

Formal Languages and Automata Theory (CS303)

Mid-Semester Examination

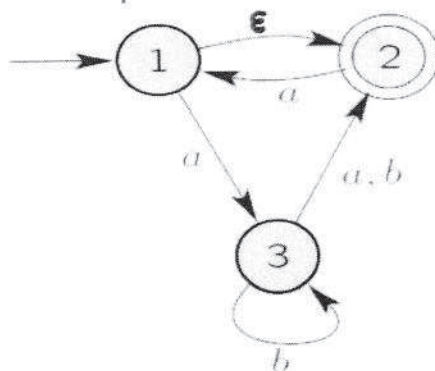
Indian Institute of Technology, Patna

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Full marks- 60, Duration- 120 min

1. Answer all the questions: [4X5]

- (a) Prove that, if M is a DFA that recognizes language B , swapping the accept and non-accept states in M yields a new DFA that recognizes \overline{B} , the complement of B . Conclude that the class of regular languages is closed under complement.
- (b) We say that a DFA M for a language A is minimal if there does not exist another DFA M' for A such that M' has strictly fewer states than M . Suppose that $M = (Q, \Sigma, \delta, q_0, F)$ is a minimal DFA for A . Using M , we construct a DFA \overline{M} for the complement \overline{A} as $\overline{M} = (Q, \Sigma, \delta, q_0, Q-F)$. Prove that \overline{M} is a minimal DFA for \overline{A} .
- (c) Prove or disprove, if M is an NFA that recognizes language C , swapping the accept and non-accept states in M doesn't necessarily yield a new NFA that recognizes \overline{C} .
- (d) Convert the following NFA N into an equivalent DFA.



2. Answer all the questions: [4+8+8]

- (a) Prove that if we remove a finite set of strings from a regular language, the result is a regular language.
- (b) Prove that the languages $A1 = \{www \mid w \in \{a, b\}^*\}$ and $A2 = \{w \in \{a, b\}^* \mid w = w^R\}$ are non-regular using pumping lemma.
- (c) Give context free grammars for following two languages: $L1 = \{a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i = j \text{ or } i = k\}$ and $L2 = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i+j = k\}$. Please write the logic behind the production rules of your grammars precisely.

3. Answer all the questions: (6+8+6)

- (a) Show that the intersection of a context-free language with a regular language is always context-free.
- (b) Give PDAs for $L1$ from Q. 2.c with at most eight states and for $L2$ from Q. 2.c with at most five states. Also write concise descriptions for your PDAs.
- (c) Prove that PDAs have more power than DPDAs.