

# Discussion 2c Week 7 Notes Yiqi Chen

## 1. Complexity analysis. Big-O notation.

for (int i=1; i≤n; t+i)  $O(n)$

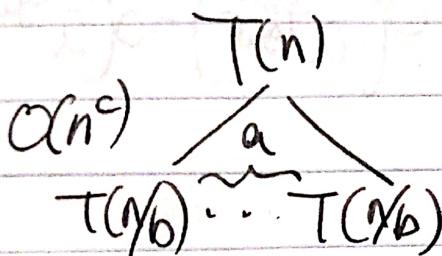
for (int i=1; i≤n; i+=t)  $O(n)$  if t is a constant

for (int i=1; i≤n; i\*=2)  $O(\log n)$

for (int i=1; i≤n; n/=2)  $O(\lg n)$ .

## Recursive Relations.

$$T(n) = a(T(n/b)) + O(n^c)$$



$$O(a(n/b^d)) / \dots \backslash / n$$

$$O(a^d \cdot (n/b^d)^c) : - - -$$

unless  $n/b^d \approx 1$ .

$$d \approx O(\log n)$$

$$\begin{aligned} T(n) &= O\left(n^c + a(n/b)^c + \dots + a^d(n/b^d)^c\right) \\ &= O\left(n^c \left(1 + \frac{a}{b^c} + \left(\frac{a}{b^c}\right)^2 + \dots + \left(\frac{a}{b^c}\right)^d\right)\right). \end{aligned}$$

Case 1:  $\frac{a}{b^c} < 1$ .  $T(n) = O(n^c)$

Case 2:  $\frac{a}{b^c} = 1$   $T(n) = O(n^c \log n)$ .

Case 3:  $\frac{a}{b^c} > 1$   $T(n) = O\left(n^c \cdot \left(\frac{a}{b^c}\right)^{\log_b n}\right)$ .

$$= O\left(n^c \cdot a^{\log_b n} \cdot b^{-c \log_b n}\right)$$

$$= O\left(n^c (b^{\log_b n})^c \cdot a^{\log_b n}\right)$$

$$= O\left(n^c \cdot n^{c - \log_b n} \cdot a^{\log_b n}\right) = O(a^{\log_b n})$$

$$= O\left(a^{\frac{\log n}{\log ab}}\right) = O\left(a^{\frac{\log n \log_b a}{\log ab}}\right) = O\left(n^{\log_b a}\right)$$

Master Theorem

## 2 | Sorty algorithms .

- Merge Sort .  $T(n) = 2T(n/2) + O(n)$ .

$O(n \log n)$

- Insertion Sort :  $O(n^2)$ . Locality helps  $O(cn)$

- Bubble Sort :  $O(n^2)$ . Best  $O(n)$

- Selection Sort :  $O(n^2)$

- Quick Sort: Average  $O(n \log n)$  Worst  $O(n^2)$ .

Quick Sort with Insertion Sort  $O(n \log n)$   
binary search  $O(\log n)$

Insertion Sort with binary Search

$O(n^2)$ .