SQL on Data Streams

Flink SQL Training

https://github.com/ververica/sql-training



Streams & Dynamic Tables

SQL Was Not (Originally) Designed for Streaming Data

Table are bounded multisets.

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Streams are infinite sequences.

DBMS queries can access all data.



Streaming queries receive data over time.

DBMS queries return a finite result.



Streaming queries continuously emit results and never complete.



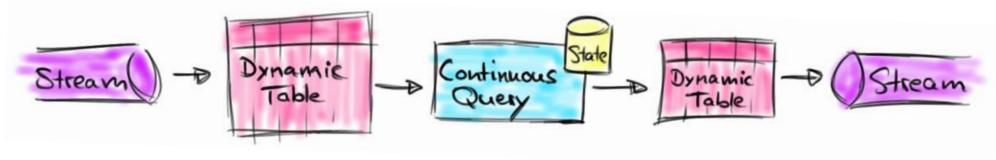
Database Systems Run Queries on Streams

- Materialized views (MV) are similar to regular views, but persisted to disk or memory
 - Used to speed-up analytical queries
 - MVs need to be updated when the base tables change
- MV maintenance is very similar to SQL on streams
 - Base table updates are a stream of DML statements
 - MV definition query is evaluated on that stream
 - MV is query result and continuously updated



Continuous Queries in Flink

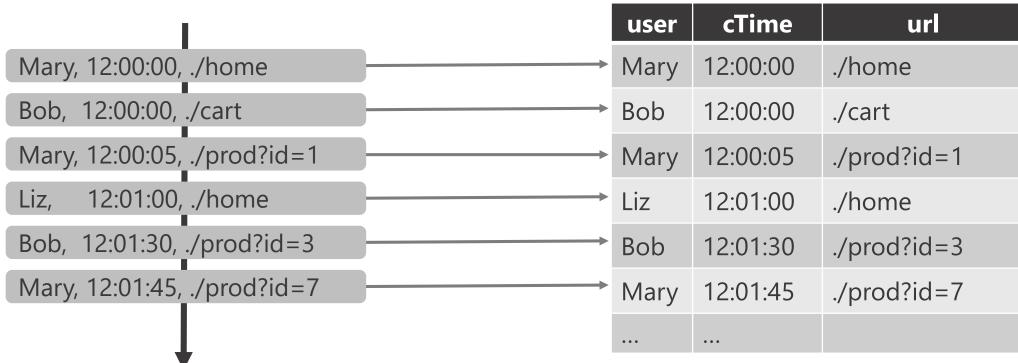
- Core concept is a "Dynamic Table"
 - Dynamic tables are changing over time
- Queries on dynamic tables
 - produce new dynamic tables (which are updated based on input)
 - do not terminate
- Stream
 → Dynamic table conversions





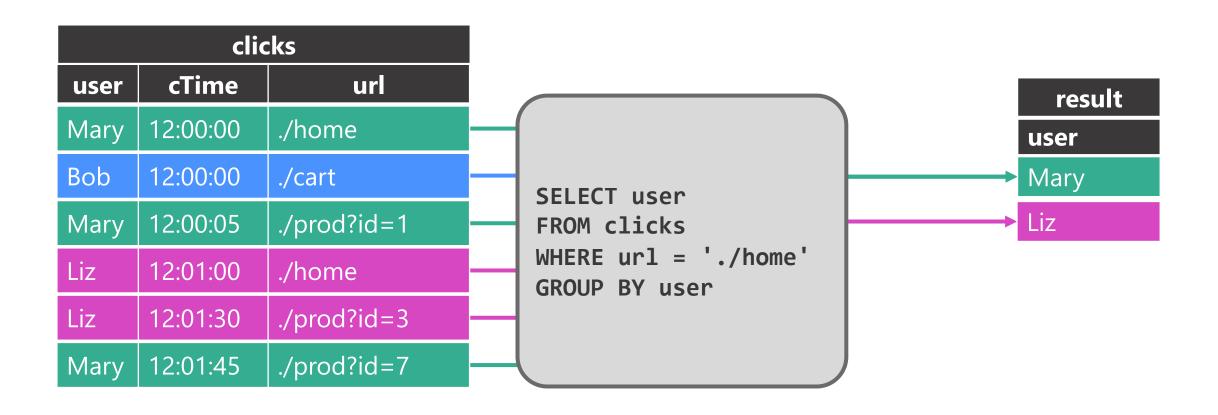
Stream → **Dynamic Table: Append**

- Append mode
 - Stream records are appended to table
 - Table grows as more data arrives





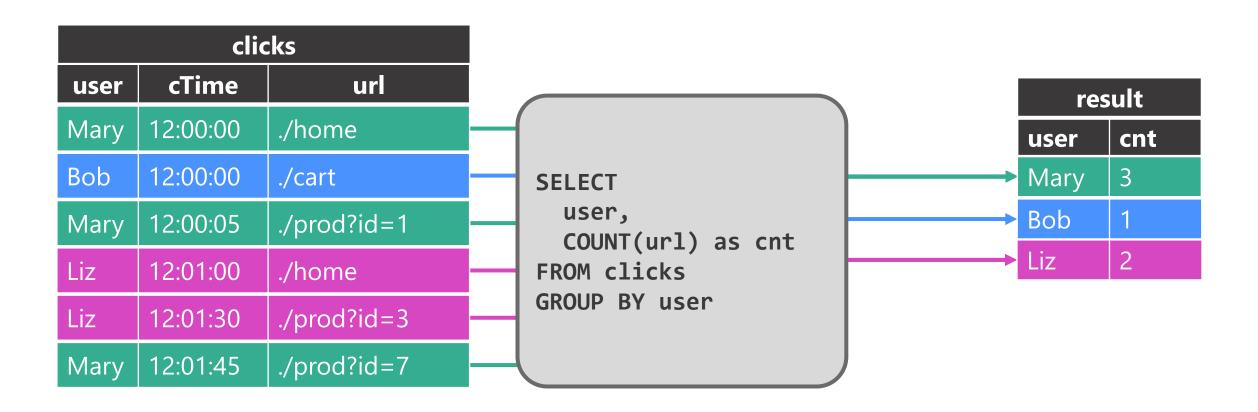
Querying a Dynamic Table



Rows of result table are appended.



Querying a Dynamic Table



Rows of result table are updated.

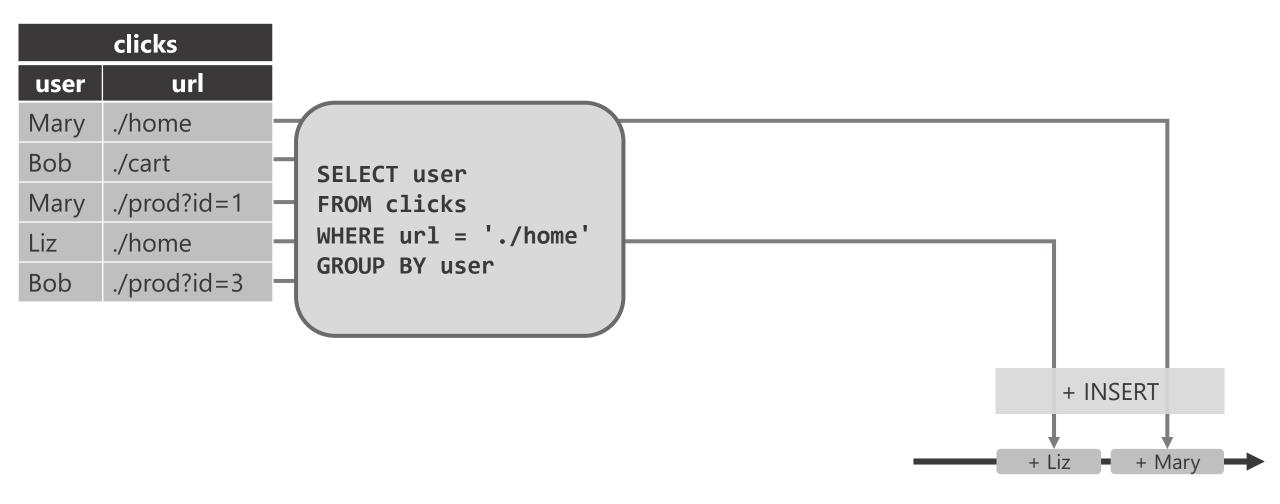


Dynamic Table → **Stream**

- Converting a dynamic table into a stream
- Dynamic tables might update or delete existing rows
- Updates must be encoded in outgoing stream

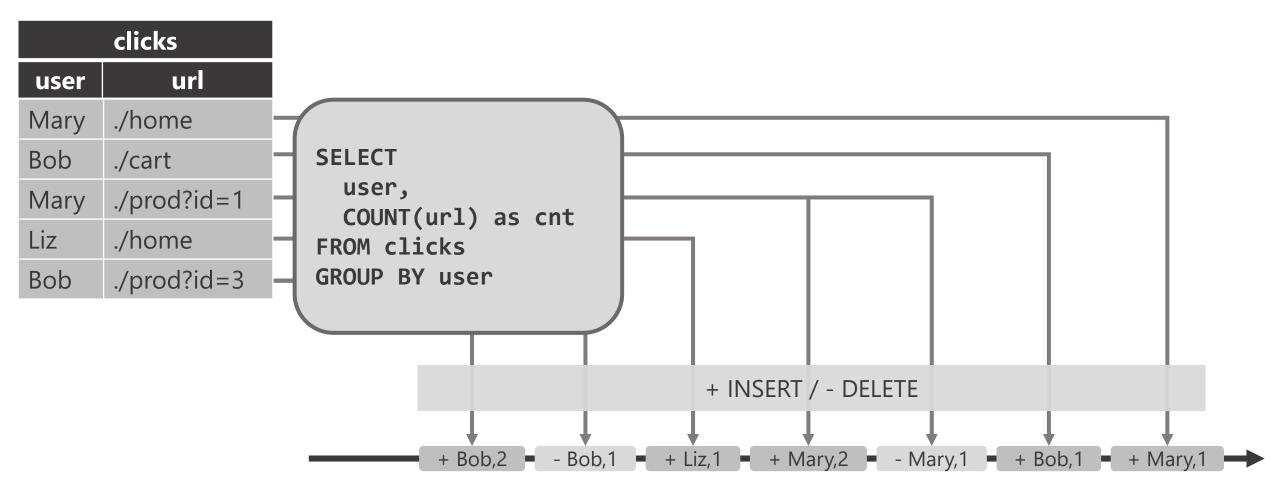


Dynamic Table → **Stream: Append-only**



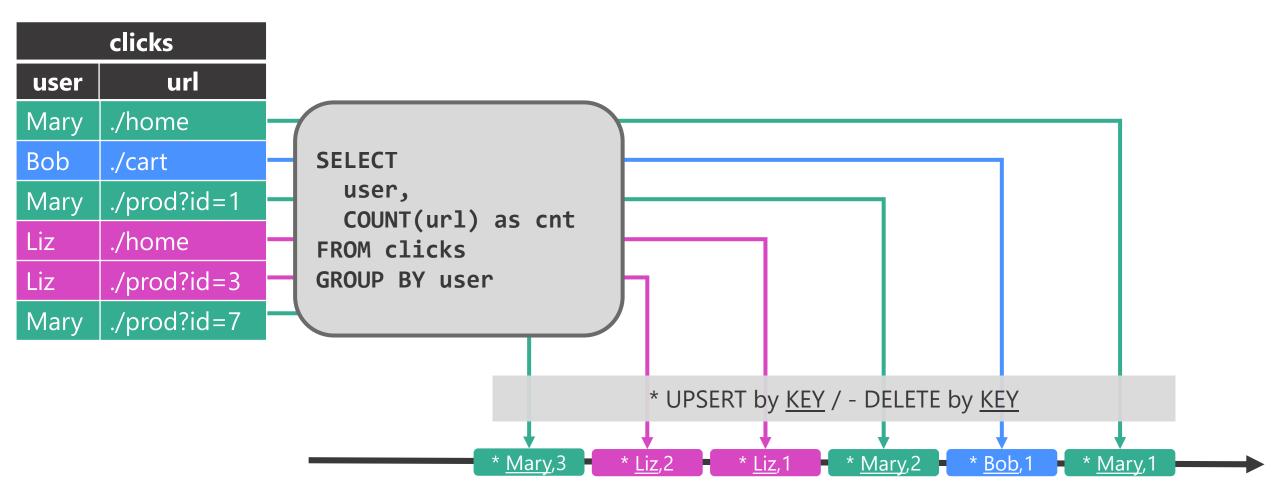


Dynamic Table → **Stream: Retraction**





Dynamic Table → **Stream: Upsert**





Summary

- Streams are interpreted as changelog for a Table
 - Flink 1.7 supports append-only stream to table conversion
 - Upsert and Insert/Update/Delete conversions are on the roadmap
- SQL queries on dynamic tables yield another dynamic table
 - Input and query determines whether resulting dynamic table is append-only or updating
- Dynamic tables can be converted back into streams
 - Append-only, Upsert, or Retraction



Event & Processing Time

How Is Time Handled in Flink SQL?

- Flink SQL supports Event and Processing Time
- Tables may include Time Attributes
 - Time Attributes provide access to event or processing time
 - Time Attributes are (mostly) treated as regular attributes
 - Time Attributes have special type extended from SQL TIMESTAMP
- Time Attributes are declared with the table schema



Event Time Attribute

- Event time attributes carry an actual timestamp
- Timestamps are extracted during table scan
 - Typically taken from an existing field
- Watermarks are generated based on the timestamps
 - Different strategies available
- Event time attributes can be used like a regular TIMESTAMP attribute
 - Event time property (and watermark alignment) are lost when it is modified



Processing Time Attribute

- Processing time attributes are virtual and do not hold data
 - Local machine time is queried on attribute access
- A processing time attribute can be used like a regular TIMESTAMP
 - Loses its processing time property when being modified



How Are Time Attributes Used?

- Some operations require time attributes in certain clauses
 - GROUP BY windows
 - OVER windows
 - Time-windowed joins
 - Joins with temporal tables
- The clicks table is used for the following examples
 - cTime (clickTime) is an event time attribute.

user	cTime	url
Mary	12:00:00	./home
Bob	12:00:00	./cart
Mary	12:00:05	./prod?id=1
•••	•••	•••



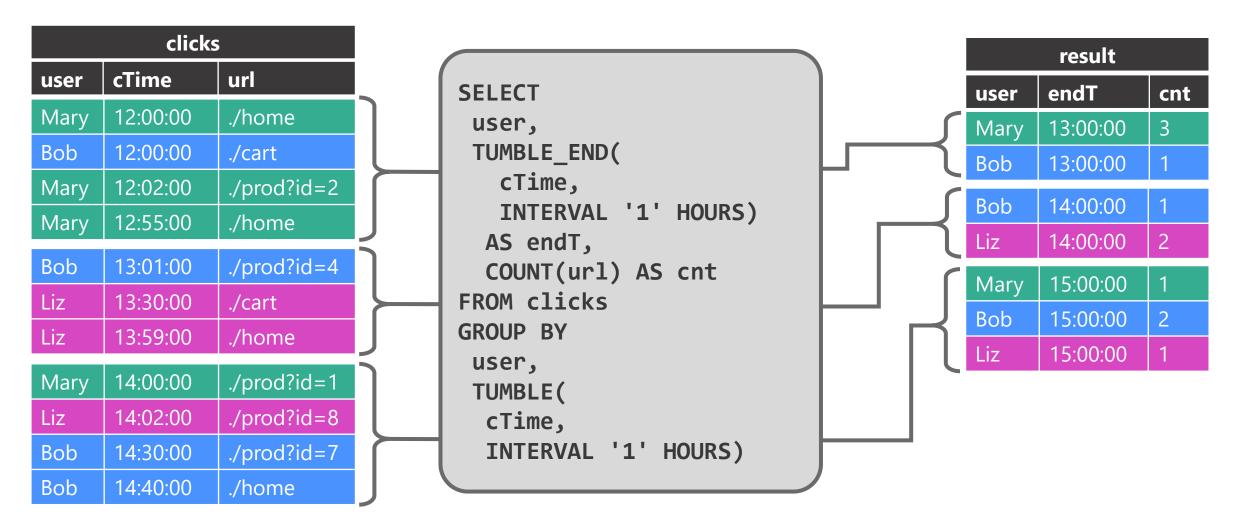
GROUP BY Window Aggregation

Compute number of clicks per hour and user

```
Time attribute
SELECT user,
       TUMBLE END(cTime, INTERVAL '1' HOURS) AS endT,
       COUNT(url) AS cnt
FROM clicks
GROUP BY TUMBLE(cTime, INTERVAL '1' HOURS),
         user
                    Time attribute
```



GROUP BY Window Aggregation





OVER Window Aggregation

Compute for each click how often the URL was clicked in the previous 10 minutes

```
SELECT
  user,
  url,
  COUNT(*) OVER (
    PARTITION BY url
    ORDER BY cTime
    RANGE BETWEEN INTERVAL '10' MINUTE PRECEDING AND CURRENT ROW)
FROM clicks
```



Why Are Time Attributes Special?

- Time attribute values are (quasi) monotonously increasing
- Operators know when rows are no longer needed
 - Watermarks (event time) or wall-clock time (processing time)
- Time-based operators automatically prune their state
 - Only hold the relevant "tail" of the stream is kept in state



Time Attributes and Non-Windowed Operators

- Non-windowed aggregations and joins do not forward time attributes
 - Window operators cannot be applied on their results!
- Non-windowed join does not preserve timestamp order
 - Any row in the state could be join with any new arriving row
 - Order of time attributes is not maintained
 - Non-windowed joins must not emit time attributes.
- Non-windowed aggregation does not forward time attributes
 - Time attributes are converted to regular TIMESTAMP attributes and lose their property
 - SELECT cTime, COUNT(*) FROM clicks GROUP BY cTime
 - cTime is regular TIMESTAMP
 - SELECT user, MAX(cTime) AS cTime FROM clicks GROUP BY user
 - cTime is regular TIMESTAMP



Time Attributes and Batch Processing

- Porting SQL queries from streaming to batch or vice versa
- Queries on event time attributes
 - Batch table requires time attribute of type TIMESTAMP
- Queries on processing time attributes
 - Not supported by batch queries
 - What would be the semantics anyway?



Queries & State

Stateless and Stateful Operators

- Stateless operators
 - Filter
 - Projection
- Stateful streaming operators
 - Window Aggregation (GROUP BY, OVER)
 - Window Joins
- Stateful materializing operators
 - Aggregation
 - Joins



Stateful Materializing Operators

- Materializing operators may never remove state
- The aggregation needs to maintain a count for every user forever.
 - Every user could click at any point in time
- The aggregation state is continuously growing
 - Unless key domain is bounded
- Joins need to materialize all input rows in state

```
SELECT
   user,
   COUNT(url) as cnt
FROM clicks
GROUP BY user
```



Managing State Size

- Query state might grow indefinitely
- Slowly growing state can be addressed by scaling the query
 - SELECT user, COUNT(*) FROM logins GROUP BY user;
- State can be automatically pruned
 - SELECT session, COUNT(*) FROM clicks GROUP BY session;
 - Rows and persisted results can be removed after an idle timeout



Idle State Clean UP

- The query result is not updated when state is removed
- Query result remains consistent if removed state is not needed again
- Query result becomes inconsistent if query needs to access state that was removed!
- Trade the accuracy of the result for size of state

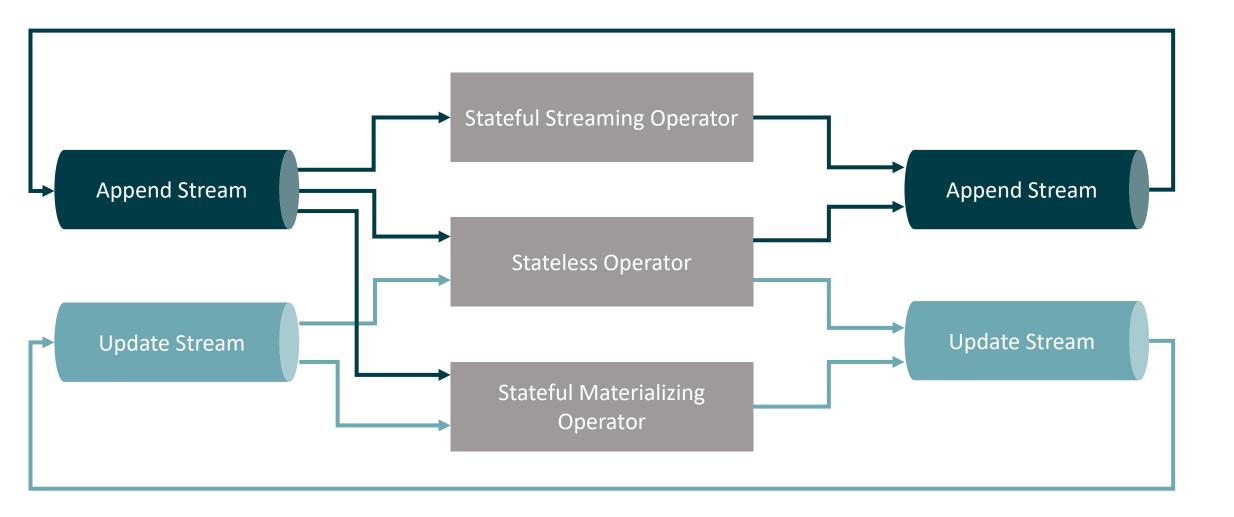


Appending and Updating Input and Output

- Stateful materializing operators may produce updates
 - All required state is available
 - Previous input can be updated
 - Results can be updated
- Stateful streaming operators cannot handle updates
 - State is evicted based on time
 - Results cannot be updated because state might be gone



Append And Update Input and Output





Summary

Summary

- Query Evaluation on Dynamic Tables
 - Stream -> Dynamic Table -> Stream conversions
 - append-only, update, retraction
- Event-Time and Processing-Time in SQL
 - Time attributes & windowed operators
- Queries and State
 - State management and automatic clean up



Hands On Exercises

Running Queries on Streams

Continue with Session 4
"Running Queries on Streams"
in the SQL training wiki:

https://github.com/ververica/sql-training/wiki

We are here to help!





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