

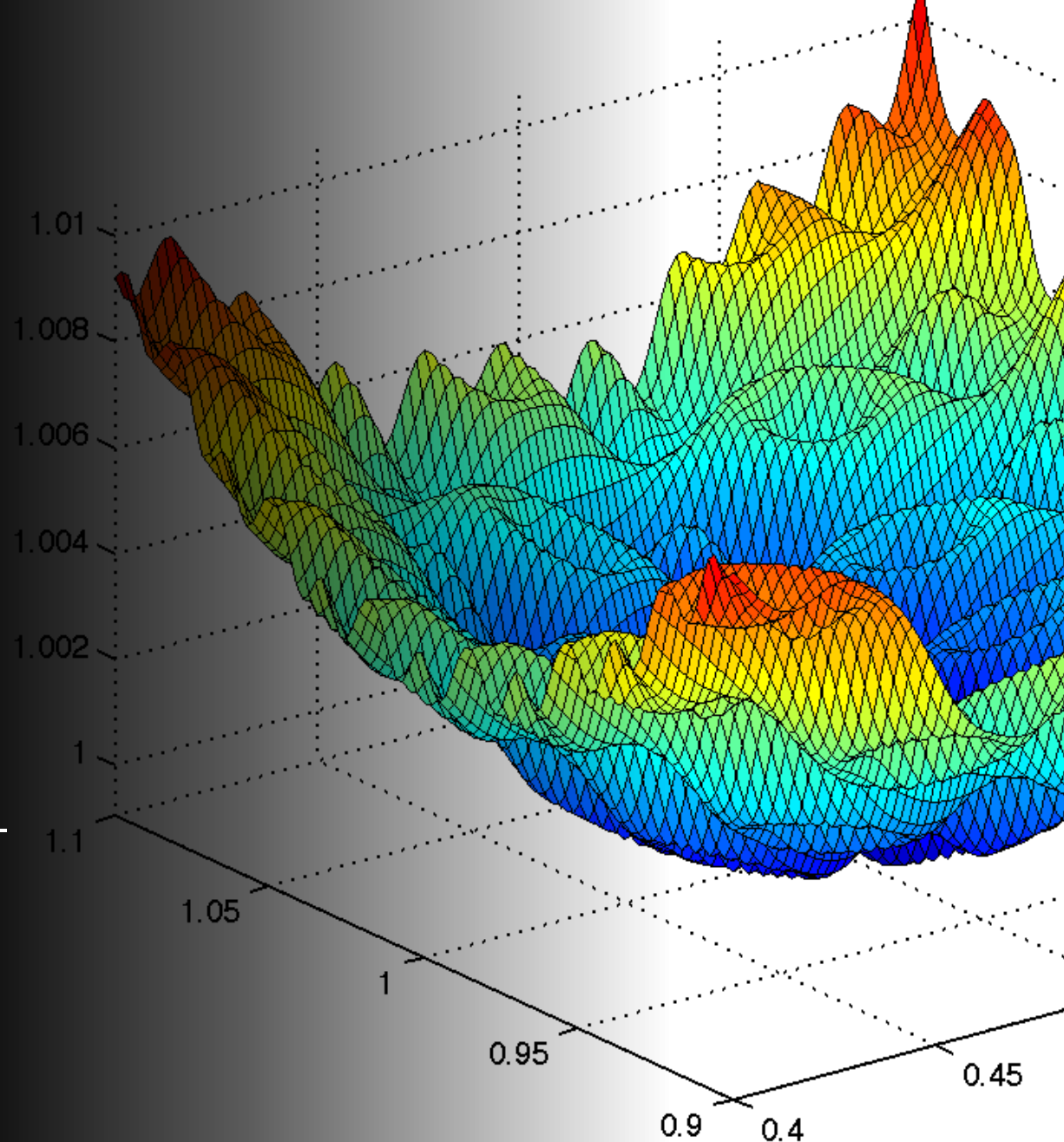
# CS 5/7320

## Artificial Intelligence

### Local Search

AIMA Chapters 4.1 & 4.2

Slides by Michael Hahsler  
based on slides by Svetlana Lazepnik  
with figures from the AIMA textbook.

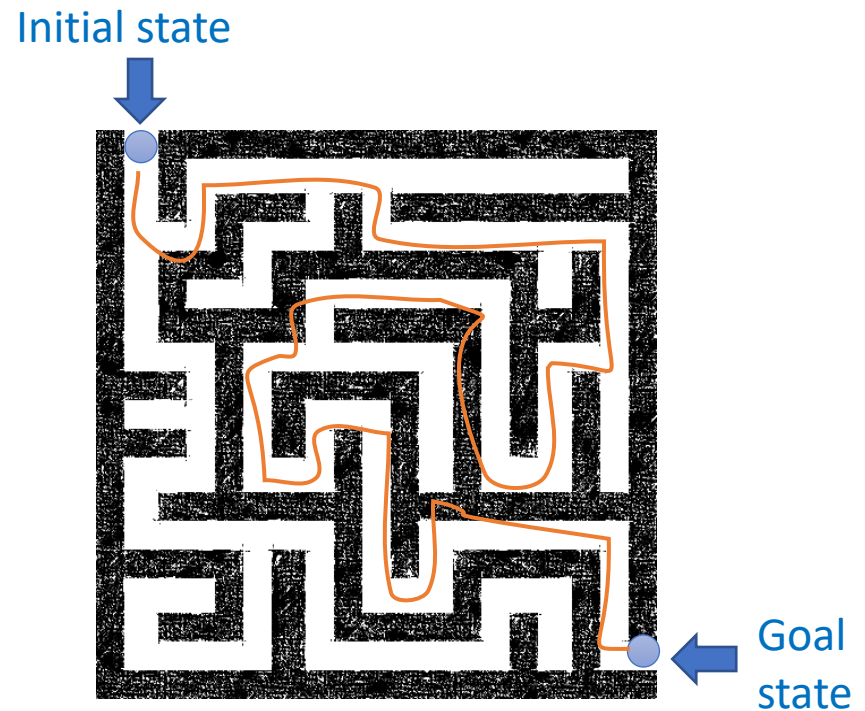


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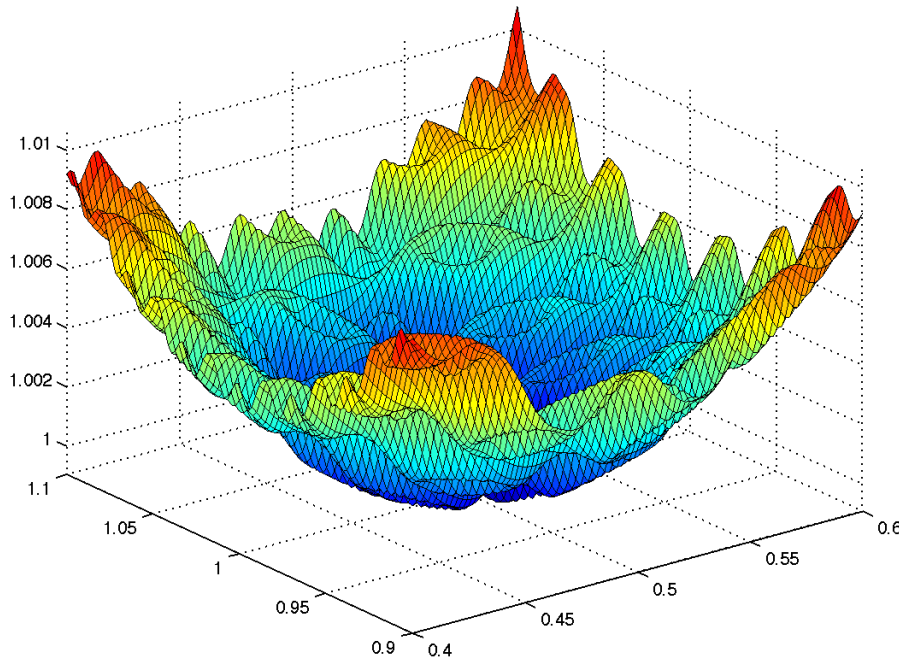
# Recap: Uninformed Search/informed search

Tries to find the  
**best path**  
from a  
**given initial state**  
to a  
**given goal state.**

- Typically searches a large portion of the search space (needs time and memory).
- Often comes with optimality guarantees.



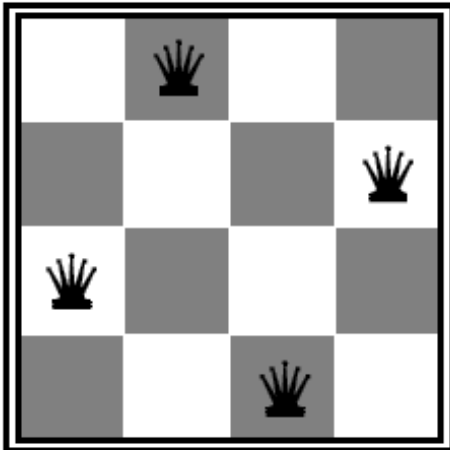
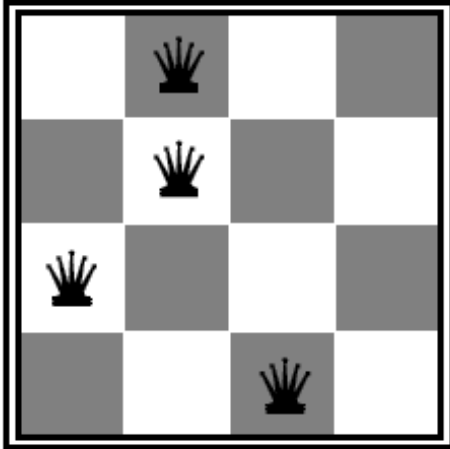
# Local search algorithms



- **Goal:** Search only a small portion of the search space to identify a **good state**.
  - a) Often no initial state
  - b) Path to solution and path cost are often not important
- We need an **objective function** over the states that defines what “good” means → **optimization problem**.
- **Idea:** search neighboring states is fast and needs little memory.

## Example applications in AI:

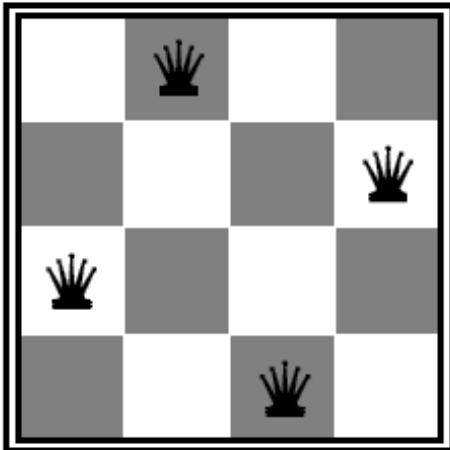
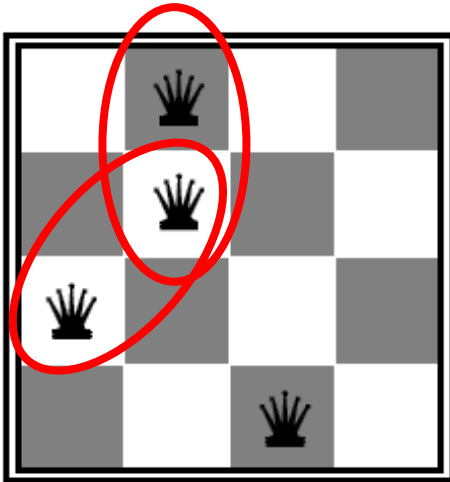
- Use for effective search in continuous space (with an infinite state space).
- Identify a good goal state (objective function might be utility).
- Each state might encode a complete plan (a solution).



# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal
- **State space:** all possible  $n$ -queen configurations. **How many are there?**
- What is the **objective function**?

2 conflicts



0 conflicts

# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal
- **State space:** all possible  $n$ -queen configurations  
 $16 \times 15 \times 14 \times 13 = 43,680$
- What is the **objective function**?  
**Number of pairwise conflicts**






## Example: Traveling salesman problem

- **Goal:** Find the shortest tour connecting a given set of cities
- **State space:** all possible tours (states are not individual cities!)
- **Objective function:** length of tour

**Note:** We have solved a different problem with uninformed/informed search! Each city was defined as a state and the path was the solution.





# Hill-Climbing Search a.k.a Greedy Local Search

**Idea:** keep a single “current” state and try to find better neighboring states.

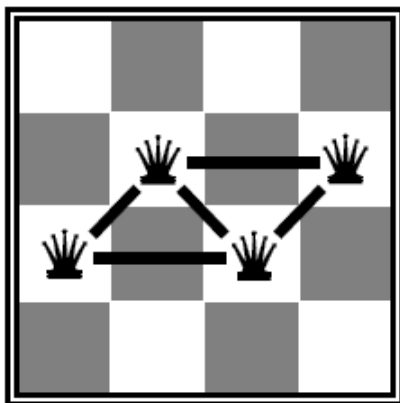
# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal.
- **State space:** all possible  $n$ -queen configurations. We can restrict the state space: Only one queen per column.
- **Objective function:** minimize the number of pairwise conflicts.

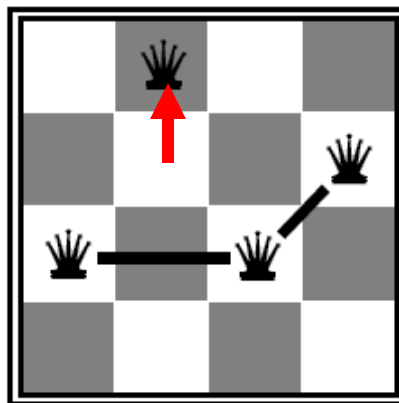
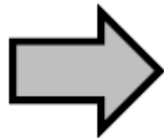
State space is reduced to  $4^4 = 256$

What is a possible local improvement strategy?

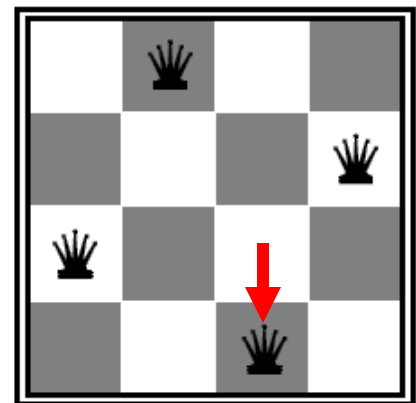
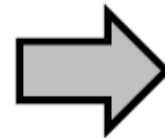
- Move one queen within its column to reduce conflicts



$h = 5$



$h = 2$



$h = 0$

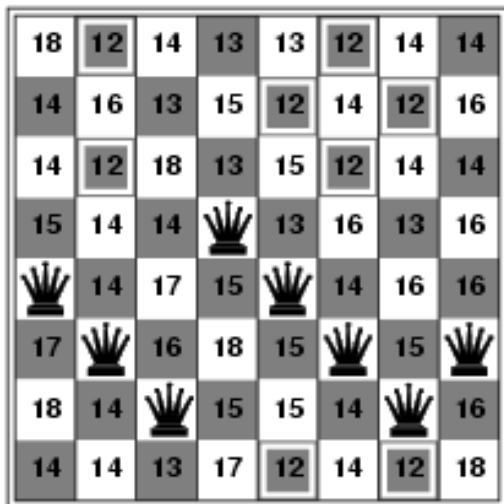


# Example: $n$ -queens problem

- **Goal:** Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal.
- **State space:** all possible  $n$ -queen configurations. We can restrict the state space: Only one queen per column.
- **Objective function:** minimize the number of pairwise conflicts.

What is a possible local improvement strategy?

- Move one queen within its column to reduce conflicts



18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	♙	13	16	13	16
♙	14	17	15	♙	14	16	16
17	♙	16	18	15	♙	15	♙
18	14	♙	15	15	14	♙	16
14	14	13	17	12	14	12	18

$h = 17$

best local improvement has  $h = 12$

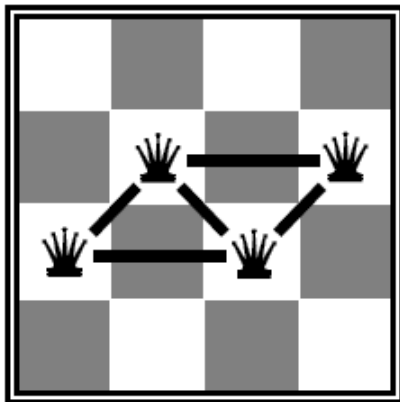
# Example: $n$ -queens problem

Optimization problem: find the best arrangement  $a$

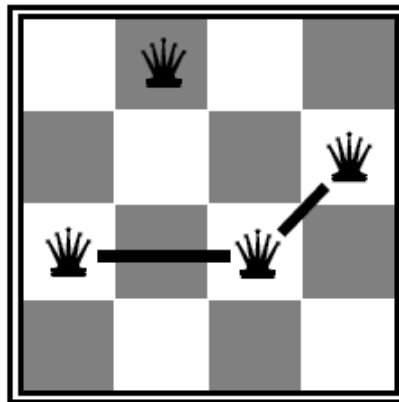
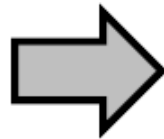
$$\operatorname{argmin}_a (\text{conflicts}(a))$$

s.t.  $a$  has one queen per column

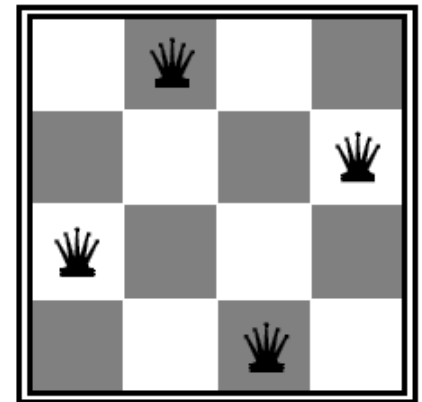
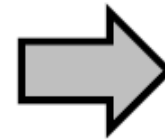
This makes the problem easier.



$h = 5$



$h = 2$



$h = 0$

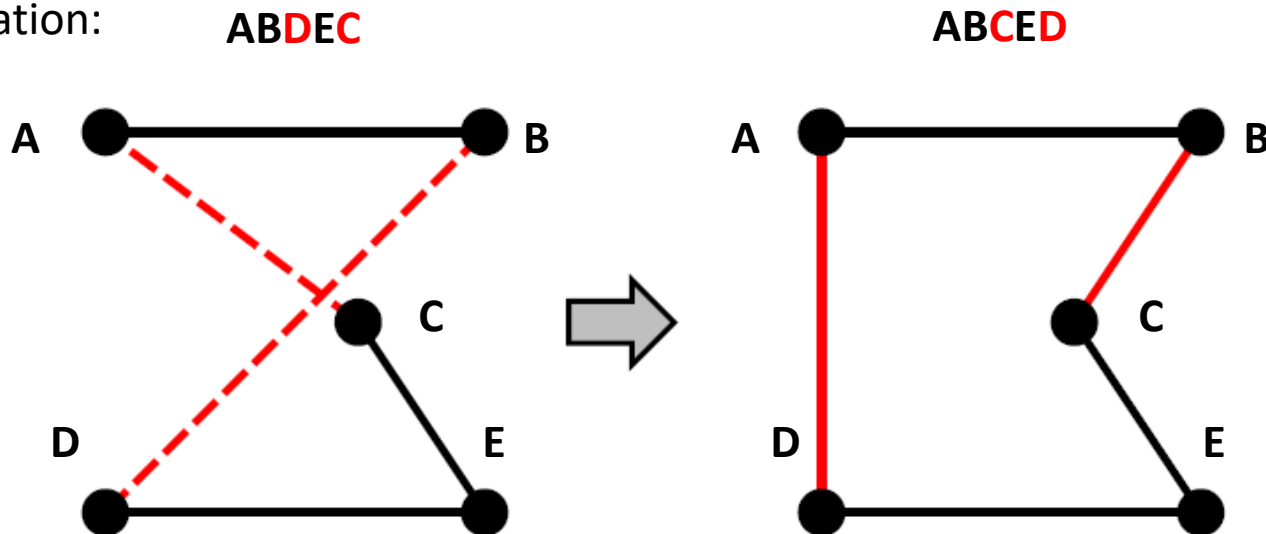
# Example: Traveling Salesman Problem

- **Goal:** Find the shortest tour connecting  $n$  cities
- **State space:** all possible tours
- **Objective function:** length of tour

What's a possible local improvement strategy?

- Start with any complete tour, perform pairwise exchanges.

Permutation:



# Example: Traveling Salesman Problem

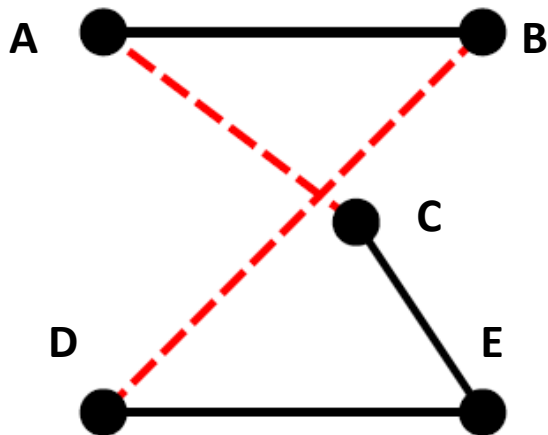
Optimization problem: Find the best tour  $\pi$

$$\operatorname{argmin}_{\pi} (\operatorname{tourLength}(\pi))$$

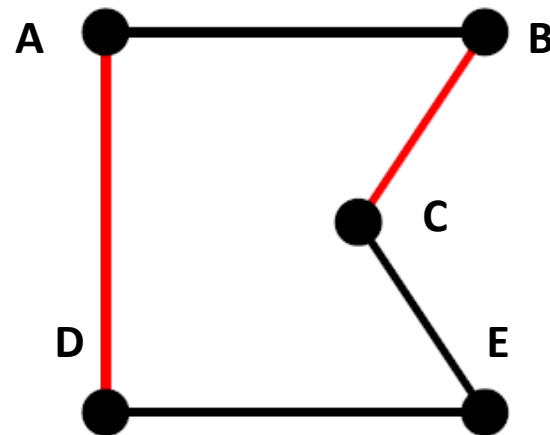
s.t.  $\pi$  is a valid permutation

Permutation:

**ABDEC**



**ABCED**





# Hill-climbing search (= Greedy local search)

Many variants

- **Steepest-ascend hill climbing**

```
function HILL-CLIMBING(problem) returns a state that is a local maximum  
  current  $\leftarrow$  problem.INITIAL  
  while true do  
    neighbor  $\leftarrow$  a highest-valued successor state of current  
    if VALUE(neighbor)  $\leq$  VALUE(current) then return current  
    current  $\leftarrow$  neighbor
```

- **Stochastic hill climbing**

- choose randomly among all uphill moves, or
- generate randomly new states until a better one is found (first-choice hill climbing)

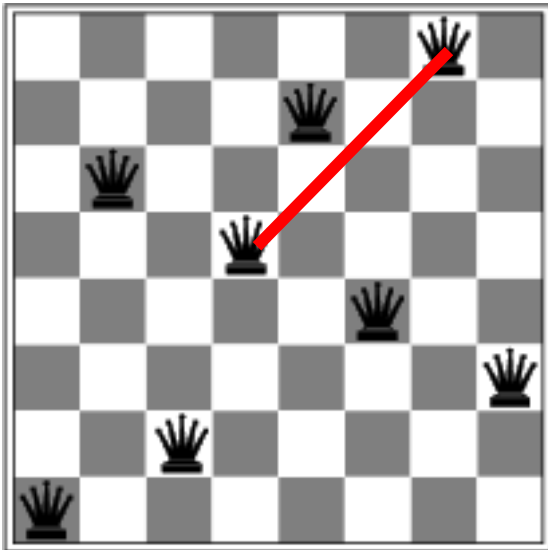
- **Random-restart hill climbing** – to deal with local optima

# Hill-climbing search

Hill-climbing search is similar to a best-first greedy search without backtracking.

Is it complete/optimal?

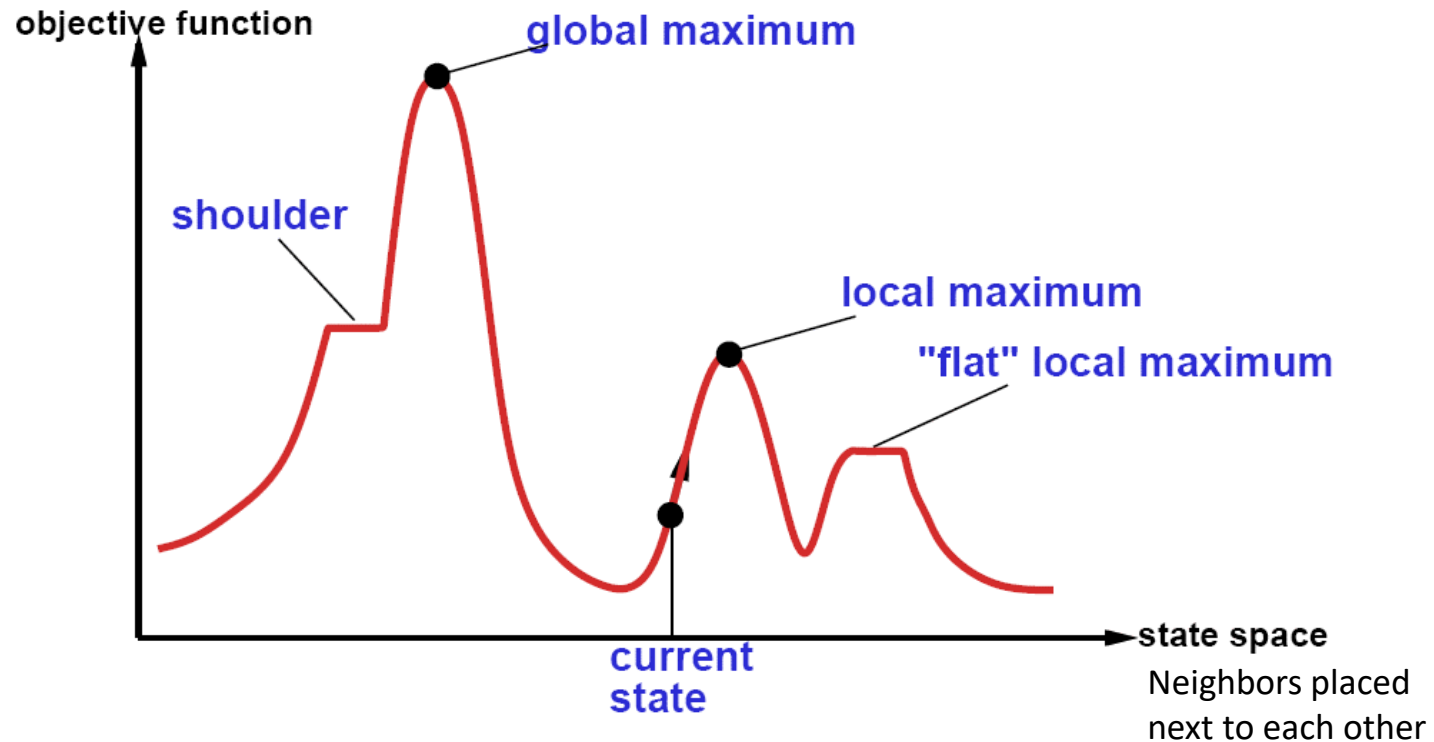
- No – can get stuck in local optima



Example: local optimum for the 8-queens problem. No single queen can be moved to improve the objective function.

$h = 1$

# The state space “landscape”



How to escape local maxima?

→ Random restart hill-climbing can help.

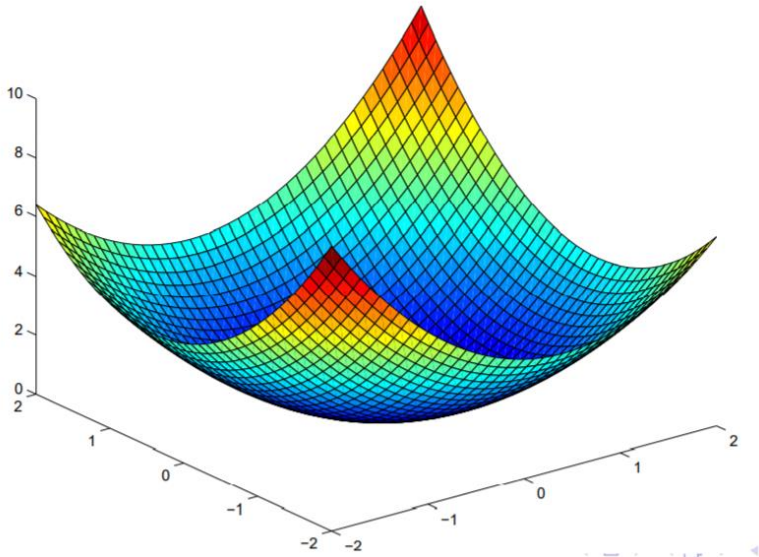
What about “shoulders” (“ridges” in higher dimensional space)?

What about “plateaux”?

→ Allow sideways moves.

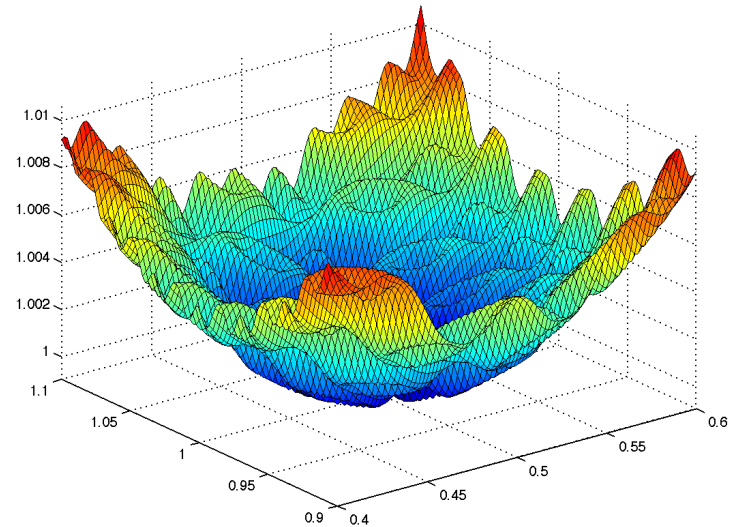
# Non-convex/convex Optimization Problems

Convex Problem



One global optimum +  
smooth function → easy

Non-convex Problem



Many local optima → hard

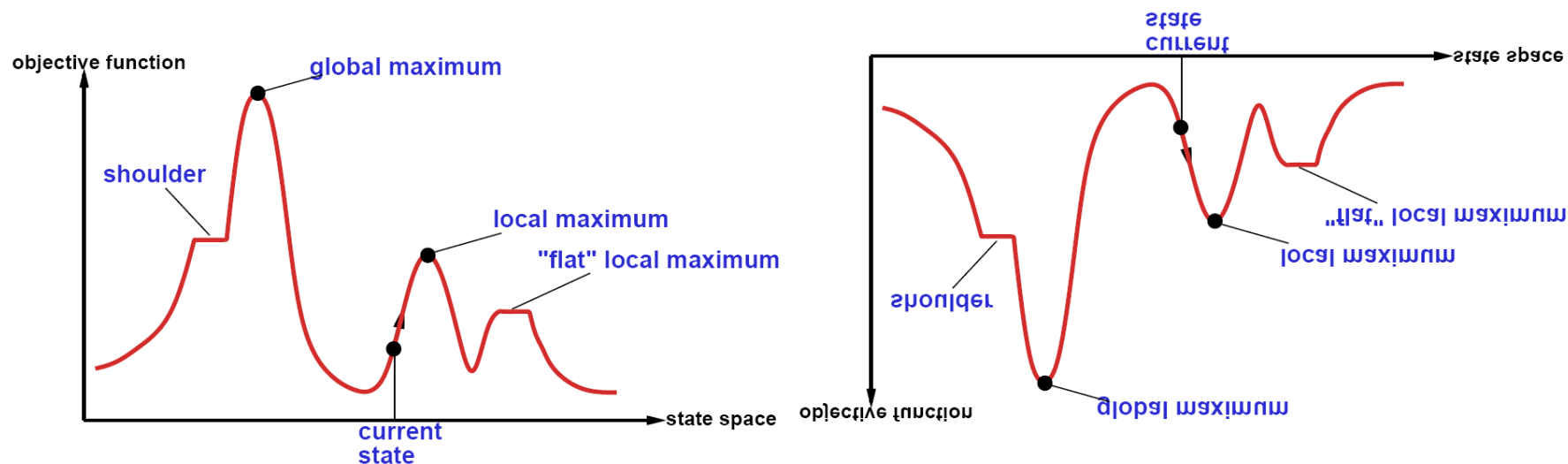
Many discrete optimization  
problems are like this.



# A Note on Minimization vs. Maximization

- Hill climbing implies **maximizing a function**.
- Optimizers like to state problems as **minimization problems**.
- Both types of problems are equivalent:

$$\max(f(x)) \Leftrightarrow \min(-f(x))$$



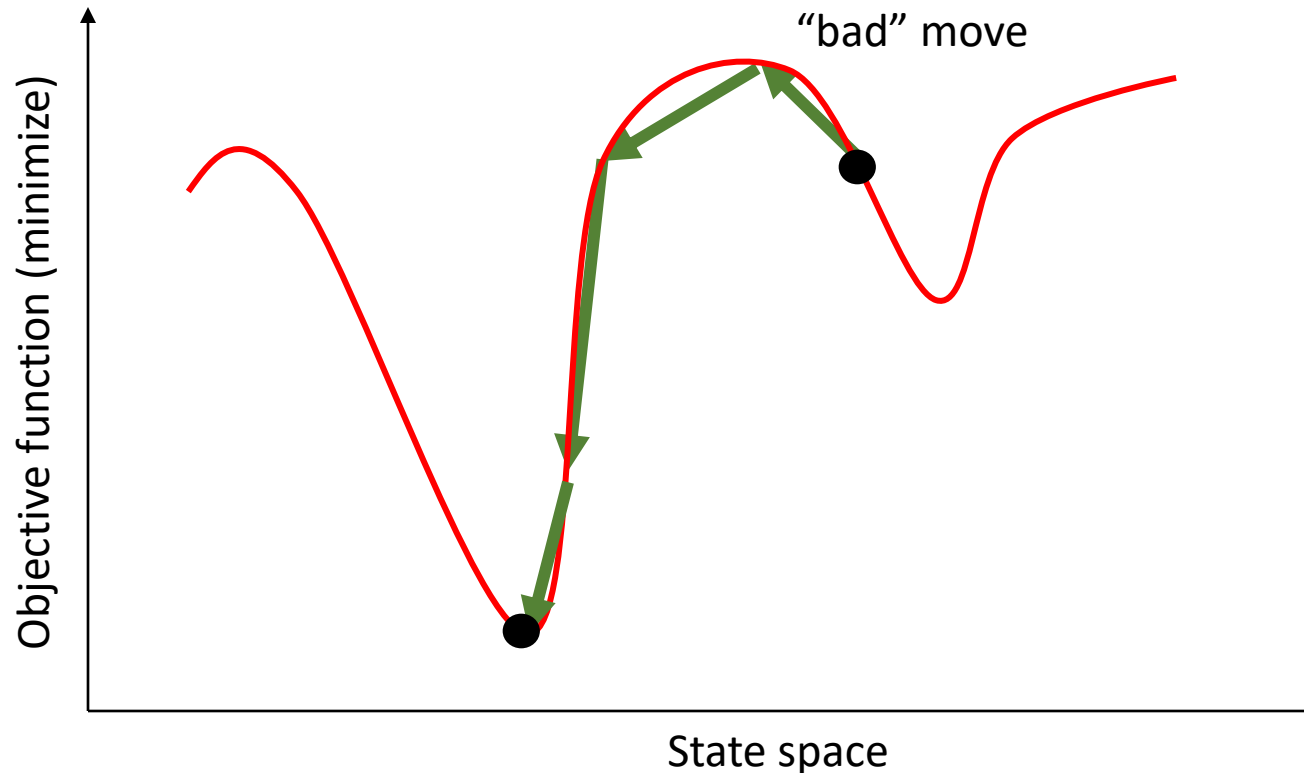
A glowing yellow-orange metal bar is being processed in a dark industrial machine. The bar is long and rectangular, with a bright, intense glow along its length, suggesting it is at a high temperature. It is positioned horizontally, and the machine's components, including rollers and structural frames, are visible in the background and foreground. The lighting is dramatic, with the bright glow of the metal contrasting sharply with the dark, shadowed machinery.

# Simulated Annealing

To escape local maxima

# Simulated annealing

- **Idea:** First-choice stochastic hill climbing + escape local minima by allowing some “bad” moves but gradually decrease their frequency.
- Inspired by the process of tempering or hardening metals by decreasing the temperature gradually.



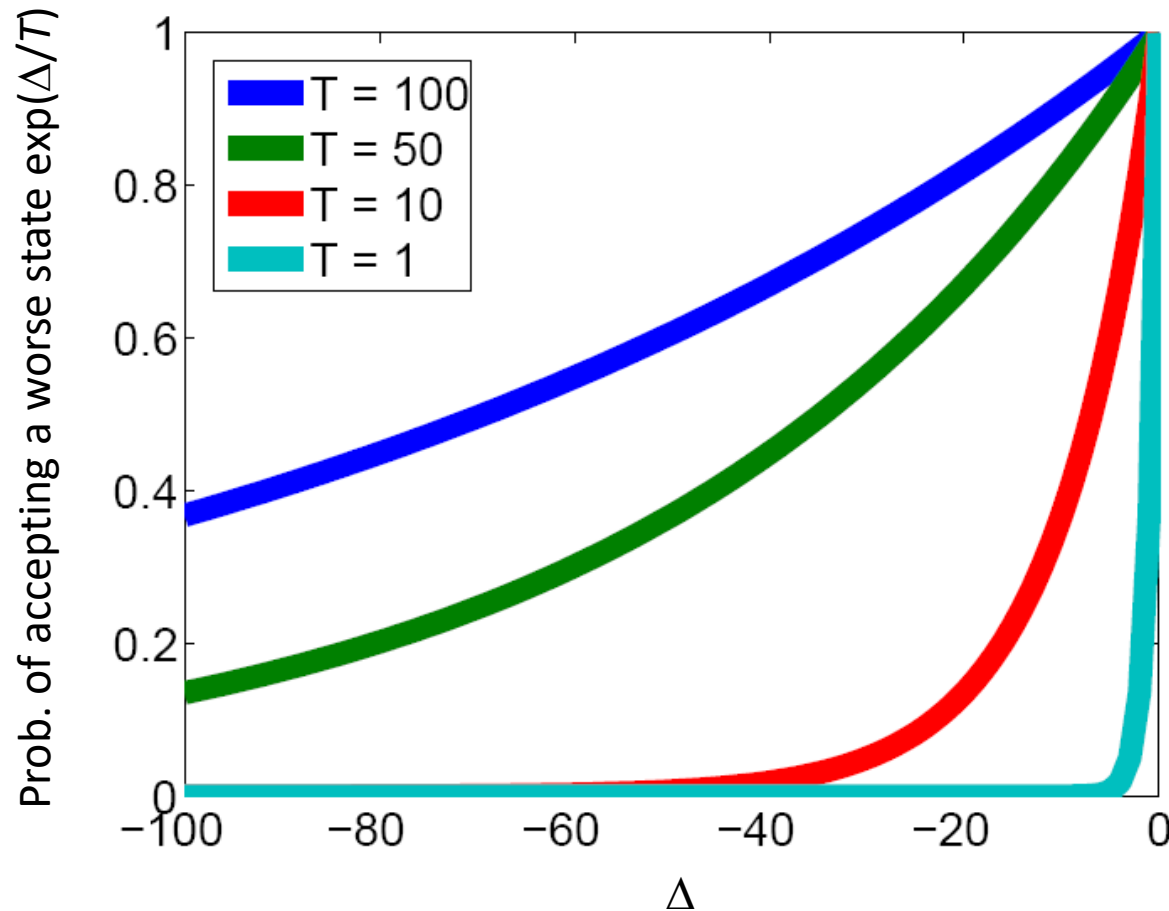
# Simulated annealing

- **Idea:** First-choice stochastic hill climbing + escape local minima by allowing some “bad” moves but gradually decrease their frequency.
- Inspired by the process of tempering or hardening metals by decreasing the temperature gradually.
- The probability of accepting “bad” moves follows **an annealing schedule** that reduces the temperature  $T$  over time  $t$ .

```
function SIMULATED-ANNEALING(problem, schedule) returns a solution state
  current  $\leftarrow$  problem.INITIAL
  for  $t = 1$  to  $\infty$  do
     $T \leftarrow$  schedule( $t$ )
    if  $T = 0$  then return current
    next  $\leftarrow$  a randomly selected successor of current
     $\Delta E \leftarrow$  VALUE(next) – Value(current)
    if  $\Delta E < 0$  then current  $\leftarrow$  next
    else current  $\leftarrow$  next only with probability  $e^{-\Delta E/T}$ 
```



# Effect of temperature



The lower the temperature, the less likely the algorithm will accept a worse state.

# Cooling Schedule

The cooling schedule is very important.

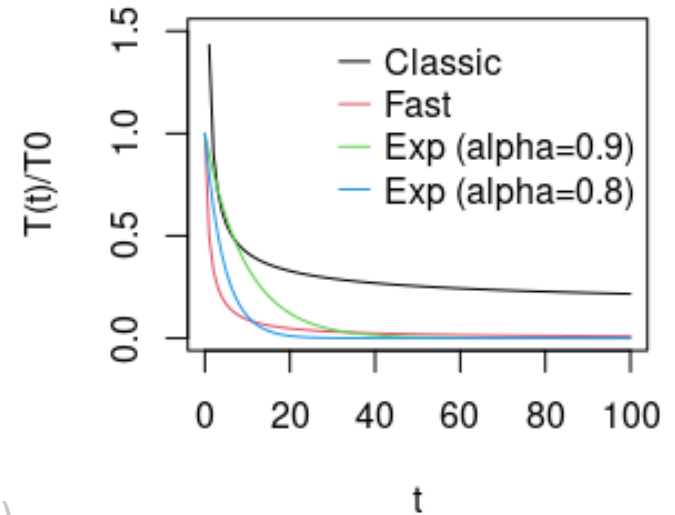
Popular schedules for the temperature at time  $t$ :

- **Classic simulated annealing:**  $T_t = T_0 \frac{1}{\log(1+t)}$
- **Fast simulated annealing** (Szy and Hartley; 1987)

$$T_t = T_0 \frac{1}{1+t}$$

- **Exponential cooling** (Kirkpatrick, Gelatt and Vecchi; 1983)

$$T_t = T_0 \alpha^t \quad \text{for } 0.8 < \alpha < 1$$

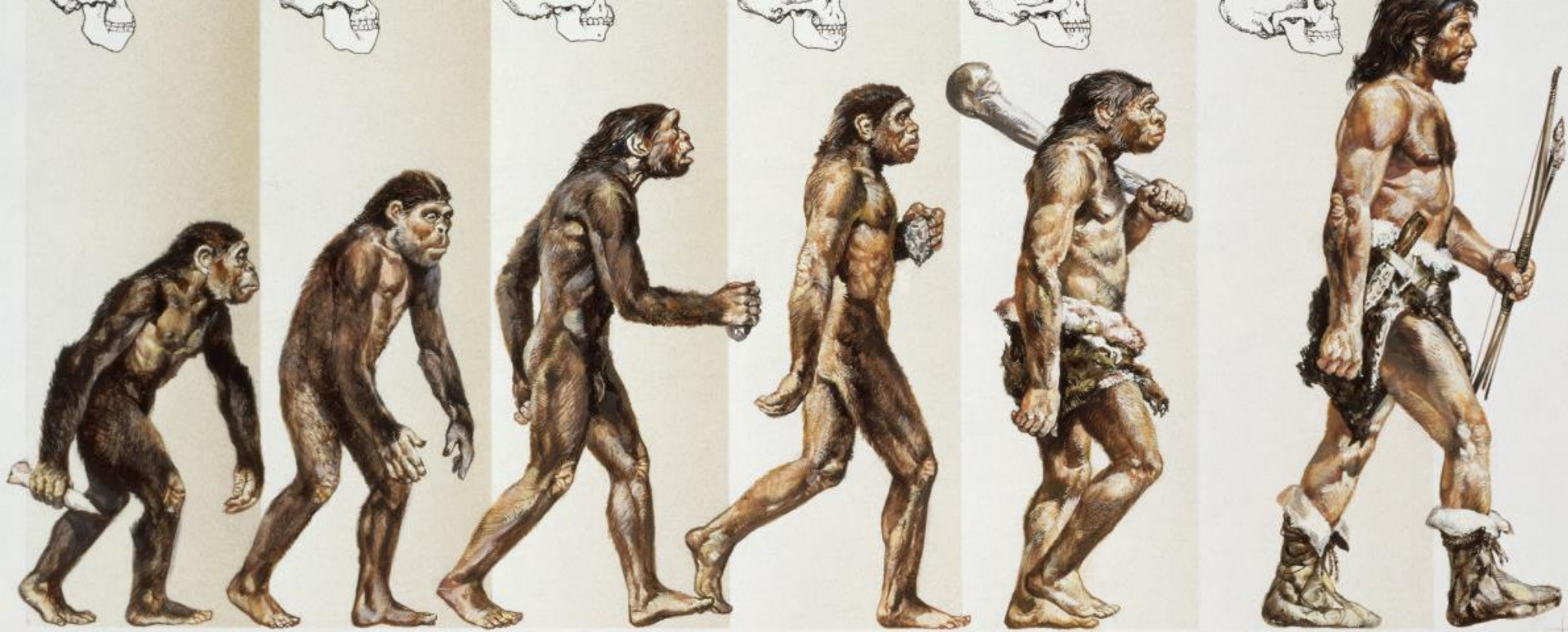


Notes:

- The best schedule is typically determined by trial-and-error.
- Choose  $T_0$  to provide a high probability that any move will be accepted at time  $t = 0$ .
- $T_t$  will not be come 0 but very small. Stop when  $T < \epsilon$  ( $\epsilon$  is a very small constant).

# Simulated annealing search

- **Guarantee:** If temperature decreases slowly enough, then simulated annealing search will find a global optimum with probability approaching one.
- However:
  - This usually takes impractically long
  - The more downhill steps you need to escape a local optimum, the less likely you are to make all of them in a row.
- More modern techniques: general family of *Markov Chain Monte Carlo* (MCMC) algorithms for exploring complicated state spaces.



# Evolutionary Algorithms

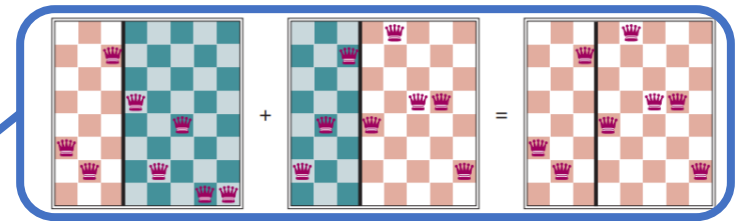
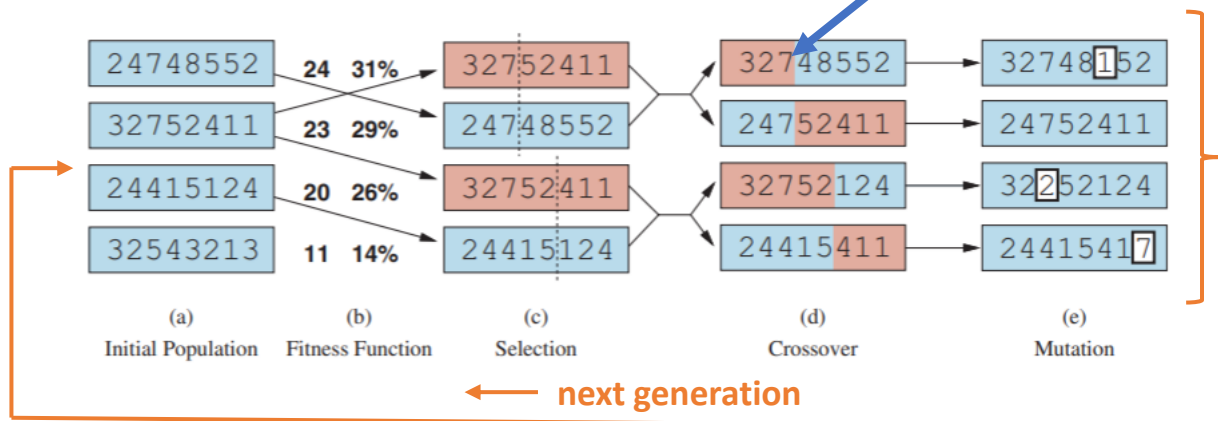
A Population-based Metaheuristics



# Evolutionary algorithms / Genetic Algorithms

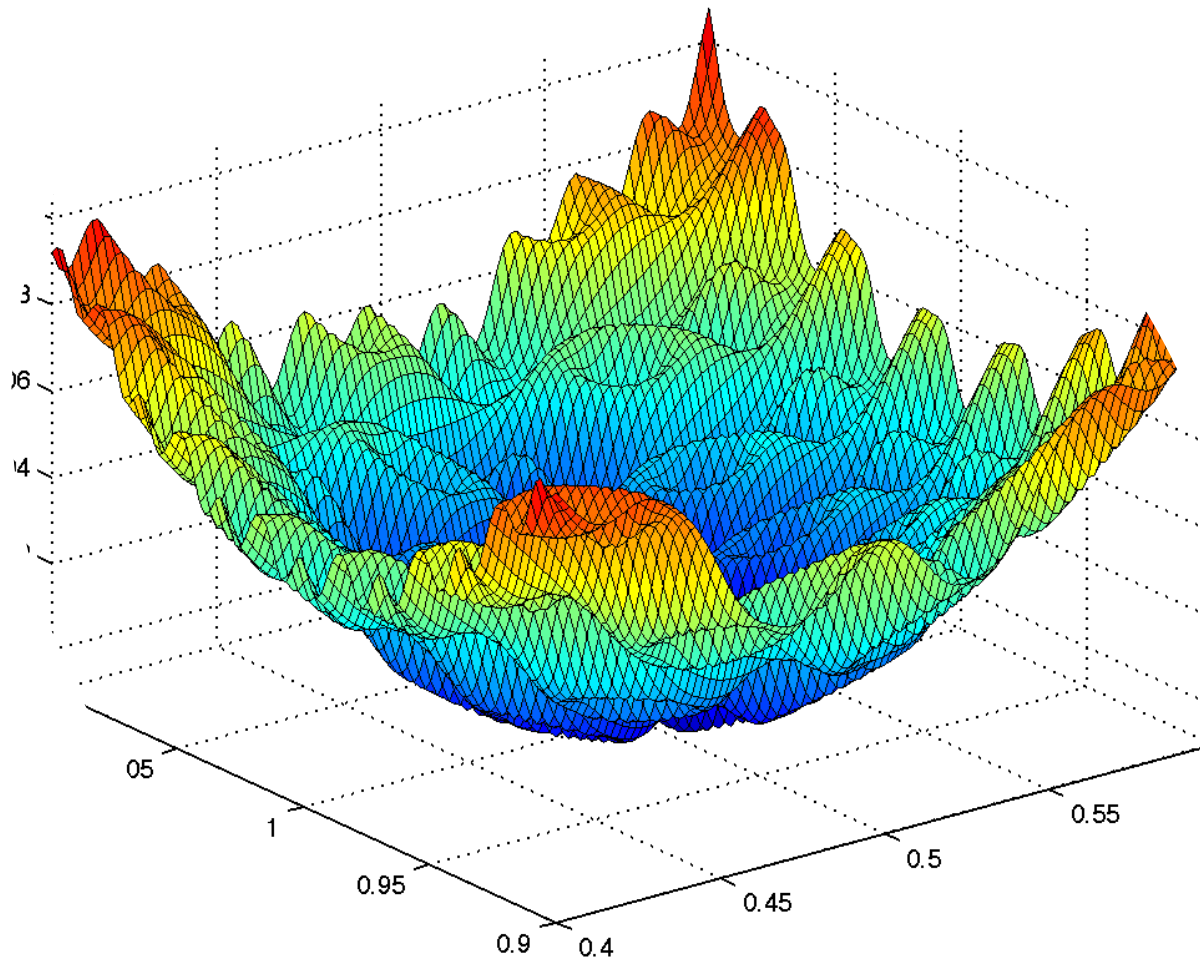
- A **population**-based metaheuristic optimization algorithm.
- Uses mechanisms inspired by biological evolution (genetics):
  - Reproduction: Random selection with probability based on a **fitness** function.
  - Random recombination (crossover)
  - Random mutation
  - Repeated for many generations

- Example: 8-queens states



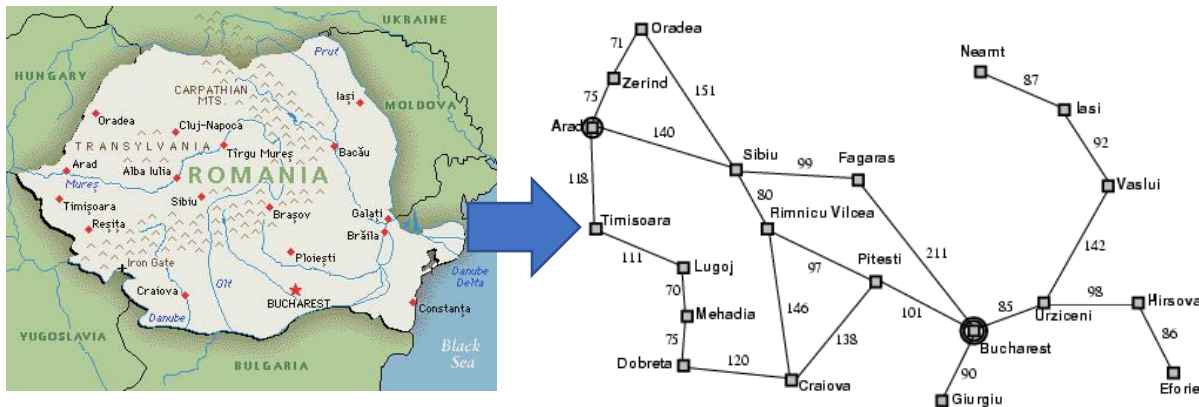
Encoding: position of the queen in each column

# Search in Continuous Spaces



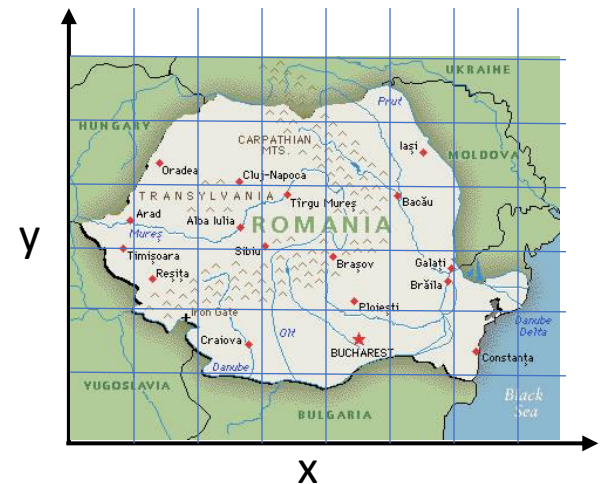
# Discretization of the continuous space

- Use atomic states



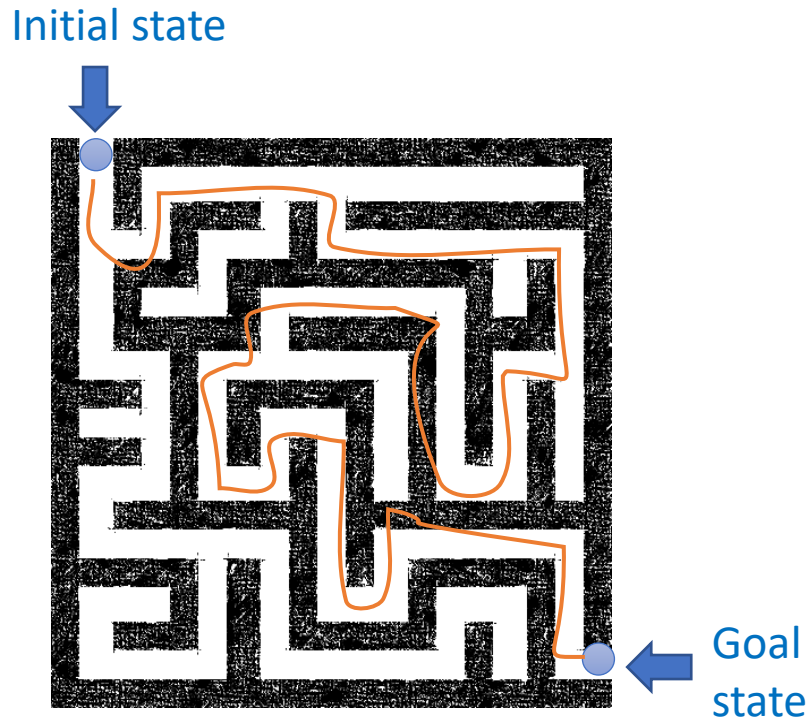
- Use a grid with spacing of size  $\delta$   
Note: You probably need a way finer grid!

Note: You probably need a way finer grid!



# Discretization of the continuous space

## How did we discretize this space?



# Search in continuous spaces: Gradient

Maximize  $f(\mathbf{x}) = f(x_1, x_2, \dots, x_k)$

Gradient at point  $\mathbf{x}$ :  $\nabla f(\mathbf{x}) = \left( \frac{\partial f(\mathbf{x})}{\partial x_1}, \frac{\partial f(\mathbf{x})}{\partial x_2}, \dots, \frac{\partial f(\mathbf{x})}{\partial x_k} \right)$

Find maximum by solving:  $\nabla f(\mathbf{x}) = 0$

- **Steepest-ascend hill climbing** (“gradient descend” for minimization) with step size  $\alpha$

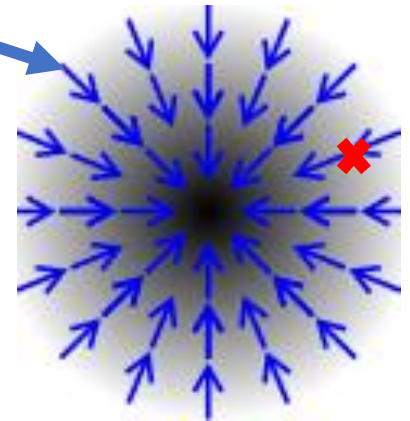
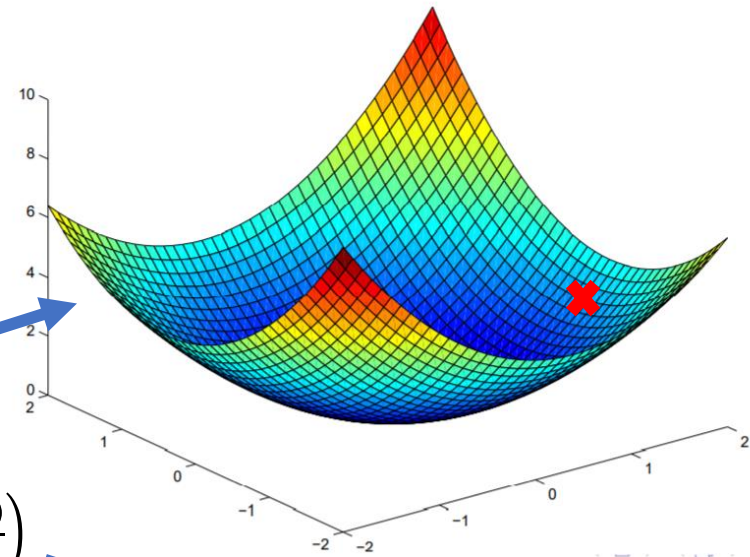
$$\mathbf{x} \leftarrow \mathbf{x} + \alpha \nabla f(\mathbf{x})$$

- **Newton-Raphson method**

uses the inverse of the Hessian matrix of second derivative  $H_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j}$  for the step size  $\alpha$

$$\mathbf{x} \leftarrow \mathbf{x} + \mathbf{H}_f^{-1}(\mathbf{x}) \nabla f(\mathbf{x})$$

May get stuck in a local optimum if the search space is non-convex! Use simulated annealing.



# Search in continuous spaces: Empirical Gradient Methods

- What if the mathematical formulation of the objective function is not known?
  - We may have objective values at fixed points, called the training data.
  - In this case we can use **empirical gradient search**. This is related to steepest ascend hill climbing in the discretized state space.
- We will talk more about search in continuous spaces with loss functions using gradient descend when we talk about **parameter learning for machine learning**.