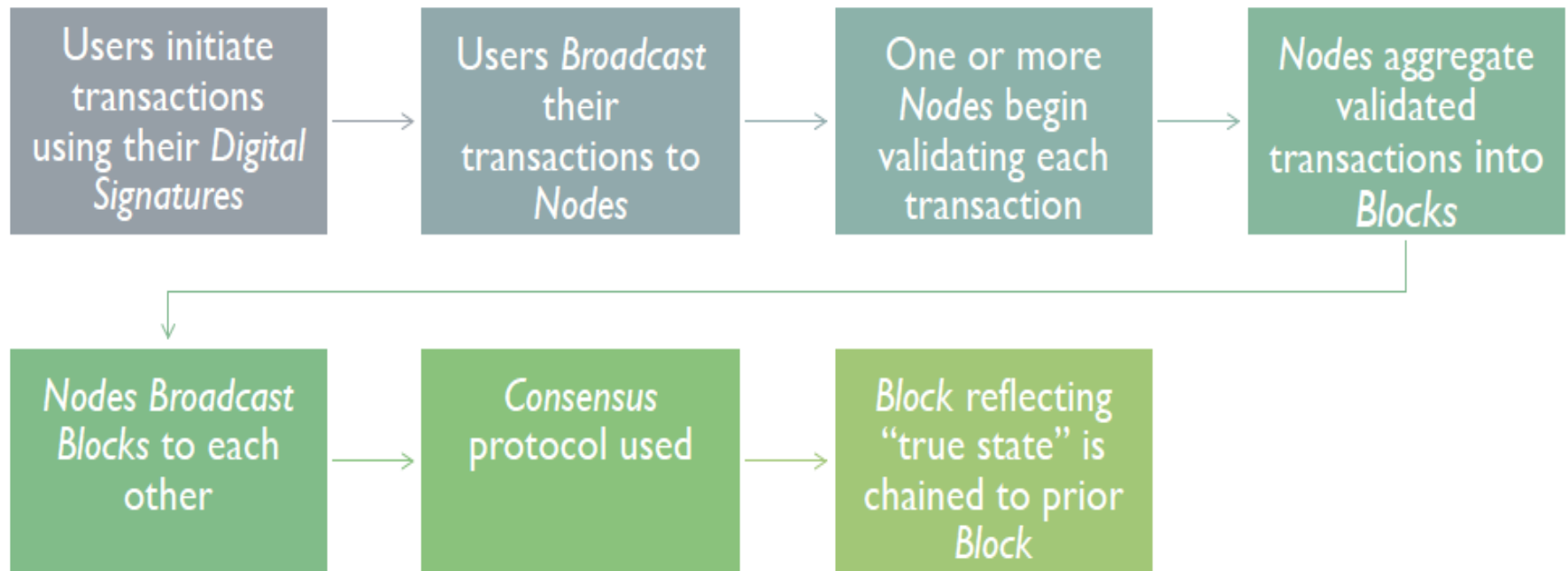


# CS577: Introduction to Blockchain and Cryptocurrency

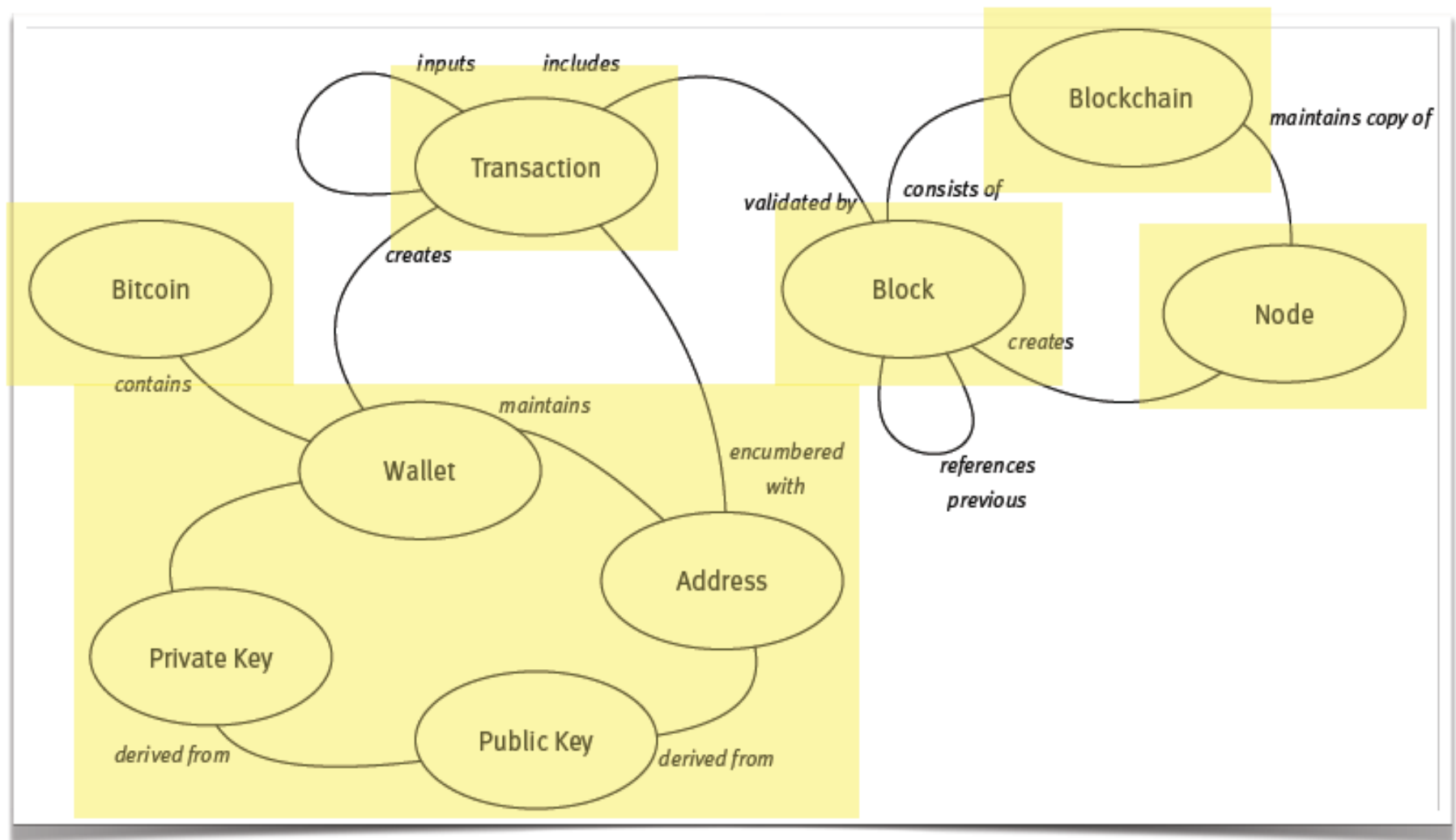
## **Consensus**

**Dr. Raju Halder**

# HOW MIGHT A DISTRIBUTED LEDGER WORK?



# Blockchain and Bitcoin



# *Key characteristics of blockchain*

- *Decentralisation.*
  - *Peer to peer to network*
- *Persistency.*
  - *Transactions stored persistently*
- *Anonymity.*
  - *Avoid Identity exposure*
- *Auditability.*
  - *Easily verifiable and traceable*

**Blockchain challenges and opportunities: a survey - by  
Zibin Zheng et al., 2018**

# Categories

- Public blockchain
  - Anybody can join anytime
- Private blockchain
  - Fully controlled by one organization who could determine the final consensus
- Consortium blockchain
  - Only a selected set of nodes are responsible for validating the block

# Comparison

**Table 1** Comparisons among *public blockchain*, *consortium blockchain* and *private blockchain*

| <i>Property</i>         | <i>Public blockchain</i>    | <i>Consortium blockchain</i>  | <i>Private blockchain</i>     |
|-------------------------|-----------------------------|-------------------------------|-------------------------------|
| Consensus determination | All miners                  | Selected set of nodes         | One organisation              |
| Read permission         | Public                      | Could be public or restricted | Could be public or restricted |
| Immutability            | Nearly impossible to tamper | Could be tampered             | Could be tampered             |
| Efficiency              | Low                         | High                          | High                          |
| Centralised             | No                          | Partial                       | Yes                           |
| Consensus process       | Permissionless              | Permissioned                  | Permissioned                  |

**In order to secure a blockchain ... it's estimated that both Bitcoin and Ethereum burn over \$1 million worth of electricity and hardware costs per day as part of their consensus mechanism.**

**- VITALIK BUTERIN**



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EDITOR'S PICK | 23,388 views | Apr 19, 2018, 11:09pm

# Bitcoin's Energy Consumption Can Power An Entire Country -- But EOS Is Trying To Fix That

**Sherman Lee** Contributor ⓘ*I write about deep tech, crypto, and artificial intelligence.*





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## Bitcoins Energy Consumption An Unsustainable Protocol That Must Evolve?



By john illic

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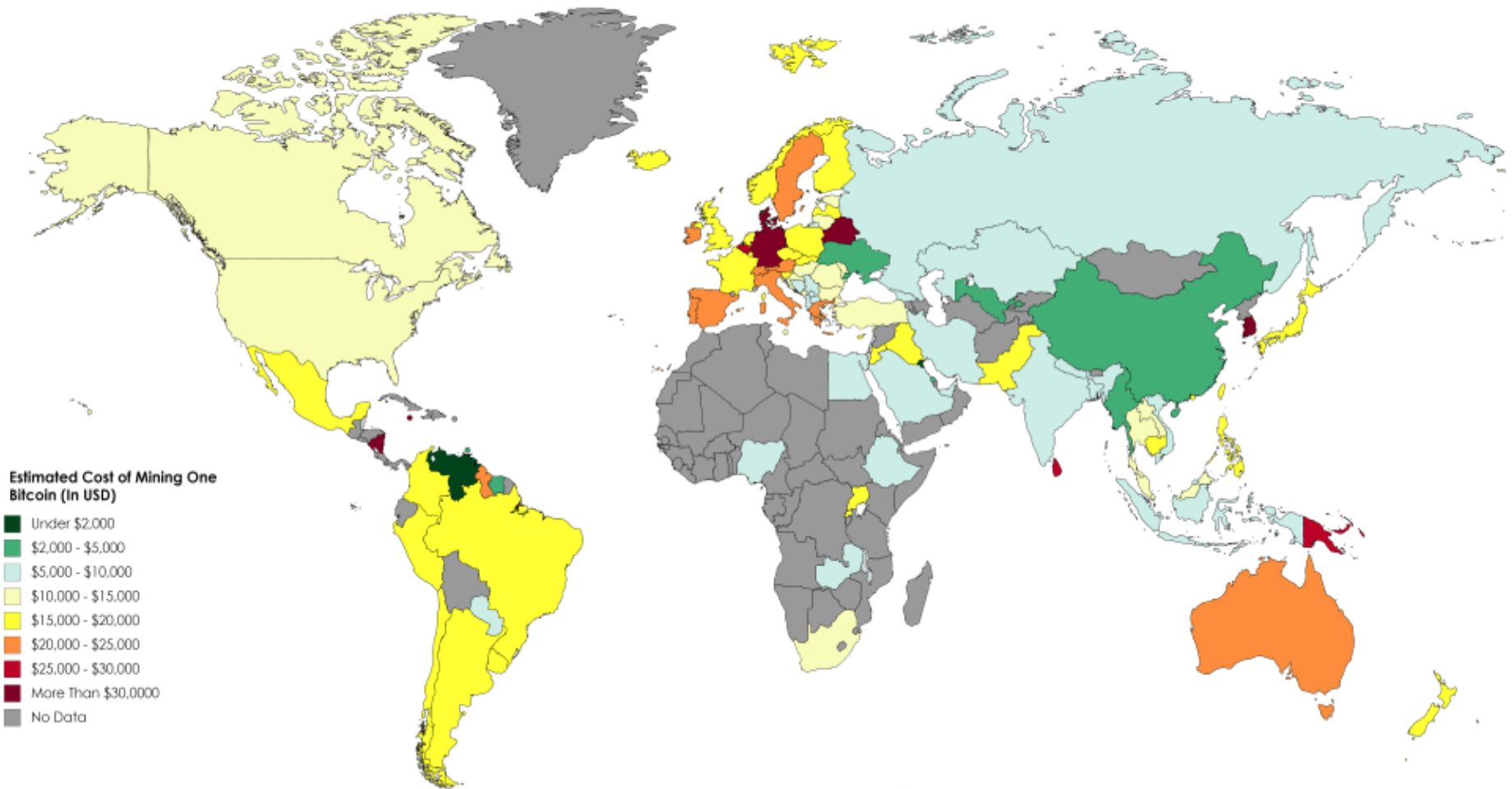


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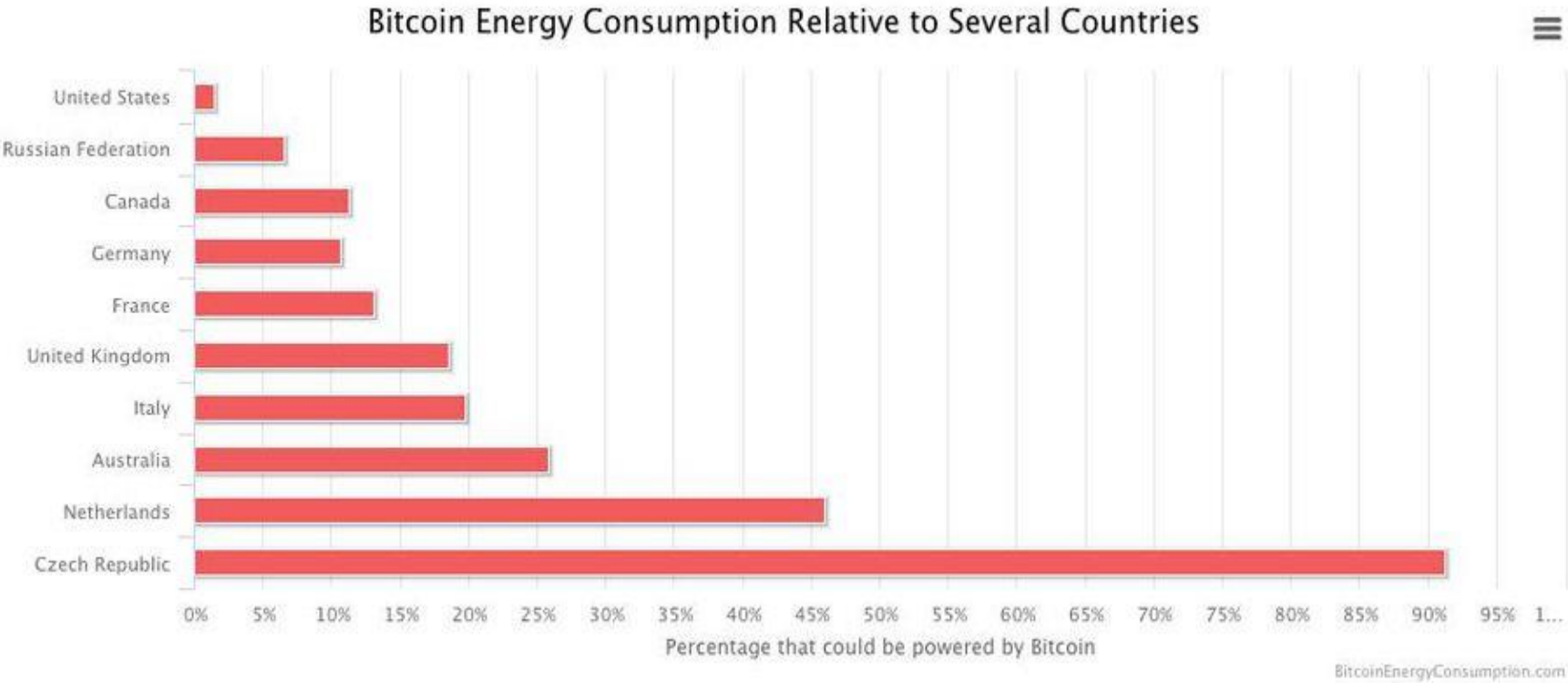
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# Estimated Electricity Cost Of Mining One Bitcoin By Country



Source: <https://powercompare.co.uk/bitcoin-electricity-cost/>

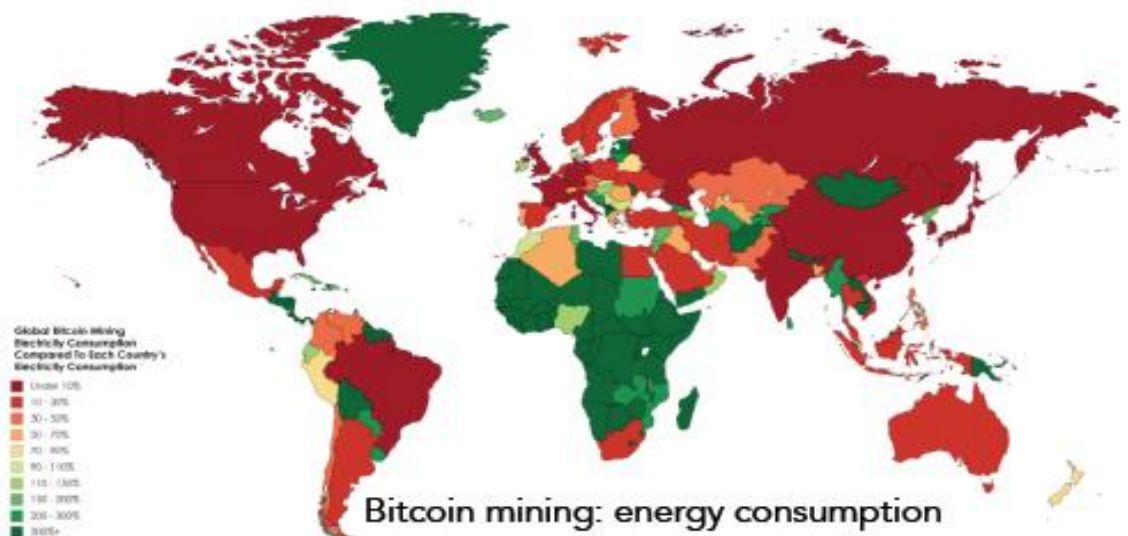
**The Bitcoin POW mechanism is so costly that it consumes the same amount of electricity it takes to power a country like Switzerland in one year. Bitcoin’s current estimated annual electricity consumption is 61.4 TWh, which is also equivalent to 1.5% of the electricity consumed in the United States.**



## Proof of X

## Proof of Stake

- And others: Burn, Elapsed time, Capacity



# Proof-of-X

- Proof-of-X (PoX) schemes is an umbrella term for systems that replace PoW with more useful and energy-efficient alternatives to Proof-of-Work (PoW).

# Proof-of-Stake

Miner/Mining Vs. Validator/Minting or forged

- POS requires people to prove the ownership of a certain amount of currency
  - It is believed that people with more currencies would be less likely to attack the network.
  - If richest person attacks, currency value falls and it may be a loss for the attackers!
- Many blockchains adopt PoW at the beginning and transform to PoS gradually.
  - For instance, Ethereum is planning to move from Ethash (a kind of PoW) (Wood, 2014) to Casper (a kind of PoS) (Zamfir, 2015).

# Proof-of-Stake

- PoS alternatives consume less energy and reach higher transactions per second.
- But they have also still to prove their attack-resistance in real open public settings like PoW so far.
- Challenge for proof-of-stake systems is to keep track of the changing stakes of the stakeholders.

# Proof-of-Stake

- Selection by account balance would result in undesirable centralization because the single richest member would have a permanent advantage as it gets richer.
- Different versions: random selection, age-based stake selection (number of coins stake multiply by the time they have been staked, when selected, time reset to 0)...



# Proof-of-Stake: Randomization

- Blackcoin (Vasin, 2014) uses randomization to predict the next generator.
- It uses a formula that looks for the lowest hash value in combination with the size of the stake.

# Proof-of-Stake: Coin age

- Peercoin (King and Nadal, 2012) favours coin age-based selection.
- In Peercoin, older and larger sets of coins have a greater probability of mining the next block.
- Once a user has forged a block, their coin age is reset to zero and then they must wait at least 30 days again before they can sign another block.

# Proof-of-Capacity

- Sometimes stake could be other things.
- For example, proof of capacity (burstcoin, 2014).
- In proof-of-capacity, participants vote on new blocks weighted by their capacity to allocate a non-trivial amount of disk space.
- Other Examples: PermaCoin, SpaceMint

# Proof-of-Capacity

- PermaCoin repurposes Bitcoin's PoW with a more broadly useful task: providing a robust, distributed storage.
- SpaceMint employs a consensus protocol based on a non-interactive variant of proof-of-capacity (called proof-of-space).

# Proof-of-Deposit

- Miners ‘lock’ a certain amount of coins, which they cannot spend for the duration of their mining.
- One such system is Tendermint, where a miner’s voting power is proportional to the amount of coins they have locked.
- Deposit could be revoked if they misbehaved.

# Proof-of-Activity

- To combine the benefits of POW and POS, proof of activity (Bentov et al., 2014) is proposed.
- In proof of activity, a mined block (based on PoW) needs to be signed by N validators (PoS) to be valid.
- In that way, if some owner of 50% of all coins exists, he/she cannot control the creation of new blocks on his/her own.
- Since POA marries POW and POS, it draws criticism for its partial use of both.

# Proof of Authority

- leverages identity instead of coins
- the PoA consensus algorithm is usually reliant upon:
  - valid and trustworthy identities: validators need to confirm their real identities.
  - difficulty to become a validator: a candidate must be willing to invest money and put his reputation at stake. A tough process reduces the risks of selecting questionable validators and incentivize a long-term commitment.
  - a standard for validator approval: the method for selecting validators must be equal to all candidates.
- Kovan and Rinkeby, the two Ethereum testnets, also use PoA as a consensus mechanism. Microsoft Azure is another example where the PoA is being implemented.

# Delegated Proof-of-Stake

- In Delegated PoS (DPOS), stake-holders don't vote on the validity of the blocks themselves, but vote (proportionately weighted based on the stake) to elect delegates to do the validation on their behalf.
- The major difference between POS and DPOS is that POS is a direct democratic while DPOS is representative democratic.
- Users can also delegate their voting power to another user who will vote on their behalf.



# Delegated Proof-of-Stake

- Higher Throughput: With significantly fewer nodes to validate the block, the block could be confirmed quickly, making the transactions confirmed quickly.
- Dishonest delegates could be voted out easily.
- Examples: Steem and BitShares

# Proof-of-Burn

- Method for distributed consensus and an alternative to Proof of Work and Proof of Stake
- Miners prove that they have destroyed a quantity of coins, for example by sending them to a verifiably unspendable address.
- Slimcode implemente this approach in 2014 but has recently been discontinued.

# Proof-of-Elapsed-Time

- Often used on the permissioned blockchain networks.
- Each node in the blockchain network generates a random wait time and goes to sleep for that specified duration.
- The one to wake up first – that is, the one with the shortest wait time – wakes up and commits a new block to the blockchain, broadcasting the necessary information to the whole peer network
- The same process then repeats for the discovery of the next block.

# Proof-of-Elapsed-Time

- The POET network consensus mechanism needs to ensure two important factors:
  - First, that the participating nodes genuinely select a time that is indeed random and not a shorter duration chosen purposely by the participants in order to win, and
  - Second, the winner has indeed completed the waiting time.

# Proof-of-Elapsed-Time

- The POET concept was invented during early 2016 by Intel.
- It offers a readymade high tech tool to solve the computing problem of "random leader election."

# Hyperledger Fabric : PBFT

- Practical byzantine fault tolerance (PBFT) is a replication algorithm to tolerate byzantine faults (Miguel and Barbara, 1999).
- Hyperledger Fabric (hyperledger, 2015) utilises the PBFT as its consensus algorithm since PBFT could handle up to  $1/3$  malicious byzantine replicas.

# Ripple

- Ripple (Schwartz et al., 2014) is a consensus algorithm that utilises collectively-trusted subnetworks within the larger network.
- In the network, nodes are divided into two types: **server** for participating consensus process and **client** for only transferring funds.
- In contrast to that PBFT nodes have to ask every node in the network, each Ripple server has a Unique Node List (UNL) to query.

# Ripple

- UNL is important to the server. When determining whether to put a transaction into the ledger, the server would query the nodes in UNL.
- If the received agreements have reached 80%, the transaction would be packed into the ledger.
- For a node, the ledger will remain correct as long as the percentage of faulty nodes in UNL is less than 20%.



# Consensus: A Comparison

**Table 2** Typical consensus algorithms comparison

| <i>Property</i>          | <i>PoW</i>      | <i>PoS</i> | <i>PBFT</i>        | <i>DPOS</i> | <i>Ripple</i>       | <i>Tendermint</i>      |
|--------------------------|-----------------|------------|--------------------|-------------|---------------------|------------------------|
| Node identity management | Open            | Open       | Permissioned       | Open        | Open                | Permissioned           |
| Energy saving            | No              | Partial    | Yes                | Partial     | Yes                 | Yes                    |
| Tolerated                | < 25%           | < 51%      | < 33.3%            | < 51%       | < 20%               | < 33.3%                |
| power of adversary       | computing power | stake      | faulty replicas    | validators  | faulty nodes in UNL | byzantine voting power |
| Example                  | Bitcoin         | Peercoin   | Hyperledger Fabric | Bitshares   | Ripple              | Tendermint             |

# Distributed ledger technologies

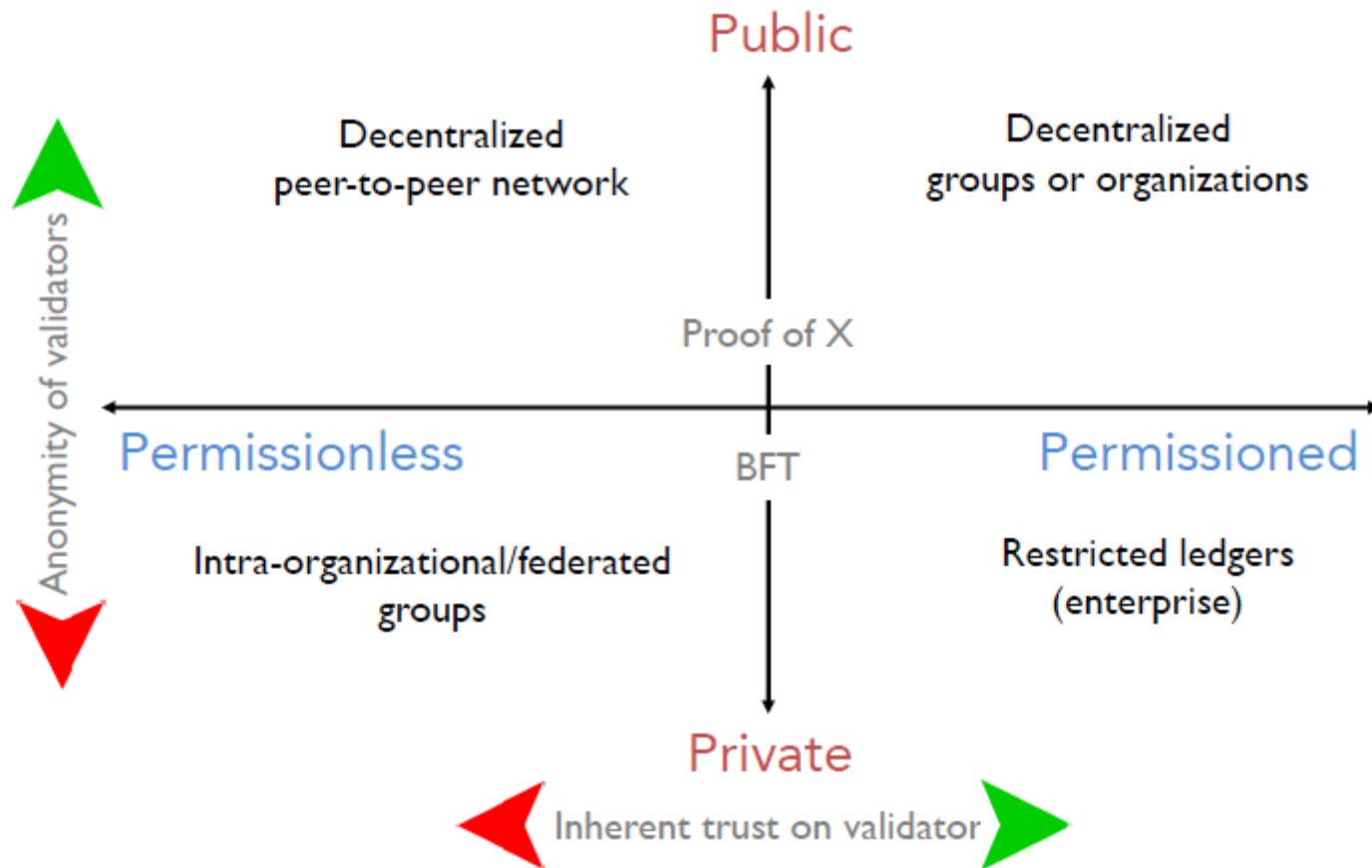


TABLE I  
A BRIEF COMPARISON OF SOME WELL-KNOWN BLOCKCHAIN SYSTEMS.

| Blockchain system   | Data structure | Permissioned | Consensus   | Smart contract language | Turing complete |
|---------------------|----------------|--------------|-------------|-------------------------|-----------------|
| Bitcoin [11]        | blockchain     | No           | PoW         | Golang, C++             | No              |
| Litecoin [55]       | blockchain     | No           | PoW         | Golang, C++             | No              |
| Ripple [56]         | blockchain     | Yes          | Ripple      | Golang, C++             | No              |
| ZCash [57]          | blockchain     | No           | PoW         | C++                     | No              |
| Hyperledger [54]    | blockchain     | Yes          | PBFT        | Golang, Java            | Yes             |
| Sawtooth Lake [58]  | blockchain     | No           | PoET        | Python                  | Yes             |
| Ethereum [52], [53] | blockchain     | No           | PoW/PoS     | Solidity, Serpent, LLL  | Yes             |
| Quorum [59]         | blockchain     | Yes          | QuorumChain | Golang                  | Yes             |
| Monax [60]          | blockchain     | Yes          | Tendermint  | Solidity                | Yes             |
| Tezos [61]          | blockchain     | No           | PoS         | Michelson               | No              |
| Corda [62]          | blockchain     | Yes          | BFT         | Kotlin, Java            | No              |
| Kadena [63], [64]   | blockchain     | Yes          | ScalableBFT | Pact                    | No              |
| IOTA [50]           | DAG            | No           | PoW         | Java                    | No              |
| Byteball [51]       | DAG            | Yes          | Main chain  | Node.js                 | No              |

# Proof of X: Attacks

- **nothing-at-stake attack:** A miners are incentivized to extend every potential fork. Since it is computationally cheap to extend a chain, in the case of forks, rational miners mine on top of every chain to increase the likelihood of getting their block in the right chain.
- **grinding attack:** A miner re-creates a block multiple times until it is likely that the miner can create a second block shortly afterwards.
- **long-range attack:** An attacker can bribe miners to sell their private keys. If these keys had considerable value in the past, then the adversary can mine previous blocks and re-write the entire history of the blockchain.