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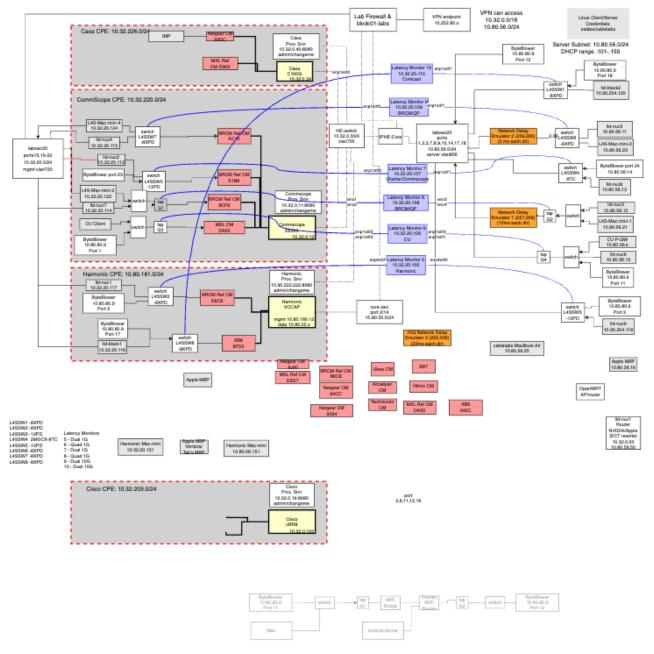


### **Intro**



- https://www.cablelabs.com/event/interoplabs-l4s-nov-2024
- Test setup with DOCSIS equipment
  - Downstream/upstream configured to 100/10Mbps
  - Linux PCs (Ubuntu 22.04 w/ Prague kernel 5.15)
  - Apple MAC Mini PCs with apple-quic (Sequoia 15.1)
- SCReAM BW test from <a href="https://github.com/EricssonResearch/scream">https://github.com/EricssonResearch/scream</a> e0af5ef4613a38a84f99b7995eef38c69a5c08fa









- Downstream test
- Max 100Mbps throughput
- RTT 12ms
  - 32 and 52ms implemented as tc qdisc netem on feedback path
- SCReAM competes with TCP Prague over the L4S queue

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## SCReAM vs up to 2 TCP prague downstream (1/2)

- RTT 12ms
- SCReAM starts
- One TCP added efter 20s
- Another TCP added after 40s
- Both TCPs terminated after 60:
- SCReAM has ~20% lower bitrathan TCP

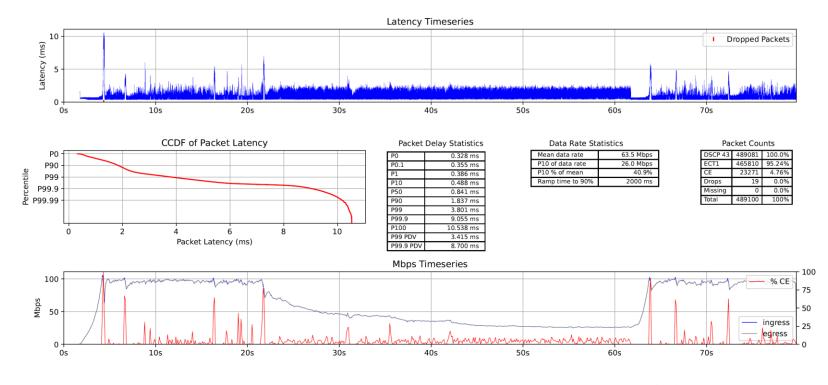
# Throughput Timeseries Downstream 80 60 Flow 1 Flow 2 Flow 3 Total

Time (s)

Aggregate thruput downstream.pdf

## SCReAM vs up to 2 TCP Prague downstream (2/2)

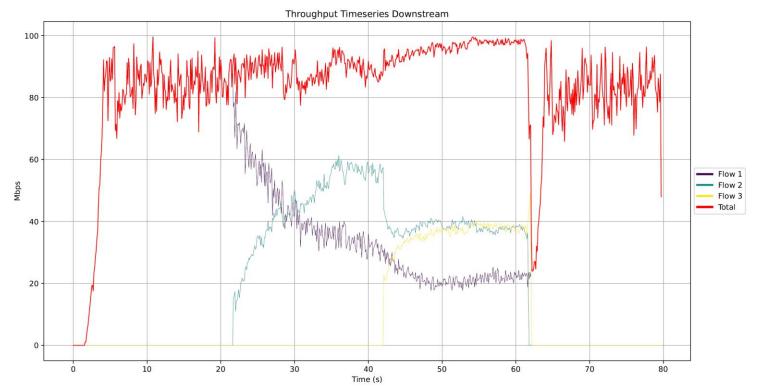
1\_downstream.pdf UDP [10.80.56.11 50000] to [10.32.226.158 50000]



# SCReAM variable video frame size vs up to 2 TCP prague downstream (1/2)

- As previous but SCReAM models variable frame sizes
  - -rand 50
- SCReAM adds extra headroom cope with varying frame sizes
   lower bitrate.
  - Intended behavior
- Slightly lower SCreAM bitrate with competing TCPs

#### Aggregate\_thruput\_downstream.pdf

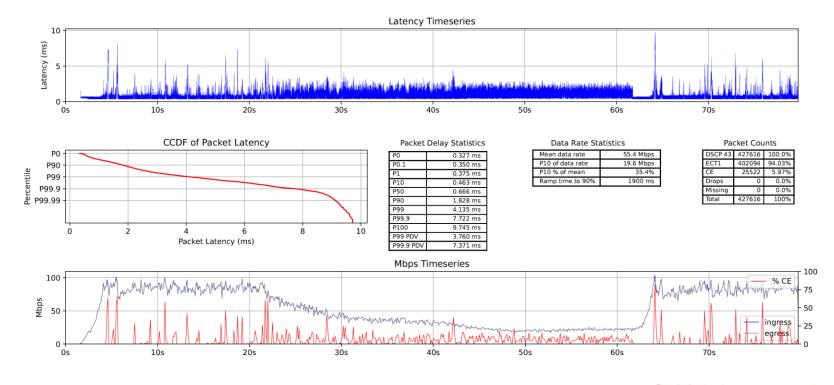




# SCReAM variable video frame size vs up to 2 TCP prague downstream (1/2)

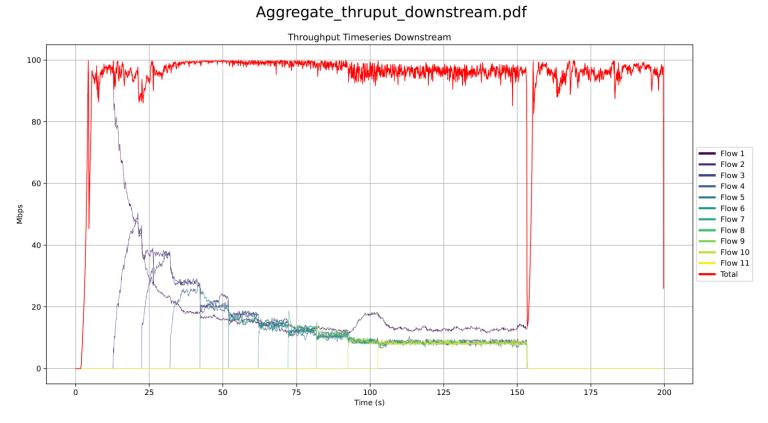
 Variable frame size gives slightly higher delay jitter, which is expected

1 downstream.pdf UDP [10.80.56.11 50000] to [10.32.226.158 50000]



## SCReAM vs up to 10 TCP prague downstream

- RTT 12ms
- SCReAM starts
- One TCP added every 10s
- All TCPs terminated after 150s
- SCReAM copes with many competing TCP Prague quite well



# SCReAM variable video frame size vs up to 10 TCP prague downstream

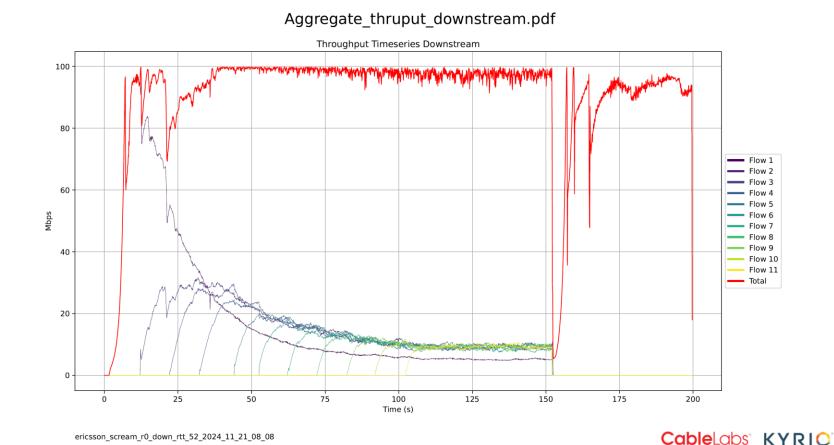
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- As previous but SCReAM models variable frame sizes
  - -rand 50
- SCReAM copes with many competing TCP Prague quite well

## Aggregate thruput downstream.pdf Throughput Timeseries Downstream Total Marie 25 125 175 200 Time (s)

## SCReAM vs up to 10 TCP prague downstream (1/4)

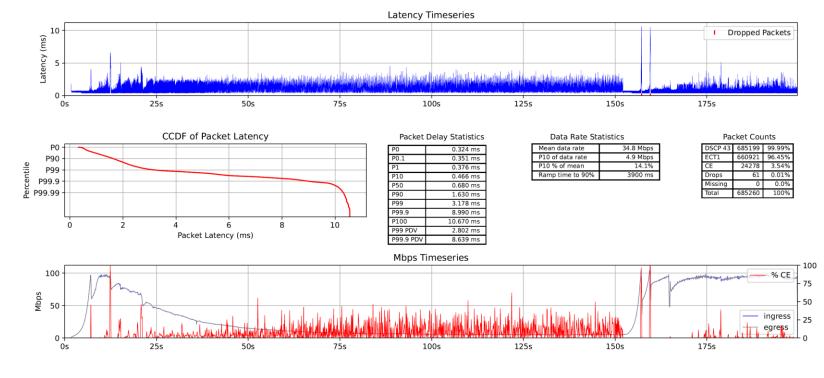
- RTT 52ms
- SCReAM starts
- One TCP added every 10s
- All TCPs terminated after 150s
- SCReAM gets roughly 40% lower bitrate
  - Possibly additional RTT compensation needed
- But SCReAM is not starved out



## SCReAM vs up to 10 TCP prague downstream (2/4)

- RTT 52ms
- Overshoot when TCP load terminated leads to packet loss
- Possibly 5%
   multiplicative
   increase is a bit
   overoptimistic

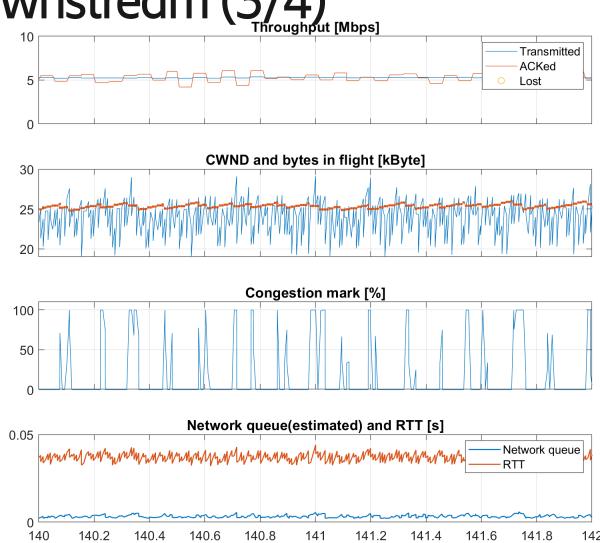
1\_downstream.pdf UDP [10.80.56.11 50000] to [10.32.226.158 50000]





vs up to 10 TCP prague downstream (3/4)

- RTT 32ms
  - no data for 52ms RTT
- SCReAM + 10 TCP Prague
- SCReAM detailed log
  - -log option
- Packet marking appears to come in bursts
  - Possible synchronization in between flows?
  - Burstiness may hit SCReAM more than Prague

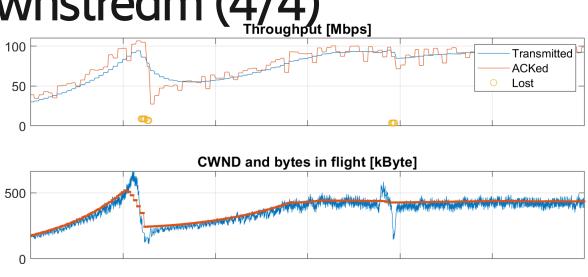


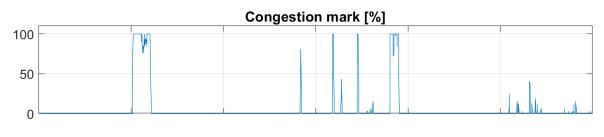
T [s]

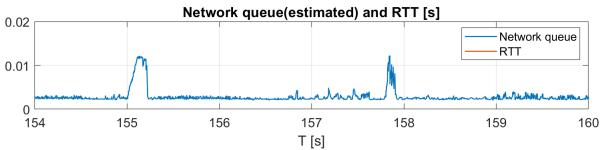


vs up to 10 TCP prague downstream (4/4)
Throughput [Mbps]

- RTT 32ms
  - no data for 52ms RTT
- SCReAM overshoot causes packet loss
- Possibly too low drop threshold for L4S queue?
- 5% multiplicative increase in SCReAM may be too optimistic









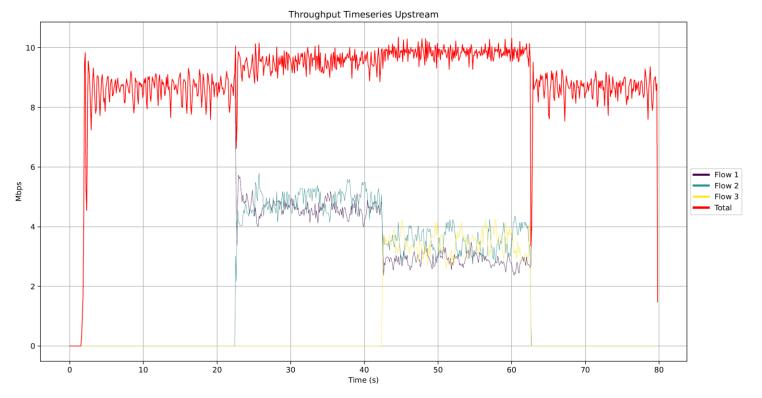


- Upstream test
- Max 10Mbps throughput
- RTT 12ms
- SCReAM competes with TCP Prague over the L4S queue

## SCReAM vs up to 2 TCP prague downstream (1/2)

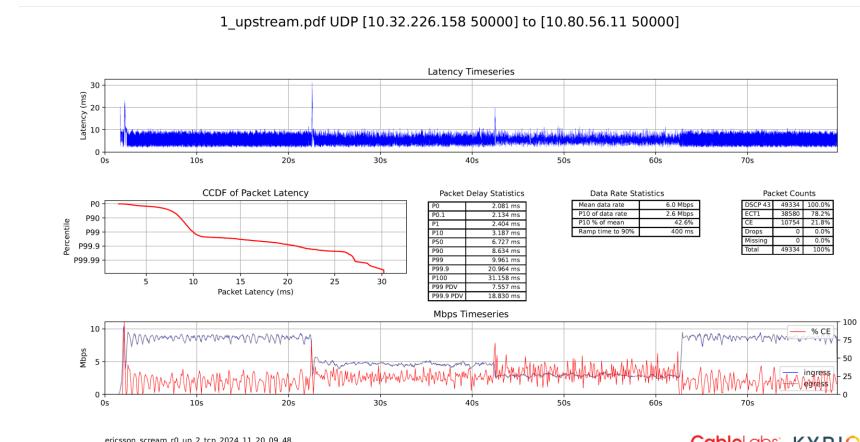
- RTT 12ms
- SCReAM starts
- One TCP added efter 20s
- Another TCP added after 40s
- Both TCPs terminated after 60
- SCreAM alone does not reach full link ulitization
  - ...but competes well with TCPrague

#### Aggregate\_thruput\_upstream.pdf





## SCReAM vs up to 2 TCP Prague downstream (2/2)



## SCReAM variable video frame size vs up to 2 TCP prague downstream (1/2)

- As previous but SCReAM models variable frame sizes
  - -rand 50
- Lower SCreAM bitrate with competing TCPs
  - ...but no starvation

# Throughput Timeseries Upstream Flow 1 Flow 2 Flow 3 Total

Time (s)

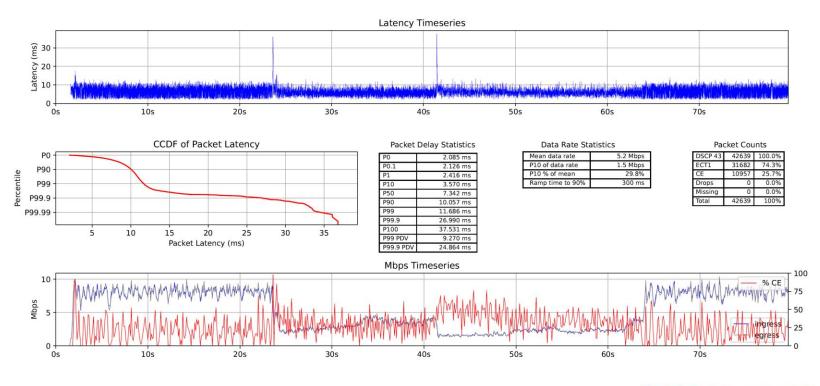
Aggregate\_thruput\_upstream.pdf



# SCReAM variable video frame size vs up to 2 TCP prague downstream (1/2)

 Variable frame size gives slightly higher delay jitter which is expected

1 upstream.pdf UDP [10.32.226.158 50000] to [10.80.56.11 50000]





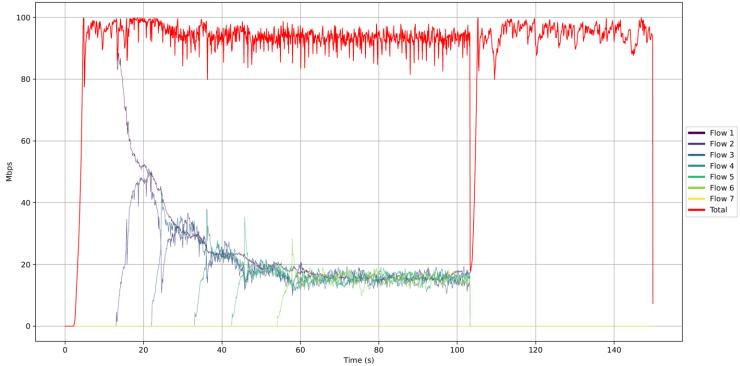


- Downstream test
- Max 100Mbps throughput
- RTT 12ms
- SCReAM competes with up to 5 Apple QUIC over the L4S queue
  - iperf3-darwin -i 0 -t 200 --apple-quic --apple-l4s -p PORT -c IP\_ADDRESS

## SCReAM vs up to 5 Apple QUIC downstream (1/2)

- RTT 12ms
- SCReAM starts
- One QUIC added every 10s
- All QUICs terminated after 100s
- Near perfect sharing of capacity
- QUIC has more uneven rate than TCP Prague

## Aggregate\_thruput\_downstream.pdf Throughput Timeseries Downstream

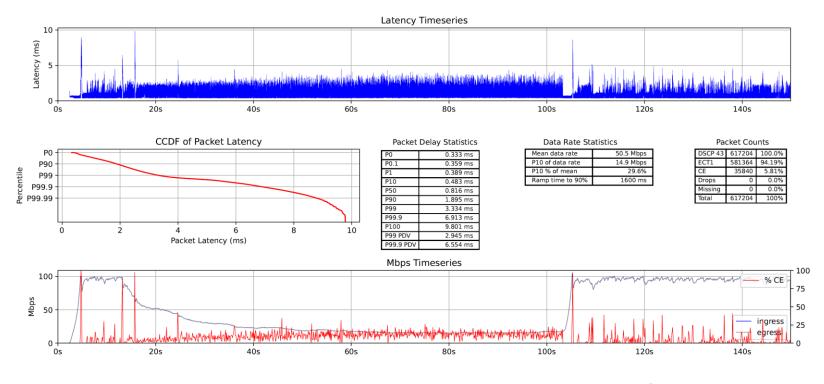






## SCReAM vs up to 5 Apple QUIC downstream (1/2)

1\_downstream.pdf UDP [10.80.56.11 50000] to [10.32.226.158 50000]







- Previous issues with SCReAM becoming starved out by TCP Prague appears to have been resolved
- SCReAM gets a lower share than TCP Prague when RTT is higher
  - Issue is not seen in <a href="https://github.com/EricssonResearch/scream/blob/master/test-record.md">https://github.com/EricssonResearch/scream/blob/master/test-record.md</a>
  - Issue is perhaps related to synchronization between flows (bursty marking)?
- SCReAM overshoot at increased capacity can give packet loss
  - Perhaps increase drop threshold in L4S queue ?.
  - 5% multiplicative increase in SCReAM may be overoptimistic.
- SCReAM performs fine with competing Apple QUIC flows.
- Over all very good test environment.

