</code>	
Write	<a href="#">progbuf2</a>	<code>ebreak</code>	
Write	<a href="#">data0</a>	address	
Write	<a href="#">command</a>	<code>write, 0x1008</code>	Write <code>s0</code>
Write	<a href="#">data0</a>	value0	
Write	<a href="#">command</a>	<code>write, postexec, 0x1009</code>	Write <code>s1</code> , then execute program buffer
Write	<a href="#">abstractauto</a>	<code>autoexecdata [0]</code>	Set <code>autoexecdata [0]</code>
Write	<a href="#">data0</a>	value1	
...	...	...	...
Write	<a href="#">data0</a>	valueN	
Write	<a href="#">abstractauto</a>	0	Clear <code>autoexecdata [0]</code>

## B.9 Handling Exceptions

Generally the debugger can avoid exceptions by being careful with the programs it writes. Sometimes they are unavoidable though, eg. if the user asks to access memory or a CSR that is not implemented. A typical debugger will not know enough about the platform to know what's going to happen, and must attempt the access to determine the outcome.

When an exception occurs while executing the Program Buffer, `cmderr` becomes set. The debugger can check this field to see whether a program encountered an exception. If there was an exception, it's left to the debugger to know what must have caused it.

## B.10 Quick Access

Halt the hart for a minimum amount of time to perform a single memory write.

There are a variety of instructions to transfer data between GPRs and the `data` registers. They are either loads/stores or CSR reads/writes. The specific addresses also vary. This is all specified in [hartinfo](#). The example here uses the pseudo-op `transfer dest, src` to represent all these options.

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	transfer arg2, s0	Save <b>s0</b>
Write	<a href="#">progbuf1</a>	transfer s0, arg0	Read first argument (address)
Write	<a href="#">progbuf2</a>	transfer arg0, s1	Save <b>s1</b>
Write	<a href="#">progbuf3</a>	transfer s1, arg1	Read second argument (data)
Write	<a href="#">progbuf4</a>	sw s1, 0(s0)	
Write	<a href="#">progbuf5</a>	transfer s1, arg0	Restore <b>s1</b>
Write	<a href="#">progbuf6</a>	transfer s0, arg2	Restore <b>s0</b>
Write	<a href="#">progbuf7</a>	ebreak	
Write	<a href="#">data0</a>	address	
Write	<a href="#">data1</a>	data	
Write	<a href="#">command</a>	0x10000000	Perform quick access



# Appendix C

## Future Ideas

**All items in this section are future ideas and should not be considered part of the specification.**

Some future version of this spec may implement some of the following features.

1. The spec defines several additions to the Device Tree which enable a debugger to discover hart IDs and supported triggers for all the cores in the system.
2. DTMs can function as general bus slaves, so they would look like regular RAM to bus masters.
3. Harts can be divided into groups. All the harts in the same group can be halted/run/stepped simultaneously. When a hart hits a breakpoint, all the other harts in the same group also halt within a few clock cycles.
4. DTMs are specified for protocols like USB, I2C, SPI, and SWD.
5. Core registers can be read without halting the processor.
6. The debugger can communicate with the power manager to power cores up or down, and to query their status.
7. Serial ports can raise an interrupt when a send/receive queue becomes full/empty.
8. The debug interrupt can be masked by running code. If the interrupt is asserted, then deasserted, and then asserted again the debug interrupt happens anyway. This mechanism can be used to eg. read/write memory with minimal interruption, making sure never to interrupt during a critical piece of code.
9. The debugger can non-intrusively sample a recent PC value from any running hart.
10. The Debug Module can include a serial interface for re-using the DTM interface as a generic communication interface.

## C.1 Serial Ports

The Debug Module may implement up to 8 serial ports. They support basic flow control and full duplex data transfer between a component and the debugger, essentially allowing the Debug Transport to be used to communicate with a debug monitor running on a hart, or more generally emulate devices which aren't present. All these uses require software support, and are not further specified here. Only the DMI side of the Debug Module serial registers are defined in this specification as the core side interface should look like a peripheral device.

Table C.1: Debug Module Debug Bus Registers

Address	Name	Page
0x34	Serial Control and Status	<a href="#">64</a>
0x35	Serial TX Data	<a href="#">65</a>
0x36	Serial RX Data	<a href="#">65</a>

### C.1.1 Serial Control and Status (`sercs`, at 0x34)

If `serialcount` is 0, this register is not present.

31	28	27	26	24	23	22	21	20	19	18
serialcount	0	serial	error7	valid7	full7	error6	valid6	full6		
4	1	3	1	1	1	1	1	1	1	
17	16	15	14	13	12	11	10	9		
error5	valid5	full5	error4	valid4	full4	error3	valid3	full3		
1	1	1	1	1	1	1	1	1		
8	7	6	5	4	3	2	1	0		
error2	valid2	full2	error1	valid1	full1	error0	valid0	full0		
1	1	1	1	1	1	1	1	1		

Field	Description	Access	Reset
<code>serialcount</code>	Number of supported serial ports.	R	Preset
<code>serial</code>	Select which serial port is accessed by <code>serrx</code> and <code>sertx</code> .	R/W	0
<code>error0</code>	1 when the debugger-to-core queue for serial port 0 has over or underflowed. This bit will remain set until it is reset by writing 1 to this bit.	R/W1C	0
<code>valid0</code>	1 when the core-to-debugger queue for serial port 0 is not empty.	R	0
Continued on next page			

full0	1 when the debugger-to-core queue for serial port 0 is full.	R	0
-------	--	---	---

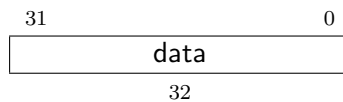
### C.1.2 Serial TX Data (sertx, at 0x35)

If serialcount is 0, this register is not present.

This register provides access to the write data queue of the serial port selected by serial in [sercs](#).

If the **error** bit is not set and the queue is not full, a write to this register adds the written data to the core-to-debugger queue. Otherwise the **error** bit is set and the write returns error.

A read to this register returns the last data written.



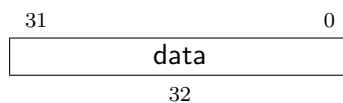
### C.1.3 Serial RX Data (serrx, at 0x36)

If serialcount is 0, this register is not present.

This register provides access to the read data queues of the serial port selected by serial in [sercs](#).

If the **error** bit is not set and the queue is not empty, a read from this register reads the oldest entry in the debugger-to-core queue, and removes that entry from the queue. Otherwise the **error** bit is set and the read returns error.

This entire register is read-only.







# Appendix D

## Change Log

Revision	Date	Author(s)	Description
6cab794	2017-10-12	Tim Newsome	Merge pull request #162 from riscv/impebreak
ecdaf0a	2017-10-11	Megan Wachs	Merge branch '0.13' into impebreak
83f9faf	2017-10-11	Tim Newsome	Clarify what <b>impebreak</b> does.
58fdc0e	2017-10-11	Megan Wachs	Merge pull request #164 from riscv/legend_on_fig
5f3933f	2017-10-11	Tim Newsome	Merge branch 'impebreak' of github.com:riscv/riscv-debug-spec into impebreak
78082b5	2017-10-11	Tim Newsome	Mention <b>impebreak</b> in Program Buffer description.
0378324	2017-10-11	mwachs5	Add legend and update some transitions on the Abstract Command State Machine diagram
a418a25	2017-10-11	Megan Wachs	Merge branch '0.13' into impebreak
d1c895a	2017-10-11	Megan Wachs	Merge pull request #161 from riscv/no_h_mode_fig
f9981e8	2017-10-11	Megan Wachs	Merge branch '0.13' into no_h_mode_fig
ae30fc1	2017-10-11	Megan Wachs	Merge pull request #163 from riscv/fix_build
fa2b600	2017-10-11	Megan Wachs	add missing period
0610630	2017-10-11	Megan Wachs	Just do simple hmode -i dmode replacement
16e11f3	2017-10-11	Tim Newsome	Remove hmode reference, to fix build.
84b9a6a	2017-10-11	Tim Newsome	Add <b>impebreak</b> , to support of implicit ebreak.
cc90b77	2017-10-11	mwachs5	Remove reference to 'H' mode from the figure
cc6a9de	2017-10-11	Megan Wachs	Change old reference to 'hmode' to 'dmode'
48f437b	2017-10-06	Megan Wachs	Merge pull request #159 from riscv/unsupported_access_size
24f0494	2017-10-06	Tim Newsome	Merge branch '0.13' into unsupported_access_size
812686d	2017-10-06	Tim Newsome	Merge pull request #157 from riscv/reset
486ecc6	2017-10-05	Tim Newsome	Refuse unsupported bus accesses.
c9cdf9e	2017-10-05	Tim Newsome	Merge branch '0.13' into reset
2be57af	2017-10-04	Tim Newsome	Merge pull request #128 from riscv/connector
d16b390	2017-10-04	Tim Newsome	Merge branch '0.13' into connector
d4118ab	2017-09-30	Tim Newsome	ndmreset can't reset logic required to access DM.
d569387	2017-09-29	Tim Newsome	Merge pull request #154 from riscv/nikhil
2a47bd1	2017-09-29	Tim Newsome	Merge branch '0.13' into nikhil
6269d9d	2017-09-29	Tim Newsome	Merge pull request #132 from riscv/progbuFSIZE

c6bd8d1	2017-09-29	Tim Newsome	and -j or
58c2441	2017-09-29	Tim Newsome	Mention <b>stepie</b> in Single Step
94c5f78	2017-09-29	Tim Newsome	Clarify ndmreset.
12810b4	2017-09-29	Tim Newsome	Clarify that sbaddress is physical.
5862fdf	2017-09-29	Tim Newsome	Unify M mode and mprv comment.
aea1bd5	2017-09-29	Tim Newsome	Define behavior when haltreq and resumereq are set
fe76d39	2017-09-28	Megan Wachs	Merge branch '0.13' into progbufsize
146b348	2017-09-28	Megan Wachs	remove superfluous 'an'
a5d16c4	2017-09-28	Megan Wachs	remove superfluous 'a'
052a8ab	2017-09-28	Tim Newsome	Clarify that a debugger can lose hart control.
cc52cff	2017-09-28	Tim Newsome	Add <b>dmerr</b> .
25685eb	2017-09-28	Tim Newsome	Explain that bus master or progbuf is required.
f75ee7d	2017-09-28	Tim Newsome	Clarify debugger can discover "almost" everything
71e6788	2017-09-27	Tim Newsome	Remove description of manual stepping.
9aea347	2017-09-27	Tim Newsome	Move Running/Single Step near Halting.
2090d9b	2017-09-27	Tim Newsome	data0 should be sbdata0 in this table.
5858cfe	2017-09-27	Tim Newsome	Clarify why <b>priv</b> exists.
bc3c2aa	2017-09-27	Tim Newsome	Mention where priv encoding comes from.
ef77cc4	2017-09-27	Tim Newsome	One more attempt to clarify DPC after single step.
80a288e	2017-09-27	Tim Newsome	Clarify instret not incrementing on ebreak.
4359b78	2017-09-27	Tim Newsome	Merge pull request #152 from riscv/nikhil
c163d22	2017-09-20	Tim Newsome	Remove ebreakh.
9971075	2017-09-20	Tim Newsome	Clarify we're talking about privilege
3fbe495	2017-09-20	Tim Newsome	Clarify that we're talking about *implementation*
3684854	2017-09-20	Tim Newsome	Use steps environment in sbdata0.
d4eda18	2017-09-20	Tim Newsome	Explain that only sbdata0 has side effects.
ae781c6	2017-09-20	Tim Newsome	Don't refer to internal system bus registers.
875922e	2017-09-20	Tim Newsome	Explain sbdata0 being stale a bit more.
cd44fd5	2017-09-20	Tim Newsome	Clarify autoread
194484b	2017-09-20	Tim Newsome	Clarify hawindow.
02f1aac	2017-09-20	Tim Newsome	Clarify that <b>dataaddr</b> is relative to <b>zero</b> .
0e9b6ae	2017-09-20	Tim Newsome	Clarify nonexistent vs unavailable.
b55ff41	2017-09-20	Tim Newsome	Fix devtreevalid.
4325ef8	2017-09-20	Megan Wachs	Merge branch '0.13' into progbufsize
2eccb86	2017-09-20	Tim Newsome	Explicitly state which registers are read-only.
4af505c	2017-09-20	Tim Newsome	Show section numbers for registers.
cbd5573	2017-09-20	Tim Newsome	Thank Nikhil
19c206f	2017-09-20	Tim Newsome	Clarify how to determine whether progbuf is RAM
0651f7d	2017-09-20	Tim Newsome	Explain what happens if ebreak is missing.
e889dae	2017-09-20	Tim Newsome	Move figure of states into its own section.
cff7b80	2017-09-20	Tim Newsome	Explain when <b>transfer</b> might be used.
6b2ee61	2017-09-20	Tim Newsome	Explain where <b>size</b> encoding came from.
900d8ab	2017-09-20	Tim Newsome	Merge pull request #145 from riscv/nikhil
c9f3b73	2017-09-14	Tim Newsome	Fix typo.
4b25400	2017-09-13	Tim Newsome	Mention dpc in CSRs abstract register numbers.
c3ee426	2017-09-13	Tim Newsome	Move abstract regno table closer to its reference.
111b9a3	2017-09-13	Tim Newsome	cycle -j operation
994afdc	2017-09-13	Tim Newsome	Account for multiple selected harts.

aa4a297	2017-09-13	Tim Newsome	Halt Control -> Run Control
e97c821	2017-09-13	Tim Newsome	continuous -> contiguous
97f73ff	2017-09-13	Tim Newsome	Clarify ndmreset behavior.
6078220	2017-09-13	Tim Newsome	Explain ndmreset
a3d4f30	2017-09-13	Tim Newsome	Describe 'halt region'
272b3d9	2017-09-13	Tim Newsome	Clarify accessing unimplemented DM DMI regs
3e91f1b	2017-09-13	Tim Newsome	Clarify either Prog Buf or Sys Bus Acc is required
e8a6145	2017-09-13	Tim Newsome	Clarify CSR access; remove serial port
ce20766	2017-09-13	Tim Newsome	Remove section referencing itself.
1195a61	2017-09-18	Tim Newsome	Generate constants to be unsigned for clang.
ba200ab	2017-08-18	Megan Wachs	Merge branch '0.13' into progbufsize
8967b0a	2017-08-16	Megan Wachs	Compressed instructions are c.foo, not foo.c
b5698a9	2017-08-16	Megan Wachs	clarify progbufsize description
d221bab	2017-08-16	Megan Wachs	Remove progbufsize enums from register description
d232d64	2017-08-16	Megan Wachs	Merge pull request #134 from riscv/sw-examples-cleanup
0498102	2017-08-16	Megan Wachs	appendix: Use standard assembly format for sw
6e20373	2017-08-15	Megan Wachs	Merge pull request #131 from riscv/devtree
50ea40c	2017-08-15	Megan Wachs	Merge branch '0.13' into devtree
4e51a25	2017-08-10	Tim Newsome	Merge pull request #130 from riscv/trigsign
4456d99	2017-08-09	Tim Newsome	Rename progsiz to progbufsize.
55d5b66	2017-08-09	Tim Newsome	Clarify that trigger comparisons are unsigned.
21e35ef	2017-08-09	Tim Newsome	Configuration String -> Device Tree
dc52f28	2017-08-03	Megan Wachs	Merge pull request #127 from riscv/cmdtype
f044f45	2017-08-02	Tim Newsome	Don't require a target to provide 25mA on VCC.
c883943	2017-08-02	Tim Newsome	Add table of Abstract Command Types
d6b8148	2017-08-02	Tim Newsome	Merge pull request #123 from riscv/lists
71f5cb2	2017-08-02	Tim Newsome	Merge branch '0.13' into lists
b83af70	2017-08-02	Megan Wachs	Merge pull request #125 from riscv/no_dmi_error
2a41bd8	2017-08-02	Megan Wachs	Merge branch '0.13' into no_dmi_error
9c73ce8	2017-08-02	Megan Wachs	Merge pull request #111 from riscv/dpc
a814400	2017-08-02	Tim Newsome	Merge branch '0.13' into dpc
8bdc5cd	2017-08-02	Tim Newsome	Merge pull request #126 from riscv/build
985a3df	2017-08-02	Tim Newsome	Fix and speed up build.
95b9108	2017-08-02	mwachs5	DTM: Clarify that there are no cases when DMI would actually return an error.
9c9e0c0	2017-08-02	mwachs5	SystemBus: No longer returns error. So DMI has no 'error' return code.
ae1e9e4	2017-07-28	Tim Newsome	Merge branch '0.13' into dpc
5ba18f9	2017-07-27	Tim Newsome	Fix more typos.
c6fef98	2017-07-26	Tim Newsome	Merge pull request #122 from riscv/version
dbc65bf	2017-07-26	Tim Newsome	Fix typos.
bba0ad9	2017-07-26	Tim Newsome	Tighten up introduction lists.
e22d5eb	2017-07-26	Tim Newsome	Add version constants for "not compatible".
c79038e	2017-07-26	Tim Newsome	Small clarification.
9df0411	2017-07-21	Tim Newsome	Incorporate review feedback.
d67419c	2017-07-21	Tim Newsome	Clarify dpc contents.
c562898	2017-07-11	Tim Newsome	Merge pull request #109 from riscv/ll

498cdf4	2017-07-11	Megan Wachs	Merge branch '0.13' into ll
0e707f1	2017-07-11	Megan Wachs	Merge pull request #105 from riscv/quick_access_errors
2d34f65	2017-07-11	Tim Newsome	Merge branch '0.13' into ll
a56831c	2017-07-11	Megan Wachs	Merge branch '0.13' into quick_access_errors
65d596e	2017-07-11	Megan Wachs	Merge pull request #106 from riscv/error_halt_resume
9f50c05	2017-07-11	Tim Newsome	Use LL instead of L for 64-bit constant suffix.
23fd24a	2017-07-10	Megan Wachs	Cleaning up whitespaces
102ba67	2017-07-10	Megan Wachs	Merge branch '0.13' into error_halt_resume
1720505	2017-07-10	Megan Wachs	Merge pull request #107 from riscv/csr_individuality
d67f6ef	2017-07-10	Megan Wachs	Merge branch '0.13' into csr_individuality
c1e61b0	2017-07-10	Megan Wachs	Merge pull request #108 from riscv/dcsr_causes
c5ab04c	2017-07-10	Megan Wachs	Update abstract_commands.xml
6e8cdf1	2017-07-10	Megan Wachs	Update abstract_commands.xml
cf6e3f2	2017-07-10	Megan Wachs	clarify DCSR.cause
79ffb9	2017-07-10	Megan Wachs	Clarify implications of CSR read, write, halt
013e191	2017-07-10	Megan Wachs	Clarify when you would get error halt/resume
231e457	2017-07-10	Megan Wachs	Quick Access error clarification
7c760b0	2017-07-03	Megan Wachs	Merge pull request #104 from riscv/serial_to_appendix
c54c2f2	2017-07-03	mwachs5	serial: add the XML file, not the TEX file
ac77477	2017-07-03	mwachs5	serial: Fix compile errors after moving serial port to appendix
6defcb8	2017-07-03	mwachs5	serial: Move serial ports out of main spec and into Future Work appendix
3541152	2017-07-03	Megan Wachs	Merge pull request #102 from riscv/remove_trace
a28f639	2017-06-30	mwachs5	remove trace dependencies from Makefile
52a122b	2017-06-30	mwachs5	remove trace section
d9e166b	2017-06-30	mwachs5	remove trace registers
7caf4e5	2017-06-30	mwachs5	remove trace appendix
aff0c16	2017-06-30	Megan Wachs	Merge pull request #82 from riscv/intdisable
4688988	2017-06-29	mwachs5	DCSR: define a 'stepie' bit which may be hard-wired to 0.
443b3fe	2017-06-29	mwachs5	Merge remote-tracking branch 'origin/0.13' into int-disable
497ed95	2017-06-29	Megan Wachs	Merge pull request #96 from riscv/jtagdtm_non_appendix
f1488c4	2017-06-29	Megan Wachs	Merge branch '0.13' into jtagdtm_non_appendix
920ec9a	2017-06-13	Megan Wachs	Merge pull request #95 from riscv/remove_spontaneous
67fa7b0	2017-06-13	Megan Wachs	Merge branch '0.13' into remove_spontaneous
78eb65e	2017-06-13	Megan Wachs	Merge pull request #94 from riscv/anynonexistent
d97b296	2017-06-13	Megan Wachs	Merge branch '0.13' into anynonexistent
61c6d30	2017-06-13	Megan Wachs	Merge pull request #93 from riscv/define-dret-again
389ee69	2017-06-13	Megan Wachs	Merge branch '0.13' into define-dret-again
421dcf2	2017-06-13	Megan Wachs	Merge pull request #97 from riscv/implementation_deets

9a0492c	2017-06-13	Megan Wachs	Add missing period and some other small text edits
13ccdbf	2017-06-13	Megan Wachs	fix typo in ProgBuf register macro
b01f989	2017-06-13	mwachs5	implementations: be a bit more concrete about the one example implementation we have.
a7b5f83	2017-06-13	mwachs5	jtagdtm: Move it out of the appendix as it is really part of the specification
87aceb0	2017-06-13	Megan Wachs	remove "spontaneous"
50b9950	2017-06-13	Megan Wachs	Forward reference for anynonexistent
adea3e2	2017-06-13	Megan Wachs	More clarifications on dret
1b8dd0e	2017-06-13	Megan Wachs	Define DRET instruction
b4f1f43	2017-06-08	Tim Newsome	Merge pull request #79 from riscv/cleanups
09c7f6e	2017-06-08	mwachs5	Merge remote-tracking branch 'origin/0.13' into cleanups
617da4c	2017-06-08	Megan Wachs	Update description of R/W1C
de2c56b	2017-06-08	Megan Wachs	Clarify that DCSR is also not updated on ebreak
efa615d	2017-06-07	Tim Newsome	Increase xdebugver field size to 4 bits. (#92)
a0e147a	2017-06-07	Tim Newsome	Address some review comments.
c1b3e54	2017-06-07	Megan Wachs	Merge pull request #91 from riscv/ndmreset
5c7c1bb	2017-06-07	Tim Newsome	Merge branch '0.13' into cleanups
72bb874	2017-06-06	Megan Wachs	Merge branch '0.13' into ndmreset
1fbbe6e	2017-06-06	Megan Wachs	Merge pull request #90 from riscv/dpc-clarifications
89ffe50	2017-06-06	mwachs5	NDMRESET: Clarify what it may and may not do
1932da0	2017-06-06	mwachs5	DPC: Clarifications on its meaning
03bcafe	2017-06-06	Megan Wachs	Merge pull request #89 from riscv/datacount
6470fdb	2017-06-06	mwachs5	ABSTRACTCS: Correct inconsistency on the number of data words.
1f4a1fe	2017-06-06	Megan Wachs	Merge pull request #88 from riscv/W0_corrections
3ca82b4	2017-06-06	Megan Wachs	More corrections for R vs R/W1C on SERCS
9705fb8	2017-06-06	Megan Wachs	Correct a bunch of W0 registers
1058690	2017-06-05	Megan Wachs	Merge branch '0.13' into intdisable
7531c41	2017-06-05	Megan Wachs	Merge pull request #80 from riscv/issue76
1347371	2017-06-05	Tim Newsome	Add intdisable to dcsr.
850bd87	2017-06-05	Megan Wachs	Merge branch '0.13' into issue76
43307eb	2017-06-05	Megan Wachs	Merge pull request #81 from riscv/issue63
989c60d	2017-06-05	Tim Newsome	Fix language. We can only halt harts, not cores.
517a08b	2017-06-05	Tim Newsome	Incorporate review feedback.
802be28	2017-06-05	Tim Newsome	Clarify/fix Quick Access example.
dbcaec8	2017-06-02	Tim Newsome	Merge branch '0.13' into cleanups
b8cc523	2017-06-02	Tim Newsome	Add included tex files as dependencies. (#78)
d0a5959	2017-06-02	Tim Newsome	Merge pull request #77 from riscv/pageno
15f864a	2017-06-01	Tim Newsome	Language cleanups, consistency and typo fixes.
4ecae86	2017-06-01	Tim Newsome	Add page numbers to list-of-register tables.
59b3e4a	2017-05-19	Megan Wachs	Setting up a Travis regression to check for build errors (#72)
124bf44	2017-05-17	mwachs5	Debug Module: CMDERR is Write-1-to clear, not R/W0
bb6c7f0	2017-05-17	mwachs5	SW Registers file should be XML, not TEX

d360358	2017-05-10	Megan Wachs	(Temporary Acct.)	Remove virtual register from core_registers.xml
bfc64fb	2017-05-10	Megan Wachs	(Temporary Acct.)	Add missing sw_registers.tex file
0512f5d	2017-05-06	mwachs5		Move virtual 'prv' register to a separate section to make it more clear it is not a real register.
6b3c9d7	2017-05-06	mwachs5		Clarify haltreq/resumereq/resumack
0a487eb	2017-04-26	mwachs5		jtag: Change specified JTAG pinout from Coretex to AVR, to provide for TRSTn option.
93cdfaf	2017-04-26	mwachs5		DM : Clarify that DATA/PROGBUF can't be written while busy.
ef98f23	2017-04-19	mwachs5		jtag: Make it clear that a NOP is really a NOP.
a6f8efa	2017-04-17	mwachs5		single_step: Exceptions count as the 'step' completion.
bf11e9e	2017-04-17	mwachs5		resumeack: fix some LaTeX cross references
4afa081	2017-04-11	mwachs5		halt/resumereq: Clarify what setting them to 0 or 1 does
297a39b	2017-04-06	mwachs5		fix chisel build
082c499	2017-04-06	mwachs5		Rename resumed to resumeack, and add more text about what these bits mean.
909d617	2017-04-06	mwachs5		Correct some cross references after removing all the multiply listed registers
dd09914	2017-04-06	mwachs5		Add 'resumedall' and 'resumedany' bits to avoid race condition on about to resume and just halted
feb88fc	2017-04-05	mwachs5		JTAG DTM: Clarify that leading bits are 0 for more than 5-bit IR
75b96ea	2017-04-04	mwachs5		use renamed dm_registers file
9f3ec7e	2017-04-04	mwachs5		debugger_implementation: remove some old TODO and commentary.
45dd5b5	2017-04-04	mwachs5		Don't list out every single DM register for those that are just indexed versions
b8b3aa2	2017-04-04	mwachs5		remove core-side register definitions from Debug Module. Rename dm1 to dm
d979a13	2017-04-04	mwachs5		remove core-side serial port specification, as these should look like implementation-specific devices with appropriate drivers.
b56870b	2017-04-04	mwachs5		Remove the wording about 'debug exception', as it is called breakpoint exception in the RISC-V Spec.
1e9347d	2017-04-03	mwachs5		Add description of hasel
0dda84d	2017-04-03	mwachs5		JTAG DTM: Clean up TAP register descriptions
82ccde5	2017-04-03	mwachs5		JTAG DTM: Add a hard DMI bit which cancels the outstanding DMI transaction
bd2a3d1	2017-04-03	mwachs5		remove preexec
02c733a	2017-04-03	mwachs5		remove preexec from Abstract State diagram.
1e271d6	2017-04-03	mwachs5		Update Debugger implementation for DMI register access, and fix tex compile issues.

155dda4	2017-04-03	mwachs5	Rewrite HW Implementation examples to describe a pure abstract command approach, and to not rely on harts executing every instruction which is fetched from the Debug Module
556c2be	2017-04-03	mwachs5	minor wording edits about RISC-V core registers
523c64a	2017-04-03	mwachs5	Edits to the Debug Module section.
b9a371f	2017-04-03	mwachs5	add missing trace.tex file.
58b2396	2017-04-03	mwachs5	Re-order the JTAG DTM Sections
a8827e2	2017-04-03	mwachs5	Edits to the System Overview.
c5417ce	2017-04-03	mwachs5	add more sections as separate files.
287d5c6	2017-04-03	mwachs5	moving more files to separate tex files.
9e873f4	2017-04-03	mwachs5	move trigger info into separate file.
2c89a86	2017-04-03	mwachs5	move risc-v core debug info into separate file.
e676491	2017-04-03	mwachs5	Move System Overview to separate file
03df6ee	2017-04-03	mwachs5	Move Debug Module description to a separate file.
5faa430	2017-04-03	mwachs5	add back in JTAG DTM in appendix
7b28b11	2017-04-03	mwachs5	Move jtag DTM to appendix. Move some text to commentary.
cc183ba	2017-04-03	mwachs5	move introduction to a separate file. Comment out reading order.
2c83830	2017-04-03	mwachs5	Merge remote-tracking branch 'origin/0.13' into 0.13
e3cf6ab	2017-04-03	Megan Wachs	Merge pull request #18 from riscv/intro_edits
60c5a1c	2017-04-03	Megan Wachs	Merge branch '0.13' into intro_edits
f727d14	2017-04-03	mwachs5	Use Chapters vs Sections. Needs reorganization.
815951d	2017-04-03	mwachs5	Formatting updates. Make this look more like the RISC-V specs. Need to use chapter vs. section
69ffaf8	2017-03-31	mwachs5	Move XML files into a subdirectory.
b276384	2017-03-31	mwachs5	Remove debug_rom.S
112bbac	2017-03-31	mwachs5	figures: reorganize the figures into directories.
2d05746	2017-03-31	Megan Wachs	Merge pull request #50 from riscv/add_license
1e5c068	2017-03-27	Megan Wachs	Add LICENSE
0e2d08a	2017-03-22	Megan Wachs	Merge pull request #47 from poweihuang17/0.13
fc17730	2017-03-22	Po-wei Huang	Change some halt mode into debug mode.
8ccf029	2017-03-22	Po-wei Huang	All halt mode changed to debug mode to synchronize with the priv spec.
f143d9e	2017-03-21	mwachs5	Correct duplicated progbuf register names
0797ec1	2017-03-17	mwachs5	autoexec: make autoexec bits match the number of data words there really are.
8e76d93	2017-03-17	mwachs5	dm1_registers: move a few more things around. Reduce abstract data words back to 12.
f8bf292	2017-03-17	mwachs5	dm1_registers: resolve some address conflicts and inconsistencies
a74dff9	2017-03-17	mwachs5	access_register: some small bit changes
2e6b0ca	2017-03-15	mwachs5	config string: Fix LaTeX compile errors.
f83260a	2017-03-10	mwachs5	Abstract Commands: clarify that 32-bit reads should always work. This allows reading MISA.

6f9347a	2017-03-10	mwachs5	Config String: change the Abstract Command to DMI registers. Allow the same registers to be used for unspecified identifier information.
4ea10ff	2017-03-10	mwachs5	abstract: Make autoexec apply to all data and progbuf words. Make a seperate register which is optional.
5008436	2017-03-10	mwachs5	abstract: Allow up to 16 progbuf and/or data words. Inform debugger about dscratch registers available for its use.
aaa13e5	2017-03-06	mwachs5	Command: use the name 'cmdtype' not 'type' to allow easier auto-generation of Scala code.
e9bb72c	2017-03-06	mwachs5	Hart Array: Add registers for hart array.
5d17a35	2017-03-06	mwachs5	DM: Move addresses around for better seperation of functionalities in HW
25ccaa8	2017-03-06	mwachs5	CONTROL: Rename control and status registers to ___CS for consistency and to accurately reflect their functionality.
45cf6c2	2017-03-06	mwachs5	Errors: fix up the bit assignments in SERSTATUS with the addition of error bit.
38cb5a0	2017-03-06	mwachs5	Errors: Make errors write-1-to-clear.
b436d77	2017-03-03	mwachs5	triggers: Clarify that matches are against virtual addresses.
793bb85	2017-03-03	mwachs5	triggers: Add suggested timings for best user experience.
2669866	2017-03-03	mwachs5	stoptime/stopcycle: Make their functionality match their name. Allow any reset value.
c85a1cf	2017-03-01	mwachs5	config_string: Simplify the Config String Address abstract command.
a303a6b	2017-03-02	Megan Wachs	Update README.md
1951ae3	2017-03-01	Megan Wachs	Merge pull request #35 from sifive/generate_chisel
2e2dc28	2017-03-01	Megan Wachs	Merge pull request #34 from sifive/serial_addr
c087c34	2017-03-01	mwachs5	Merge remote-tracking branch 'origin/0.13' into generate_chisel
92a4923	2017-03-01	mwachs5	serial: tweak addresses.
b09f460	2017-03-01	mwachs5	serial: tweak addresses.
6477837	2017-03-01	mwachs5	chisel: tweaks to class names.
be83e3e	2017-02-28	Tim Newsome	Clarify stoptime, stopcycle.
7f94662	2017-02-27	mwachs5	Merge remote-tracking branch 'origin/0.13' into generate_chisel
c17c17c	2017-02-27	Tim Newsome	Abstract command that returns config string addr.
096dfbc	2017-02-27	Tim Newsome	Acknowledge Alex.
c0253ab	2017-02-24	Tim Newsome	Explain tdata1 type a bit more.
e43ac2e	2017-02-24	Tim Newsome	Clarify how to enumerate triggers again.
c6e3e20	2017-02-23	Tim Newsome	Revert previous commit.
ef770bf	2017-02-23	Tim Newsome	mcontrol and icount mask tdata2, not tdata1.
27806f2	2017-02-23	mwachs5	rename 'type' to 'cmdtype' purely so my auto-generation scripts work.
e46798d	2017-02-22	mwachs5	Add Abstract Commands to automatic chisel
b3bb939	2017-02-21	mwachs5	Generate Chisel headers as well for Debug Module.



3d5b6f6	2017-02-22	Tim Newsome	Merge pull request #31 from sifive/abstract_command_types
c9db98c	2017-02-22	Tim Newsome	Simplify description of op statuses.
bda39cc	2017-02-22	mwachs5	Add explicit type field to Abstract Command.
34ff1d8	2017-02-22	Tim Newsome	Merge pull request #30 from sifive/-more_ibuf_progbuf
f83a1ca	2017-02-22	mwachs5	Finish up replacement of ibuf->progbuf
ddde0a2	2017-02-22	Tim Newsome	Merge pull request #28 from sifive/inst_supply_vs_progbuf
9666e51	2017-02-22	mwachs5	IBUF->PROGBUF
5308ecd	2017-02-22	mwachs5	Remove last references to "Instruction Supply"
f6ebde9	2017-02-22	Tim Newsome	Move authentication to a serial protocol.
0f079c8	2017-02-22	Tim Newsome	Reserve bit for per-hart reset.
f2c93ac	2017-02-22	Tim Newsome	Clarify that dmactive resets authentication.
59154ac	2017-02-22	Tim Newsome	Merge pull request #27 from asb/clarify_reset
f5e7b1c	2017-02-22	Alex Bradbury	Clarify that the halt state of all harts is maintained through reset
3dfe8fd	2017-02-22	Tim Newsome	More Debug Mode -> Halt Mode.
d29fc1f	2017-02-22	Tim Newsome	Debug Mode -> Halt Mode
55d6030	2017-02-21	Tim Newsome	Generate debug_defines.h as part of normal make
b0e6a7f	2017-02-21	Tim Newsome	Minor clarifications.
0f9885c	2017-02-20	Tim Newsome	Various clarifications.
e443ab9	2017-02-15	Tim Newsome	Merge pull request #25 from sifive/ctrl_status
3b08e90	2017-02-15	Tim Newsome	Merge pull request #24 from sifive/sm_diagram_resumereq
0802d5a	2017-02-15	mwachs5	Use consistent 'Control and Status' naming for CS registers.
5acc7d	2017-02-15	Tim Newsome	Change all the "other" JTAG IRs to just reserved.
bcbd7da	2017-02-15	mwachs5	sm_diagram: Show using resumereq bit to resume.
18f6e55	2017-02-14	Tim Newsome	Introduce resumereq command, similar to haltreq.
fb40538	2017-02-14	Tim Newsome	Merge pull request #22 from sifive/sb_errors
4b62c40	2017-02-14	mwachs5	SystemBus: Clean up some formatting and error specification notes.
0f346e4	2017-02-14	Tim Newsome	Merge pull request #21 from sifive/sm_for_quick_access
bc97723	2017-02-14	mwachs5	quick-access: Update SM Diagram for Quick Access
d27066e	2017-02-14	Tim Newsome	Clarify haltreq bit.
6f8ec43	2017-02-14	Tim Newsome	Always generate long constants when required.
c6ac6bc	2017-02-13	Tim Newsome	Include field descriptions in C header file.
b849213	2017-02-13	Tim Newsome	Fix the build.
c82c62e	2017-02-12	Tim Newsome	Merge pull request #20 from sifive/jtag_ir_minimum
1cf8033	2017-02-12	mwachs5	jtag: More clarifications
6203bd6	2017-02-12	Megan Wachs	Update requirements- W GPRs Required
f2b43a7	2017-02-12	Megan Wachs	Remove double 'the'
2c64ef1	2017-02-12	Megan Wachs	Remove comma
f84abce	2017-02-12	Megan Wachs	Whitespace edits and address come comments
7246b44	2017-02-12	Tim Newsome	Merge pull request #19 from sifive/jtag_dtm_edits
23c2648	2017-02-11	mwachs5	jtag-dtm: ask for clarification on TAP sharing.

7020d23	2017-02-11	mwachs5	jtag_dtm: Clarifications, DBUS- $\mathcal{I}$ DMI
292d49c	2017-02-11	Megan Wachs	fix indentation
55ef8d6	2017-02-11	Tim Newsome	Merge pull request #17 from sifive/prog_buffer_size
b879b86	2017-02-11	Megan Wachs	Add missing period
bbe0521	2017-02-11	mwachs5	Make comments on program buffer size match the address map.
4ceaa37	2017-02-11	mwachs5	Flesh out and edit the introduction/background Add a description of use cases this spec has in mind, and what it doesn't cover.
cbf89d6	2017-02-11	Tim Newsome	Rewrite Quick Access.
9115db1	2017-02-10	Tim Newsome	Merge pull request #16 from sifive/reduce_prog_buffer_size
170bff1	2017-02-10	Megan Wachs	Allow size 4 for the program buffer
9d46077	2017-02-10	Tim Newsome	Merge pull request #15 from sifive/dmactive
c911e6e	2017-02-10	Tim Newsome	Clarify use of dmactive.
2ca296f	2017-02-09	Tim Newsome	Reserve command register space for custom use.
e49666e	2017-02-09	Tim Newsome	Clarify hart index change per Megan's comments.
84865e9	2017-02-09	Tim Newsome	Add header prefix for abstract commands.
2434f4f	2017-02-09	Tim Newsome	Select harts by index instead of hart ID.
7bf112a	2017-02-09	Tim Newsome	Generate correct headers for $\mathcal{I}$ 32-bit registers.
7f0f09a	2017-02-08	Tim Newsome	Reset dbus status to "failure" to avoid confusion.
7b1803e	2017-02-08	Tim Newsome	Merge pull request #13 from sifive/arg0_clarification
8b1c6f0	2017-02-08	Megan Wachs	Fix line wrap issue
345c33f	2017-02-08	Megan Wachs	Call out "arg0" specifically.
9f080f5	2017-02-08	Megan Wachs	Clarify "arguments" to commands
259badd	2017-02-08	Tim Newsome	Make haltsum/halt registers mandatory.
eb0f1d3	2017-02-07	Tim Newsome	Allow for early abstract command failures.
bb49bd1	2017-02-07	Tim Newsome	Clarify error handling a little.
3fc0a97	2017-02-07	Tim Newsome	Explain when abstract data regs may be clobbered.
c37167e	2017-02-07	Tim Newsome	Fix old language in description of halt registers.
6943c96	2017-02-07	Tim Newsome	Generate more useful C header files from reg defs
d7a8045	2017-02-06	Tim Newsome	Merge pull request #11 from sifive/sm_diagram
8bef40e	2017-02-05	mwachs5	Merge remote-tracking branch 'origin/0.13' into sm_diagram
98639df	2017-02-05	mwachs5	Include the SM Diagram as a figure. Also some minor capitalization fixes.
a95e4c3	2017-02-05	mwachs5	Update State Machine diagram to show uncertainty of halt bit during auto halt/resume.
ba76744	2017-02-05	Tim Newsome	Combine loabits and hiabits.
02b1d92	2017-02-05	Tim Newsome	DMI can get away with just 6 address bits.
35d6e33	2017-02-05	mwachs5	Update State machine diagram to show BUSY without HALTED
f511b05	2017-02-04	Tim Newsome	Clarify command busy bit.
a8e5ae7	2017-02-03	mwachs5	Merge remote-tracking branch 'origin/0.13' into sm_diagram
d0f8961	2017-02-03	mwachs5	Update figures
e18a68d	2017-02-03	Tim Newsome	Clarify prehalt/postresume failure.
ac3e2a9	2017-02-02	Tim Newsome	Clarify abstract command failure behavior.

ce4baee	2017-02-02	Tim Newsome	Add Quick Access section.
0490377	2017-02-02	Tim Newsome	Add prehalt and postresume to reg command.
67515bd	2017-02-02	Tim Newsome	Deal with a few minor TODOs.
96456fc	2017-02-02	Tim Newsome	Turn register names into links.
317cd98	2017-02-02	Tim Newsome	Explain what register access is required.
f3ad2f2	2017-02-01	Tim Newsome	Revert Plain Exception implementation to be simple
a0ad281	2017-02-01	Tim Newsome	execb -i, preexec, execa -i, postexec
1d4a2c3	2017-02-01	Tim Newsome	Limit Program Buffer sizes to 0, 1, 8.
cc40815	2017-02-01	Tim Newsome	Incorporate Po-wei's feedback.
c8b45d6	2017-02-01	Tim Newsome	Clarify how all autoexec bits work.
dbb1deb	2017-02-01	Tim Newsome	Remove stale TODO.
c5f8f59	2017-02-01	Tim Newsome	Explain why cmderr inhibits starting new commands.
5c69194	2017-02-01	Tim Newsome	Fix editing error.
50f7c48	2017-02-01	Tim Newsome	Remove empty hart info register.
781c68e	2017-02-01	Megan Wachs	Update README.md
f46b32e	2017-02-01	mwachs5	Add a diagram of Abstract Command flow.
633bd63	2017-02-01	Tim Newsome	Move Reading Order into About This Document
51ec4d1	2017-02-01	Tim Newsome	Add reading order section.
03d20ad	2017-02-01	Tim Newsome	autoexec0 applies to data0, not inst0.
c302353	2017-01-31	Tim Newsome	Don't rely on hart fetching instructions once.
2558c25	2017-01-31	Tim Newsome	Change how exceptions in Halt Mode are handled.
a36ddce	2017-01-31	Tim Newsome	Add size to abstract register command.
64de458	2017-01-31	Tim Newsome	Detail bus master reads.
c08486f	2017-01-31	Megan Wachs	reset: Add some comments (#5)
1558049	2017-01-30	Tim Newsome	Automate Change Log.
51525a4	2017-01-29	Tim Newsome	Update System Overview
7d39ac0	2017-01-29	Tim Newsome	Update Supported Features.
9e7cbea	2017-01-29	Tim Newsome	Update RISC-V Core section.
515188d	2017-01-29	Tim Newsome	Update Hardware Implementations section.
4b19ed8	2017-01-29	mwachs5	system.bus: be consistent and always call it 'System Bus'. Even if some dislike the name, we should be consistent and clear in the spec.
9cceef3d	2017-01-29	Tim Newsome	Fleshed out some debugger implementation.
04b9176	2017-01-28	Tim Newsome	Rename debug exception to breakpoint exception.
5ac4ea1	2017-01-27	Tim Newsome	WIP on big update on instruction supply.
2d9c3e2	2017-01-27	Tim Newsome	Reorganize dm registers.
de50ba8	2017-01-27	Tim Newsome	Abstract command support is already addressed.
27cb0da	2017-01-26	Tim Newsome	Merge pull request #4 from sifive/access_renames
5085046	2017-01-26	mwachs5	Rename registers and fields like 'access' that were confusingly the same name.
10bbf6f	2017-01-26	Tim Newsome	Fix #2: DM address space table
a05c582	2017-01-26	Tim Newsome	Add debugger inspection as a feature.
4062681	2017-01-24	Tim Newsome	Add publish target.
5c8bb83	2017-01-24	Tim Newsome	Clarify use of data registers.
1504da6	2017-01-24	Tim Newsome	Replace manual date with automatic git hash/date.
997f2a0	2017-01-23	Tim Newsome	Deal with unsupported abstract commands.
cb6f2b8	2017-01-23	Tim Newsome	Renumber registers to prevent duplicates.
8b4db96	2017-01-23	Tim Newsome	Don't print out addresses if they're not provided.

b00cd21	2017-01-23	Tim Newsome	Add an abstract command.
675b556	2017-01-23	Tim Newsome	Reorganize DM bits into functional group regs.
5fc7512	2017-01-23	Tim Newsome	Remove bits 33:32 from sbdata[23].
ceb5d66	2017-01-20	Tim Newsome	Starting point for a comprehensive spec