

## CS571: Artificial Intelligence

### End Semester Quiz

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**Ans 1:**

#### **Gini Index**

- calculates the degree of probability of a specific variable that is wrongly being classified when chosen randomly and a variation of gini coefficient.
- works on categorical variables, provides outcomes either be “successful” or “failure” and hence conducts binary splitting only.

#### **Information Gain**

- used for determining the best features/attributes that render maximum information about a class.
- follows the concept of entropy while aiming at decreasing the level of entropy, beginning from the root node to the leaf nodes.
- computes the difference between entropy before and after split and specifies the impurity in class elements.

#### **Differences of Gini Index and Information Gain:**

- Gini Index performs binary splitting to determine the efficiency of the split, Information Gain is robust.
- Reason: the Gini algorithm itself experiments with the data in a greedy fashion in-order to land at the best possible split using the concept of entropy.
- Though Information gain is computationally more heavy compared to Gini Index, we can be sure that it enables us to get the best possible split.
- However, if the data consists of large number of partitions, Information Gain will take lots of time to reach optimality. In such cases Gini Index could be preferred.
- Gini's maximum impurity is 0.5 and maximum purity is 0 while Entropy's maximum impurity is 1 and maximum purity is 0

**Counterexample:** If we have a dataset of 1000 clases, using information gain will take much time as compared to gini index.

**Ans 2:**

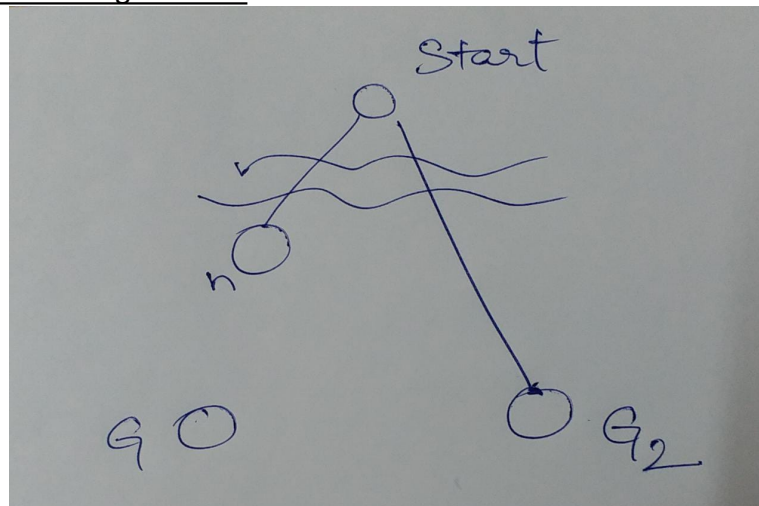
A heuristic  $h(n)$  is said to be admissible if  $h(n) \leq h^*(n)$  where  $h^*(n)$  is the true cost to the goal  $\rightarrow h(n)$  never overestimates the cost from node to goal state.

$A^*$  is guaranteed to always return the shortest path. All other algorithms return the optimal path must expand all other nodes whereas  $A^*$  does not.

So, the given statement that "If  $A^*$  knows the ACTUAL cost of the optimal path of any node to the goal, then no useless node is expanded" is **TRUE**

- We are using an admissible heuristic  $h(n)$ , so there is no overestimation of the actual cost from a node  $n$  to goal.
- Since  $A^*$  heuristic is calculated in such a way that it adds both the cost to that node and cost from that node to Goal state,  $A^*$  is always guaranteed to return the shortest path and doesn't expand to an useless node

Consider the following scenario:



Node  $n$  is an unexpanded node on the shortest path to goal ( $G$ ) and  $G_2$  is a suboptimal goal.

$$h(G_2) = 0 \Rightarrow f(G_2) = g(G_2)$$

As  $G_2$  is not optimal,  $g(G_2) > g(G)$

$$\text{so, } h(G) = 0 \Rightarrow f(G) = g(G) \Rightarrow f(G_2) > f(G)$$

as  $n$  is an unexpanded node,  $f(n) = g(n) + h(n) \leq f(G)$

So,  $f(n) \leq f(G)$ . (Since  $h$  is admissible)

This implies  $f(G_2) > f(n)$

Hence,  $A^*$  will never expand node  $G_2$ .

Therefore,  $A^*$  will never expand a useless node.

**Ans 3:**

Our knowledge base consists of the following first-order Horn clauses:

$\text{Ancestor}(\text{Mother}(x), x)$   
 $\text{Ancestor}(x, y) \wedge \text{Ancestor}(y, z) \implies \text{Ancestor}(x, z)$

Clearly, we can see that the resolution is complete

We need to prove the following:

$\sim \text{Ancestor}(\text{John}, \text{John})$

The resolution cannot prove the above statement despite being complete or finished. A careful analysis will show that the statement to be proven doesn't follow the knowledge base.

Reason: There is nothing in the knowledge base that dismisses the possibility of everything being the ancestor of everything else.