

1 Categories

Definition 1.1 (Category). A category, \mathcal{A} is defined to have each of the following,

- (i) A collection of objects, denoted by $\text{ob}(\mathcal{A})$ and written $A, B, C \in \mathcal{A}$.
Such that each object has an ‘identity’, $1_A \in \mathcal{A}(A, A), 1_B \in \mathcal{A}(B, B), 1_C \in \mathcal{A}(C, C)$
- (ii) For each pair of objects, a collection of ‘links’/morphisms between them, denoted by $\mathcal{A}(A, B)$ and written as $f \in \mathcal{A}(A, B), g \in \mathcal{A}(B, C)$. Such that,
 - (a) morphisms with matching domain,co-domain can be ‘chained’/composed $(g, f) = g \circ f$
 - (b) with this composition being associative, $(h \circ g) \circ f = h \circ (g \circ f)$
 - (c) and they are ‘fixed’ by the identity $f \circ 1_A = f = 1_B \circ f$

Example 1.1. Non-trivial Identity Consider the objects to be groups, and morphisms to be direct product between them:

- i $\text{ob}(\mathcal{A}) = \{G \mid G \text{ is a group}\}$
- ii $\mathcal{A}(A, B) := A \times B$
- iii $\mathcal{A}(B, C) \circ \mathcal{A}(A, B) \mapsto \mathcal{A}(A, C)$

So, there’s a unique morphism between any two objects i.e groups. And the identity morphism,

$$\forall A, B \in \mathcal{A}, \text{ if } f \in \mathcal{A}(A, B), \text{ then } f \circ 1_A \in \mathcal{A}(A, B) \times \mathcal{A}(A, A) \mapsto \mathcal{A}(A, B) \text{ and } 1_B \circ f \in \mathcal{A}(B, B) \times \mathcal{A}(A, B) \mapsto \mathcal{A}(A, B)$$

Thus, $\text{ob}(\mathcal{A})$ along with \circ is actually a group. And hence has a unique inverse. [But how exactly?](#)

Example 1.2. Set The objects are defined to be sets, and morphisms are the functions between them, with the usual composition law:

- i $\text{ob}(\mathcal{A}) = \{S \mid S \text{ is a set}\}$
- ii $(f : A \mapsto B) \in \mathcal{A}(A, B)$
- iii $(g \in \mathcal{A}(B, C)) \circ (f \in \mathcal{A}(A, B)) \mapsto g(f) \in \mathcal{A}(A, C)$

Example 1.3. Grp Objects are groups, with homomorphisms between them being the morphisms, and composition being as usual:

- i $\text{ob}(\mathcal{A}) = \{G \mid G \text{ is a group}\}$
- ii $\mathcal{A}(A, B) = \text{Hom}(A, B)$ i.e. all f such that $\forall x, y \in A, f((x) \cdot_A (y)) = (f(x)) \cdot_B (f(y))$
- iii composition is defined as that between two group homomorphisms

In this example, the set of all morphisms along with composition forms a group.

Example 1.4. Ring Objects are rings, and arrows are ring homomorphisms between them.

- i $\text{ob}(\mathcal{A}) = \{G \mid G \text{ is a ring}\}$
- ii $\mathcal{A}(A, B) = \text{Hom}(A, B)$
- iii composition is defined as that between two ring homomorphisms

Example 1.5. Vect_k Objects are vector spaces *over field k*, and the morphisms between them are linear transformations

- i $\text{ob}(\mathcal{A}) = \{A \mid A \text{ is a vector space}\}$
- ii $\mathcal{A}(A, B) = \mathcal{L}(A, B)$
- iii composition is defined as that of linear transformations

Definition 1.2 (Isomorphism). An isomorphism, between objects, is a morphism between them such that it's 'inverse' is also a morphism. So,

$$f : A \mapsto B \text{ is an isomorphism} \iff \exists g \in \mathcal{A}(B, A) : gf = 1_A \text{ and } fg = 1_B$$

Definition 1.3 (Product Category). Somewhat like a cartesian product of categories. Given categories \mathcal{A} and \mathcal{B} , $\mathcal{A} \times \mathcal{B}$ is defined as:

- i $ob(\mathcal{A} \times \mathcal{B}) := ob(\mathcal{A}) \times ob(\mathcal{B})$
- ii $(\mathcal{A} \times \mathcal{B})((A, B), (A', B')) := \mathcal{A}(A, A') \times \mathcal{B}(B, B')$
- iii $(f, g) \in \mathcal{A} \times \mathcal{B}((A, B), (C, D))$, $(a, b) \in \mathcal{A} \times \mathcal{B}((C, D), (E, F)) \implies (a, b) \circ (f, g) := (a \circ f, b \circ g)$
- iv $\forall (A, B) \in ob(\mathcal{A} \times \mathcal{B})$, $1_{(A, B)} := (1_A, 1_B)$

Example 1.6 (CAT). The category of all categories with morphisms being functors.

- i $ob(\mathcal{A}) = \{A | A \text{ is a category}\}$
- ii $\mathcal{A}(A, B) = F(A, B)$
- iii $F : \mathcal{A} \mapsto \mathcal{B}$, $G : \mathcal{B} \mapsto \mathcal{C} \implies G \circ F := H : \mathcal{A} \mapsto \mathcal{C}$

And thus, the identity of \mathcal{A} is the functor, $1_{\mathcal{A}} : \mathcal{A} \mapsto \mathcal{A}$

2 Functors

Definition 2.1 ((Covariant)Functor). A functor is a map between categories, written $F : \mathcal{A} \mapsto \mathcal{B}$, consists :

- (i) function taking objects of \mathcal{A} to those of \mathcal{B} i.e. $ob(\mathcal{A}) \mapsto ob(\mathcal{B})$. Written as $A \mapsto F(A)$.
- (ii) associative, identity-preserving function taking links between objects of \mathcal{A} to those for \mathcal{B} , $f \mapsto F(f)$, i.e.
 $\forall A, B \in \mathcal{A}, \mathcal{A}(A, B) \mapsto \mathcal{B}(F(A), F(B))$ such that (a) $f \in \mathcal{A}(A, B), g \in \mathcal{A}(B, C) \implies F(g \circ f) = F(g) \circ F(f) = F(g \circ f)$
(b) $A \in \mathcal{A} \implies F(1_A) = 1_{F(A)}$

Example 2.1. Forgetful Functors They essentially ignore some of the structure of the 'domain'.

- (a) $U : Grp \mapsto Set$ takes groups to their underlying set, and homomorphisms to maps between the sets. Similarly, $Ring \mapsto Set$ and $Vect_k \mapsto Set$
- (b) Let Ab be the category of abelian groups, then, $U : Ring \mapsto Ab$ takes rings to their additive group, 'forgetting' the multiplicative group. And if Mon is the category of monoids, $U : Ring \mapsto Mon$ 'forgets' the additive group.
- (c) $U : Ab \mapsto Grp$ just takes each abelian group to itself, and does the same for (homo)morphisms.

Example 2.2. Free Functors

- (a) let $F(S)$ denote the free group on a set S . Then, $U : Set \mapsto Grp$ is a 'free' functor taking sets to their free group, and thus the maps between them becomes a homomorphism between their free groups.
- (b) Same for Rings pp20
- (c) Same for Vector Spaces

Example 2.3. Let \mathcal{G}, \mathcal{H} be the one object categories of monoids G, H respectively. Then, due to composition being associative and identity preseving, possible functors are precisely the homomorphisms.

Definition 2.2 (Contravariant Functor). For categories \mathcal{A} and \mathcal{B} , $\mathcal{A}^{op} \mapsto \mathcal{B}$ is a contravariant functor from \mathcal{A} to \mathcal{B} .

Example 2.4. Let k be a field and V, W be vector spaces over it. Then fixing W ,

$$\forall f \in Hom(V, V'), \exists f^* : Hom(V', W) \mapsto Hom(V, W)$$

So, for each $V \in ob(Vect_k)$, $Hom(V, W)$ defines a contravariant functor from $Vect_k^{op}$ to $Vect_k$

3 Natural Isomorphisms

Example 3.1.