MATH230: Tutorial Two [Solutions]

Propositional Logic: Natural Deductions with Negation

Key ideas

 \bullet Write natural deduction proofs involving \neg negation.

 \bullet Write natural deductions using Ex Falso i.e. \bot elimination.

Relevant Topic: Propositional Logic

Relevant reading: L $\exists \forall N$ Chapters 3,4 and Simon section 1.1

Hand in exercises: 1a, 1b, 1c, 1i, 1k

Due Friday @ 5pm to the submission box on Learn.

Discussion Questions

1. Show $A \vdash \neg \neg A$.

2. Show $(A \lor B) \land (A \lor C) \vdash A \lor (B \land C)$.

Hints

Read these hints and suggestions before reading the solutions below. It is important that you try to solve these problems first, before reading a solution. All hints from earlier labs still apply to this lab.

In this second lab we now use propositions of the form $\neg \alpha$. However, this requires no new rules to deal with the \neg connective as we define $\neg \alpha := \alpha \to \bot$. Therefore, we need only the implication introduction and elimination rules to deal with negations.

This introduction of \neg and \bot can lead to inconsistent contexts i.e. $\Sigma \vdash \bot$ is possible for some Σ . Different logics deal with the introduction of \bot in different ways. In this tutorial you will be asked to make use of the following \bot -elimination rule called *ex falso quodlibet*.

$$\begin{array}{c} \Sigma\\ \vdots\\ \frac{\bot}{A} \text{ XF}\\ \bot \text{ Elimination (Ex Falso)} \end{array}$$

More commonly called the *principle of explosion*, this rule of inference states that if a contradiction can be derived from Σ , then in fact any proposition at all follows.

Hints and Tips

Always unpack $\neg \alpha$ for $\alpha \to \bot$ to help remember the rules that apply in eliminating and introducing \neg negations.

Tutorial Exercises

1. **Minimal Logic.** Provide natural deduction proofs of the following sequents. These deductions require only the use of minimal logic; the introduction and elimination rules for ∧ conjunction, ∨ disjunction, → implication, and the definition of ¬ negation as an implication.

(a)
$$\neg A \vdash (C \rightarrow A) \rightarrow \neg C$$

Solution:

$$\frac{\overline{C}^2 \quad \overline{C} \to \overline{A}^1}{\underline{A}} \text{ MP} \qquad \frac{A}{\neg A} \text{ MP} \qquad \frac{\bot}{\neg C} \to I, 2 \qquad \text{MP} \qquad \frac{\bot}{(C \to A) \to \neg C} \to I, 1$$

(b)
$$A \wedge \neg B \vdash \neg (A \rightarrow B)$$

Solution:

$$\frac{\frac{A \wedge \neg B}{A} \wedge E_{l} \quad \overline{A \rightarrow B}^{1} \text{ MP} \quad \frac{A \wedge \neg B}{\neg B} \wedge E_{r}}{\frac{\bot}{\neg (A \rightarrow B)} \rightarrow I, 1} \wedge E_{r}$$

(c)
$$A \to C$$
, $B \to D$, $\neg C \lor \neg D \vdash \neg A \lor \neg B$

Solution:

(d)
$$A, \neg A \vdash \neg B$$

Solution:

$$\frac{A \quad \neg A}{\frac{\bot}{\neg B} \rightarrow I} \mathsf{MP}$$

(e)
$$A \rightarrow B$$
, $A \rightarrow \neg B \vdash \neg A$

$$\frac{\overline{A}^1 \quad A \to B}{B} \text{ MP} \quad \frac{\overline{A}^1 \quad A \to \neg B}{\neg B} \text{ MP}$$

$$\frac{\bot}{\neg A} \to I, 1$$

(f) $A \rightarrow \neg B \vdash B \rightarrow \neg A$

Solution:

$$\begin{array}{cccc} \overline{A}^2 & A \to \neg B & \mathsf{MP} & \overline{B}^1 \\ \hline & \frac{\neg B}{A} & \frac{\bot}{\neg A} \to I, 2 \\ \hline & \frac{B}{A} \to \neg A} & \to I, 1 \end{array}$$

(g) $\neg (A \land B) \vdash A \rightarrow \neg B$

Solution:

$$\frac{\overline{A}^{1} \quad \overline{B}^{2}}{A \wedge B} \wedge I \quad \neg (A \wedge B) \\ \frac{\bot}{\neg B} \rightarrow I, 2 \\ A \rightarrow \neg B \rightarrow I, 1$$
 MP

(h) $A \vdash \neg \neg A$

Solution:

$$\frac{A \quad \overline{\neg A}^1}{\stackrel{\perp}{\neg \neg A} \rightarrow I, 1} \mathsf{MP}$$

(i) $\neg \neg \neg A \vdash \neg A$

Solution:

$$\frac{\overline{A}^{1}}{\neg \neg A} \text{ THM } \frac{}{\neg \neg \neg A} \text{ MP}$$

$$\frac{\bot}{\neg A} \rightarrow I, 1$$

(j) $\neg A \lor \neg B \vdash \neg (A \land B)$

Solution:

$$\frac{ \frac{\overline{A \wedge B}^1}{\underline{A} \wedge E_l} \wedge E_l \quad \frac{\overline{A}^2}{\neg A^2} \text{MP} \quad \frac{\overline{A \wedge B}^1}{\underline{B} \wedge E_r} \wedge E_r \quad \frac{\overline{B}^3}{\neg B^3} \text{MP} \\ \frac{\underline{\bot}}{\neg A \rightarrow \bot} \rightarrow I, 2 \quad \frac{\underline{\bot}}{\neg B \rightarrow \bot} \rightarrow I, 3 \\ \frac{\underline{\bot}}{\neg (A \wedge B)} \rightarrow I, 1$$

(k) $\neg A \land \neg B \vdash \neg (A \lor B)$

$$\underbrace{\frac{\neg A \wedge \neg B}{\neg A} \wedge E_{l} \quad \overline{A}^{2}}_{\substack{A \vee B^{1}} \quad MP} \underbrace{\frac{\neg A \wedge \neg B}{\neg B} \wedge E_{r} \quad \overline{B}^{3}}_{\substack{A \cup B \cup A \cup B}} MP \underbrace{\frac{\bot}{A \to \bot} \rightarrow I, 3}_{\substack{A \cup B \cup A \cup B \cup A \cup B}} \rightarrow I, 1$$

(I)
$$\neg (A \lor B) \vdash \neg A \land \neg B$$

Solution:

$$\frac{\overline{A}^1}{\frac{A \vee B}{A \vee B} \vee I_r} \vee I_r \vee$$

(m)
$$A \rightarrow \neg B \vdash \neg (A \land B)$$

Solution:

$$\frac{\overline{A \wedge B}^1}{\underline{A} \wedge E_l} \wedge E_l \xrightarrow{A \to \neg B} \mathsf{MP} \quad \frac{\overline{A \wedge B}^1}{\underline{B}} \wedge E_r$$

$$\frac{\bot}{\neg (A \wedge B)} \to I, 1$$

(n)
$$\vdash \neg \neg (A \lor \neg A)$$

$$\frac{\overline{A}^2}{\frac{A \vee \neg A}{A \vee \neg A} \vee I_r} \vee I_r \frac{\neg (A \vee \neg A)^1}{\neg (A \vee \neg A)^1} \text{ MP}$$

$$\frac{\frac{\bot}{\neg A} \to I, 2}{\frac{A \vee \neg A}{A \vee \neg A} \vee I_l} \frac{\neg (A \vee \neg A)^1}{\neg (A \vee \neg A)} \text{ MP}$$

$$\frac{\bot}{\neg \neg (A \vee \neg A)} \to I, 1$$

- Intuitionistic derivations. Provide natural deduction proofs of the following. You do not need to use the *classical* ⊥ rule for these questions, but may find that the *intuitionistic* ⊥ rule is necessary.
 - (a) $A, \neg A \vdash B$

Solution:

$$\frac{A - A}{\frac{\perp}{B}} XF$$

(b) $\neg A \vdash A \rightarrow B$

Solution:

$$\frac{\overline{A}^1 \quad \neg A}{\frac{\bot}{B} \text{ XF}} \text{ MP}$$

$$\frac{A \rightarrow B}{A \rightarrow B} \rightarrow I, 1$$

(c) $\neg A \lor B \vdash A \to B$

Solution:

$$\begin{array}{ccc} & \overline{A}^1 & \overline{\neg A}^2 \\ \frac{\bot}{B} \ \mathsf{XF} & & \\ \underline{-A \lor B} & \overline{-A \to B} & \to I, 2 & \\ & & \underline{B \to B} \\ & & \underline{A \to B} & \to I, 1 \end{array} \lor E$$

(d) $A \vee B$, $\neg A \vdash B$

Solution:

$$\begin{array}{c|c} \overline{A}^1 & \neg A \\ \hline \underline{\frac{\bot}{B}} \ \mathsf{XF} \\ A \vee B & \overline{A \to B} & \rightarrow I, 1 \\ \hline B & \\ \end{array} \lor E$$

(e) $\vdash \neg (B \rightarrow A) \rightarrow (A \rightarrow B)$

Solution:

$$\frac{\overline{A}^2}{\frac{B \to A}{B \to A} \to I} \xrightarrow{\neg (B \to A)^1} \text{MP}$$

$$\frac{\frac{\bot}{B} \text{ XF}}{\frac{A \to B}{\neg (B \to A) \to (A \to B)} \to I, 1}$$

(f) $A \to B, A \to \neg B \vdash A \to C$

$$\frac{\overline{A}^1 \quad A \to B}{B} \text{ MP} \quad \frac{\overline{A}^1 \quad A \to \neg B}{\neg B} \text{ MP}$$

$$\frac{\frac{\bot}{C} \text{ XF}}{A \to C} \to I, 1$$

(g)
$$A \lor B, \neg A, \neg B \vdash C$$