First-class Isomorphic Specialization by Staged Evaluation

Alexander Slesarenko Alexander Filippov Alexey Romanov

WGP'14, August 31, 2014, Gothenburg, Sweden

www.huawei.com



Agenda

- ☐ Isomorphic Specialization in a Nutshell
- **☐** Why it matters
- ☐ First-class Isomorphisms
- ☐ How it works

Program Specialization

Partial Evaluation

$$P(x,y)$$
, $x = const$



P'(y)

for all y: P'(y) = P(x,y)

Supercompilation

$$P(x), x \in D$$



 $P'(x) \land x \in D' \subset D$

for all $x \in D'$: P'(x) = P(x)

Isomorphic Specialization

$$P: A \rightarrow B$$
, $A \approx_A A'$, $B \approx_B B'$

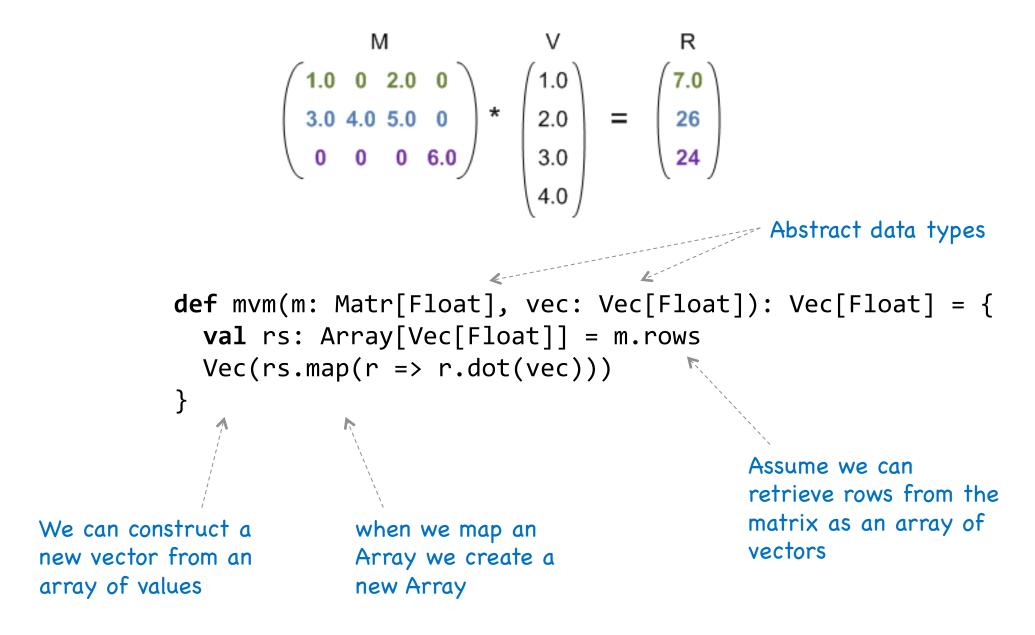


P': X' -> Y'
such that
the diagram
commutes

$$\begin{array}{ccc}
A & \xrightarrow{P} & B \\
\approx_{A} \cdot to & & \downarrow \approx_{B} \cdot from \\
A' & \xrightarrow{P'} & B'
\end{array}$$

P' is better in some sence

Example: Matrix Vector Multiplication

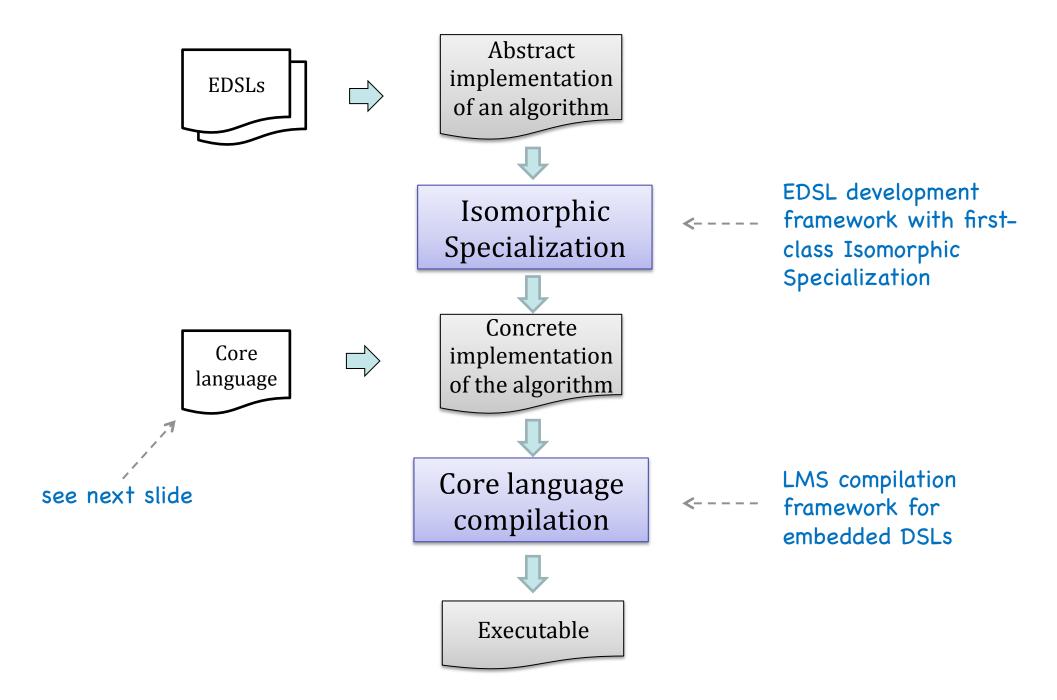


What kind of language we use here? Looks like Scala...

Embedded Domain Specific Languages

```
trait Array[T] { // abstract array interface
                                                   Examples:
  def length: Int
                                                   val average = sum(xs) / xs.length
 def map[R](f: T => R): Array[R]
                                                   val squares = xs.map(x \Rightarrow x * x)
trait Vec[T] {     // abstract vector interface
  def length: Int
  def coords: Array[T]
                                                   val x = v.coords(i)
  def dot(vec: Vec[T]): T
                                                   val s = v1.dot(v2)
trait Matr[T] {    // abstract matrix interface
  def rows: Array[Vec[T]]
                                                   m.rows.map(v => v.dot(vec))
                                            this languages can grow
def Vec(coords: Array[T]): Vec[T]
                                            as a domain specific
def sum[T](arr: Array[T]): T
                                            library in Scala (EDSL)
def mvm(m: Matr[T], vec: Vec[T]): Vec[T] = {
  val vs: Array[Vec[T]] = m.rows
 Vec(vs.map(v => v.dot(vec)))
```

Outline of the method



Core language with immutable arrays

You can select any language which works for you

In the Core language we can construct the following types

```
T = Unit | Int | Float | Boolean // base types
  (T1,T2)
                                     // pair of types
  | Either[T1,T2]
| Array[T]
                                     // sum of types
                                     // array of values of the type T
trait Array[T] {
 def length: Int
 def apply(index: Int): T // get element at index
 def apply(indices: Array[Int]): Array[T]
 def map[R](f: T => R): Array[R]
 def filter(p: T => Boolean): Array[T]
 def zip[U](other: Array[U]): Array[(T,U)]
 def |*| (other: Array[T]): Array[T] // element-wise op
def range(start: Int, len: Int): Array[Int]
def sum[T](arr: Array[T]): T
def unzip[T,U](pairs: Array[(T,U)]): (Array[T],Array[U])
def dotSV[T](
      indices1: Array[Int], values1: Array[T],
      indices2: Array[Int], values2: Array[T]): T
```

Now, if we have a compiler from the Core language into something executable

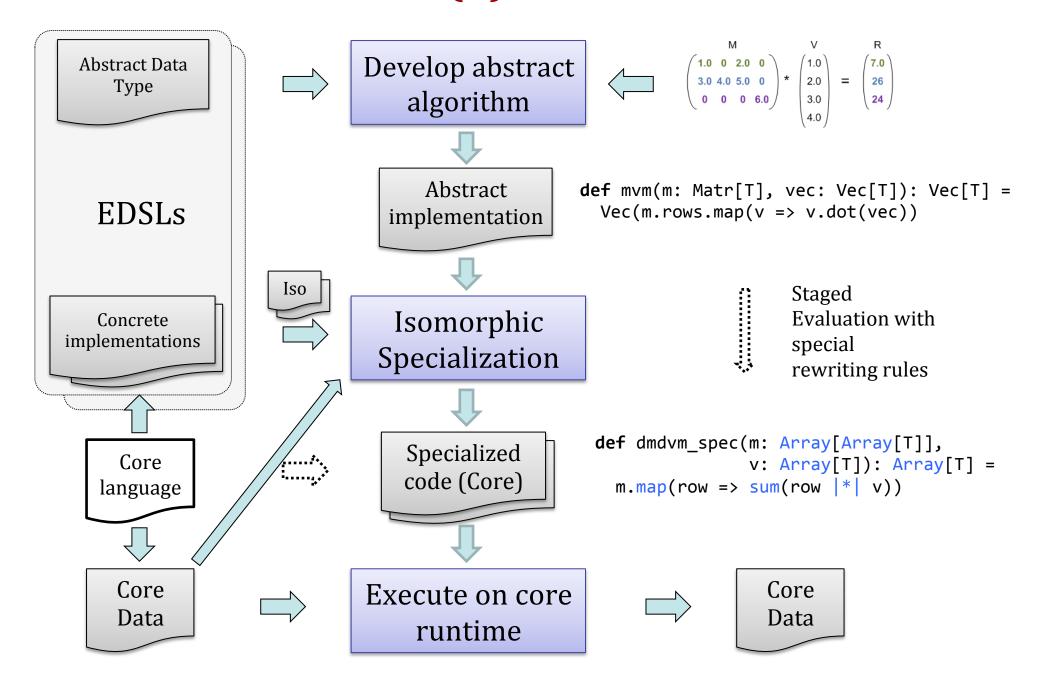
then we can implement abstract data types using the Core language.

Let's look at the implementation

MVM using dense representations

```
Implements interface
class DenseVec[T](val coords: Array[T]) extends Vec[T] { ←
  def length = coords.length
  def dot(vec: Vec[T]) = vec match {
                                                                    Use core language
    case dv: DenseVec[T] => sum(coords | * | dv.coords)
                                                                    types and primitives
    case sv: SparseVec[T] =>
      sum(sv.values |*| coords(sv.indices))
                                                                    Allow mixed types
class DenseMatr[T](val rows: Array[DenseVec[T]]) extends Matr[T]
                                                            default vector is dense
def Vec(coords: Array[T]): Vec[T] = new DenseVec(coords)
                                                                   In abstract algorithm
def mvm(m: Matr[T], vec: Vec[T]): Vec[T] =
                                                                    we don't know how
  Vec(m.rows.map(v => v.dot(vec)))
                                                                    Matr[T] and Vec[T] are
                                                                    implemented
def dmdvm(m: Array[Array[T]],
                                                                    In order to execute
          v: Array[T]): Array[T] = {
                                                                    using core data we do
  val dm = new DenseMatr(
             m.map(r => new DenseVec(r)))
                                                                    three steps:
  val dv = new DenseVec(v)
                                                                    1) wrap
                                                                    2) apply
  val v = mvm(dm, dv)
                                                                    3) extract
  v.coords
```

Outline of the method (2)



Because of staging, specialization can happen at runtime

Specialization by Staged Evaluation

Staged Evaluation with rewriting

```
val dv1 = new DenseVec[Double](v1)
                                          // class DenseVec[T](val coords: Array[T])
val dv2 = new DenseVec[Double](v2)
SE[sum(dv1.coords |*| dv2.coords)]
                                           // Ctx[new C(v_1,...,v_n).f_i] \rightarrow Ctx[v_i]
==> SE[sum(v1 |*| dv2.coords)]
==> SE[sum(v1 |*| v2)]
==> sum(v1 |*| v2)
                                                              Every field is also a
                                                              constructor argument
val m: Array[Array[T]]
SE[m.map(r => new DenseVec(r)).map(dv => dv.coords)]
               // RW[as.map(f).map(g)] \rightarrow as.map(a => g(f(a)))
               // f = (r => new DenseVec(r))
               // g = (dv => dv.coords)
==> SE[m.map(r => new DenseVec(r).coords)]
               // Ctx[new C(v_1,...,v_n).f<sub>i</sub>] \rightarrow Ctx[v_i]
==> SE[m.map(r => r)]
               // RW[as.map(a => a)] \rightarrow as
==> m
```

Take It Home

- 1. Develop an algorithm using abstract data types of embedded DSLs
- 2. Implement abstract data types using a functional core language with an efficient compile
- 3. Specialize an abstract code into the core language

Agenda

- ☐ Isomorphic Specialization in a Nutshell
- **☐** Why it matters
- ☐ First-class Isomorphisms
- ☐ How it works

Generated executable code



Core language compilation with loop fusion, deforestation etc.

```
def dmdv(m: Array[Array[Double]], v: Array[Double]): Array[Double] = {
   val nRows = m.length
   var res = new Array[Double](nRows)
   for (i <- 0 until nRows) {
      val row = m(i)
      val nCols = row.length
      var sum: Double = 0
      for (j <- 0 until nCols) {
        sum += row(j) * v(j)
      }
      res(i) = sum
   }
   res
}</pre>
```

Why it matters

All matrices: $10^4 \times 10^4$ elements $\mathbf{S_m}$ is matrix sparseness (% of zeros) $\mathbf{S_v}$ is vector sparseness

S_m	S_v	dmdv
0%	0%	11389
10%	10%	11408
50%	50%	12944
90%	90%	13093
99%	99%	13251
0%	50%	11483
50%	0%	12946
10%	90%	13907
90%	10%	14304

Execution time in milliseconds running as Scala program without abstraction elimination and optimizations

- 1) Isomorphic Specialization produces version in Core language with immutable arrays
- 2) LMS compiler optimizes purely functional array operations (deforestation, loop fusion etc.)

S_m	S_v	dmdv
0%	0%	309
10%	10%	311
50%	50%	310
90%	90%	307
99%	99%	307
0%	50%	308
50%	0%	310
10%	90%	311
90%	10%	311

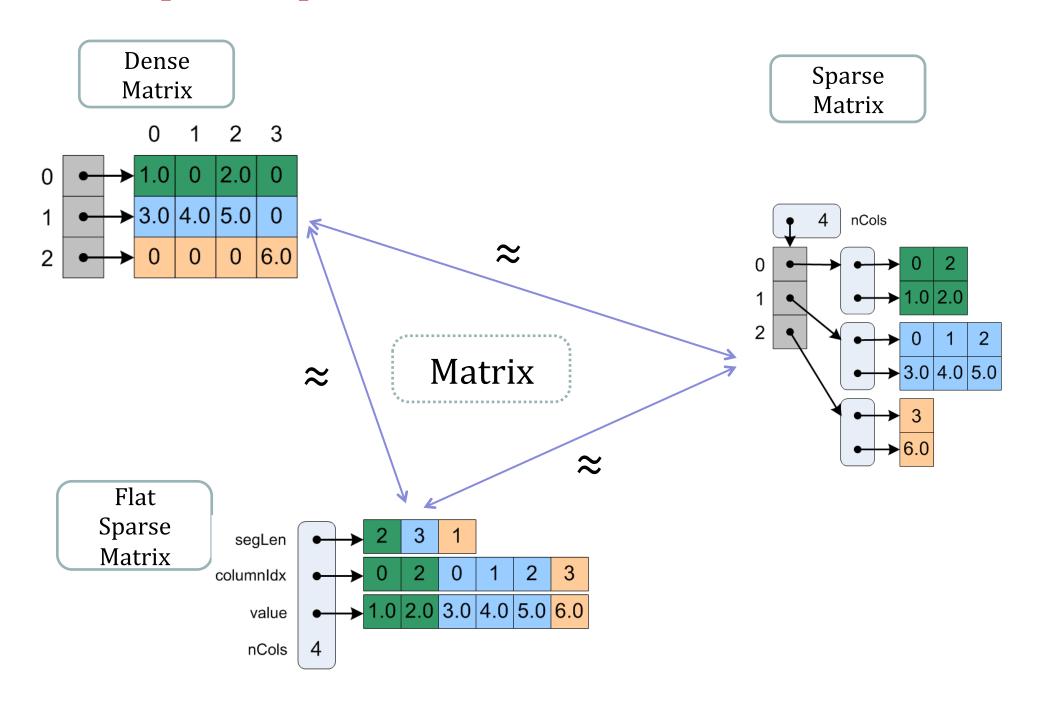
Execution time in milliseconds

~ 40x combined performance improvement

Agenda

- ☐ Isomorphic Specialization in a Nutshell
- **☐** Why it matters
- ☐ First-class Isomorphisms
- ☐ How it works

Isomorphic representations of data



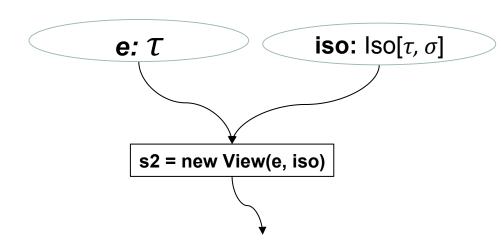
First-class Isomorphisms

```
Vectors
                   Array[T]
  DSL
                        ≈ (iso)
class DenseVec[T](val coords: Array[T])
  extends Vec[T]
{ ... }
         trait Vec[T] {
           def length: Int
           def coords: Array[T]
class SparseVec[T](
        val indices: Array[Int],
        val values: Array[T],
        val length: Int) extends Vec[T]
 (Array[Int], Array[T], Int)
```

```
Matrix
               Array[Array[T]]
 DSL
class DenseMatr[T](
   val rows: Array[DenseVec[T]])
  extends Matr[T]
{ ... }
      trait Matr[T] {
        def rows: Array[Vec[T]]
class SparseMatr[T](say about why
  val rows: Array[(Array[Int], Array[T])];
  val nCol: Int) extends Matr[T]
{ ... }
              ≈ (iso)
(Array[(Array[Int],Array[T])], Int)
```

Generic composition of isomorphisms

IDEA: Let's build an isomorphisms for each constructor of the Core language

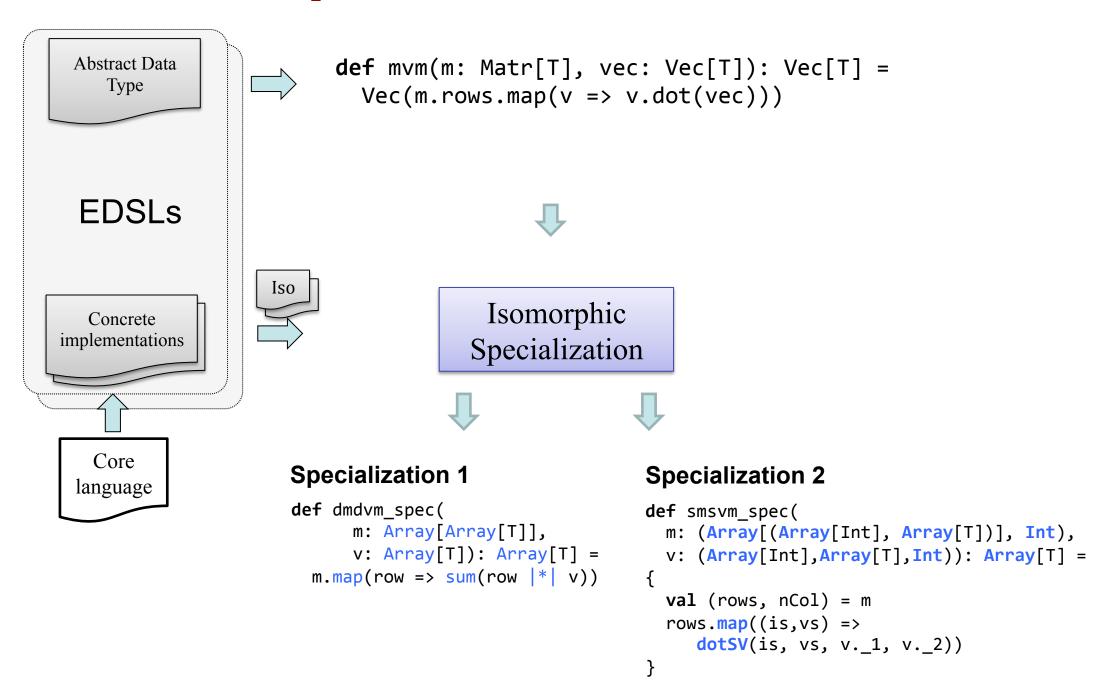


 $\Gamma \vdash e \triangleleft iso : \sigma$

```
class [so_{\times}]A_1, A_2, B_1, B_2] (
     val iso_1: Iso[A_1, B_1], val iso_2: Iso[A_2, B_2])
  extends Iso[A_1 \times A_2, B_1 \times B_2] {
  def to(a:A_1 \times A_2) = (iso_1.to(fst(a)), iso_2.to(snd(a)))
  def from(b:B_1 \times B_2) = (iso_1.from(fst(b)), iso_2.from(snd(b)))
class [so_{+}]A_{1}, A_{2}, B_{1}, B_{2}]
     val iso_1: Iso[A_1, B_1], val iso_2: Iso[A_2, B_2])
  extends Iso [A_1 + A_2, B_1 + B_2] {
  def to (a:A_1+A_2) = case a of {
    1 \cdot a_1 \rightarrow 1 \cdot iso_1 \cdot to(a_1); r \cdot a_2 \rightarrow r \cdot iso_2 \cdot to(a_2)
  def from(b:B_1 + B_2) = case b of {
     1 \cdot b_1 \rightarrow 1 \cdot iso_1 \cdot from(b_1); r \cdot b_2 \rightarrow r \cdot iso_2 \cdot from(b_2)
class (Iso_{axy}[A, B]) (val iso: Iso[A, B])
  extends Iso[Array[A],Array[B]] {
  def to(as: Array[A]) = as.map(iso.to)
  def from(bs: Array[B]) = bs.map(iso.from)
class [so_{\to}]A_1, A_2, B_1, B_2] (
     val iso_1: Iso[A_1, B_1], val iso_2: Iso[A_2, B_2])
  extends Iso[A_1 \rightarrow A_2, B_1 \rightarrow B_2] {
  def to(f:A_1 \rightarrow A_2) = b \Rightarrow iso_2.to(f(iso_1.from(b)))
  def from(g:B_1 \rightarrow B_2)= a \Rightarrow iso_2.from(g(iso_1.to(a)))
```

NOTE: These isomorphisms are defined in the same language and thus they are first class citizents in the framework

Alternative specialized versions



Why it matters

All matrices: $10^4 \times 10^4$ elements

 S_{m} is matrix sparseness (% of zeros)

 S_v is vector sparseness

S_m	S_v	dmdv	dmsv	smdv	smsv
0%	0%	11389	14740	14827	53348
10%	10%	11408	13326	13253	44376
50%	50%	12944	7443	7428	21788
90%	90%	13093	1682	1546	3548
99%	99%	13251	227	167	280
0%	50%	11483	7466	14758	27987
50%	0%	12946	14852	7462	42471
10%	90%	13907	1659	14029	9728
90%	10%	14304	14581	1608	27586

Execution time in milliseconds running as Scala program.

Without abstraction elimination and optimizations

- 1) Isomorphic Specialization produces version in Core language with immutable arrays
 2) LMS compiler optimizes purely functional
- 2) LMS compiler optimizes purely functional array operations (deforestation, loop fusion etc.)

S_m	S_v	dmdv	dmsv	smdv	smsv
0%	0%	(309)	354	366	760
10%	10%	311	323	332	1002
50%	50%	310	202	(187)	924
90%	90%	307	104	42	172 ^
99%	99%	307	18	8	18 ¦
0%	50%	308	198	373	1134
50%	0%	310	359	(187)	986
10%	90%	311	(118)	335	497
90%	10%	311	323	42	345

Execution time in milliseconds

How do you know this is bad choice?

 \sim 20x - 40x performance improvement

A Remarkable Property of Isomorphic Specialization

$$P: A \rightarrow B, A \approx_{A} A', B \approx_{B} B'$$
 $A, B - domain types$
 $A', B' - Core language types$
 $P': X' \rightarrow Y'$
 $A', B' - Core language types$
 $A', B' - Core language types$

commutes

A, B - domain types

A', B' - Core language types

$$\approx_A$$
 to - wrapper

 \approx_B from - extractor

Our conjecture (we haven't proved it formally):

For all P: A -> B,
$$A \approx_A A'$$
, $B \approx_B B'$ there exists P' : X' -> Y' such that $P' = SE_{RW}[\approx_B .from \circ P \circ \approx_A .to]$ and $P' \in Core$

It holds for many non-trivial examples.

Agenda

- ☐ Isomorphic Specialization in a Nutshell
- **☐** Why it matters
- ☐ First-class Isomorphisms
- ☐ How it works

Staged method invocation with graphs

IDEA: let's implement semantics of virtual method call at staging time using IR nodes as class instances

During staging of mvm we have the following redex val r: Array[Double]

val dvec: DenseVec

new DenseVec(r).dot(dvec)

r: Array[Double]

dvec: DenseVec

s2 = new DenseVec(r)

invoke

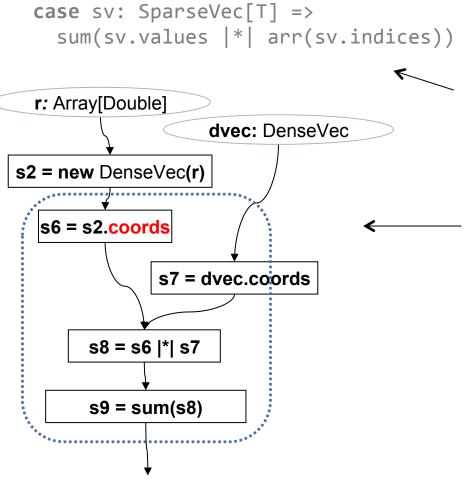
s = s2.dot(dvec)

res

Invocation is possible

Invocation is possible because s2 is defined and we can invoke the method of the concrete instance (which is the node of the graph).

def dot(vec: Vec[T]) = vec match {
 case dv: DenseVec[T] =>
 sum(coords |*| dv.coords)
 case sv: SparseVec[T] =>
 sum(sv.values |*| arr(sv.indices))



res

Method selected dynamically with respect to semantics of virtual method call of FJ language

Pattent matching on types during staged invocation. Even for undefined symbols.

Invocation is NOT possible. This method call is reified in the graph, whereas if invocation IS possible then the node is replaced with invocation result and the graph is rewired.

Take it home

- 1. Develop an algorithm using abstract data types of embedded DSLs
- 2. Identify isomorphic representations in the Domain
- 3. Implement abstract data types using a Core language and isomorphic representations
- 4. Generate representation-specific implementations in the Core Language
- 5. Use optimizing compiler of the Core language
- 6. Check it out (https://github.com/scalan)

Thank you

github.com/scalan