

Serverless Dataflows: ...

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Computer Science and Engineering

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DeclarationI declare that this document is an original work of my own authorship and that it fulfills all the requirements of the Code of Conduct and Good Practices of the Universidade de Lisboa.

Acknowledgments

I would like to thank my parents for their friendship, encouragement and caring over all these years, for always being there for me through thick and thin and without whom this project would not be possible. I would also like to thank my grandparents, aunts, uncles and cousins for their understanding and support throughout all these years.

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Last but not least, to all my friends and colleagues that helped me grow as a person and were always there for me during the good and bad times in my life. Thank you.

To each and every one of you - Thank you.

Abstract

Serverless computing has become a suitable cloud paradigm for many applications, prized for its operational ease, automatic scalability, and fine-grained pay-per-use pricing model. However, executing workflows, which are compositions of multiple tasks, in Function-as-a-Service (FaaS) environments remains inefficient. This inefficiency stems from the stateless nature of functions, and a heavy reliance on external services for intermediate data transfers and inter-function communication.

In this document, we introduce a decentralized DAG engine that leverages historical metadata to plan and influence task scheduling. Our solution encompasses metadata management, static workflow planning, and a worker-level scheduling strategy designed to drive workflow execution with minimal synchronization. We compare our scheduling approach against WUKONG, another decentralized server-less DAG engine. Our evaluation demonstrates that utilizing historical information significantly improves performance and reduces resource utilization for workflows running on serverless platforms.

Keywords

Maecenas tempus dictum libero; Donec non tortor in arcu mollis feugiat; Cras rutrum pulvinar tellus.

Resumo

A computação serverless tornou-se um paradigma de nuvem adequado para muitas aplicações, val-

orizado pela sua facilidade operacional, escalabilidade automática e modelo de preços granular baseado

na utilização. Contudo, a execução de workflows, que são composições de múltiplas tarefas, em ambi-

entes Function-as-a-Service (FaaS) permanece ineficiente. Esta ineficiência resulta da natureza state-

less (sem estado) destas funções e de uma forte dependência de serviços externos para transferências

de dados intermédios e comunicação entre funções.

workflows executados em plataformas serverless.

Neste documento, apresentamos um motor de workflows serverless descentralizado que utiliza

métricas recolhidas durante a execução para planear e influenciar o scheduling de tarefas. A nossa

solução abrange a gestão de metadados, o planeamento estático de workflows e uma estratégia de

scheduling ao nível dos workers concebida para conduzir a execução de workflows de uma forma de-

scentralizada e com sincronização mínima. Comparamos a nossa abordagem com o WUKONG, outro

motor de workflows serverless descentralizado. A nossa avaliação demonstra que a utilização de in-

formação histórica melhora significativamente o desempenho e reduz a utilização de recursos para

Palavras Chave

Cloud Computing; Serverless; FaaS; Serverless Workflows; Serverless DAGs; Metadata; Workflow

Prediction

ν

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Acronyms



1

Introduction

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Function-as-a-Service (FaaS) represents a serverless cloud computing paradigm that simplifies application deployment by abstracting away infrastructure management. It provides automatic, elastic scalability—potentially without limit—along with a fine-grained, pay-per-use pricing model. This has led to its widespread adoption for event-driven systems, microservices, and web services on platforms like AWS Lambda¹, Azure Functions², and Google Cloud Functions³. These applications typically benefit the most from FaaS because they are lightweight, stateless, and characterized by highly variable or unpredictable workloads, allowing them to leverage serverless platforms' on-demand scalability and cost-efficiency.

This paradigm is also increasingly used to execute complex scientific and data processing workflows,

¹https://aws.amazon.com/pt/lambda/

²https://azure.microsoft.com/en-us/products/functions

³https://cloud.google.com/functions

such as the Cybershake [1] seismic hazard analysis or Montage [2], an astronomy image mosaicking workflow. These applications are structured as workflows—formally represented as Directed Acyclic Graphs (DAGs) of interdependent tasks. However, efficiently executing these complex workflows on serverless platforms remains a significant challenge.

Despite their advantages, serverless platforms present several limitations that complicate the execution of complex workflows. Since these platforms allow scaling down to zero resources to save costs, they can also introduce unpredictable latency, known as *cold starts* [3], particularly for short-lived functions, affecting overall workflow performance. The lack of *direct inter-function communication* [4] means that tasks often have to rely on external services, such as message brokers or databases to exchange intermediate data, which can increase overhead and reduce efficiency. Interoperability between platforms is further limited by the use of platform-specific workflow definition languages, which restricts the portability of workflows across different serverless environments. Additionally, while statelessness simplifies scaling and management, it can introduce overhead and complexity for applications that require continuity or coordination across multiple function invocations. Finally, developers have limited control over the underlying infrastructure, restricting the ability to optimize resource usage or tune performance for specific workloads.

Several solutions have emerged to address the limitations of serverless platforms. Stateful functions (e.g., AWS Step Functions⁴, Azure Durable Functions⁵, and Google Cloud Workflows⁶) expand the range of applications that can run on serverless platforms by maintaining state across multiple function invocations, coordinating complex workflows, and providing built-in fault tolerance. Other approaches tackle limitations at the runtime level, proposing extensions to FaaS platforms (e.g., Faa\$T [5], Palette [6], Lambdata [7]) or entirely new serverless architectures (e.g., Apache OpenWhisk [8]). Finally, some workflow-focused solutions (e.g., WUKONG [9], Unum [10], DEWEv3 [11]) employ scheduling strategies and workflow-level optimizations to enhance efficiency, primarily by improving data locality to bring computation closer to the data and minimize reliance on external services.

These workflow-focused approaches, however, often use uniform resources for workers and rely on "one-step scheduling," making decisions based solely on the immediate workflow stage without considering the broader context or the downstream effects of their decisions. This combination of homogeneous worker configurations and limited scheduling foresight can lead to inefficient use of resources when tasks have diverse computational or memory requirements. Furthermore, the heuristic-based approaches used by other solutions can be inefficient in certain scenarios, as they lack mechanisms to adapt worker resource allocations to the specific needs of individual tasks. Moreover, we found no prior work that leverages metadata or historical metrics to inform scheduling decisions across an entire

⁴https://aws.amazon.com/pt/step-functions/

⁵https://learn.microsoft.com/en-us/azure/azure-functions/durable/durable-functions-overview?tabs=in-process%2Cnodejs-v3%2Cv1-model&pivots=csharp

⁶https://cloud.google.com/workflows

serverless workflow.

This limitation motivates the central research question of this work: if we have knowledge of the computation steps, collect sufficient metrics on their behavior, and understand how they are composed to form the full workflow, can we leverage this information to make smarter scheduling decisions that minimize makespan and maximize resource efficiency in a FaaS environment?

To answer this research question, we propose a decentralized serverless workflow execution engine that leverages historical metadata from previous workflow runs to generate informed task allocation plans, which are then executed by FaaS workers in a choreographed manner, without needing a central scheduler. By relying on such planning, our approach aims to minimize the usage of external cloud storage services, which are often employed by similar solutions for intermediate data exchange and synchronization, while also avoiding the inefficiencies of homogeneous worker resource allocations.

The main contributions of this work are as follows:

- Analysis of the serverless workflow orchestration research landscape;
- Propose a decentralized serverless workflow execution engine that overcomes the "one-step scheduling" and uniform-resource limitations of existing workflow-focused solutions by leveraging historical metadata to generate informed execution plans;
- Demonstrate how incorporating historical execution data can improve task allocation, reduce reliance on external cloud storage services, and enhance overall workflow efficiency on FaaS platforms.

1.1 Problem/Motivation

1.2 Gaps in prior work

1.3 Proposed Solution

1.4 Document Organization

2

Related Work

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Architecture

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Algorithm 1 Worker Assignment Algorithm

```
Require: nodes, predictions, base_rc, SLA, MAX_CLUSTERING
 1: assigned \leftarrow \emptyset
                                                                                             ▷ nodes are topologically sorted
 2: for all n \in nodes do
       if n \in assigned then
 4:
           continue
 5:
        if n.upstream = \emptyset then
 6:
                                                                                                                   7:
           roots \leftarrow \{r \in nodes \mid r.upstream = \emptyset \land r \notin assigned\}
 8:
           ASSIGNGROUP(null, roots)
 9:
        else if |n.upstream| = 1 then
                                                                                                                \triangleright 1 \rightarrow 1 or 1 \rightarrow N
10:
           u \leftarrow n.upstream[0]
           if |u.downstream| = 1 then
11:
               AssignWorker([n], u.worker)
                                                                                                                > reuse worker
12:
                                                                                                                         > 1→N
13:
                fanout \leftarrow \{d \in u.downstream \mid d \notin assigned\}
14.
15:
               AssignGroup(u.worker, fanout)
16:
           end if
17:
        else
                                                                                                                         \triangleright N \rightarrow 1
18:
           outputs \leftarrow \{u.worker : predictions.output\_size(u) \mid u \in n.upstream\}
19:
           best \leftarrow \arg\max_{w \in outputs} outputs[w]
20:
            ASSIGNWORKER([n], best)
21:
        end if
22: end for
23: function ASSIGNGROUP(up\_worker, tasks)
        if tasks = \emptyset then return
24:
25:
        end if
26:
        exec \ t \leftarrow \{t : predictions.exec \ time(t) \mid t \in tasks\}
        out \ sz \leftarrow \{t: predictions.output\_size(t) \mid t \in tasks\}
27:
        median \leftarrow \mathsf{MEDIAN}(exec\ t.values())
28:
29:
        longs \leftarrow \{t \in tasks \mid exec\_t[t] > median\}
30:
        shorts \leftarrow \mathsf{SORTLARGEROUTPUTFIRST}(\{t \in tasks \mid exec\_t[t] \leq median\})
                                                          ▷ 1) cluster short tasks with bigger outputs on upstream worker
31:
32:
        if up\_worker \neq null \land shorts \neq \emptyset then
33:
            cluster \leftarrow shorts[0:MAX\_CLUSTERING]
           {\tt ASSIGNWORKER}(cluster,\,up\_worker)
34:
35:
            shorts \leftarrow shorts[MAX\_CLUSTERING:]
36:
        end if
37:
                                                        ▷ 2) pair long tasks with remaining short tasks (1 long per group)
38:
        while longs \neq \emptyset \land shorts \neq \emptyset do
            cluster \leftarrow [longs[0]] + shorts[0:MAX\_CLUSTERING-1]
39:
           worker id \leftarrow \mathsf{NEWWORKERID}
40:
41:
           ASSIGNWORKER(cluster, worker_id)
42.
           longs \leftarrow longs[1:]
            shorts \leftarrow shorts[MAX\_CLUSTERING - 1:]
43:
44:
        end while
45:

⊳ 3) group remaining short tasks

46:
        while shorts \neq \emptyset do
47:
           worker\_id \leftarrow \mathsf{NEWWORKERID}
           AssignWorker(shorts[0:MAX\_CLUSTERING], worker\_id)
48:
            shorts \leftarrow shorts[MAX\_CLUSTERING:]
49:
        end while
50:
51.
                                                                                       52:
        half \leftarrow \max(1, \lfloor MAX\_CLUSTERING/2 \rfloor)
53:
        while longs \neq \emptyset do
            worker\_id \leftarrow \mathsf{NEWWORKERID}
54:
            {\sf ASSIGNWORKER}(longs[0:half],worker\_id)
55:
            longs \leftarrow longs[half:]
56:
57:
        end while
58: end function
```

Algorithm 2 Resource Downgrading Algorithm

```
Require: dag, nodes, critical_path_ids, original_cp_time, configs, predictions
 1: workers\_outside \leftarrow \emptyset
 2:
                                                                           3: for all n \in nodes do
                                                                                        > nodes are topologically sorted
 4.
       wid \leftarrow n.worker\_id
 5:
       if n.id \notin critical\_path\_ids \land \forall cp \in dag.critical\_path\_nodes : wid \neq cp.worker\_id then
 6:
           workers\_outside \leftarrow workers\_outside \cup \{wid\}
 7:
 8: end for
 9: nodes\_outside\_cp \leftarrow \{n \in nodes \mid n.id \notin critical\_path\_ids\}
                                                         > 2) Attempt downgrade for each worker outside critical path
10:
11: for all wid \in workers\_outside do
       last\_good\_rc \leftarrow \{n.id : n.config \mid n \in nodes\_outside\_cp \land n.worker\_id = wid\}
12:
                                                    ▷ Iterate through weaker configurations (skip strongest at index 0)
13:
14:
       for i \leftarrow 1 to |configs| - 1 do
15:
           trial \leftarrow configs[i].\mathsf{CLONE}(wid)
16:
                                                                   > Apply trial configuration to all nodes of this worker
17:
           for all n \in nodes outside cp do
18:
              if n.worker\_id = wid then
19:
                  n.config \leftarrow trial
              end if
20:
           end for
21:
22:
                                                                         ▷ Recompute workflow timing with predictions
23:
           cp\_time \leftarrow \mathsf{SIMULATECRITICALPATHTIME}(dag)
24:
           if cp\_time = original\_cp\_time then
25:
                                                                    Downgrade acceptable, record as last good state
26:
              for all n \in nodes outside cp do
27:
                  if n.worker\_id = wid then
28:
                      last\_good\_rc[n.id] \leftarrow n.config
                  end if
29:
              end for
30:
           else
31:
32:
                                          Downgrade increases critical path, revert and move on to the next worker
33:
              for all n \in nodes outside cp do
                  if n.worker\_id = wid then
34:
35:
                      n.config \leftarrow last\_good\_rc[n.id]
36:
                  end if
37:
              end for
              break
                                                                                                  38:
           end if
39:
40:
       end for
41: end for
```

4

Evaluation

11

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4.2 Proin ornare dignissim lacus
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4.1 Maecenas vitae nulla consequat

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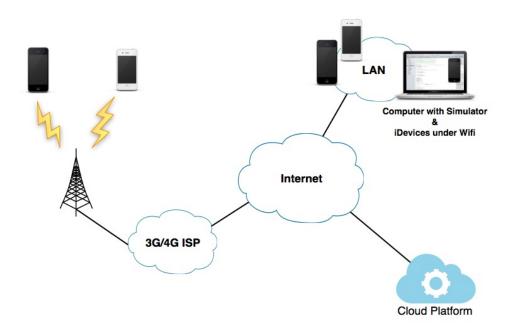


Figure 4.1: Test Environment

Aliquam aliquet, est a ullamcorper condimentum, tellus nulla fringilla elit, a iaculis nulla turpis sed wisi. Fusce volutpat. Etiam sodales ante id nunc. Proin ornare dignissim lacus. Nunc porttitor nunc a sem. Sed sollicitudin velit eu magna. Aliquam erat volutpat. Vivamus egestas. Nunc tempor diam vehicula mauris. Nullam sapien eros, facilisis vel, eleifend non, auctor dapibus, pede 4.1 used in the tests. The Network Link Conditioner allows to force/simulate fluctuations in fixed network segments.

Table 4.1: Network Link Conditioner Profiles

Network Profile	Bandwidth	Packets Droped	Delay
Wifi	40 mbps	0%	1 ms
3G	780 kbps	0%	100 ms
Edge	240 kbps	0%	400 ms

Aliquam aliquet, est a ullamcorper condimentum, tellus nulla fringilla elit, a iaculis nulla turpis sed wisi. Fusce volutpat. Etiam sodales ante id nunc. Proin ornare dignissim lacus. Nunc porttitor nunc a sem. Sed sollicitudin velit eu magna. Aliquam erat volutpat. Vivamus ornare est non wisi. Proin vel quam. Vivamus egestas. Nunc tempor diam vehicula mauris. Nullam sapien eros, facilisis vel, eleifend non, auctor dapibus, pede.

4.2 Proin ornare dignissim lacus

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 N_j Is the number of times peer j has been optimistically unchoked.

 n_j Among the N_j unchokes, the number of times that peer j responded with unchoke or supplied segments to peer p.

 $C_{r[j]}$ The cooperation ratio of peer j. If peer j never supplied peer p, the information of $C_{r[j]}$ may not be available.

 $C_{r(max)}$ The maximum cooperation ratio of peer p's neighbors, i.e., $C_{r(max)} = max(C_r)$.

$$G_j = \begin{cases} \frac{n_j C_{r[j]}}{N_j} & \text{if } n_j > 0\\ \frac{C_{r(max)}}{N_j + 1} & \text{if } n_j = 0 \end{cases} \tag{4.1}$$

Cursus $C_{r(max)}$ conubia nostra, per inceptos hymenaeos j gadipiscing mollis massa $N_j=0$, unc ut dui eget nulla venenatis aliquet $G_j=C_{r(max)}$.

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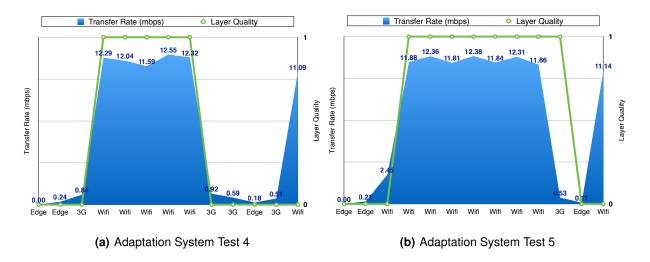


Figure 4.2: Adaptation System Behavior Test

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Conclusion

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5.1 **Conclusions**

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Rui Cruz You should always start a Chapter with an in-

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5.2 System Limitations and Future Work

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Code of Project

Nulla dui purus, eleifend vel, consequat non, dictum porta, nulla. Duis ante mi, laoreet ut, commodo eleifend, cursus nec, lorem. Aenean eu est. Etiam imperdiet turpis. Praesent nec augue. Curabitur ligula quam, rutrum id, tempor sed, consequat ac, dui. Vestibulum accumsan eros nec magna. Vestibulum vitae dui. Vestibulum nec ligula et lorem consequat ullamcorper.

Listing A.1: Example of a XML file.

Etiam imperdiet turpis. Praesent nec augue. Curabitur ligula quam, rutrum id, tempor sed, consequat ac, dui. Maecenas tincidunt velit quis orci. Sed in dui. Nullam ut mauris eu mi mollis luctus. Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos. Sed cursus cursus velit. Sed a massa. Duis dignissim euismod quam.

Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos. Phasellus eget nisl ut elit porta ullamcorper. Maecenas tincidunt velit quis orci. Sed in dui. Nullam ut mauris eu mi mollis luctus. Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos.

This inline MATLAB code for i=1:3, disp('cool'); end; uses the \mcode{} command.1

Nullam ut mauris eu mi mollis luctus. Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos. Sed cursus cursus velit. Sed a massa. Duis dignissim euismod quam. Nullam euismod metus ut orci.

Listing A.2: Matlab Function

Nullam ut mauris eu mi mollis luctus. Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos. Sed cursus cursus velit. Sed a massa. Duis dignissim euismod quam. Nullam euismod metus ut orci.

¹MATLAB Works also in footnotes: for i=1:3, disp('cool'); end;

Listing A.3: function.m

```
Copyright 2010 The MathWorks, Inc.
2 function ObjTrack(position)
3 % #codegen
4 % First, setup the figure
5 numPts = 300;
                           % Process and plot 300 samples
6 figure; hold; grid;
                      % Prepare plot window
7 % Main loop
8 for idx = 1: numPts
      z = position(:,idx); % Get the input data
      y = kalmanfilter(z); % Call Kalman filter to estimate the position
      plot_trajectory(z,y); % Plot the results
12 end
13 hold;
14 end
      % of the function
```

Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos. Phasellus eget nisl ut elit porta ullamcorper. Maecenas tincidunt velit quis orci. Sed in dui. Nullam ut mauris eu mi mollis luctus. Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos hymenaeos. Sed cursus cursus velit. Sed a massa. Duis dignissim euismod quam. Nullam euismod metus ut orci. Vestibulum erat libero, scelerisque et, porttitor et, varius a, leo.

Listing A.4: HTML with CSS Code

```
1 <!DOCTYPE html>
2 <html>
    <head>
      <title>Listings Style Test</title>
      <meta charset="UTF-8">
      <style>
        /* CSS Test */
         * {
8
           padding: 0;
           border: 0;
10
           margin: 0;
        }
12
      </style>
13
      <link rel="stylesheet" href="css/style.css" />
14
    </head>
15
```

```
<header> hey </header>
    <article> this is a article </article>
    <body>
      <!-- Paragraphs are fine -->
19
      <div id="box">
20
        >
21
          Hello World
22
        23
        Hello World
24
        Hello World
        </div>
27
      <div>Test</div>
28
      <!-- HTML script is not consistent -->
29
      <script src="js/benchmark.js"></script>
30
      <script>
31
        function createSquare(x, y) {
          // This is a comment.
          var square = document.createElement('div');
           square.style.width = square.style.height = '50px';
35
           square.style.backgroundColor = 'blue';
37
          /*
           * This is another comment.
           */
           square.style.position = 'absolute';
41
           square.style.left = x + 'px';
42
          square.style.top = y + 'px';
43
          var body = document.getElementsByTagName('body')[0];
          body.appendChild(square);
        };
47
48
        // Please take a look at +=
49
        window.addEventListener('mousedown', function(event) {
50
          // German umlaut test: Berührungspunkt ermitteln
51
          var x = event.touches[0].pageX;
          var y = event.touches[0].pageY;
```

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Listing A.5: HTML CSS Javascript Code

```
@media only screen and (min-width: 768px) and (max-width: 991px) {
2
    \# main {
       width: 712px;
       padding: 100px 28px 120px;
    }
    /* .mono {
      font-size: 90%;
10
    } */
11
    .cssbtn a {
13
       margin-top: 10px;
14
       margin-bottom: 10px;
15
       width: 60px;
16
       height: 60px;
17
       font-size: 28px;
18
       line-height: 62px;
    }
20
```

Nulla dui purus, eleifend vel, consequat non, dictum porta, nulla. Duis ante mi, laoreet ut, commodo eleifend, cursus nec, lorem. Aenean eu est. Etiam imperdiet turpis. Praesent nec augue. Curabitur ligula quam, rutrum id, tempor sed, consequat ac, dui. Vestibulum accumsan eros nec magna. Vestibulum vitae dui. Vestibulum nec ligula et lorem consequat ullamcorper.

Listing A.6: PYTHON Code

```
1 class TelgramRequestHandler(object):
2   def handle(self):
3    addr = self.client_address[0]  # Client IP-adress
4   telgram = self.request.recv(1024)  # Recieve telgram
5   print "From: %s, Received: %s" % (addr, telgram)
6   return
```

B

A Large Table

Aliquam et nisl vel ligula consectetuer suscipit. Morbi euismod enim eget neque. Donec sagittis massa. Vestibulum quis augue sit amet ipsum laoreet pretium. Nulla facilisi. Duis tincidunt, felis et luctus placerat, ipsum libero vestibulum sem, vitae elementum wisi ipsum a metus. Nulla a enim sed dui hendrerit lobortis. Donec lacinia vulputate magna. Vivamus suscipit lectus at quam. In lectus est, viverra a, ultricies ut, pulvinar vitae, tellus. Donec et lectus et sem rutrum sodales. Morbi cursus. Aliquam a odio. Sed tortor velit, convallis eget, porta interdum, convallis sed, tortor. Phasellus ac libero a lorem auctor mattis. Lorem ipsum dolor sit amet, consectetuer adipiscing elit.

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As B.1 shows, the data can be inserted from a file, in the case of a somehow complex structure. Notice the Table footnotes.

Table B.1: Example table

Danahmanler	#1	#Nets	#Nladaa*	Critical	Latency
Benchmark:	#Layers		#Nodes*	path	(T_{iter})
ANN	(1)	(2)	$(3) = 8 \cdot (1) \cdot (2)$	$(4) = 4 \cdot (1)$	(5)
A1	3–1501	1	24-12008	12-6004	4
A2	501	1	4008 2004		2–2000
A3	10	2-1024	160-81920	40	60 [†]
A4	10	50	4000	40	80–1200
Benchmark: FFT	FFT size [‡]	#Inputs		Critical	Latency
			#Nodes*	path	(T_{iter})
	(1)	$(2) = 2^{(1)}$	$(3) = 10 \cdot (1) \cdot (2)$	$(4) = 4 \cdot (1)$	(5)
F1	1–10	2-1024	20-102400	4–40	6-60 [†]
F2	5	32	1600	20	40 – 1500
Benchmark:	<u> </u>			Critical	Latency
Random	#Types	#Nodes	#Networks	path	(T_{iter})
networks	(1)	(2)	(3)	(4)	(5)
R1	3	10–2000	500	variable	(4)
R2	3	50	500	variable	$(4) \times [1; \cdots; 20]$

^{*} Excluding constant nodes.

Values in bold indicate the parameter being varied.

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And now an example (??) of a table that extends to more than one page. Notice the repetition of the Caption (with indication that is continued) and of the Header, as well as the continuation text at the bottom.

An example of a large Table that autofits the size to the page margins is illustrated in B.2. Please notice the text size that is shrunken in order for the table to adjust to the page:

Table B.2: Sample Table.

URL	First Time Visit	Last Time Visit	URL Counts	Value	Reference
https://web.facebook.com/	1521241972	1522351859	177	56640	[facebook-2021]
http://localhost/phpmyadmin/	1518413861	1522075694	24	39312	database-management
https://mail.google.com/mail/u/	1516596003	1522352010	36	33264	Google-Gmail-2021
https://github.com/shawon100	1517215489	1522352266	37	27528	Code-Repository
https://www.youtube.com/	1517229227	1521978502	24	14792	Youtube-video-2021

[†] Value kept proportional to the critical path: (5) = (4) * 1.5.

 $^{^{\}ddagger}$ A size of x corresponds to a 2^x point FFT.