

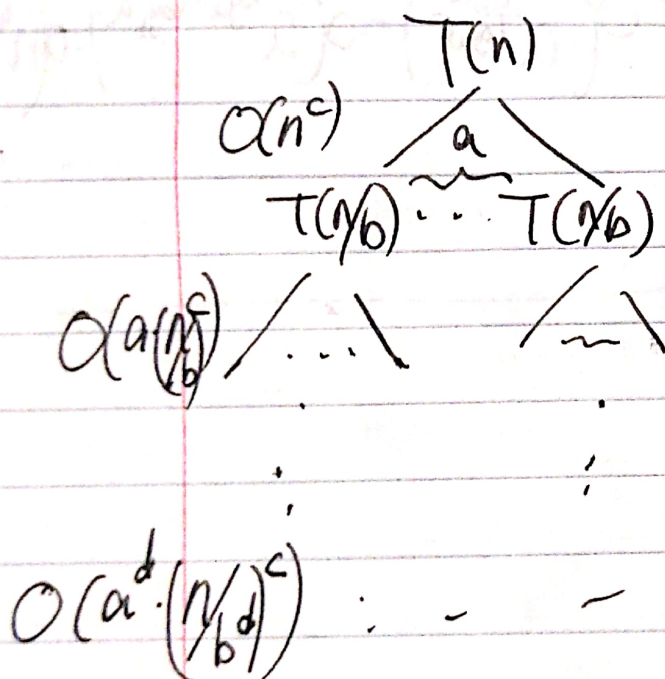
# Discussion 2c Week 7 notes Yijia Chen

## 1. Complexity analysis. Big-O notation.

$\text{for (int } i=1; i \leq n; i++)$   $O(n)$   
 $\text{for (int } i=1; i \leq n; i+=t)$   $O(n)$  if  $t$  is a constant  
 $\text{for (int } i=1; i \leq n; i*=2)$   $O(\log n)$   
 $\text{for (int } i=1; i \leq n; i/=2)$   $O(\log n)$ .

## Recursive Relations.

$$T(n) = a(T(n/b)) + O(n^c)$$



ensures  $n/b^d \approx 1$ .  
 $d \approx O(\log n)$

$$\begin{aligned}
 T(n) &= O\left(n^c + a(n/b)^c + \dots + a^d(n/b^d)^c\right) \\
 &= O\left(n^c \left(1 + \frac{a}{b^c} + \left(\frac{a}{b^c}\right)^2 + \dots + \left(\frac{a}{b^c}\right)^d\right)\right).
 \end{aligned}$$

Case 1:  $\frac{a}{b^c} < 1$ .  $T(n) = O(n^c)$

Case 2:  $\frac{a}{b^c} = 1$ .  $T(n) = O(n^c \log_b n)$ .

Case 3:  $\frac{a}{b^c} > 1$ .  $T(n) = O\left(n^c \cdot \left(\frac{a}{b^c}\right)^{\log_b n}\right)$ .

$$= O\left(n^c \cdot a^{\log_b n} \cdot b^{-c \log_b n}\right)$$

$$= O\left(n^c (b^{\log_b n})^c \cdot a^{\log_b n}\right)$$

$$= O\left(n^c \cdot n^c \cdot a^{\log_b n}\right) = O\left(a^{\log_b n}\right)$$

$$= O\left(a^{\frac{\log_a n}{\log_a b}}\right) = O\left(a^{\log_a n \log_b a}\right) = O\left(n^{\log_b a}\right)$$

Master Theorem



## 2. Sorting algorithms .

- Merge Sort:  $T(n) = 2T(n/2) + O(n)$ .

$$O(n \log n)$$

- Insertion Sort:  $O(n^2)$ . Locality helps  $O(n)$

- Bubble Sort:  $O(n^2)$ . Best  $O(n)$

- Selection Sort:  $O(n^2)$

- Quick Sort: Average  $O(n \log n)$  worst  $O(n^2)$ .  
(Quick Sort with Insertion Sort  $O(n \log n)$ )  
binary search  $O(\log n)$

Insertion Sort with binary search  $O(n^2)$ .