

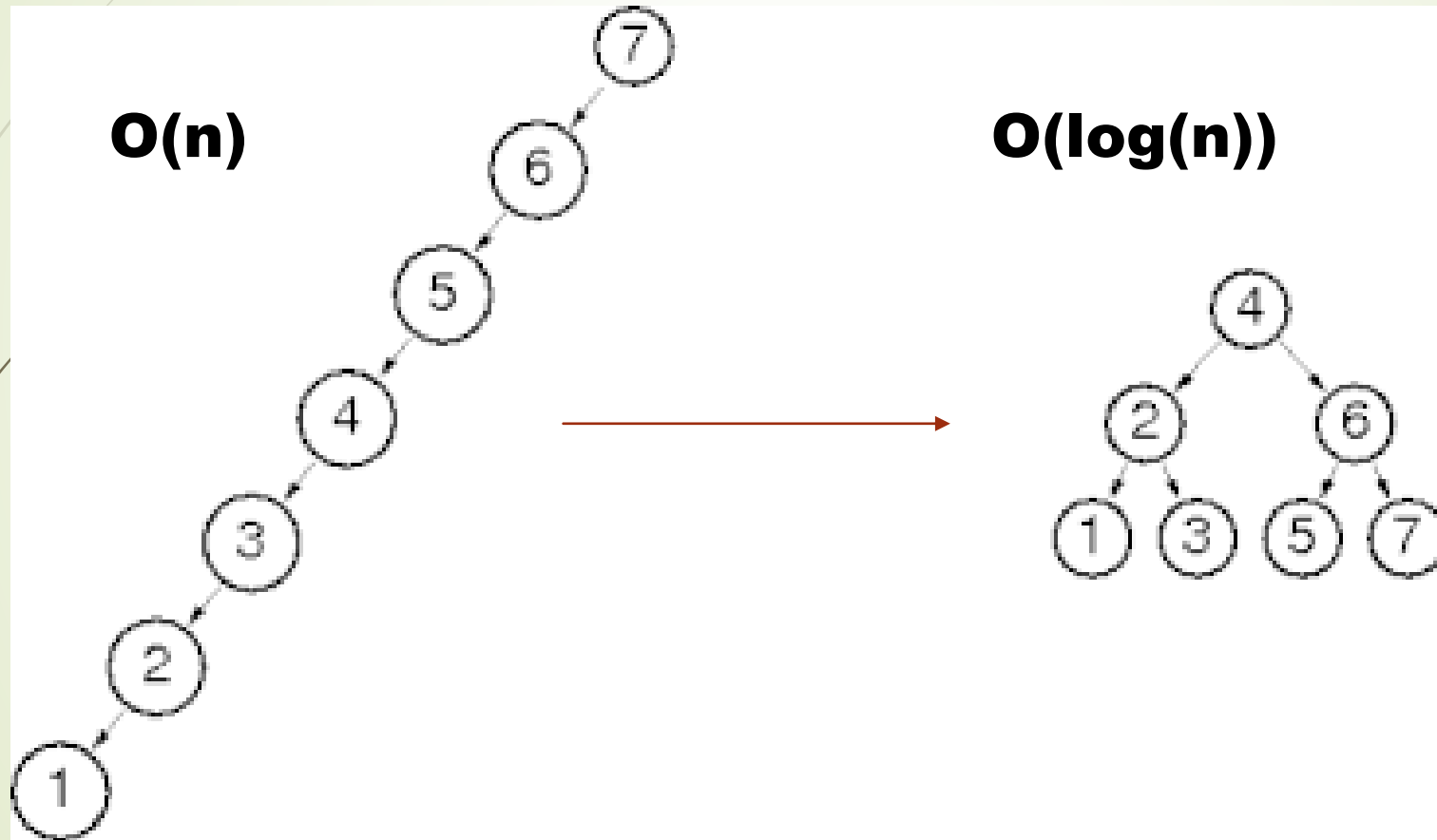


AVL Tree

Pittsford Sutherland Programming Club

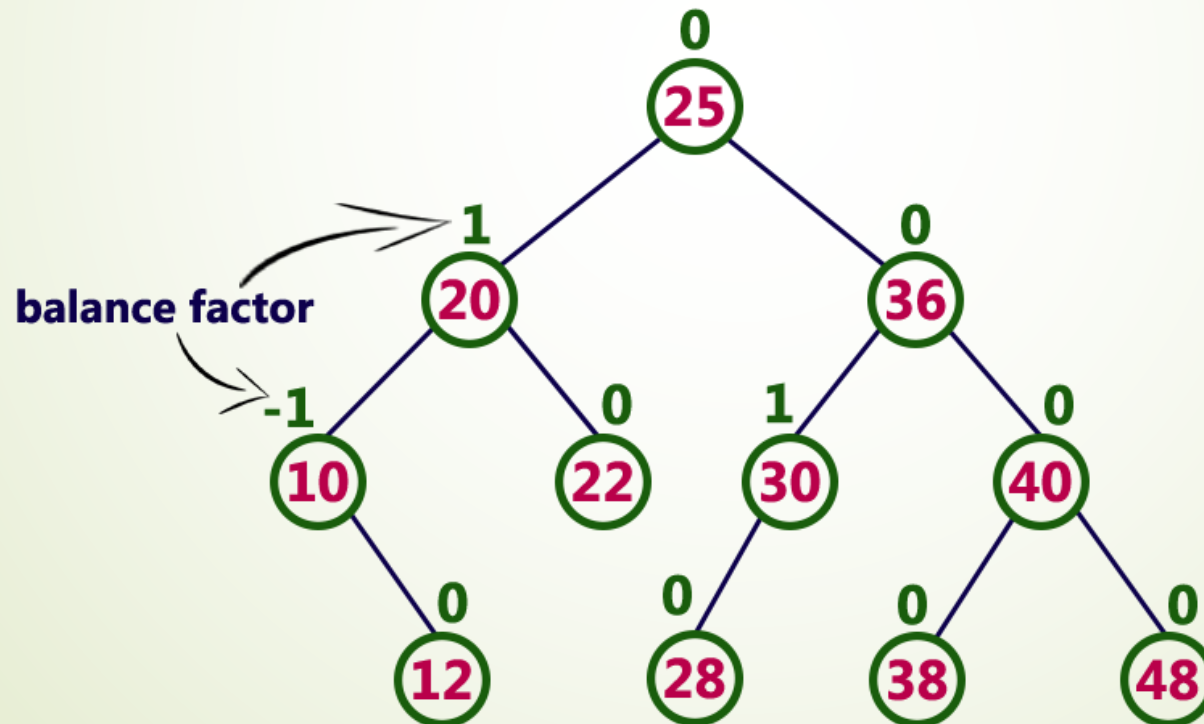
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What is "Balance"?



AVL Tree

- A binary search tree with a balance condition.
- The height of left subtree and the height of right subtree differs at most 1.
- The height of an empty tree is defined to be -1.
- Balance factor saves the difference between the height of two subtrees

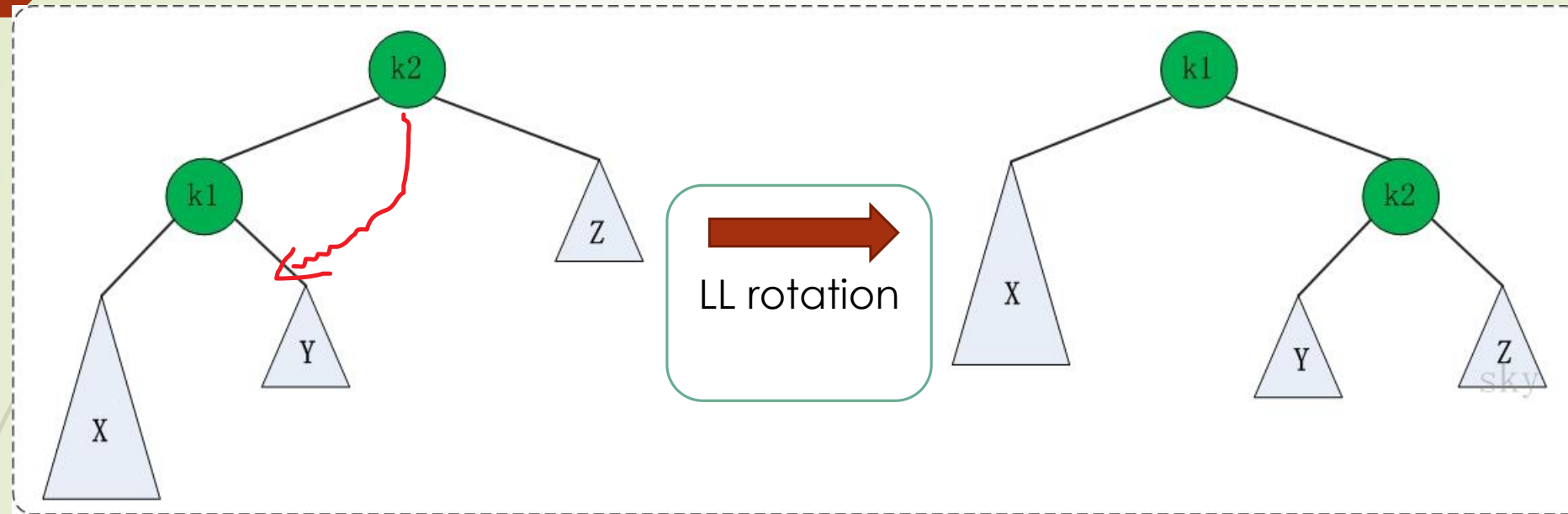




How to keep it “Balanced”?

Use Rotations!

Left-Left(LL) Rotation

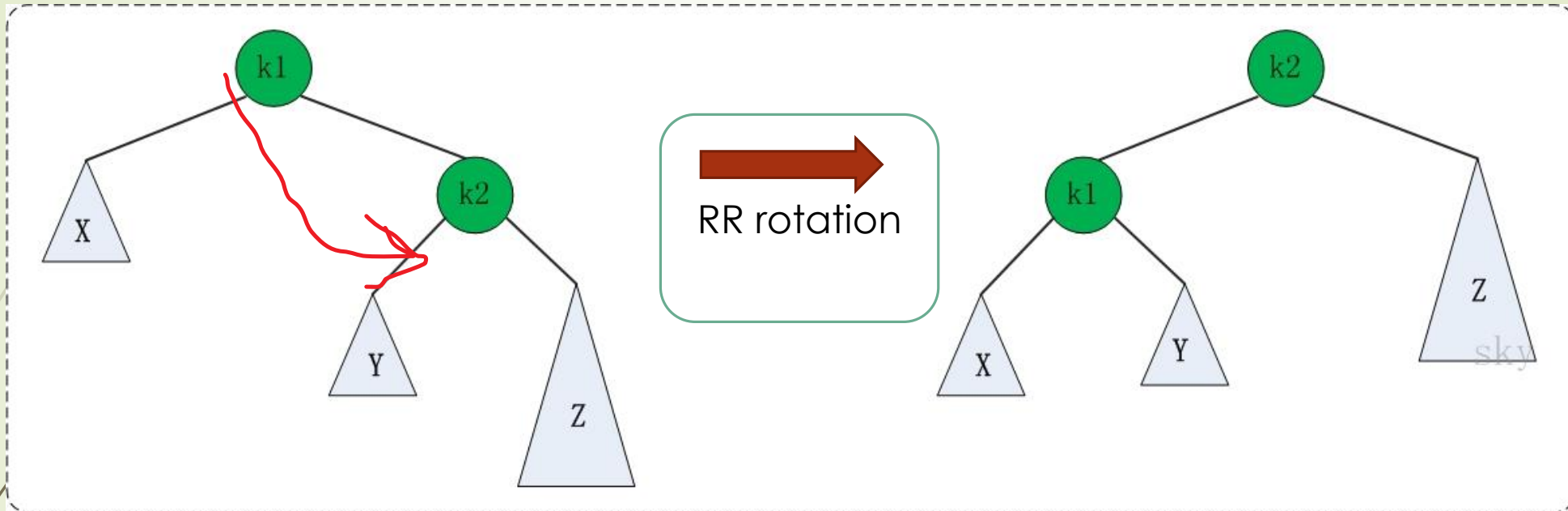


Condition: When root's left subtree is greater than its right subtree, and its left son's left subtree is greater than its left son's right subtree. ($H(k1) > H(Z)$, $H(X) > H(Y)$)

How it works: We want to rotate $k2$ to be the right son of $k1$ and make $k1$ the root in order to make the tree more balanced.

We cannot directly make $k2$ the right son of $k1$ since $k1$ has already had its right subtree Y . According to the property of binary search tree, all the elements in subtree Y must be less than $k2$, so we can make Y be the left subtree of $k2$. Also, all the elements in Y , $k2$, and Z are greater than $k1$, so we can put all these elements in the right subtree of $k1$, and remain X unchanged.

Right-Right(RR) Rotation

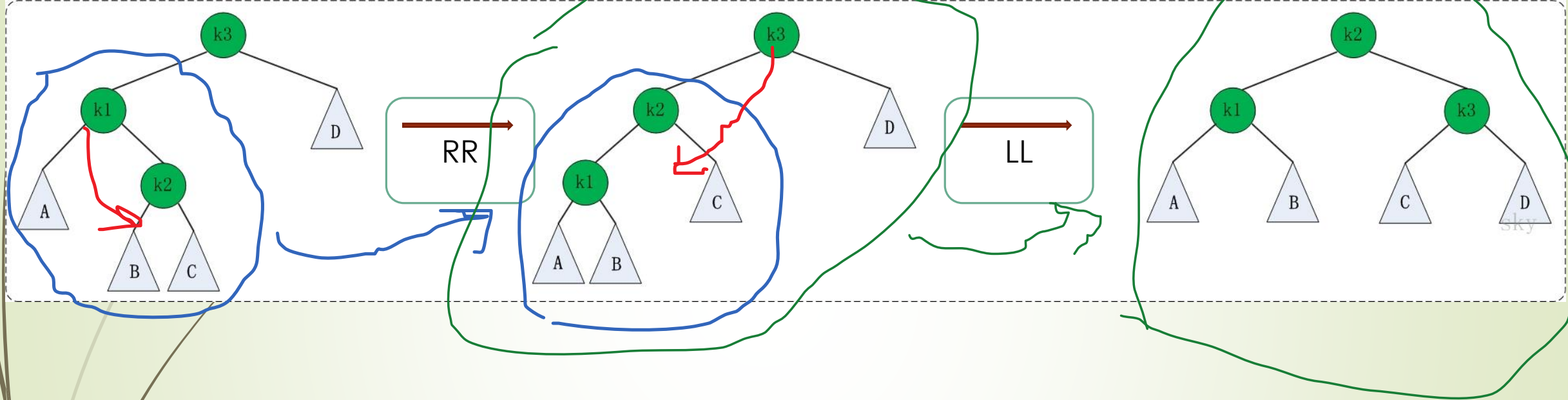


Condition: When root's right subtree is greater than its left subtree, and its right son's right subtree is greater than its right son's left subtree. ($H(k2) > H(X)$, $H(Z) > H(Y)$)

How it works: We want to rotate $k1$ to be the left son of $k2$ and make $k2$ the root in order to make the tree more balanced.

We cannot directly make $k1$ the left son of $k2$ since $k2$ has already had its left subtree Y . According to the property of binary search tree, all the elements in subtree Y must be greater than $k1$, so we can make Y be the right subtree of $k1$. Also, all the elements in Y , $k1$, and X are less than $k2$, so we can put all these elements in the left subtree of $k2$, and remain Z unchanged.

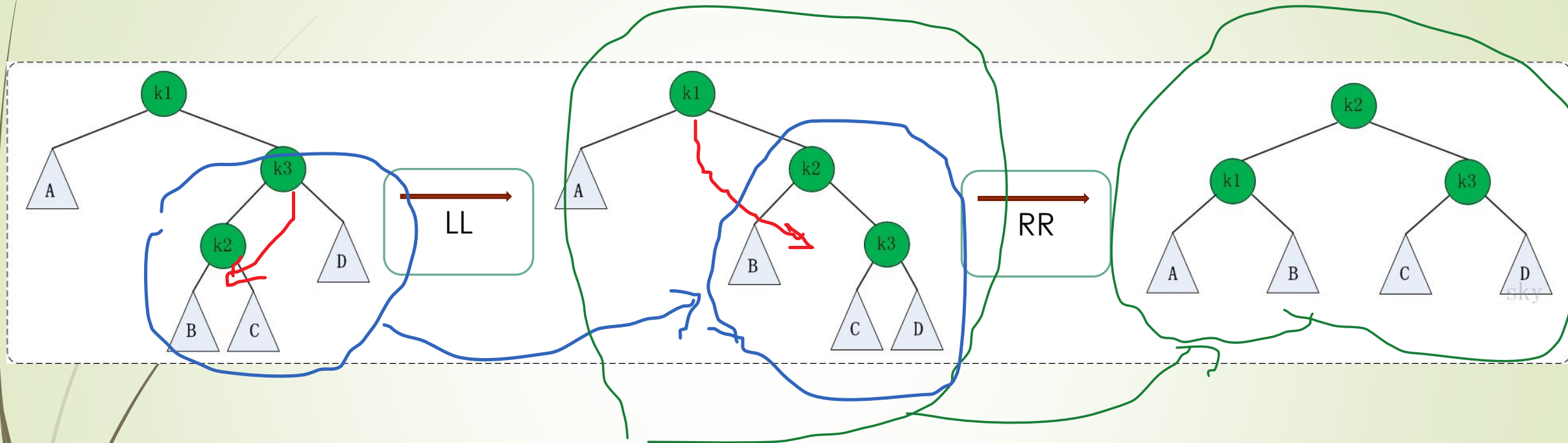
Left-Right(LR) Rotation



Condition: When root's left subtree is greater than its right subtree, and its left son's right subtree is greater than its left son's left subtree. ($H(k2) > H(D)$, $H(k2) > H(A)$)

How it works: First make root's left subtree balance by using RR rotation.
Then make whole tree balance by using LL rotation.

Right-Left(RL) Rotation





Condition: When root's right subtree is greater than its left subtree, and its right son's left subtree is greater than its right son's right subtree. ($H(k3) > H(A)$, $H(k2) > H(D)$)

How it works: First make root's right subtree balance by using LL rotation.
Then make whole tree balance by using RR rotation.



Build an AVL Tree

```
Struct NODE{  
    Tree right;  
    Tree Left;  
    int element;  
    int height; // -1 if it's an empty node, otherwise the height;  
}
```



ALL the operations(Insert, remove, search) are the same as the binary search tree; Only need to add the balance process to each operation. Make sure to look at the C++ code.