SQL ON DATA STREAMS

FABIAN HUESKE, SOFTWARE ENGINEER

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STREAMS & DYNAMIC TABLES



SQL WAS NOT DESIGNED FOR STREAMS

 Table are bounded multisets.



Streams are infinite sequences.

 DBMS queries can access all data.



Streaming queries receive data over time.

• DBMS queries return a finite result.



Streaming queries continuously emit results and never complete.



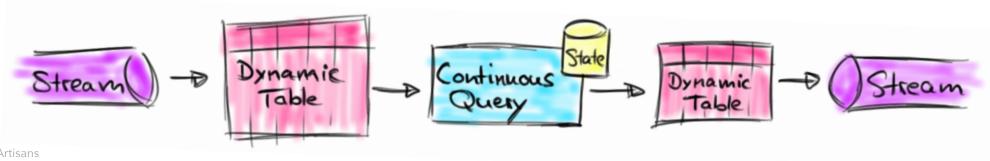
DATABASE SYSTEMS RUN QUERIES ON STREAMS

- Materialized views (MV) are similar to regular views, but persisted to disk or memory
 - Used to speed-up analytical queries
 - MVs need to be updated when the base tables change
- MV maintenance is very similar to SQL on streams
 - Base table updates are a stream of DML statements
 - MV definition query is evaluated on that stream
 - MV is query result and continuously updated



CONTINUOUS QUERIES IN FLINK

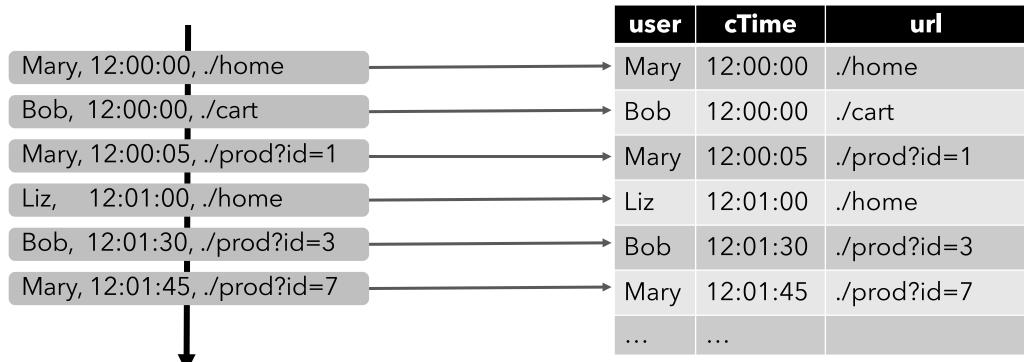
- Core concept is a "Dynamic Table"
 - Dynamic tables are changing over time
- Queries on dynamic tables
 - produce new dynamic tables (which are updated based on input)
 - do not terminate
- Stream → Dynamic table conversions





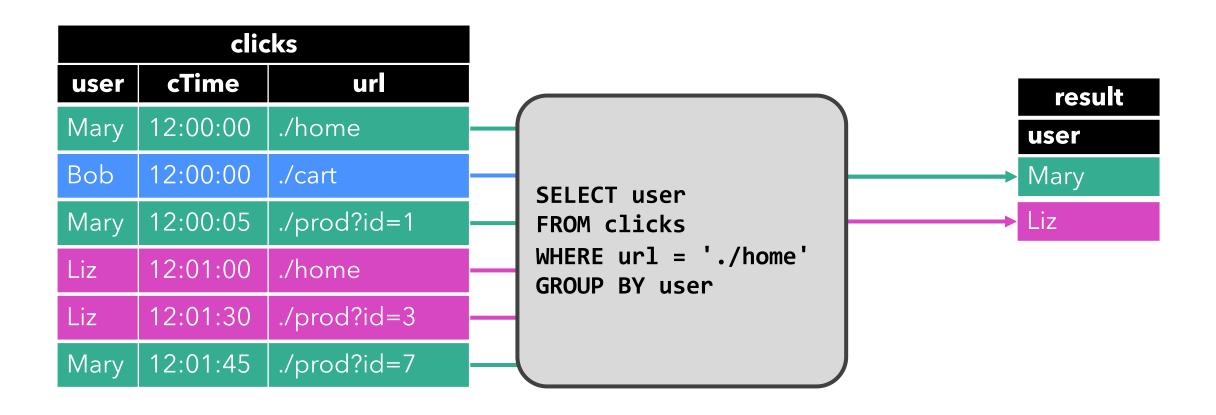
STREAM → DYNAMIC TABLE: APPEND

- Append mode
 - Stream records are appended to table
 - Table grows as more data arrives





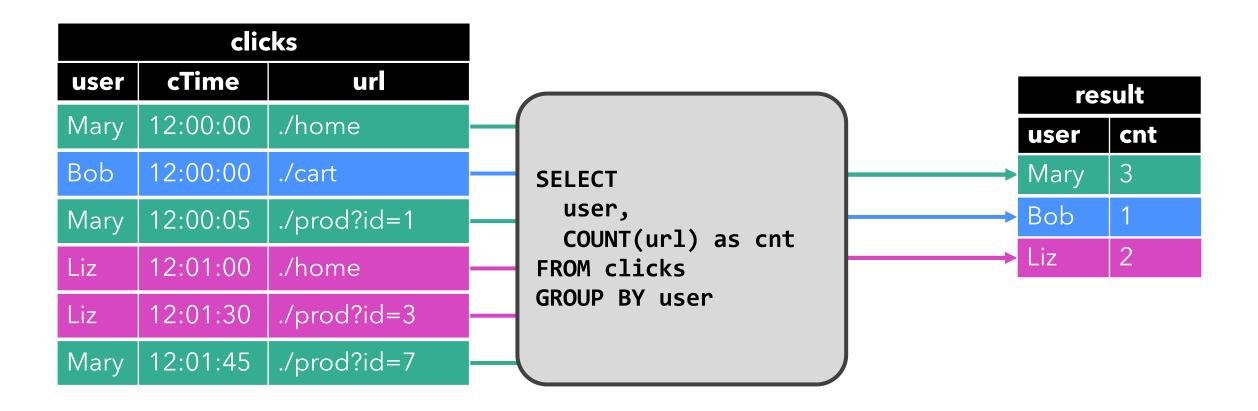
QUERYING A DYNAMIC TABLE



Rows of result table are appended.



QUERYING A DYNAMIC TABLE



Rows of result table are updated.

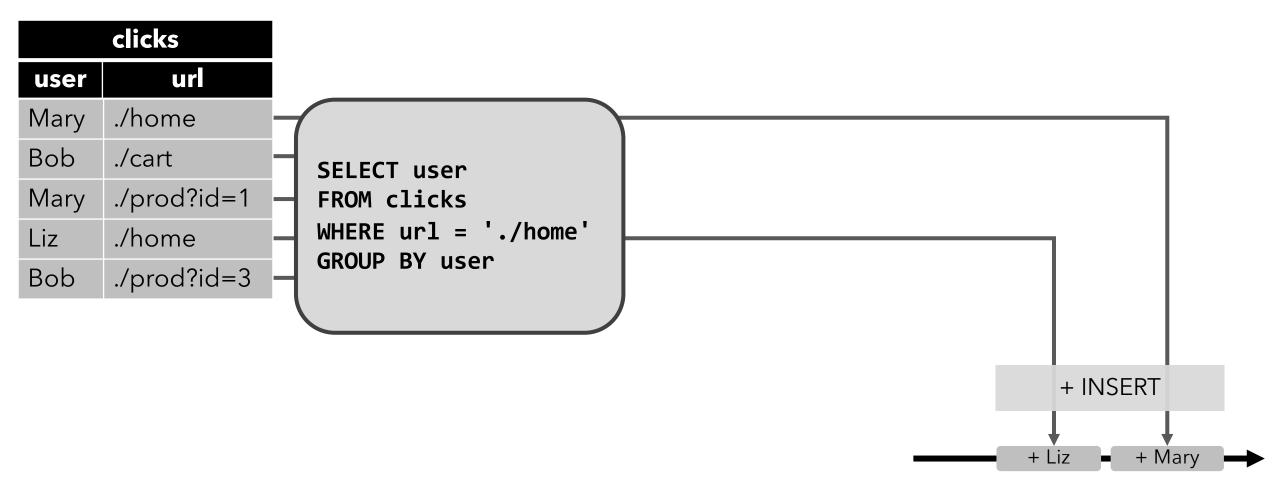


DYNAMIC TABLE → **STREAM**

- Converting a dynamic table into a stream
 - Dynamic tables might update or delete existing rows
 - Updates must be encoded in outgoing stream

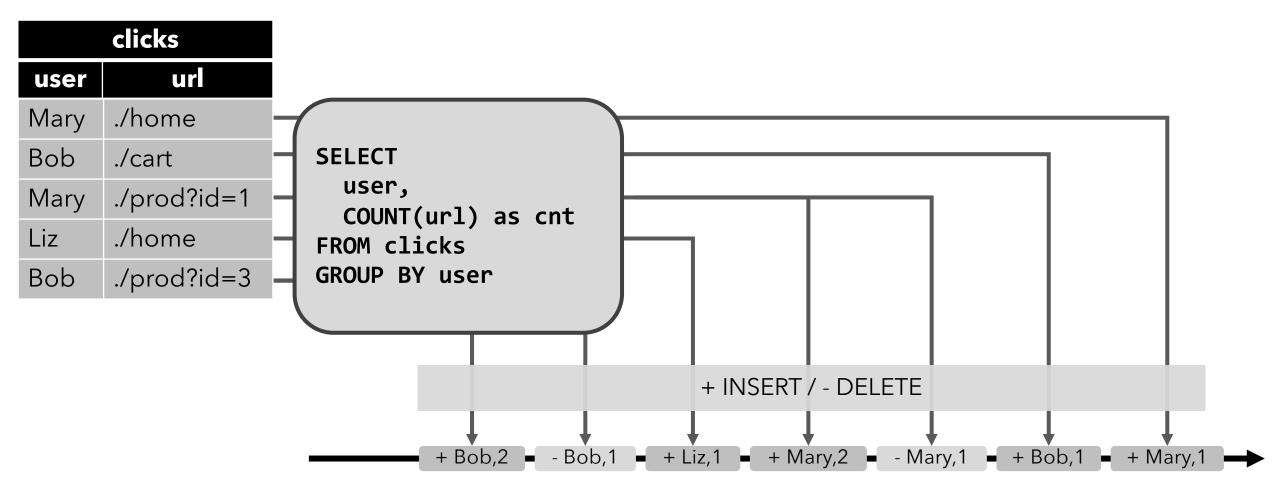


DYNAMIC TABLE \rightarrow STREAM: APPEND-ONLY



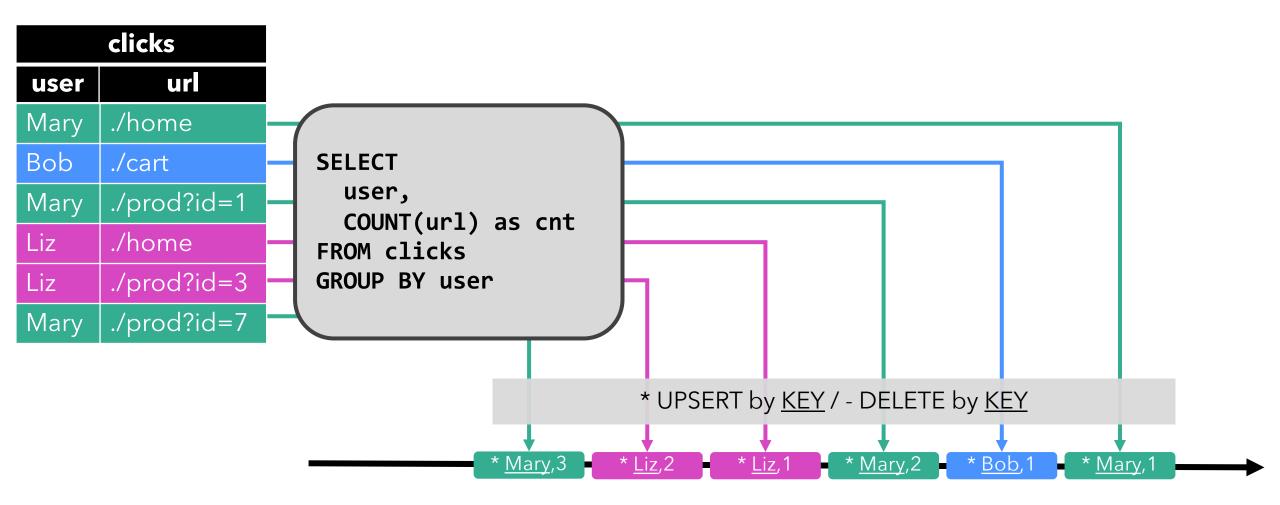


DYNAMIC TABLE \rightarrow STREAM: RETRACTION





DYNAMIC TABLE \rightarrow STREAM: UPSERT





SUMMARY

- Streams are interpreted as changelog for a Table
 - Flink 1.6.0 supports append-only stream to table conversion
 - Upsert and Insert/Update/Delete conversions are on the roadmap
- SQL queries on dynamic tables yield another dynamic table
 - Input and query determines whether resulting dynamic table is append-only or updating
- Dynamic tables can be converted back into streams
 - Append-only, Upsert, or Retraction



EVENT & PROCESSING TIME



HOW IS TIME HANDLED IN FLINK SQL?

- Flink SQL supports Event and Processing Time
- Tables may include *Time Attributes*
 - Time Attributes provide access to event or processing time
 - Time Attributes are (mostly) treated as regular attributes
 - Time Attributes have special type extended from SQL TIMESTAMP

Time Attributes are declared with the table schema



EVENT TIME ATTRIBUTE

- Event time attributes carry an actual timestamp
- Timestamps are extracted during table scan
 - Typically taken from an existing field
- Watermarks are generated based on the timestamps
 - Different strategies available
- Event time attributes can be used like a regular TIMESTAMP
 - Event time property (and watermark alignment) are lost when being modified



PROCESSING TIME ATTRIBUTE

- Processing time attributes are virtual and do not hold data
 - Local machine time is queried on attribute access
- A proc time attribute be used like a regular TIMESTAMP
 - Loses its processing time property when being modified



HOW ARE TIME ATTRIBUTES USED?

- Some operators require time attributes in certain clauses
 - GROUP BY windows
 - OVER windows
 - Time-windowed joins
- The clicks table is used for the following examples
 - -cTime (clickTime) is an event time attribute.

user	cTime	url
Mary	12:00:00	./home
Bob	12:00:00	./cart
Mary	12:00:05	./prod?id=1
•••	• • •	•••



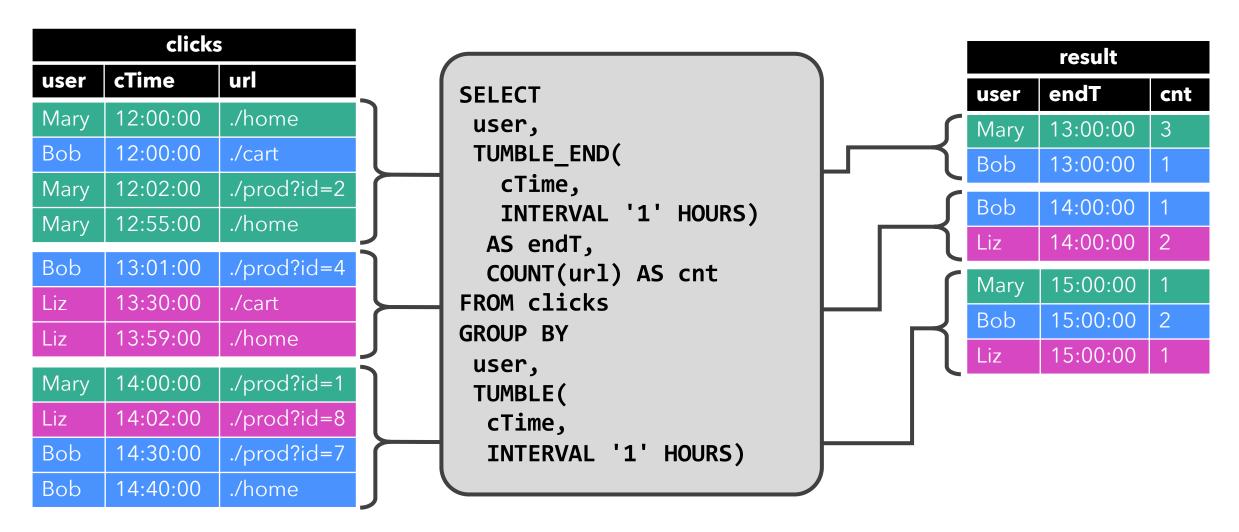
GROUP BY WINDOW AGGREGATION

Compute number of clicks per hour and user

```
Time attribute
SELECT user,
       TUMBLE END(cTime, INTERVAL '1' HOURS) AS endT,
       COUNT(url) AS cnt
FROM clicks
GROUP BY TUMBLE(cTime, INTERVAL '1' HOURS),
         user
                  Time attribute
```



GROUP BY WINDOW AGGREGATION





OVER WINDOW AGGREGATION

Compute for each click
 how often the URL was clicked in the previous 10 minutes

```
user,
url,
COUNT(*) OVER (
PARTITION BY url
ORDER BY cTime
RANGE BETWEEN INTERVAL '10' MINUTE PRECEDING AND CURRENT ROW)
FROM clicks
```



TIME-WINDOWED JOIN

- Find all product URLs that were clicked by a user less than 5 seconds after being served.
 - Assuming a table serves with an event time attribute sTime

```
SELECT url
FROM clicks c, serves s
WHERE s.url LIKE './prod%' AND
    s.url = c.url AND
    s.user = c.user AND
```

url	sTime	user
./home	12:00:00	Mary
./cart	12:00:00	Liz
./prod?id=1	12:00:05	Bob
•••	•••	•••

c.cTime BETWEEN s.sTime AND s.sTime + INTERVAL '5' SECOND





WHY ARE TIME ATTRIBUTES SPECIAL?

- Time attribute values are (quasi) monotonously increasing
- Operators know when rows are no longer needed
 - -Watermarks (event time) or wall-clock time (processing time)
- Time-based operators automatically prune their state
 - -Only hold the relevant "tail" of the stream is kept in state



TIME ATTRIBUTES AND NON-WINDOWED OPERATORS

- Non-windowed aggregations and joins do not forward time attributes
 - Window operators cannot be applied on their results!
- Non-windowed join does not preserve timestamp order
 - Any row in the state could be join with any new arriving row
 - Order of time attributes is not maintained
 - Non-windowed joins must not emit time attributes.
- Non-windowed aggregation does not forward time attributes
 - Time attributes are converted to regular TIMESTAMP attributes and lose their property
 - SELECT cTime, COUNT(*) FROM clicks GROUP BY cTime
 - cTime is regular TIMESTAMP
 - SELECT user, MAX(cTime) AS cTime FROM clicks GROUP BY user
 - cTime is regular TIMESTAMP



TIME ATTRIBUTES AND BATCH PROCESSING

- Porting SQL queries from streaming to batch or vice versa
- Queries on event time attributes
 - -Batch table requires time attribute of type TIMESTAMP
- Queries on processing time attributes
 - Not supported by batch queries
 - -What would be the semantics anyway?



QUERIES AND STATE



STATELESS AND STATEFUL OPERATORS

- Stateless operators
 - Filter
 - Projection
- Stateful streaming operators
 - Window Aggregation (GROUP BY, OVER)
 - -Window Joins
- Stateful materializing operators
 - Aggregation
 - Joins



STATEFUL MATERIALIZING OPERATORS

- Materializing operators may never remove state
- The aggregation needs to maintain a count for every user forever.
 - -Every user could click at any point in time

```
SELECT
  user,
  COUNT(url) as cnt
FROM clicks
GROUP BY user
```

- The aggregation state is continuously growing
 - -Unless key domain is bounded
- Joins need to materialize all input rows in state



MANAGING STATE SIZE

- Query state might grow indefinitely
- Slowly growing state can be addressed by scaling the query
 - -SELECT user, COUNT(*) FROM logins GROUP BY user;
- State can be automatically pruned
 - -SELECT session, COUNT(*) FROM clicks GROUP BY session;
 - -Rows and persisted results can be removed after an idle timeout



IDLE STATE CLEAN UP

- The query result is not updated when state is removed
- Query result remains consistent if removed state is not needed again
- Query result becomes inconsistent if query needs to access state that was removed!
- Trade the accuracy of the result for size of state



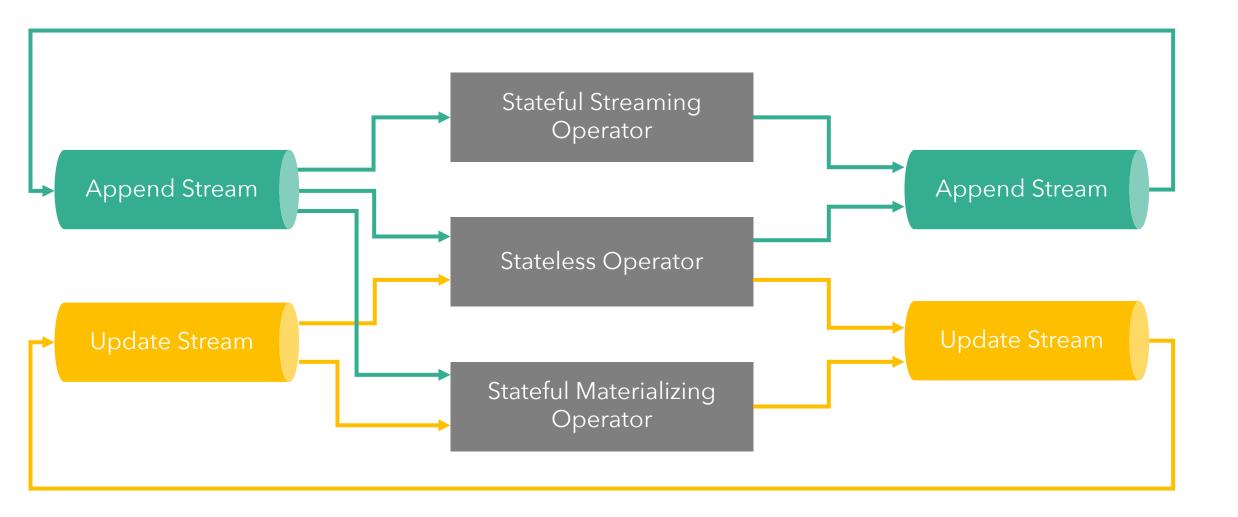
APPEND AND UPDATE INPUT AND OUTPUT

- Stateful materializing operators may produce updates
 - -All required state is available
 - -Previous input can be updated
 - -Results can be updated

- Stateful streaming operators cannot handle updates
 - -State is evicted based on time
 - -Results cannot be updated because state might be gone



APPEND AND UPDATE INPUT AND OUTPUT





OTHER CONSTRAINTS?

 A change of an input table may only trigger a partial re-computation of the result table

SELECT user, RANK() OVER (ORDER BY lastLogin) FROM users;



SUMMARY



SUMMARY

- Query evaluation on Dynamic Tables
 - -Stream -> Dynamic Table -> Stream conversions
 - -append-only, update, retraction
- Event-time and processing-time in SQL
 - -Time attributes & windowed operators
- Queries and State
 - -State management and automatic clean up



HANDS ON



RUNNING QUERIES ON STREAMS

Continue with Session 4 "Running Queries on Streams" in the SQL training wiki:

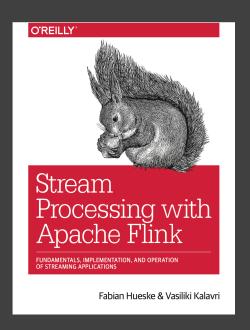
https://github.com/dataArtisans/sql-training/wiki

We are here to help!



THANK YOU!

@fhueske@dataArtisans@ApacheFlink



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