

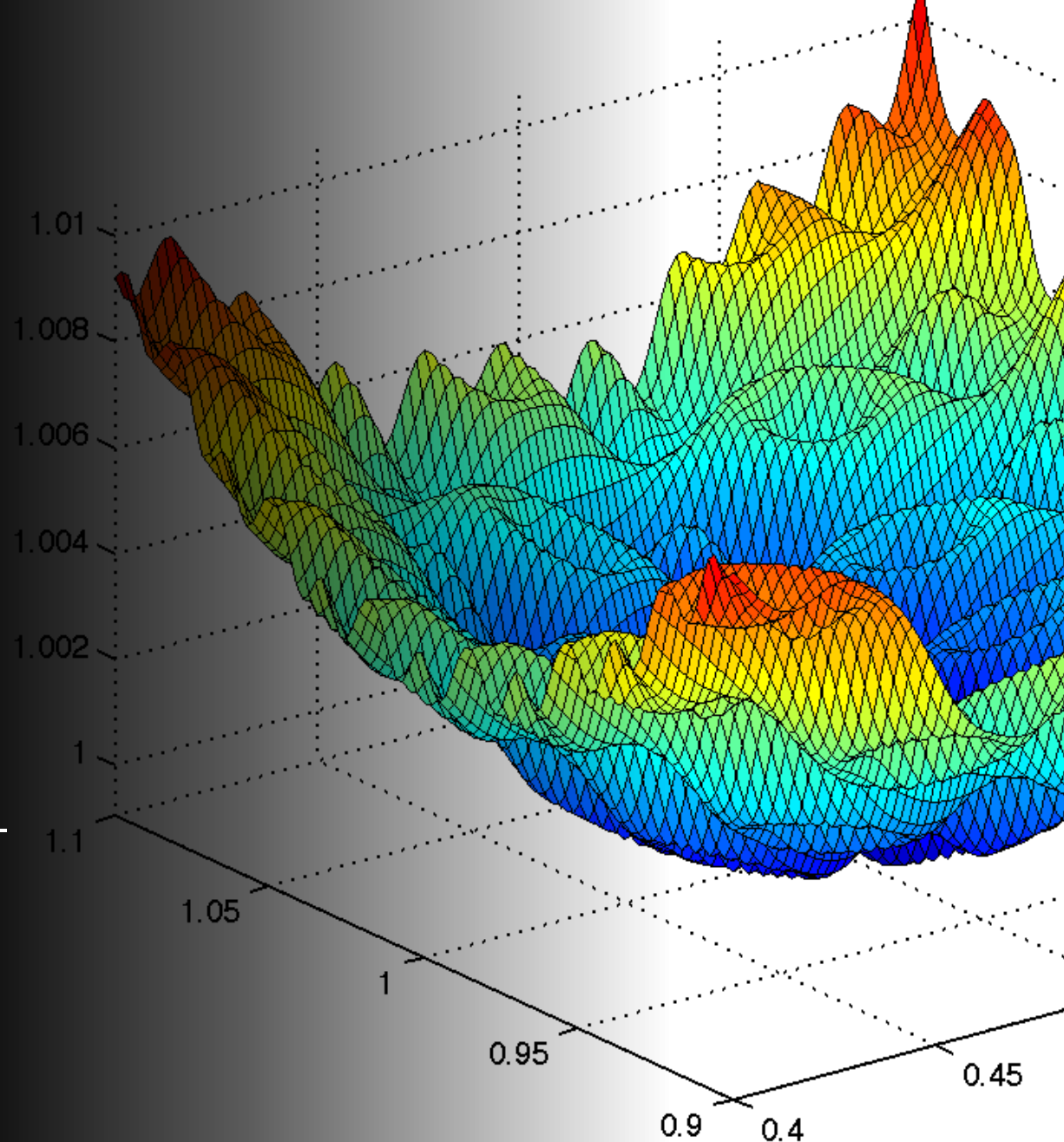
# CS 5/7320

## Artificial Intelligence

### Local Search

AIMA Chapters 4.1 & 4.2

Slides by Michael Hahsler  
based on slides by Svetlana Lazepnik  
with figures from the AIMA textbook.

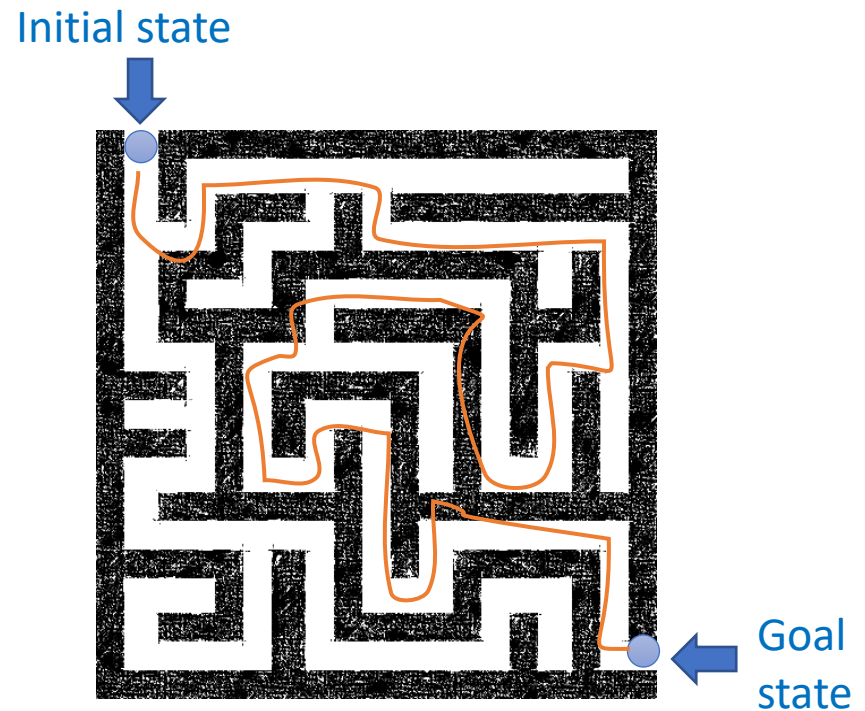


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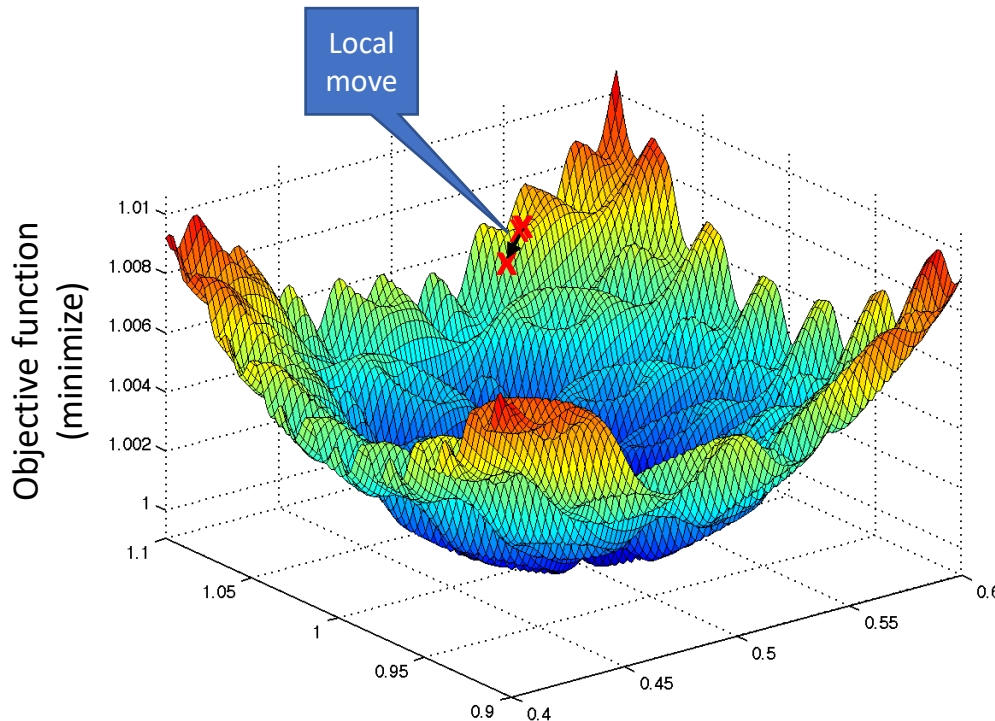
# Recap: Uninformed Search/informed search

Tries to find the  
**best path**  
from a  
**given initial state**  
to a  
**given goal state.**

- Typically searches a large portion of the search space (needs time and memory).
- Often comes with optimality guarantees.



# Local search algorithms



- **Goal:** A fast and memory-efficient way to find a **good state**. That means to search only a small portion of the search space.
- We need an **objective function** over the states that defines what “good” means → **optimization problem**.

## Idea:

1. Formulate a solution as the state.
2. Improve the solution by moving to a neighboring state (i.e., search locally). This is fast and needs little memory (no search tree).

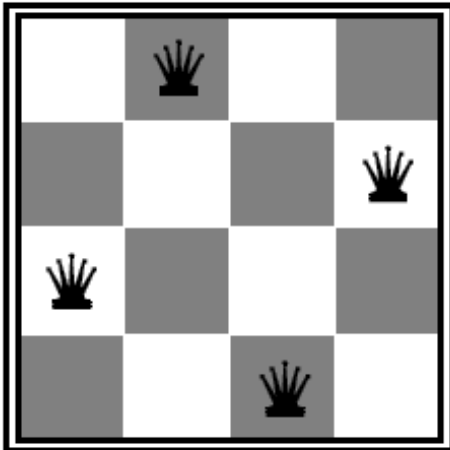
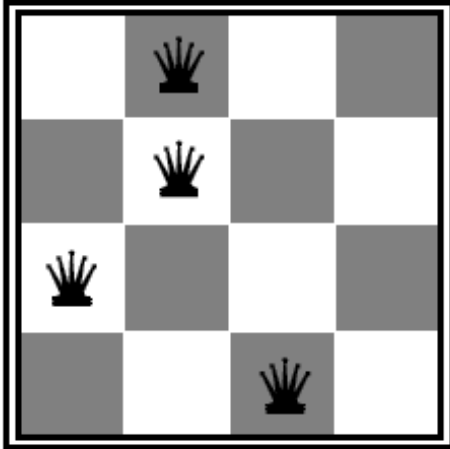
# Local search algorithms

## Difference to search from the previous chapter:

- a) Goal state is unknown and needs to be identified.
- b) Often no explicit initial state + path to goal and path cost are not important.
- c) No search tree structure. Just stores the current state.

## Use in AI

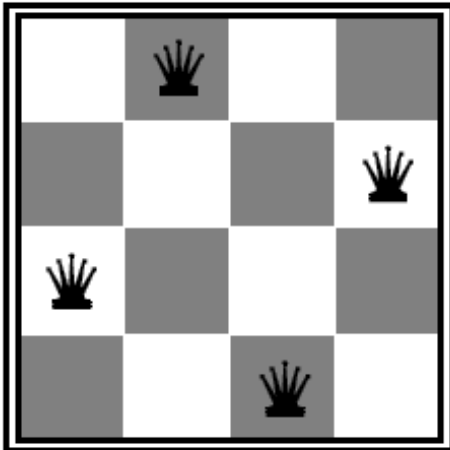
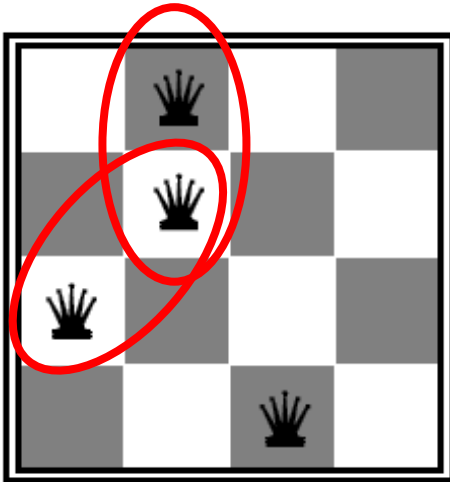
- **Utility-based agent:** Use utility as the objective function and always move to higher utility states. A greedy method used for complicated/large state spaces or online search.
- **Goal-based agent:** Identify a good goal state with a good objective function value before planning the path to that state.
- **General optimization:** Use for effective heuristic search in large or continuous spaces (with an infinite state space). E.g., learn neural networks.



# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal.
- **State space:** All possible  $n$ -queen configurations. **How many are there?**
- What is a possible **objective function**?

2 conflicts



0 conflicts

# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal
- **State space:** all possible  $n$ -queen configurations:  
4-queens problem:  $\binom{16}{4} = 1820$
- What is a possible **objective function**?  
**Minimize the number of pairwise conflicts**  
**Note:** this can be seen as a heuristic used in informed search.






## Example: Traveling salesman problem

- **Goal:** Find the shortest tour connecting a given set of cities
- **State space:** all possible tours (states are not individual cities!)
- **Objective function:** minimize the length of the tour

**Note:** We have solved a different problem with uninformed/informed search! Each city was defined as a state and the path was the solution.



A motorcyclist wearing a red and white racing suit and a helmet is climbing a steep, rocky hill. The rider is leaning forward, and the motorcycle is angled upwards. The background shows a rugged, rocky landscape with some sparse vegetation. The text "Hill-Climbing Search aka Greedy Local Search" is overlaid on the image in a large, white, sans-serif font. Below the title, there is a smaller line of text in a white, sans-serif font. In the bottom right corner, the website "MotorCycleUSA.com" is visible in a small, white, sans-serif font.

# Hill-Climbing Search aka Greedy Local Search

**Idea:** keep a single “current” state and try to find better neighboring states.



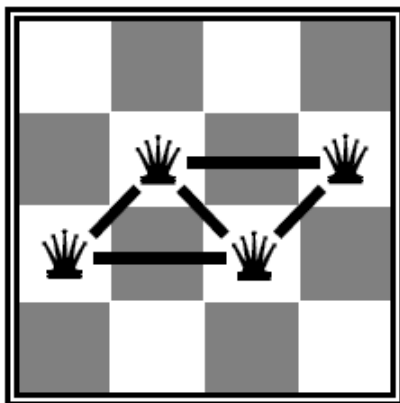
# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal.
- **State space:** all possible  $n$ -queen configurations. We can restrict the state space: Only one queen per column.
- **Objective function:** minimize the number of pairwise conflicts.

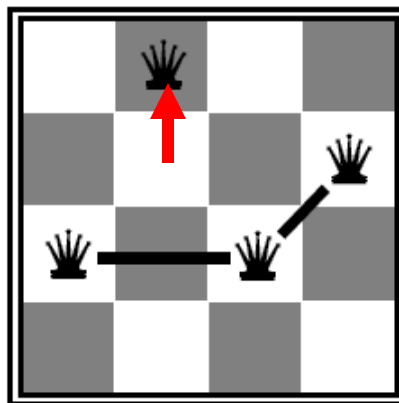
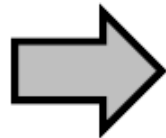
State space is reduced from 1820 to  $4^4 = 256$

What is a possible local improvement strategy?

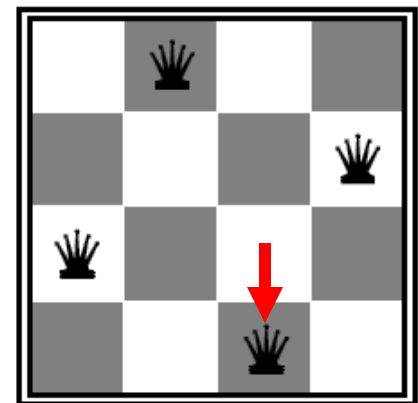
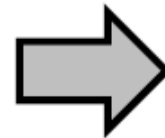
- Move one queen within its column to reduce conflicts



$h = 5$



$h = 2$



$h = 0$

# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal.
- **State space:** all possible  $n$ -queen configurations. We can restrict the state space: Only one queen per column.
- **Objective function:** minimize the number of pairwise conflicts.

What is a possible local improvement strategy?

- Move one queen within its column to reduce conflicts

18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	♙	13	16	13	16
♙	14	17	15	♙	14	16	16
17	♙	16	18	15	♙	15	♙
18	14	♙	15	15	14	♙	16
14	14	13	17	12	14	12	18

$h = 17$

best local improvement has  $h = 12$

Note that there are many options, and we have to choose one!

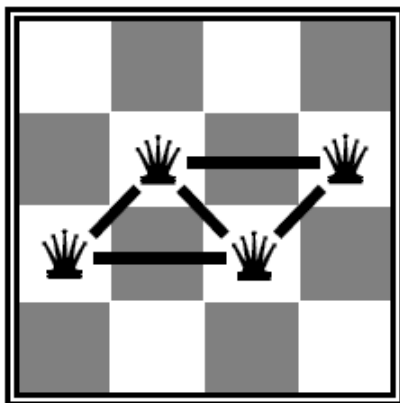
# Example: $n$ -queens problem

Optimization problem: find the best arrangement  $a$

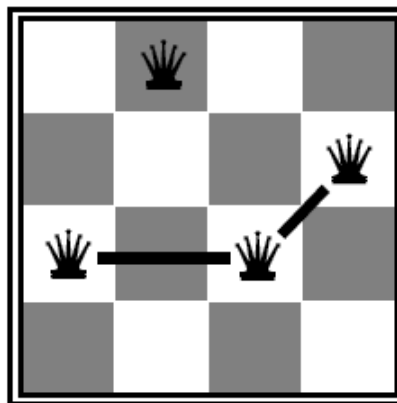
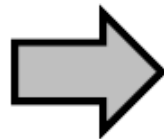
$$a^* = \operatorname{argmin}_a \operatorname{conflicts}(a)$$

s.t.  $a$  has one queen per column

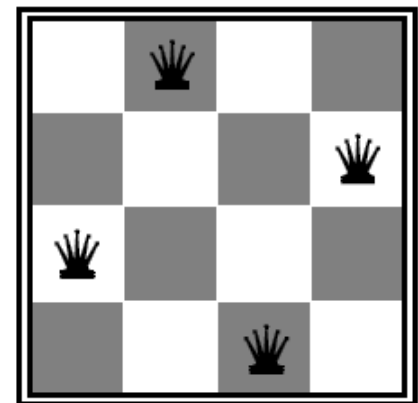
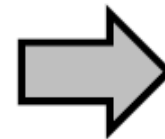
This makes the problem a lot easier. The search space is now:  $4! = 24$



$h = 5$



$h = 2$



$h = 0$

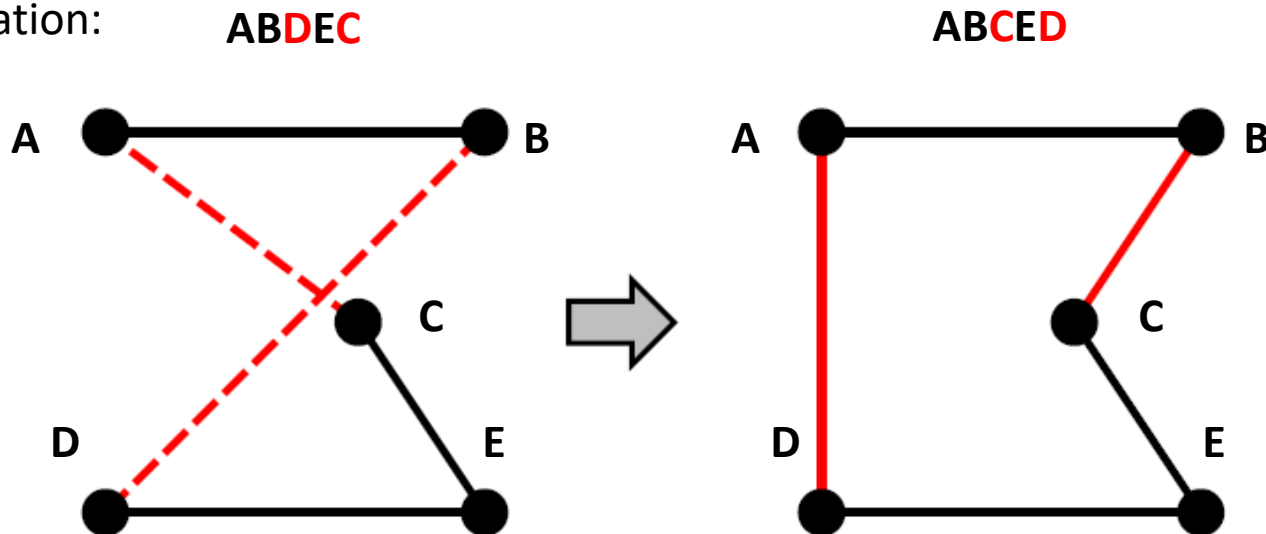
# Example: Traveling Salesman Problem

- **Goal:** Find the shortest tour connecting  $n$  cities
- **State space:** all possible tours
- **Objective function:** length of tour

What's a possible local improvement strategy?

- Start with any complete tour, perform pairwise exchanges.

Permutation:





# Example: Traveling Salesman Problem

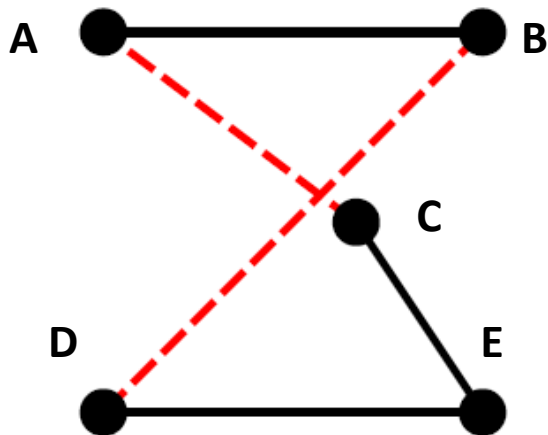
Optimization problem: Find the best tour  $\pi$

$$\pi^* = \operatorname{argmin}_{\pi} \operatorname{tourLength}(\pi)$$

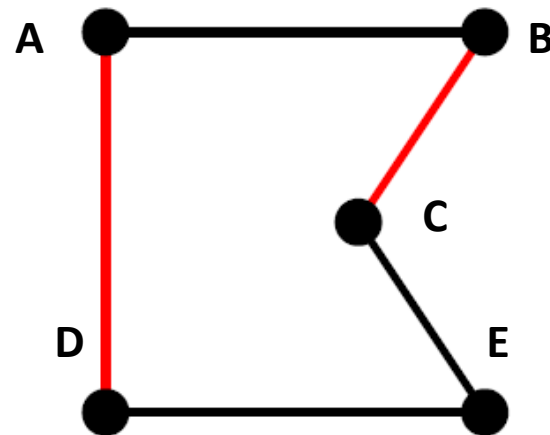
s.t.  $\pi$  is a valid permutation (i.e., sub-tour elimination)

Permutation:

**ABDEC**



**ABCED**



# Hill-climbing search (= Greedy local search)

```
function HILL-CLIMBING(problem) returns a state that is a local maximum
  current  $\leftarrow$  problem.INITIAL
  while true do
    neighbor  $\leftarrow$  a highest-valued successor state of current
    if VALUE(neighbor)  $\leq$  VALUE(current) then return current
    current  $\leftarrow$  neighbor
```

Variants:

- **Steepest-ascend hill climbing**
  - Check all possible successors and choose the highest-valued successors.
- **Stochastic hill climbing**
  - choose randomly among all uphill moves, or
  - generate randomly one new successor at a time until a better one is found = first-choice hill climbing – the most popular variant, this is what people often mean when they say “stochastic hill climbing”

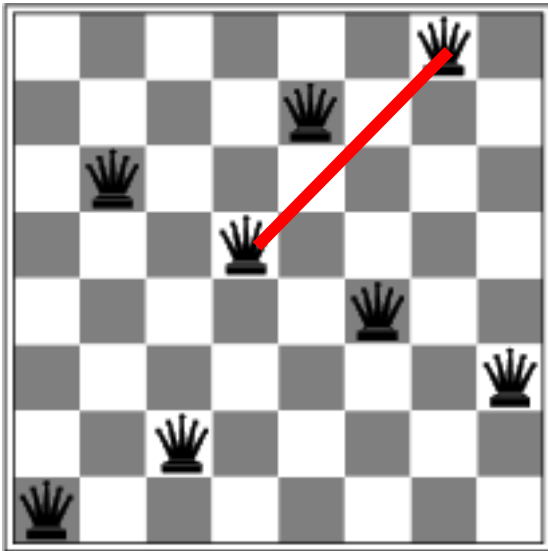
**Random-restart hill climbing** – to deal with local optima

# Hill-climbing search

Hill-climbing search is similar to greedy best-first search without backtracking.

Is it complete/optimal?

- No – can get stuck in local optima

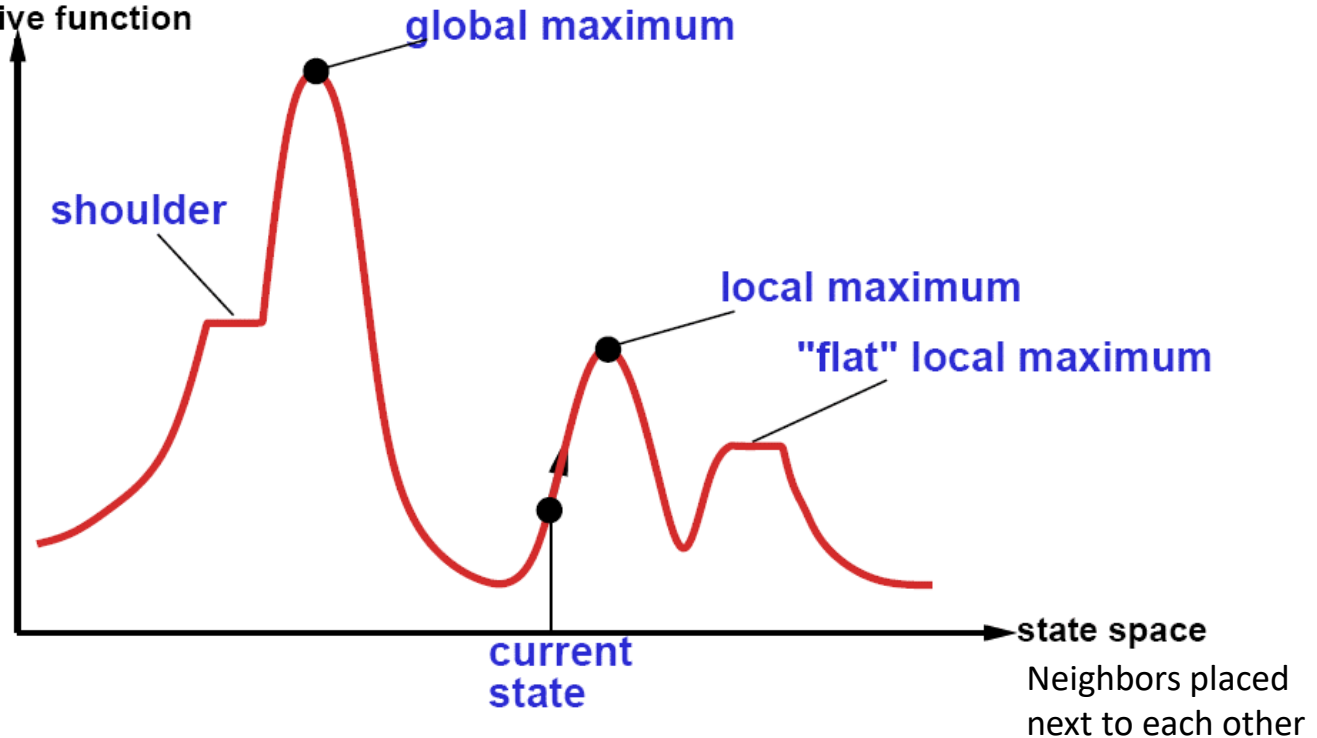


Example: local optimum for the 8-queens problem. No single queen can be moved within its column to improve the objective function.

$$h = 1$$

# The state space “landscape”

Maximization problem  
objective function



How to escape local maxima?

→ Random restart hill-climbing can help.

What about “shoulders” (called “ridges” in higher dimensional space)?

What about “plateaus”?

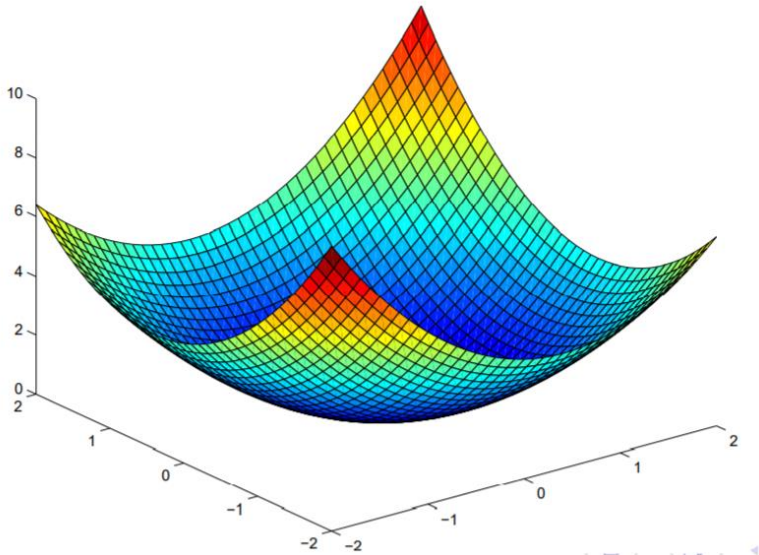
→ Allow sideways moves.



# Non-convex/convex Optimization Problems

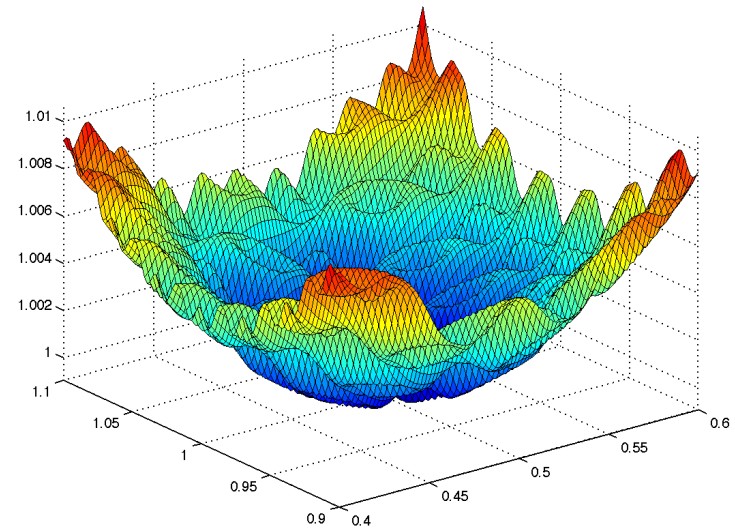
Minimization problem

Convex Problem



One global optimum +  
smooth function → easy

Non-convex Problem



Many local optima → hard

Many discrete optimization  
problems are like this.

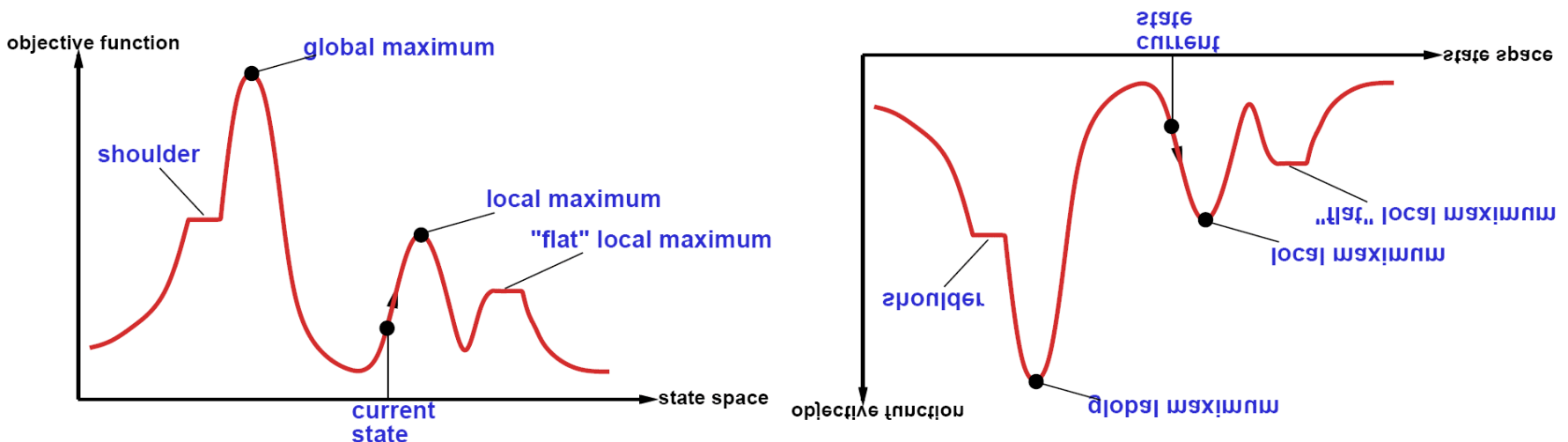
# A Note on Minimization vs. Maximization

- The name **hill climbing** implies **maximizing a function**.
- Optimizers like to state problems as **minimization problems** (e.g., gradient descent).
- Both types of problems are equivalent:

$$\max(f(x))$$



$$\min(-f(x))$$



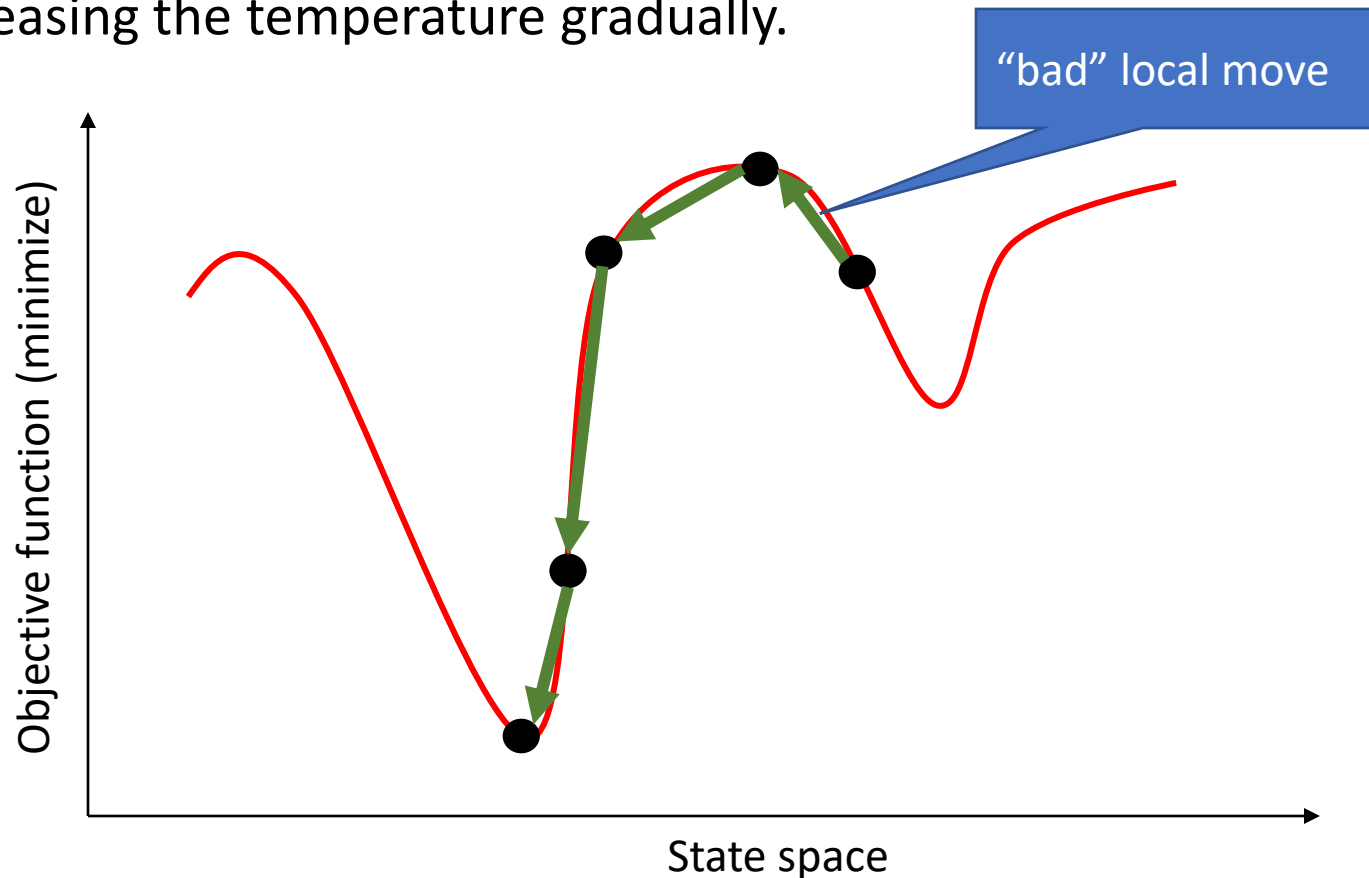
A glowing yellow-orange metal bar is being processed in a dark industrial machine. The bar is long and rectangular, with a bright, intense glow along its length, suggesting it is being heated or melted. The machine components are dark and metallic, with various rollers and guides visible. The background is dark, emphasizing the bright light of the metal.

# Simulated Annealing

Use heat to escape local optima...

# Simulated annealing

- **Idea:** First-choice stochastic hill climbing + escape local minima by **allowing some “bad” moves** but gradually decrease their frequency.
- Inspired by the process of tempering or hardening metals by decreasing the temperature gradually.





# Simulated annealing

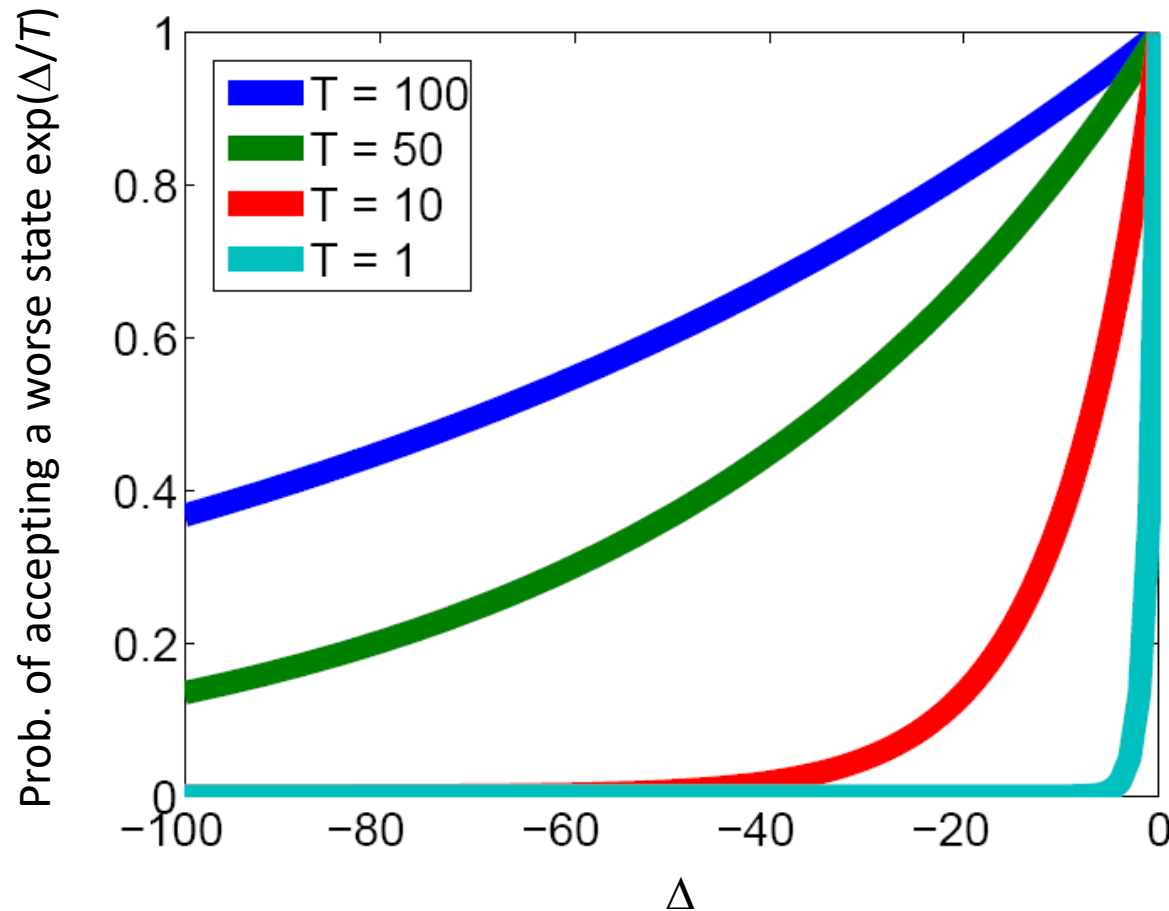
- **Idea:** First-choice stochastic hill climbing + escape local minima by allowing some “bad” moves but gradually decrease their frequency.
- Inspired by the process of tempering or hardening metals by decreasing the temperature gradually.
- The probability of accepting “bad” moves follows **an annealing schedule** that reduces the temperature  $T$  over time  $t$ .

```
function SIMULATED-ANNEALING(problem, schedule) returns a solution state  
  current  $\leftarrow$  problem.INITIAL  
  for  $t = 1$  to  $\infty$  do  
     $T \leftarrow$  schedule( $t$ )  
    if  $T = 0$  then return current  
    next  $\leftarrow$  a randomly selected successor of current  
     $\Delta E \leftarrow$  VALUE(next) – Value(current)  
    if  $\Delta E < 0$  then current  $\leftarrow$  next  
    else current  $\leftarrow$  next only with probability  $e^{-\Delta E/T}$ 
```

Always do good moves

Uses the Metropolis acceptance criterion to accept “bad” moves

# Effect of temperature



The lower the temperature, the less likely the algorithm will accept a worse state.

# Cooling Schedule

The cooling schedule is very important.

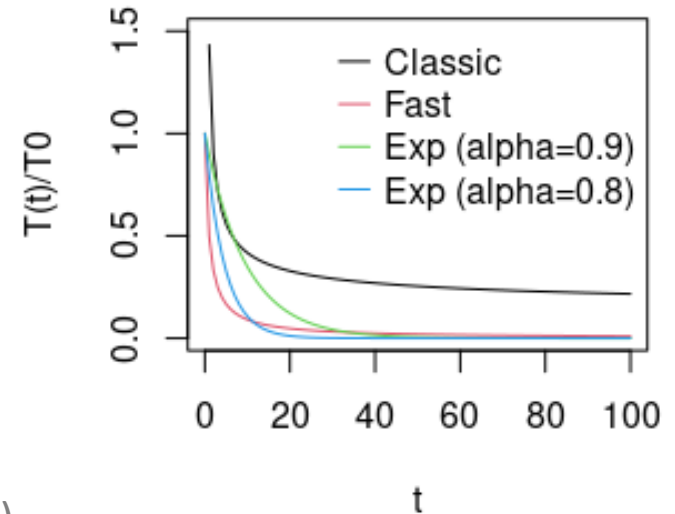
Popular schedules for the temperature at time  $t$ :

- **Classic simulated annealing:**  $T_t = T_0 \frac{1}{\log(1+t)}$
- **Fast simulated annealing** (Szy and Hartley; 1987)

$$T_t = T_0 \frac{1}{1+t}$$

- **Exponential cooling** (Kirkpatrick, Gelatt and Vecchi; 1983)

$$T_t = T_0 \alpha^t \quad \text{for } 0.8 < \alpha < 1$$



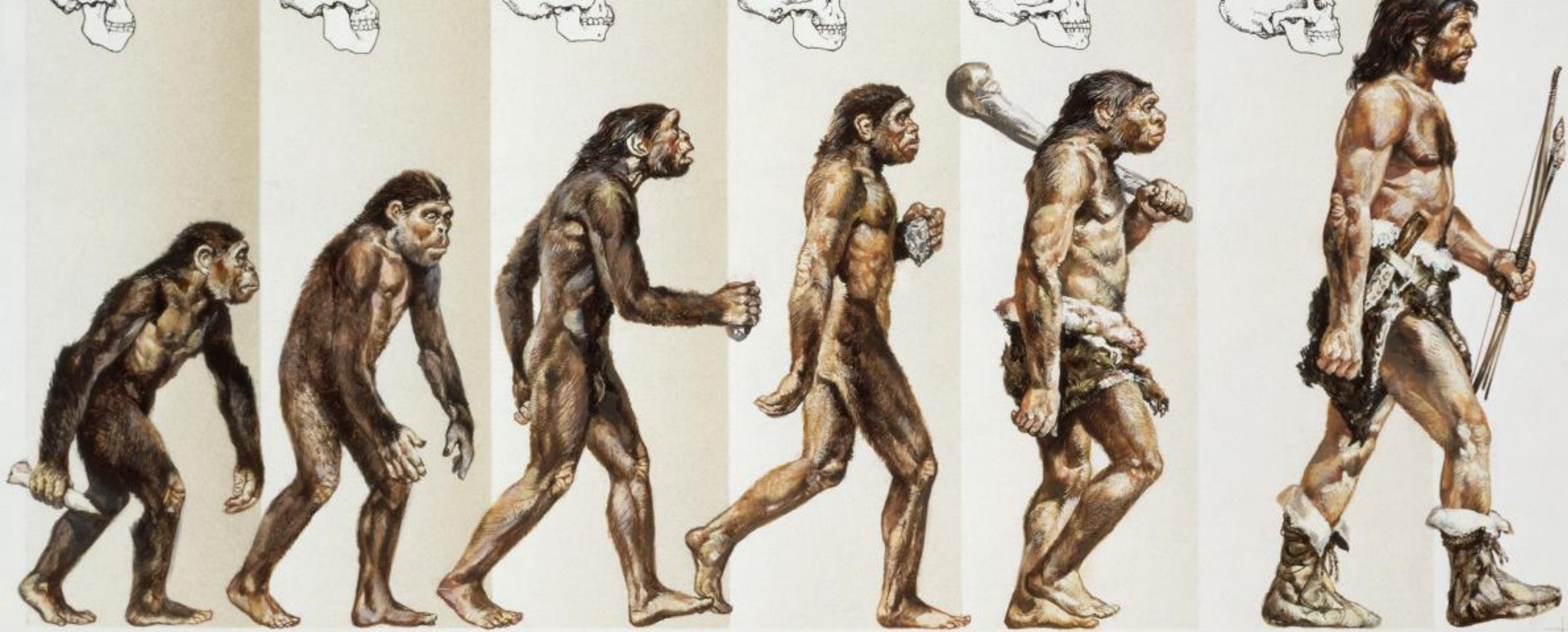
Notes:

- The best schedule is typically determined by trial-and-error.
- Choose  $T_0$  to provide a high probability that any move will be accepted at time  $t = 0$ .
- $T_t$  will not be come 0 but very small. Stop when  $T < \epsilon$  ( $\epsilon$  is a very small constant).

# Simulated annealing search

- **Guarantee:** If temperature decreases slowly enough, then simulated annealing search will find a global optimum with probability approaching one.
- However:
  - This usually takes impractically long.
  - The more downhill steps you need to escape a local optimum, the less likely you are to make all of them in a row.
- The related **Markov Chain Monte Carlo (MCMC)** method is a general family of randomized algorithms for exploring complicated state spaces.





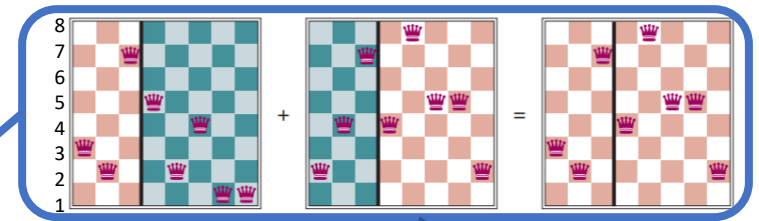
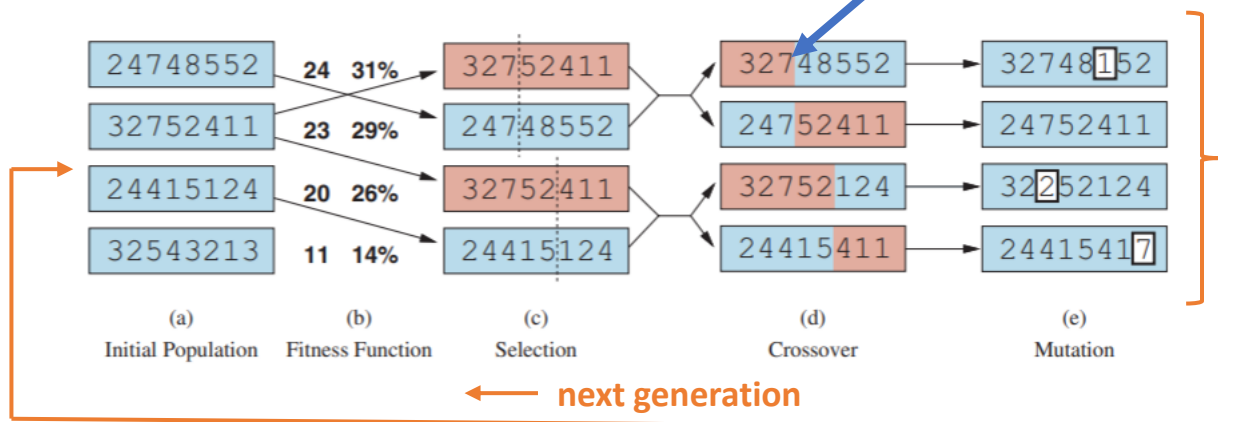
# Evolutionary Algorithms

A Population-based Metaheuristics

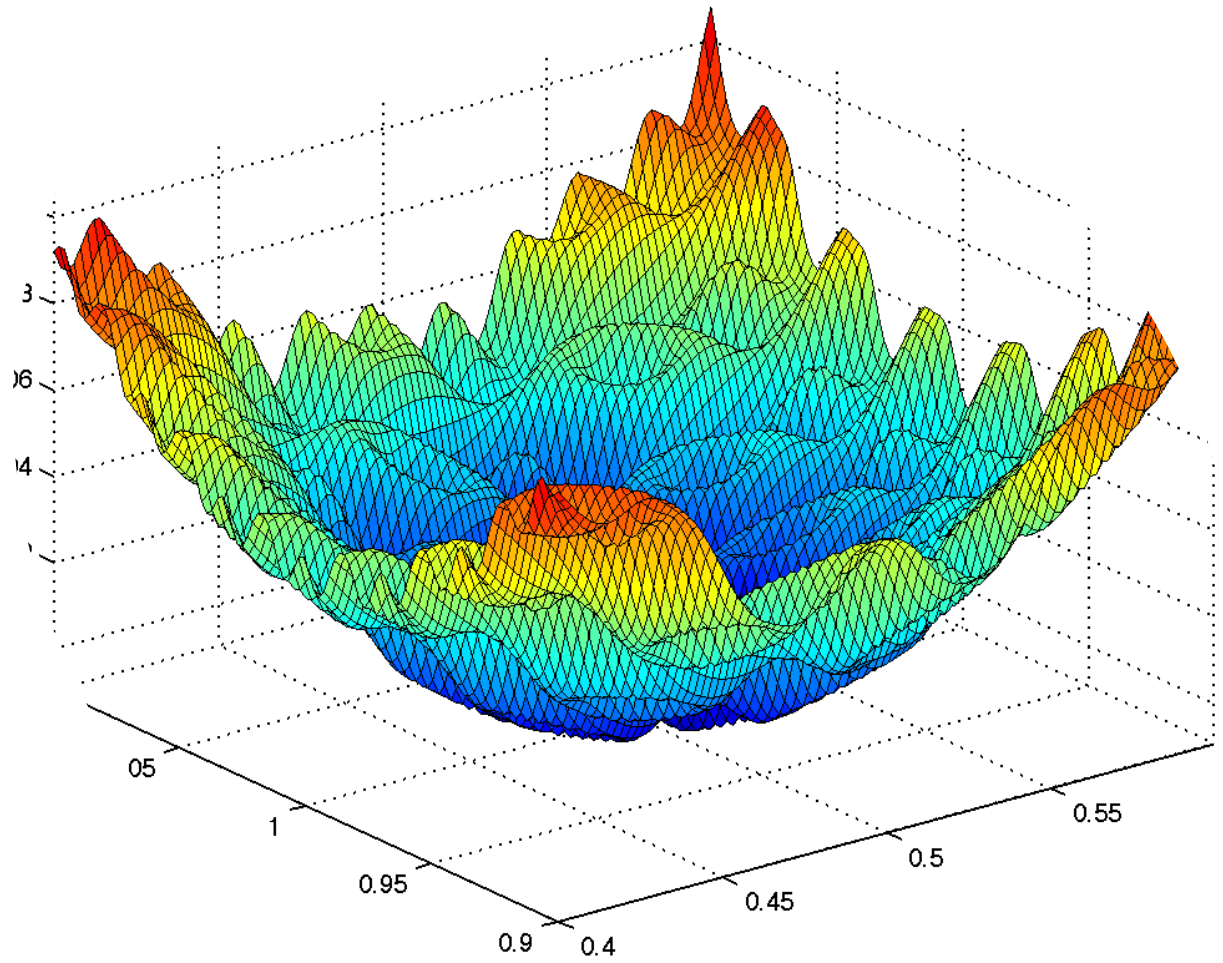
# Evolutionary algorithms / Genetic Algorithms

- A metaheuristic for **population**-based optimization.
- Uses mechanisms inspired by biological evolution (genetics):
  - Reproduction: Random selection with probability based on a **fitness** function.
  - Random recombination (crossover)
  - Random mutation
  - Repeated for many generations

- Example: 8-queens problem



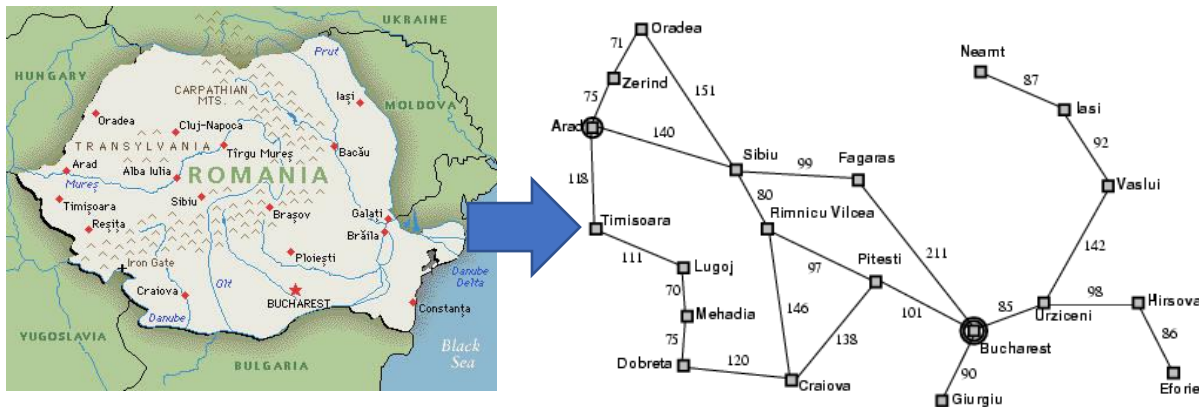
Individual = state  
Encoding: row of  
the queen in each  
column



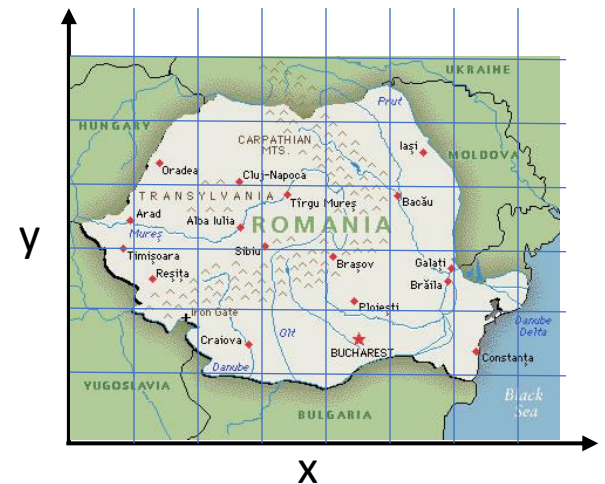
# Search in Continuous Spaces

# Discretization of the continuous space

- Use atomic states to create a graph

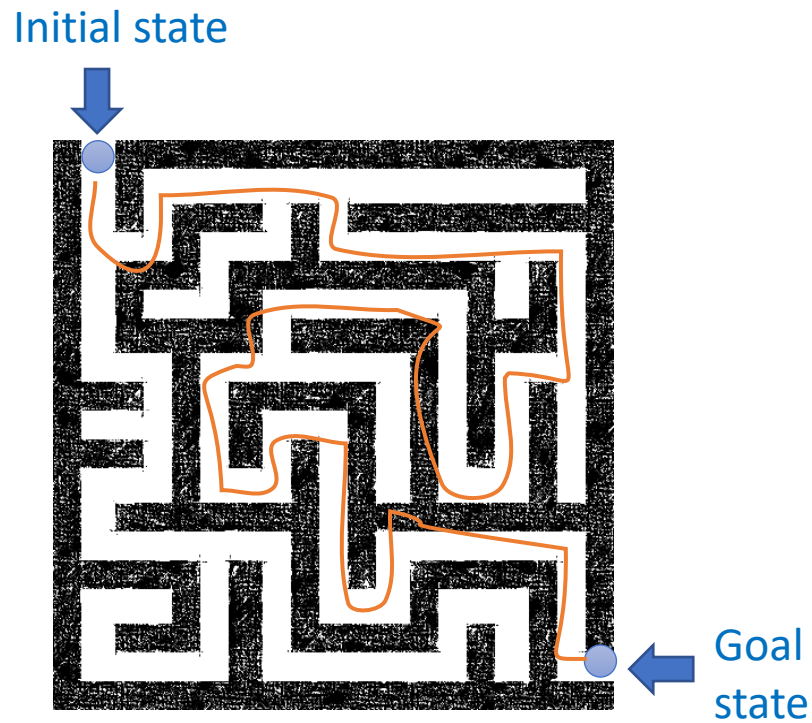


- Use a grid with spacing of size  $\delta$   
Note: You probably need a way finer grid!



# Discretization of the continuous space

How did we discretize this space?





# Search in continuous spaces: Gradient

Maximize  $f(\mathbf{x}) = f(x_1, x_2, \dots, x_k)$

Gradient at point  $\mathbf{x}$ :  $\nabla f(\mathbf{x}) = \left( \frac{\partial f(\mathbf{x})}{\partial x_1}, \frac{\partial f(\mathbf{x})}{\partial x_2}, \dots, \frac{\partial f(\mathbf{x})}{\partial x_k} \right)$

Find maximum by solving:  $\nabla f(\mathbf{x}) = 0$

- **Gradient descent (= Steepest-ascent hill climbing for minimization)** with step size  $\alpha$

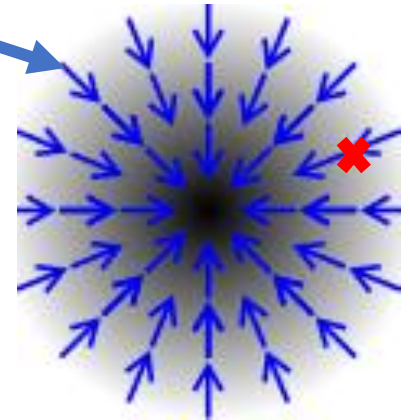
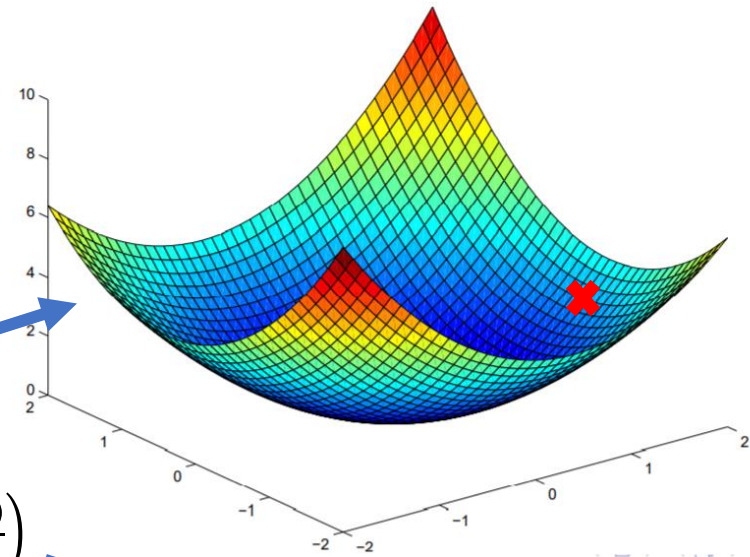
$$\mathbf{x} \leftarrow \mathbf{x} - \alpha \nabla f(\mathbf{x})$$

- **Newton-Raphson method**

uses the inverse of the Hessian matrix of the second derivative  $H_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j}$  for the step size  $\alpha$

$$\mathbf{x} \leftarrow \mathbf{x} - \mathbf{H}_f^{-1}(\mathbf{x}) \nabla f(\mathbf{x})$$

May get stuck in a local optimum if the search space is non-convex! Use simulated annealing.



# Search in continuous spaces: Empirical Gradient Methods

- What if the mathematical formulation of the objective function is not known?
  - We may have objective values at fixed points, called the training data.
  - In this case we can use **empirical gradient search**. This is related to steepest ascend hill climbing in the discretized state space.
- We will talk more about search in continuous spaces with loss functions using gradient descend when we talk about **parameter learning for machine learning**.