Introduction to Functional Programming, or:

How I Learned to Stop Worrying and Love Referential Transparency

George Wilson

Ephox

george.wilson@ephox.com

March 22, 2017

Me

- ► Graduated from Griffith in 2014
- ▶ Working at Ephox ever since
- Assistant organiser of Brisbane FP Group



This talk

- ▶ Beginner talk you should be able to understand it!
- ► Ask lots of questions
- ► Let's hang out after the talk

This talk

- ► You won't learn functional programming tonight
- ► You should learn what it is and hopefully why you should care
- ▶ The aim is to motivate you and provide the tools to teach yourself

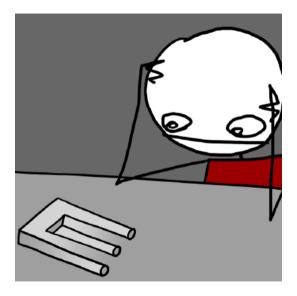


Programming is hard

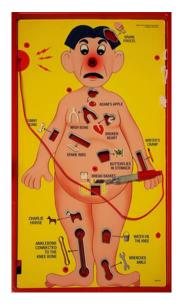
Programming is hard

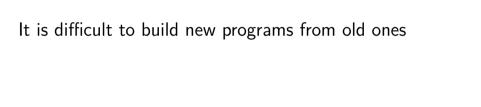
Why?

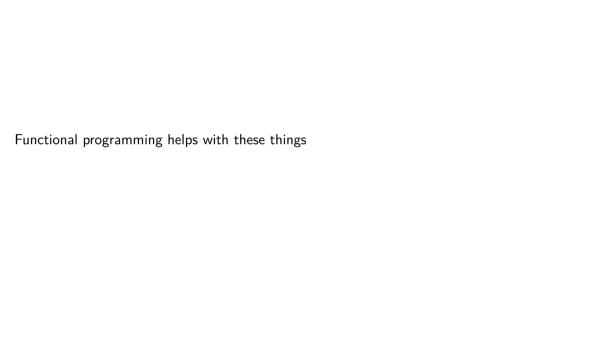
Programs are difficult to understand



Programs are difficult to modify







What is functional programming?





 $Functional\ programming\ is\ programming\ with\ functions.$



A function relates its input to a result and does nothing else.

FUNCTION f: OUTPUT f(x)

$$f(x) = x + 5$$

By definition, a function can not:

- ► Modify or depend on external state
- ▶ Perform I/O (such as reading a file or printing to the screen)
- ► Throw an exception

We call those things side effects.

Some people call functions pure functions	
to differentiate them from "functions" with side effects	

Useful properties of functions:

▶ Passing the same input *always* gives the same output

Useful properties of functions:

- ▶ Passing the same input *always* gives the same output
- ➤ You do not need to understand the environment in which the function is run to understand its result

Useful properties of functions:

- ▶ Passing the same input *always* gives the same output
- ➤ You do not need to understand the environment in which the function is run to understand its result
- Functions can be composed to build bigger functions

Examples of functions

```
def add5(x):
    return x + 5
```

Examples of functions

```
def add5(x):
    return x + 5

def greet(name):
    return f"Hello, {name}"
```

Example of not a function

```
count = 0

def countedGreet(name):
   global count
   count = count + 1
   s = f"Hello, {name}. I've greeted {count} people!"
   return s
```

Example of not a function

```
count = 0
def countedGreet(name):
  global count
  count = count + 1
  s = f"Hello, {name}. I've greeted {count} people!"
  return s
g = countedGreet("George")
h = countedGreet("George")
# g = "Hello, George. I've greeted 1 people!"
# h = "Hello, George, I've greeted 2 people!"
```

Turning it into a function

```
def countedGreet2(name, count):
   newCount = count + 1
   s = f"Hello, {name}. I've greeted {count} people!"
   return (s, newCount)
```

Turning it into a function

```
def countedGreet2(name, count):
  newCount = count + 1
  s = f"Hello, {name}. I've greeted {count} people!"
  return (s. newCount)
(g,n) = countedGreet2("George", 0)
(h,m) = countedGreet2("George", 0)
# g = "Hello, George. I've greeted 1 people!"
# h = "Hello, George. I've greeted 1 people!"
```

Turning it into a function

```
def countedGreet2(name, count):
  newCount = count + 1
  s = f"Hello, {name}. I've greeted {count} people!"
  return (s. newCount)
(g,n) = countedGreet2("George", 0)
(h,m) = countedGreet2("George", n)
# g = "Hello, George. I've greeted 1 people!"
# h = "Hello, George. I've greeted 2 people!"
```

Removing side e	ffects makes our	types reveal m	ore information

countedGreet2 :: (String, Int) -> (String, Int)

countedGreet :: String -> String

You've seen how to remove mutable state by passing more parameters Can we remove exceptions somehow? What about null?	

Say we're writing a function to parse a String into an Int

parse("34") ==> 34

Say we're writing a function to parse a String into an $\ensuremath{\mathsf{Int}}$

parse("34") ==> 34

F-----

parse("nope") ==> ?

Say we're writing a function to parse a String into an Int

parse("34") ==> 34

parse("nope") ==> ?

This function takes a String and either returns an Int or throws an exception

```
So the type is
String -> Int

let x: Int = parse("nope")
x + 5
```

Instead,	let's use a type that is either an error or an Int
String	-> Result <int></int>

Instead, let's use a type that is either an error or an ${\tt Int}$

```
enum Result < A > {
  case Error(String)
  case Success(A)
```

String -> Result<Int>

let r: Result < Int > = parse("34")

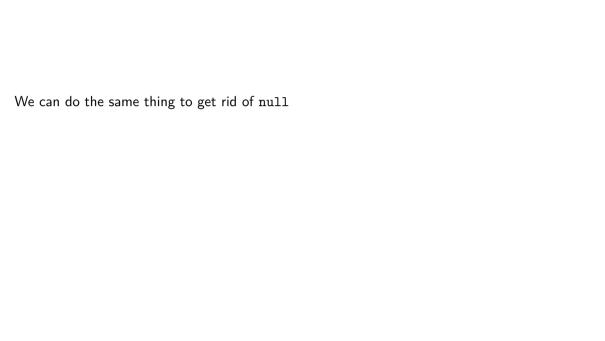
```
let r: Result < Int > = parse("34")
```

That doesn't compile! We can't forget to check for the exception!

r + 5

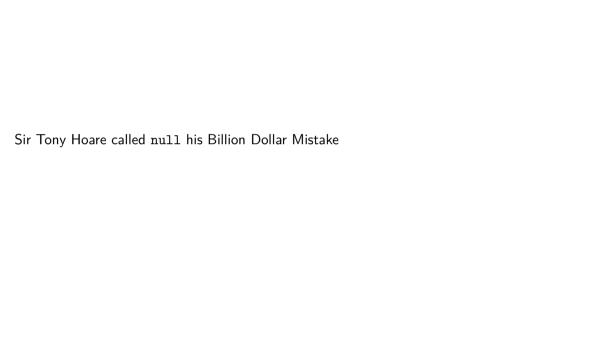
```
let r: Result < Int > = parse("34")

switch r {
   case Error(message):
     // handle the message somehow
   case Success(n):
     n + 5
```

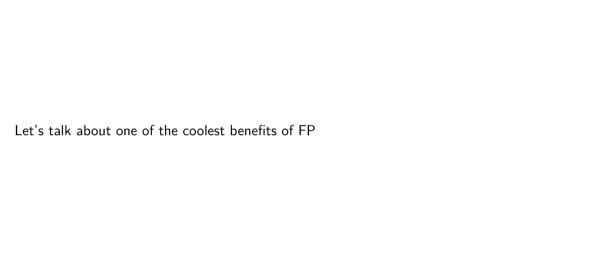


We can do the same thing to get rid of null

```
enum Optional <A> {
  case Nothing
  case Something(A)
```







```
def p():
   proc(expression, expression)
```

```
def p():
    x = expression
    proc(x, x)
```

```
def print2(s, t):
    print(s)
    print(t)

def strpopthen():
    s = ['a','b','c','d','e','f']
    print2(s[5],s[5])

# 'f'
# 'f'
```

```
def print2(s, t):
    print(s)
    print(t)

def strpopthen():
    s = ['a','b','c','d','e','f']
    x = s[5]
    print2(x,x)
```

'f'

```
def print2(s, t):
    print(s)
    print(t)

def listpopthen():
    s = ['a','b','c','d','e','f']
```

print2(s.pop(), s.pop())

'f' # 'e'

```
def print2(s, t):
    print(s)
    print(t)

def listpopthen():
    s = ['a','b','c','d','e','f']
    x = s.pop()
    print2(x,x)
```

'f' # 'f'

Referential Transparency

An expression e in a program is $referentially\ transparent$ if and only if replacing all occurrences of e with its value does not change the observable behaviour of the program

Referential Transparency

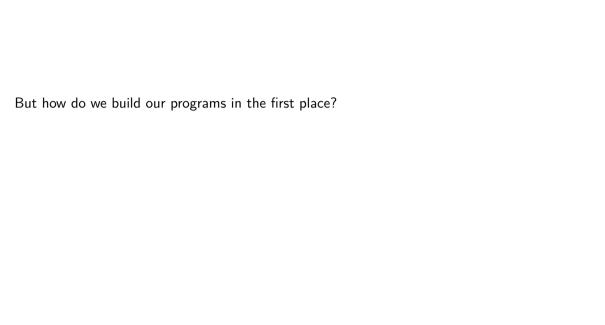
Referential transparency enables equational reasoning

Referential Transparency

Referential transparency enables equational reasoning

Equational reasoning is the process of substituting equals for equals and knowing that you are not altering the result of the program.

We can fearlessly modify our programs!



But how do we build our programs in the first place?	
No define data types (like Regult (AN)	

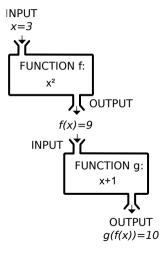
We define data-types (like Result<A>) and define functions on them

and then just combine the functions

Higher-order functions let us build new functions from others

- A higher-order function:
 - takes a function as input
 - produces a function as output
 - or both!

An example is function composition



```
def compose(f,g):
    return lambda x: f(g(x))
```

```
def compose(f,g):
    return lambda x: f(g(x))
```

def excited(s):
 return s + "!"

def loud(s):

return s.upper()

```
def excited(s):
    return s + "!"

def loud(s):
    return s.upper()
```

talkingAboutFunctionalProgramming = compose(excited, loud)

def compose(f,g):

return lambda x: f(g(x))

Another higher-order function is map

map runs a function on every element of a list

Another higher-order function is map map runs a function on every element of a list

```
def embiggen(x):
    return x * 10

bigNumbers = map(embiggen, numbers)
# [10,20,30,40,50]
```

numbers = [1,2,3,4,5]

Higher-order functions let us build larger programs from smaller ones

We call this *modularity*

Summary

Pure functions give benefits

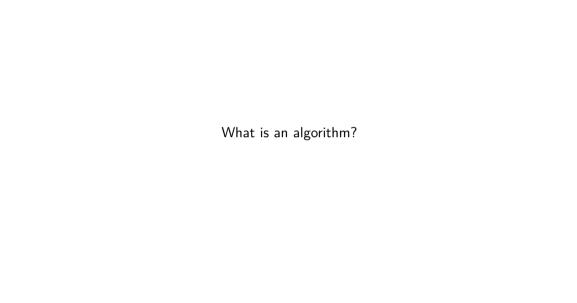
- ▶ It's easier to understand what a function does
- It's easier to refactor functions (and not break the program)
- Pure functions can always be run in parallel
- ▶ The type system helps you more

Referential transparency and higher-order functions give rise to modularity and reuse.

Necessary evil?

"we only use side effects when they are necessary" $% \left(1\right) =\left(1\right) \left(1\right) \left($

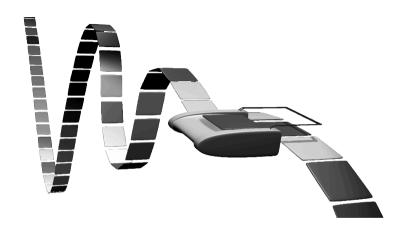
"the algorithm we're using is inherently stateful"



In the 1930's, **Alan Turing** invented Turing Machines

to answer this question



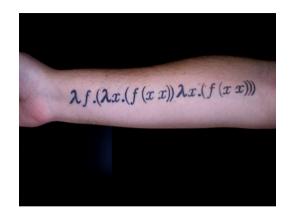


Turing machines compute by reading and writing to a tape

Computation is done by this *modification* of the tape

At the same time, Alonzo Church invented $\lambda \ {\rm calculus}$





Lambda calculus is pure functions and nothing else

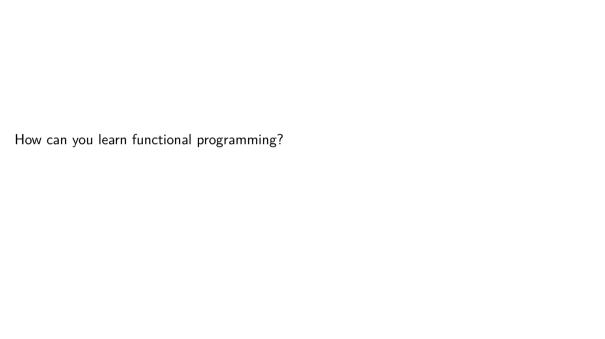
Computation is done by function application

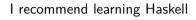




Church and Turing discovered that λ calculus and Turing Machines $are \ equivalent$

Anything computable with side effects can be computed with only pure functions







Haskell lectures
https://www.seas.upenn.edu/%7Ecis194/spring13
https://github.com/bfpg/cis194-yorgey-lectures

Books

http://www.haskellbook.com http://http://learnyouahaskell.com/

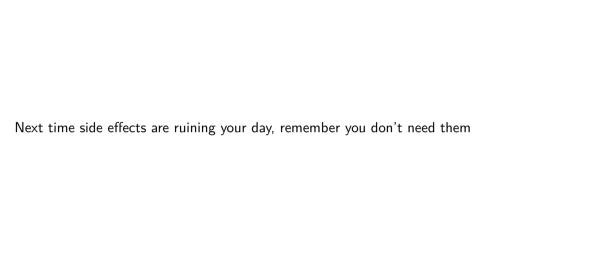
Nothing in this talk is specific to Haskell!

Other languages with excellent support for pure functional programming:

- Purescript
- ▶ Idris

and other languages with support for functional programming:

- OCaml
 - Scala
 - Rust
 - Swift
 - ► C#
 - Java
 - ► Many others!





"If you are scared $[\dots]$ go to Church!" - Ice Cube



Thanks for listening!