# The Exponential Mechanism for Medians

## April 27, 2020

## 1 THE EXPONENTIAL MECHANISM

Sometimes, the global sensitivity of a function is too great, so the Laplace mechanism will not produce meaningful results. The median is one such function. In many cases, the *Exponential mechanism* is an alternate approach that gives reasonable utility. Introduced in 2007 by McSherry and Talwar, the exponential mechanism posits that for a given database, users prefer some outputs over others. That those preferences may be encapsulated with a utility score, where a high utility score indicates a higher preference for that output. The exponential mechanism releases outputs with probability proportional (in the exponent) to the utility score and the sensitivity of the utility function.

**Definition 1.** Let  $\mathcal{X}$  be a space of databases and let [m,M] be an arbitrary range. Let  $u: \mathcal{X} \times [m,M] \to \mathbb{R}$  be a utility function, which maps pairs of databases and outputs to a utility score. Let  $\Delta u$  be the sensitivity of u with respect to the database argument. The exponential mechanism outputs  $r \in [m,M]$  with probability proportional to  $\exp\left(\frac{\varepsilon u(x,r)}{2\Delta u}\right)$   $[MT07,DR^+14].^2$ 

**Theorem 1.** The exponential mechanism preserves  $(\varepsilon, 0)$ -differential privacy [MT07, DR<sup>+</sup>14].<sup>3</sup> Note that the exponential mechanism may not be tractable in many cases, as it assumes the existence of a utility function, and even if one exists it may not be efficiently computable.

## 2 AN EXPONENTIAL MECHANISM FOR A MEDIAN

# 2.1 Defining a sensible utility function

$$u(x,r) = \max(\alpha, (1-\alpha))N - |(1-\alpha)\#(Z < x) - \alpha\#(Z > x)|$$
 (2.1)

<sup>&</sup>lt;sup>1</sup>This is not the *only* advantage of the exponential mechanism. It is a way to compute differentially private queries on non-numeric data, unlike the Laplace mechanism it does not assume that the probability of outputting a response ought to be symmetric about the true response, etc.

<sup>&</sup>lt;sup>2</sup>The original definition is from [MT07], but here we state the version rewritten in [DR<sup>+</sup>14] as it is slightly clearer.

<sup>&</sup>lt;sup>3</sup>As written in [MT07], the mechanism actually preserves  $(2\varepsilon\Delta u, 0)$ -differential privacy; the main difference in the [DR<sup>+</sup>14] version is that it has the extra factor of  $2\Delta u$  to avoid these extra terms.

#### 2.2 Sensitivity of the utility function

#### 2.2.1 Neighboring Definition: Change One

Let  $C_1 = \#(Z < x), C_2 = \#(Z > x)$ . Worst case,  $C_1$  increases by 1 and  $C_2$  decreases by 1. Then

$$\Delta u = ||(1 - \alpha)(C_1 + 1) - \alpha(C_2 - 1)| - |(1 - \alpha)C_1 - \alpha C_2||$$

$$\leq |(1 - \alpha)(C_1 + 1) - \alpha(C_2 - 1) - (1 - \alpha)C_1 + \alpha C_2|$$

$$\leq |C_1 + 1 - \alpha C_1 - \alpha - \alpha C_2 + \alpha - C_1 + \alpha C_2 + \alpha C_2|$$

$$= 1$$

If instead  $C_2$  increases by 1 and  $C_1$  decreases by 1, the result is identical (negative sign falls out).

## 2.2.2 Neighboring Definition: Add/Drop One

In the add/remove-1 model, utility function has sensitivity  $\max(1 - \alpha, \alpha)$ : If we add 1, there are two worse-cases:  $C_2$  increases by 1, nothing happens to  $C_2$ .

$$\Delta u = |(1 - \alpha)(C_1 + 1) - \alpha(C_2)| - |(1 - \alpha)C_1 - \alpha C_2|$$

$$\leq |(1 - \alpha)(C_1 + 1) - \alpha(C_2) - (1 - \alpha)C_1 + \alpha C_2|$$

$$= |C_1 + 1 - \alpha C_1 - \alpha - \alpha C_2 - C_1 + \alpha C_1 + \alpha C_2|$$

$$= 1 - \alpha$$

b. nothing happens to  $C_1, C_2$  increases by 1

$$\Delta u = |(1 - \alpha)(C_1) - \alpha(C_2 + 1)| - |(1 - \alpha)C_1 - \alpha C_2|$$

$$\leq |C_1 - \alpha C_1 - \alpha C_2 - \alpha - C_1 + \alpha C_1 + \alpha C_2|$$

$$= \alpha$$

Subtracting a point gives you same thing but with some negative signs inside the absolute values that come out in the wash.

#### 2.3 The Normalization Factor

#### References

- [DFM<sup>+</sup>20] Wenxin Du, Canyon Foot, Monica Moniot, Andrew Bray, and Adam Groce. Differentially private confidence intervals. arXiv preprint arXiv:2001.02285, 2020.
- [DR<sup>+</sup>14] Cynthia Dwork, Aaron Roth, et al. The algorithmic foundations of differential privacy. Foundations and Trends® in Theoretical Computer Science, 9(3–4):211–407, 2014.
- [MT07] Frank McSherry and Kunal Talwar. Mechanism design via differential privacy. In 48th Annual IEEE Symposium on Foundations of Computer Science (FOCS'07), pages 94–103. IEEE, 2007.
- [Smi11] Adam Smith. Privacy-preserving statistical estimation with optimal convergence rates. In *Proceedings of the forty-third annual ACM symposium on Theory of computing*, pages 813–822, 2011.