ULFM

Обзор функциональных возможностей

User Level Failure Mitigation

The User Level Failure Mitigation (ULFM) proposal was developed by the MPI Forum's Fault Tolerance Working Group [1].

Designing the mechanism that users would use to manage failures was built around three concepts:

- 1. simplicity, the API should be easy to understand and use in most common scenarios;
- 2. **flexibility**, the API should allow varied fault tolerant models to be built as external libraries and;
- **3. absence of deadlock**, no MPI call (point-to-point or collective) can block indefinitely after a failure, but must either succeed or raise an MPI error.
- [1] http://fault-tolerance.org/

User Level Failure Mitigation

To use this ULFM implementation, an MPI application must change the default error handler on (at least)

- □ MPI_COMM_WORLD from MPI_ERRORS_ARE_FATAL
- to either
- ☐ MPI_ERRORS_RETURN or a custom MPI Errorhandler

User Level Failure Mitigation

Предложение по смягчению последствий на уровне пользователя (ULFM) было разработано рабочей группой MPI Forum Fault Tolerance [1].

Проектирование механизма, который пользователи будут использовать для управления отказами, строится вокруг трех концепций:

- **1. простота,** АРІ должен быть легко понят и использоваться в большинстве распространенных сценариев;
- **2. гибкость,** API должен позволять создавать различные отказоустойчивые модели как внешние библиотеки;
- **3. отсутствие взаимоблокировки,** никакой вызов MPI (точка-точка или коллективный) не может блокироваться неограниченно после сбоя, он должен либо корректно завершиться, либо сообщить об ошибке.
- [1] http://fault-tolerance.org/

Overview

- ☐ Code Repository: https://bitbucket.org/icldistcomp/ulfm
- ☐ Bulding&Running: http://fault-tolerance.org/ulfm/ulfm-setup/
- ☐ ULFM proposal ticket: https://github.com/mpi-forum/mpi-issues/20
- □ MPI4.0 ???
- ☐ ULFM is prototype implementation in OpenMPI

Сборка&Запуск&Тестирование

- http://fault-tolerance.org/ulfm/ulfm-setup/
- ☐ hg clone ssh://hg@bitbucket.org/icldistcomp/ulfm
- ☐ ./configure && make && make install
- □ Сборка и установка происходит также как и ОрепМРІ

You will need to add one additional header to get access to the new functions, just after the include for mpi.h:

For C/C++ this looks like:

```
#include <mpi.h>
#include <mpi-ext.h>
```

Error-handler sample

```
static void verbose errhandler (MPI Comm* comm, int* err, ...) {
   int rank, size, len;
   char errstr[MPI MAX ERROR STRING];
   MPI Comm rank (MPI COMM WORLD, &rank);
   MPI Comm size (MPI COMM WORLD, &size);
   MPI Error string( *err, errstr, &len );
   printf("Rank %d / %d: Notified of error %s\n", rank, size, errstr);
int main(int argc, char *argv[]) {
   int rank, size;
   MPI Errhandler errh;
   MPI Init(NULL, NULL);
   MPI Comm rank (MPI COMM WORLD, &rank); MPI Comm size (MPI COMM WORLD, &size);
   MPI Comm create errhandler(verbose errhandler, &errh); // создаем обработчик
   MPI Comm set errhandler (MPI COMM WORLD, errh);
   if (rank == (size-1)) { raise(SIGKILL); } // Убиваем процесс
   MPI Barrier (MPI COMM WORLD);// Точка синхронизации
   printf("Rank %d / %d: Stayin' alive!\n", rank, size);
   MPI Finalize();
   return 0;
```

Respawn sample

```
static int MPIX Comm replace(MPI Comm comm, MPI Comm *newcomm) {
        MPIX Comm shrink(comm, &scomm); MPI Comm size(scomm, &ns); MPI Comm size(comm, &nc);
        MPI Comm set errhandler( scomm, MPI ERRORS RETURN );
        rc = MPI Comm spawn(...);
static int app needs repair(MPI Comm comm) {
        if (comm == world) {
                 worldi = (worldi + 1) % 2;
                 if (MPI COMM NULL != world) { MPI Comm free(&world); }
                 MPIX Comm replace (comm, &world);
                 if (world == comm) {
                          return false; // Ok, we repaired nothing, no need to redo any work
        return true; // We have repaired the world, we need to reexecute
static void errhandler respawn (MPI Comm* pcomm, int* errcode, ...) {
        if (MPIX ERR PROC FAILED != eclass && MPIX ERR REVOKED != eclass) {
                 MPI Abort (MPI COMM WORLD, *errcode);
        MPIX Comm revoke (*pcomm);
        app needs repair(*pcomm);
int main( int argc, char* argv[] ) {
        MPI Comm create errhandler(&errhandler respawn, &errh);
        while(iteration < max iterations || app needs repair(world)) {</pre>
                 if (iteration == error iteration && victim) raise( SIGKILL );
                 rc = MPI Bcast(array, COUNT, MPI DOUBLE, 0, world);
```

Базовые примеры

http://fault-tolerance.org/ulfm/usage-guide/

- ☐ Master/Worker The example below presents a master code that handles failures by ignoring failed processes and resubmitting requests. It demonstrates the different failure cases that may occur when posting receptions from MPI_ANY_SOURCE as discussed in the advice to users in the proposal.
- ☐ Iterative Refinement The example below demonstrates a method of fault-tolerance to detect and handle failures. At each iteration, the algorithm checks the return code of the MPI_ALLREDUCE. If the return code indicates a process failure for at least one process, the algorithm revokes the communicator, agrees on the presence of failures, and later shrinks it to create a new communicator.

```
int master(void) {
    MPI Comm set errhandler (comm, MPI ERRORS RETURN);
   MPI Comm size(comm, &size);
    // ... submit the initial work requests ...
    MPI Irecv (buffer, 1, MPI INT, MPI ANY SOURCE, tag, comm, &req);
   // Progress engine: Get answers, send new requests,
    // and handle process failures
    while((active workers > 0) && work available) {
        rc = MPI Wait(&req, &status); // ТОЧКА СИНХРОНИЗАЦИИ!
        if((MPI ERR PROC FAILED == rc) | (MPI ERR PENDING == rc)) {
            MPI Comm failure ack(comm); MPI Comm failure get acked(comm, &g);
            MPI Group size(q, &gsize);
            // ... find the lost work and requeue it ...
            active workers = size - gsize - 1;
            MPI Group free (&g);
            // repost the request if it matched the failed process
            if(rc == MPI ERR PROC FAILED) {
                MPI Irecv (buffer, 1, MPI INT, MPI ANY SOURCE, tag, comm, &req);
            continue;
        // ... process the answer and update work available ...
        MPI Irecv(buffer, 1, MPI INT, MPI ANY SOURCE, tag, comm, &req);
     // ... cancel request and cleanup ...
```

FAULT TOLERANT MPI APPLICATIONS WITH ULFM

George Bosilca, UTK
Keita Teranishi, SNL
Marc Gamell, Rutgers
Tsutomu Ikegami, AIST
Sara Salem Hamouda, ANU

And many more

SC'15 BoF Austin, TX, USA



ULFM-based Applications

- ORNL: Molecular Dynamic simulation, C/R in memory with Shrink
- **UAB**: transactional FT programming model
- Tsukuba: Phalanx Master-worker framework
- Georgia University: Wang Landau Polymer Freezing and Collapse, localized subdomain C/R restart
- Sandia, INRIA, Cray: PDE sparse solver

Programming models: resilient X10

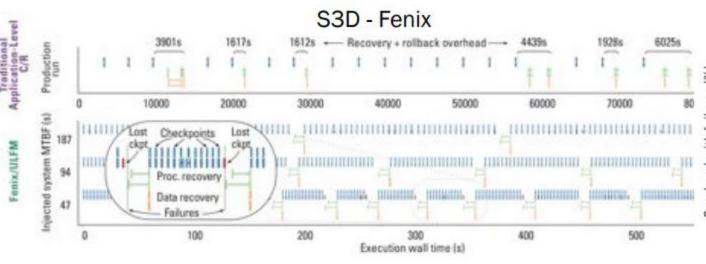
- Cray: CREST miniapps, PDE solver Schwartz, PPStee (Mesh, automotive), HemeLB (Lattice Boltzmann)
- ETH Zurich: Monte-Carlo, on failure the global communicator (that contains spares) is shrunk, ranks reordered to recreate the same domain decomposition

 $E(\rho) [kg/m^2]$ mean of rho at t=0.06 mean of rho at t=0.06 17.5 15.0 15.0 10.0 7.5 (a) failure-free (b) few failures (c) many failures

Figure 5. Results of the FT-MLMC implementation for three different failure scenarios.

FRAMEWORKS USING ULFM

LFLR, FENIX, FTLA, Falanx



In Proceedings of SC '14

35 Traditional Application-Level C/R 30 25 Overhead with failures (%) Image courtesy of the authors, M.Gamell, D.Katz, H.Kolla, J.Chen, S.Klasky, and M.P. Exploring automatic, online failure recovery for scientific applications at extreme s 9600 MITREIN

Credits: ETH Zurich

User activities

- ORNL: Molecular Dynamic simulation
 - Employs coordinated user-level C/R, in place restart with Shrink
- UAB: transactional FT programming model
- Tsukuba: Phalanx Master-worker framework
- Georgia University: Wang Landau Polymer Freezing and Collapse
 - Employs two-level communication scheme with group checkpoints
 - Upon failure, the tightly coupled group restarts from checkpoint, the other distant groups continue undisturbed
- Sandia: PDE sparse solver
- INRIA: Sparse PDE solver

- Cray: CREST miniapps, PDE solver Schwartz, PPStee (Mesh, automotive), HemeLB (Lattice Boltzmann)
- UTK: FTLA (dense Linear Algebra)
 - Employs ABFT
 - FTQR returns an error to the app, App calls new BLACS repair constructs (spawn new processes with MPI_COMM_SPAWN), and re-enters FTQR to resume (ABFT recovery embedded)
- ETH Zurich: Monte-Carlo
 - Upon failure, shrink the global communicator (that contains spares) to recreate the same domain decomposition, restart MC with same rank mapping as before

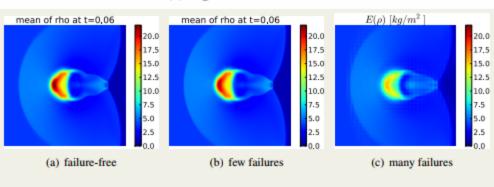
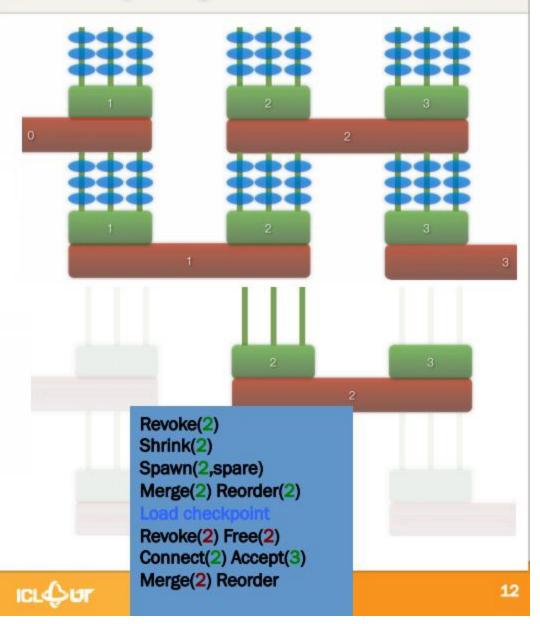


Figure 5. Results of the FT-MLMC implementation for three different failure scenarios.

Credits: ETH Zurich

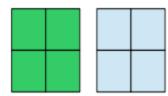
U-GA: Wang-Landau polymer Freeze

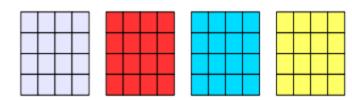
- Long independent computation on each processor
 - Dataset protected by small, cheap checkpoints (stored on neighbors)
- Periodically, an AllReduce on the communicator of the Energy window
- Immediately after, a Scatter and many pt2pt on the communicator linking neighboring energy windows



ETH-Zurich: Monte-Carlo PDE

- X is the solution to a stochastic PDE
- Each sample X^i is computed with a FVM solver
- MC error is determined by
 - stochastic error (depends on M)
 - discretization error (depends on the mesh-width h)
- A more accurate MC approximation requires more samples M and a finer mesh h



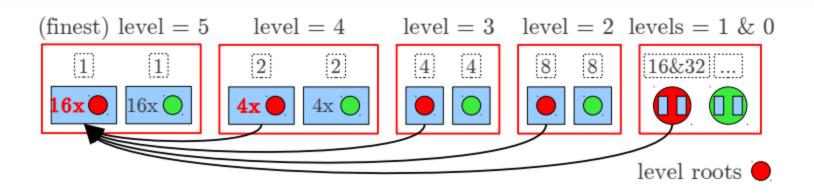


Fault Tolerant Monte Carlo:

ullet The number of samples M turns into a random variable \hat{M}

•
$$\|\mathbb{E}[X] - E_{\hat{M}}[X]\| \leq \mathbb{E}\left[\frac{1}{\sqrt{\hat{M}}}\right] \|X\|$$

ETH-Zurich: Monte-Carlo PDE



- Try to collect the mean as in fault-free ALSVID-UQ
- Call MPI_BARRIER on MPI_COMM_WORLD at the end to discover failed processes
- non-uniform success of MPI_BARRIER: MPI_BARRIER is followed by MPI_COMM_AGREE
- In case of failure: (Re)assign the level roots and repeat the collection of the means