

## 1 Propositional Practice

Note 1

Convert the following English sentences into propositional logic and the following propositions into English. State whether or not each statement is true with brief justification.

- (a) There is a real number which is not rational.
- (b) All integers are natural numbers or are negative, but not both.
- (c) If a natural number is divisible by 6, it is divisible by 2 or it is divisible by 3.
- (d)  $(\forall x \in \mathbb{Z}) (x \in \mathbb{Q})$
- (e)  $(\forall x \in \mathbb{Z}) (((2 \mid x) \vee (3 \mid x)) \implies (6 \mid x))$
- (f)  $(\forall x \in \mathbb{N}) ((x > 7) \implies ((\exists a, b \in \mathbb{N}) (a + b = x)))$

**Solution:**

- (a)  $(\exists x \in \mathbb{R}) (x \notin \mathbb{Q})$ , or equivalently  $(\exists x \in \mathbb{R}) \neg(x \in \mathbb{Q})$ . This is true, and we can use  $\pi$  as an example to prove it.
- (b)  $(\forall x \in \mathbb{Z}) (((x \in \mathbb{N}) \vee (x < 0)) \wedge \neg((x \in \mathbb{N}) \wedge (x < 0)))$ . This is true, since we define the naturals to contain all integers which are not negative.
- (c)  $(\forall x \in \mathbb{N}) ((6 \mid x) \implies ((2 \mid x) \vee (3 \mid x)))$ . This is true, since any number divisible by 6 can be written as  $6k = (2 \cdot 3)k = 2(3k)$ , meaning it must also be divisible by 2.
- (d) All integers are rational numbers. This is true, since any integer number  $n$  can be written as  $n/1$ .
- (e) Any integer that is divisible by 2 or 3 is also divisible by 6. This is false—2 provides the easiest counterexample. Note that this statement is false even though its converse (part c) is true.
- (f) If a natural number is larger than 7, it can be written as the sum of two other natural numbers. This is trivially true, since we can take  $a = x$  and  $b = 0$ .

(Aside: this is a reference to the very weak Goldbach Conjecture (<https://xkcd.com/1310/>).)

## 2 Truth Tables

**Note 1** Determine whether the following equivalences hold, by writing out truth tables. Clearly state whether or not each pair is equivalent.

(a)  $P \wedge (Q \vee P) \equiv P \wedge Q$

(b)  $(P \vee Q) \wedge R \equiv (P \wedge R) \vee (Q \wedge R)$

(c)  $(P \wedge Q) \vee R \equiv (P \vee R) \wedge (Q \vee R)$

**Solution:**

(a) Not equivalent.

$P$	$Q$	$P \wedge (Q \vee P)$	$P \wedge Q$
T	T	T	T
T	F	T	F
F	T	F	F
F	F	F	F

(b) Equivalent.

$P$	$Q$	$R$	$(P \vee Q) \wedge R$	$(P \wedge R) \vee (Q \wedge R)$
T	T	T	T	T
T	T	F	F	F
T	F	T	T	T
T	F	F	F	F
F	T	T	T	T
F	T	F	F	F
F	F	T	F	F
F	F	F	F	F

(c) Equivalent.

$P$	$Q$	$R$	$(P \wedge Q) \vee R$	$(P \vee R) \wedge (Q \vee R)$
T	T	T	T	T
T	T	F	T	T
T	F	T	T	T
T	F	F	F	F
F	T	T	T	T
F	T	F	F	F
F	F	T	T	T
F	F	F	F	F

### 3 Implication

Note 0  
Note 1

Which of the following implications are always true, regardless of  $P$ ? Give a counterexample for each false assertion (i.e. come up with a statement  $P(x,y)$  that would make the implication false).

- (a)  $\forall x \forall y P(x,y) \implies \forall y \forall x P(x,y)$ .
- (b)  $\forall x \exists y P(x,y) \implies \exists y \forall x P(x,y)$ .
- (c)  $\exists x \forall y P(x,y) \implies \forall y \exists x P(x,y)$ .

#### Solution:

- (a) True. For all can be switched if they are adjacent; since  $\forall x, \forall y$  and  $\forall y, \forall x$  means for all  $x$  and  $y$  in our universe.
- (b) False. Let  $P(x,y)$  be  $x < y$ , and the universe for  $x$  and  $y$  be the integers. Or let  $P(x,y)$  be  $x = y$  and the universe be any set with at least two elements. In both cases, the antecedent is true and the consequent is false, thus the entire implication statement is false.
- (c) True. The first statement says that there is an  $x$ , say  $x'$  where for every  $y$ ,  $P(x,y)$  is true. Thus, one can choose  $x = x'$  for the second statement and that statement will be true again for every  $y$ .