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# Count Sensitivity Proofs

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**Definition 1.** Let  $\mathcal{X}$  be the universe of possible rows (individuals) and let  $I : \mathcal{X} \rightarrow \{0, 1\}$  be a predicate on rows. Let  $x \in \mathcal{X}^n$  be a dataset. Then a count over  $x$  is defined as

$$q(x) = \sum_{i=1}^n I(x_i).$$

**Definition 2.** Let  $q_1, \dots, q_k$  be a series of counts with predicates  $I_1, \dots, I_k$ . These counts are disjoint for every row in the database if only one of them evaluates to 1. In other words, they are disjoint if  $\forall x_i \in \mathcal{X}$ ,

$$\sum_{j=1}^k I_j(x_i) = 1.$$

## 1 NEIGHBORING DEFINITION: CHANGE ONE

### 1.1 $\ell_1$ -sensitivity

**Theorem 1.** A single count query has  $\ell_1$ -sensitivity 1 in the change-one model. A series of  $k$  disjoint counts has  $\ell_1$ -sensitivity 2 in the change-one model.

*Proof.* Let  $q$  be a count query with predicate  $I$ , and let  $x'$  be equal to  $x$  with point  $x_i$  changed to  $x'_i$ . Then,

$$\begin{aligned} |q(x') - q(x)| &= \left| \sum_{j=1}^n I(x'_j) - \sum_{j=1}^n I(x_j) \right| \\ &= \left| \sum_{\{i \in [n] | i \neq j\}} I(x_j) + I(x'_i) - \sum_{\{i \in [n] | i \neq j\}} I(x_j) - I(x_i) \right| \\ &= |I(x'_i) - I(x_i)| \\ &\leq 1. \end{aligned}$$

Consider a series of  $k$  disjoint count queries  $\mathbf{q} = \{q_1, \dots, q_k\}$  on the same databases  $x$  and  $x'$ . Note that since the counts are disjoint,  $x_i$  and  $x'_i$  can at most each increment a single one of the  $k$  counts by 1. Call the count that  $x_i$  impacts  $q_i$ , and the count that  $x'_i$  impacts  $q_j$ . Then,

$$\begin{aligned} |\mathbf{q}(x) - \mathbf{q}(x')| &= |(q_1(x) - q_1(x')) + \dots + (q_k(x) - q_k(x'))| \\ &= |(q_i(x) - q_i(x')) + (q_j(x) - q_j(x'))| \\ &\leq 2 \end{aligned}$$

□

## 1.2 $\ell_2$ -sensitivity

**Theorem 2.** *A single count query has  $\ell_2$ -sensitivity 1 in the change-one model. A series of  $k$  disjoint counts has  $\ell_2$ -sensitivity 2 in the change-one model.*

*Proof.* From the proof of Theorem 1, the difference between counts on two neighboring databases is at most 1. Squaring this gives the same value. For a series of  $k$  disjoint counts,

$$\begin{aligned} |\mathbf{q}(x) - \mathbf{q}(x')|_2 &= \left| (q_1(x) - q_1(x'))^2 + \dots + (q_k(x) - q_k(x'))^2 \right| \\ &\leq |1^2 + 1^2| \\ &= 2. \end{aligned}$$

□

## 2 NEIGHBORING DEFINITION: ADD/DROP ONE

### 2.1 $\ell_1$ -sensitivity

**Theorem 3.** *A single count query has  $\ell_1$ -sensitivity 1 in the add/drop-one model. A series of  $k$  disjoint counts also has  $\ell_1$ -sensitivity 1 in the add/drop-one model.*

*Proof.* Let  $q$  be a count query with predicate  $I$ , and let  $x'$  be equal to database  $x$  with point  $x_i$  removed. Then

$$\begin{aligned} |q(x) - q(x')| &= \left| \sum_{j=1}^n I(x_j) - \sum_{\{i \in [n] | i \neq j\}} I(x_j) \right| \\ &= \left| \sum_{\{i \in [n] | i \neq j\}} I(x_j) + I(x_i) - \sum_{\{i \in [n] | i \neq j\}} I(x_j) \right| \\ &\leq 1. \end{aligned}$$

A nearly identical argument holds for adding a point.

Consider a series of  $k$  disjoint count queries  $\mathbf{q} = \{q_1, \dots, q_k\}$  and consider database  $x'$  equal to database  $x$  with point  $x_i$  removed. Note that only a single one of the  $k$  queries will be affected by the change from  $x$  to  $x'$ , so

$$\begin{aligned} |\mathbf{q}(x) - \mathbf{q}(x')| &= |(q_1(x) - q_1(x')) + \dots + (q_k(x) - q_k(x'))| \\ &\leq 1. \end{aligned}$$

The same argument holds for  $x'$  equal to  $x$  with a single point added.

□

## 2.2 $\ell_2$ -sensitivity

**Theorem 4.** *A single count query has  $\ell_2$ -sensitivity 1 in the add/drop-one model. A series of  $k$  disjoint counts also has  $\ell_2$ -sensitivity 1 in the add/drop-one model.*

*Proof.* Squaring the sensitivity bounds from Theorem 3 gives 1 as an upper bound on the  $\ell_2$  sensitivity for both a single count and a series of  $k$  disjoint counts.  $\square$