

# CS 547: Foundation of Computer Security

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# *Previous Class*

- Access Control
  - Discretionary Access Control

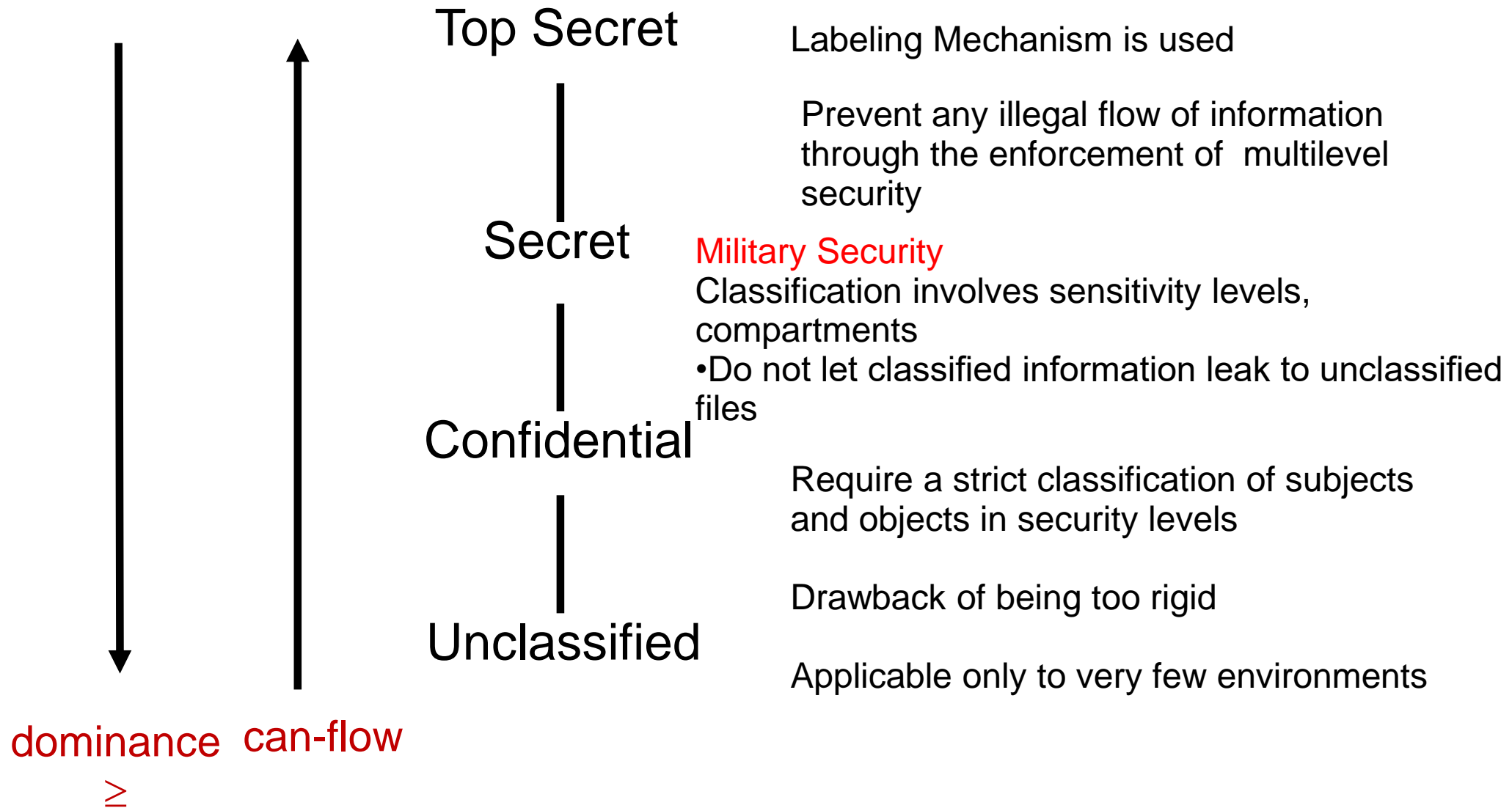
# Present class

- Access Control
  - Mandatory Access Control
  - Role-Based Access Control

# Mandatory Access Control (MAC)

- Defined by three major properties:
  - Administratively-defined security policy
  - Control over all subjects (process) and objects (files, sockets, network interfaces)
  - Decisions based on all security-relevant info
- **MAC**
  - by assigning security levels to users and objects'
  - Access to an object is granted only if the security levels of the subject and the object satisfy certain constraints.
- The MAC pattern is also known as multilevel security model and lattice-based access control.

# Mandatory Access Control (MAC)



# Bell-LaPadula Model: Multi-level Security

- Introduced in 1973
- Air Force was concerned with security in time-sharing systems
  - Many OS bugs
  - Accidental misuse
- Main Objective:
  - Enable one to formally show that a computer system can securely process classified information

# The BLP Security Policy

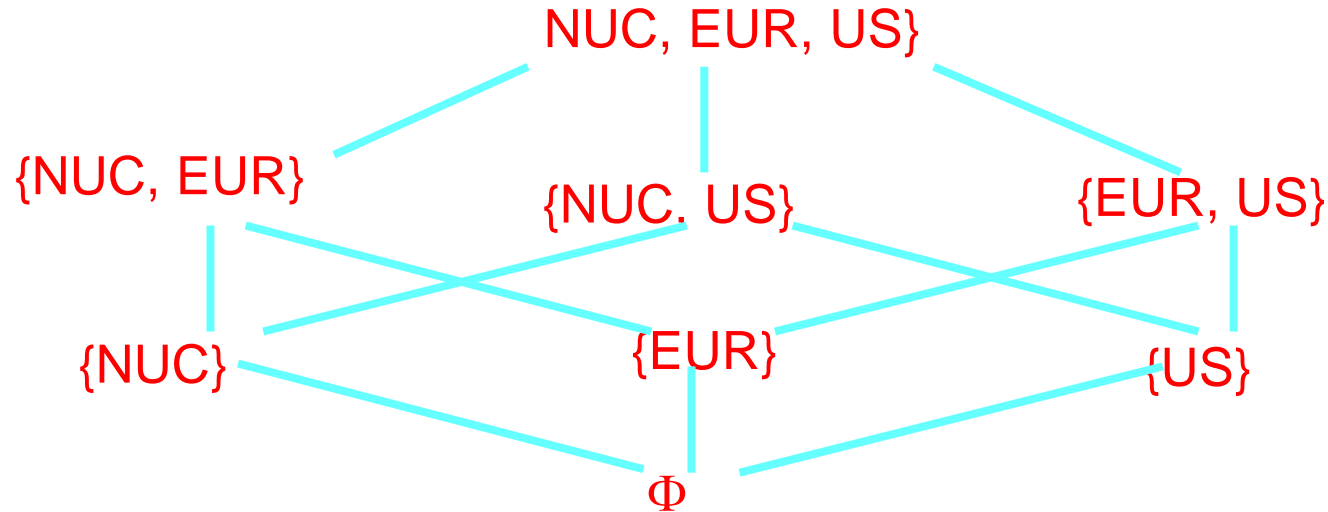
- A state is secure if it satisfies
  - Simple Security Condition (no read up):
    - $S$  can read  $O$  iff  $L_m(S) \geq L(O)$
  - The Star Property (no write down): for any  $S$  that is not trusted
    - $S$  can read  $O$  iff  $L_c(S) \geq L(O)$  (no read up)
    - $S$  can write  $O$  iff  $L_c(S) \leq L(O)$  (no write down)
  - Discretionary-security property
    - every access is allowed by the access matrix
- A system is secure if and only if every reachable state is secure.

# Categories and Need to Know Principle

- Expand the model by adding a set of categories.
  - Each category describe a kind of information.
  - These category arise from the “need to know” principle
    - “ no subject should be able to read objects unless reading them is necessary for that subject to perform its function.”
- Example: three categories: NUC, EUR, US.
- Each security level and category form a security level or compartment.
- Subjects *have clearance at* (are cleared into, or are in) a security level.
- Objects are *at the level of* (or are in) a security level.



# Security Lattice



- William may be cleared into level  $(S, \{EUR\})$
- George into level  $(TS, \{NUC, US\})$ .
- A document may be classified as  $(C, \{EUR\})$
- Someone with clearance at  $(TS, \{NUC, US\})$  will be denied access to document with category EUR.

# Dominate (dom) Relation

- The security level  $(L, C)$  dominates the security level  $(L', C')$  if and only if  $L' \leq L$  and  $C' \subseteq C$
- $\neg \text{Dom} \rightarrow$  dominate relation is false.
- George is cleared into security level  $(S, \{\text{NUC}, \text{EUR}\})$
- DocA is classified as  $(C, \{\text{NUC}\})$
- DocB is classified as  $(S, \{\text{EUR}, \text{US}\})$
- DocC is classified as  $(S, \{\text{EUR}\})$
- George dom DocA
- George  $\neg$  dom DocB
- George dom DocC

# New Security Condition and \*-Property

- Let  $C(S)$  be the category set of subject  $S$ .
- Let  $C(O)$  be the category set of object  $O$ .
- Simple Security Condition (**not read up**):  
     $S$  can read  $O$  if and only if  $S \text{ dom } O$  and  
     $S$  has discretionary read access to  $O$ .
- \*-Property (**not write down**):  
     $S$  can write to  $O$  if and only if  $O \text{ dom } S$  and  
     $S$  has discretionary write access to  $O$ .
- Basic Security Theorem:  
    Let  $\Sigma$  be a system with secure initial state  $\sigma_0$   
    Let  $T$  be the set of state transformations.  
    If every element of  $T$  preserves the simple security condition,  
    preliminary version, and the \*-property, preliminary version,  
    Then every state  $\sigma_i, i \geq 0$ , is secure.

# Allow Write Down?

- Bell-LaPadula allows higher-level subject to write into lower level object that low level subject can read.
- A subject has a maximum security level and a current security level.
  - maximum security level must dominate current security level.
- A subject may (effectively) decrease its security level from the maximum in order to communicate with entities at lower security levels.
- Colonel's maximum security level is  $(S, \{NUC, EUR\})$ . She changes her current security level to  $(S, \{EUR\})$ . Now she can create document at Major is clearance level  $(S, \{EUR\})$ .

# Limitations with BLP

- Deal only with confidentiality, does not deal with integrity at all
  - Addressed by integrity models (such as Biba, Clark-Wilson)
- Does not deal with information flow through covert channels

# What is integrity?

- Attempt 1: Critical data do not change.
- Attempt 2: Critical data changed only in “correct ways”
  - E.g., in DB, integrity constraints are used for consistency
- Attempt 3: Critical data changed only through certain “trusted programs”
- Attempt 4: Critical data changed only as intended by authorized users.

# The Biba Model

- Kenneth J. Biba: "Integrity Considerations for Secure Computer Systems", MTR-3153, The Mitre Corporation, April 1977.
- Motivated by the fact that BLP does not deal with integrity

# Biba: Integrity Levels

- Each subject (program) has an integrity level
- Each object has an integrity level
- Integrity levels are totally ordered
- Integrity levels different from security levels in confidentiality protection
  - a highly sensitive data may have low integrity



# Five Mandatory Policies in Biba

- Strict integrity policy
- Subject low-water mark policy
- Object low-water mark policy
- Low-water mark Integrity audit policy
- Ring policy

# Strict Integrity Policy (BLP reversed)

- Rules:
  - $s$  can read  $o$  iff  $i(s) \leq i(o)$ 
    - no read down
    - stops indirect sabotage by contaminated data
  - $s$  can write to  $o$  iff  $i(s) \geq i(o)$ 
    - no write up
    - stops directly malicious modification
- Fixed integrity levels

- Thanks