

#### Università degli Studi Roma Tre Dipartimento di Ingegneria Computer Networks Research Group

### kathara lab

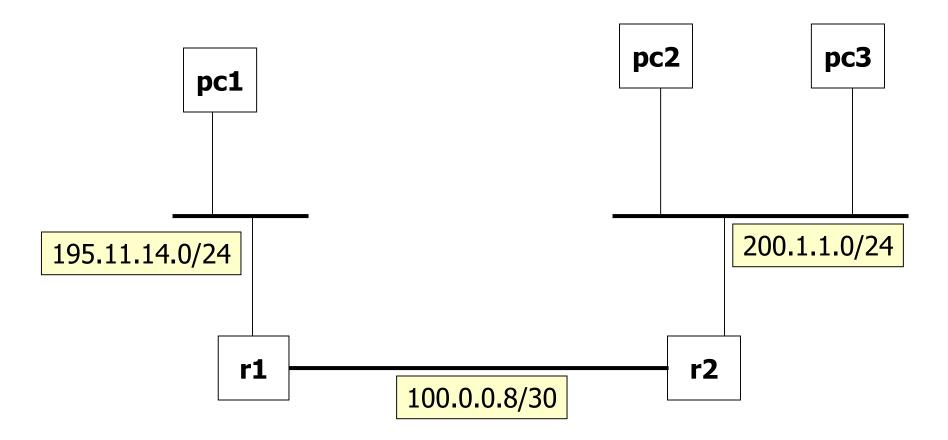
#### arp

Version	1.0
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Web	http://www.kathara.org/
Description	using the address resolution protocol for local and non local traffic – kathara version of the netkit lab on arp version 2.2

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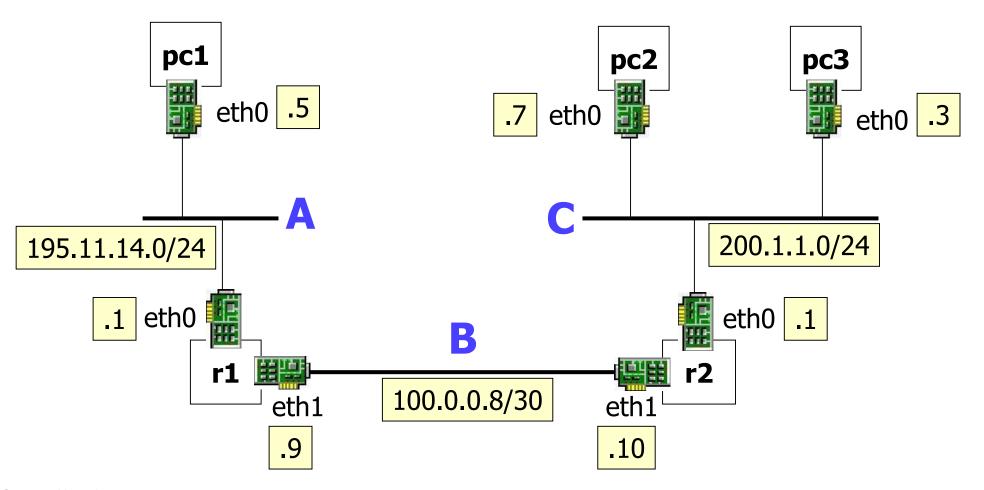
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### step 1 – network topology high level view



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# step 1 – network topology configuration details



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### step 2 – a quick look at the lab

# lab.conf r1[0]="A" r1[1]="B" r2[0]="C" r2[1]="B" pc1[0]="A" pc2[0]="C" pc3[0]="C"

```
pc1.startup

ifconfig eth0 195.11.14.5 up
route add default gw 195.11.14.1

pc2.startup

ifconfig eth0 200.1.1.7 up
route add default gw 200.1.1.1
```

```
pc3.startup
ifconfig eth0 200.1.1.3 up
route add default gw 200.1.1.1
```

### step 2 – a quick look at the lab

#### r1.startup

```
ifconfig eth0 195.11.14.1 up
ifconfig eth1 100.0.0.9 netmask 255.255.255.252 broadcast 100.0.0.11 up
route add -net 200.1.1.0 netmask 255.255.255.0 gw 100.0.0.10 dev eth1
```

#### r2.startup

```
ifconfig eth0 200.1.1.1 up
ifconfig eth1 100.0.0.10 netmask 255.255.255.252 broadcast 100.0.0.11 up
route add -net 195.11.14.0 netmask 255.255.255.0 gw 100.0.0.9 dev eth1
```

#### start the lab

### host machine user@localhost:~\$ cd kathara-lab\_arp user@localhost:~/kathara-lab\_arp\$ lstart ■

### step 3 – inspecting the arp cache

```
ARP (8)
                                                                         ARP (8)
                           Linux Programmer's Manual
NAME
       arp - manipulate the system ARP cache
SYNOPSIS
       arp [-vn] [-H type] [-i if] -a [hostname]
       arp [-v] [-i if] -d hostname [pub]
       arp [-v] [-H type] [-i if] -s hostname hw addr [temp]
       arp [-v] [-H type] [-i if] -s hostname hw addr [netmask nm] pub
       arp [-v] [-H type] [-i if] -Ds hostname ifa [netmask nm] pub
       arp [-vnD] [-H type] [-i if] -f [filename]
DESCRIPTION
       Arp manipulates the kernel's ARP cache in various ways. The primary
```

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dump of the ARP cache.

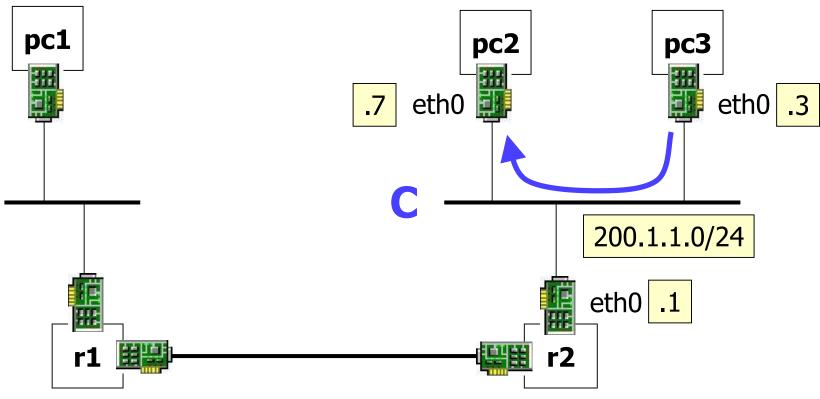
kathara – [ lab: arp ]

options are clearing an address mapping entry and manually setting up

debugging purposes, the arp program also allows a complete

# step 3 – inspecting the arp cache (local traffic)

traffic within the same network does not traverse routers



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# step 3 – inspecting the arp cache (local traffic)

the arp cache is sending packets to initially empty v pc3 ▲× 200.1.1.7 requires address resolution pc3:~# arp pc3:~# ping 200.1.1.7 = PING 200.1.1.7 (200.1.1.7) 56(84) bytes of data. 64 bytes from 200.1.1.7: icmp\_seq=1 ttl=64 time=1.39 ms 64 bytes from 200.1.1.7: icmp\_seq=2 ttl=64 time=0.542 ms --- 200.1.1.7 ping statistics ---2 packets transmitted, 2 received, 0% packet loss, time 1022ms rtt min/avg/max/mdev = 0.542/0.969/1.396/0.427 mspc3:~# arp -n HWtype HWaddress Address Iface Flags Mask ether FE:FD:C8:01:01:07 200.1.1.7 eth0 pc3:~#

address resolution results are stored in the arp cache

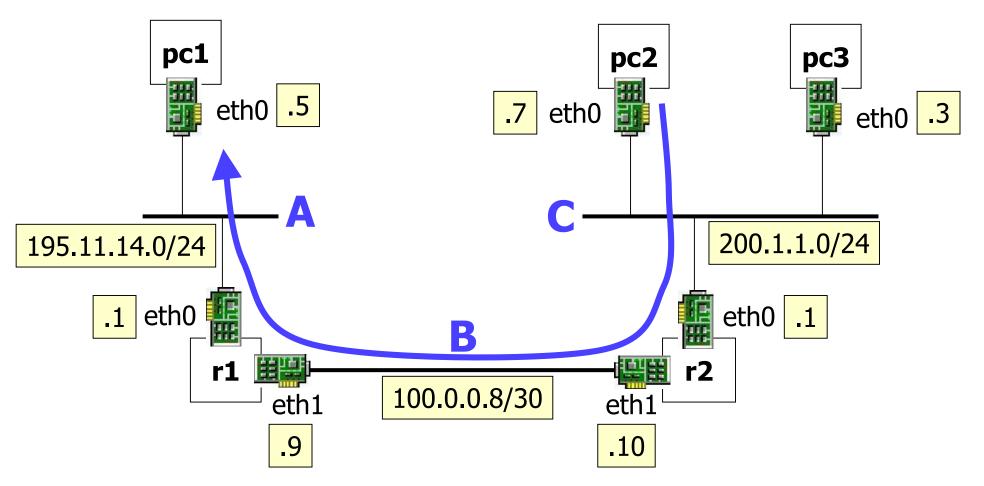
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# step 3 – inspecting the arp cache (local traffic)

- communications are usually bi-directional
- the receiver of the arp request learns the mac address of the other party, to avoid a new arp in opposite direction (standard behavior, see rfc 826)



# step 4 – inspecting the arp cache (non local traffic)



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kathara – [ lab: arp ]

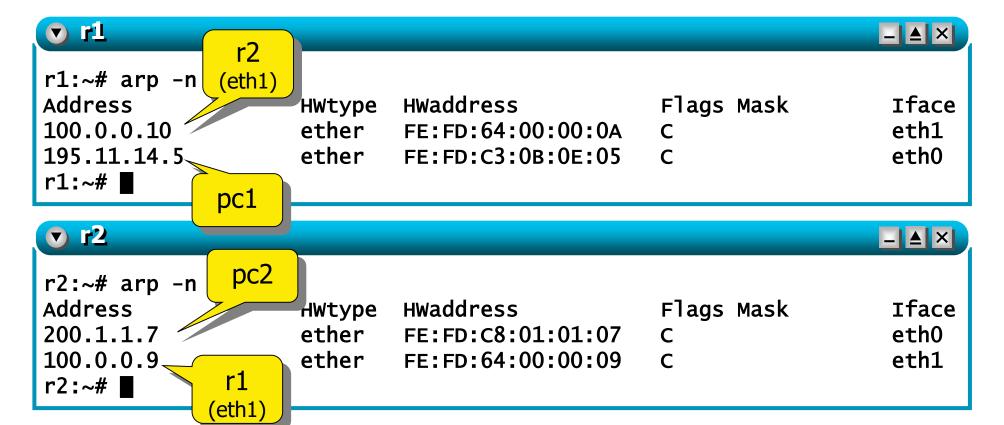
# step 4 – inspecting the arp cache (non local traffic)

- when ip traffic is addressed outside the local network, the sender needs the mac address of the router
- arp requests can get replies only within the local network

```
v pc2
                                                                  _ ≜ ×
pc2:~# ping 195.11.14.5
PING 195.11.14.5 (195.11.14.5) 56(84) bytes of data.
64 bytes from 195.11.14.5: icmp_seq=1 ttl=62 time=30.4 ms
64 bytes from 195.11.14.5: icmp_seq=2 ttl=62 time=1.02 ms
--- 195.11.14.5 ping statistics ---
2 packets transmitted, 2 received, 0% packet loss, time 1013ms
rtt min/avg/max/mdev = 1.024/15.731/30.438/14.707 ms
pc2:~# arp -n
Address
                                                Flags Mask
                            HWaddress
                                                                   Iface
                    HWtype
200.1.1.1
                    ether
                            FE:FD:C8:01:01:01
                                                                   eth0
                                                                   etnu
                    etner
                            FE:FD:C8:U1:U1:U3
pc2:~#
```

### step 4 – inspecting the arp cache (non local traffic)

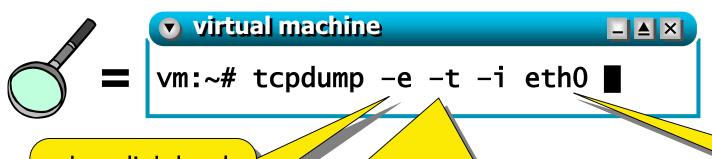
- what about routers?
- routers perform arp too (hence have arp caches) anytime they have to send ip packets on an ethernet lan



restart the lab in order to clear arp caches

```
v host machine
user@localhost:~/kathara-lab_arp$ lclean
user@localhost:~/kathara-lab_arp$ lstart ■
```

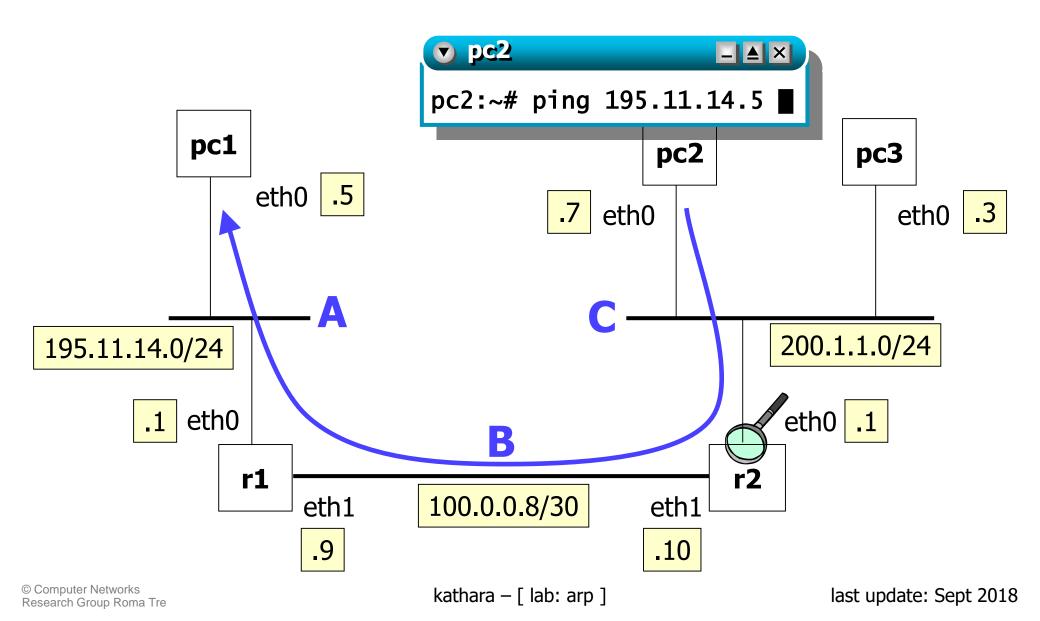
get ready to sniff



show link-level headers (mac addresses)

suppress timestamps (kathara is not a simulation system, hence timestamps are not meaningful) sniff on this interface

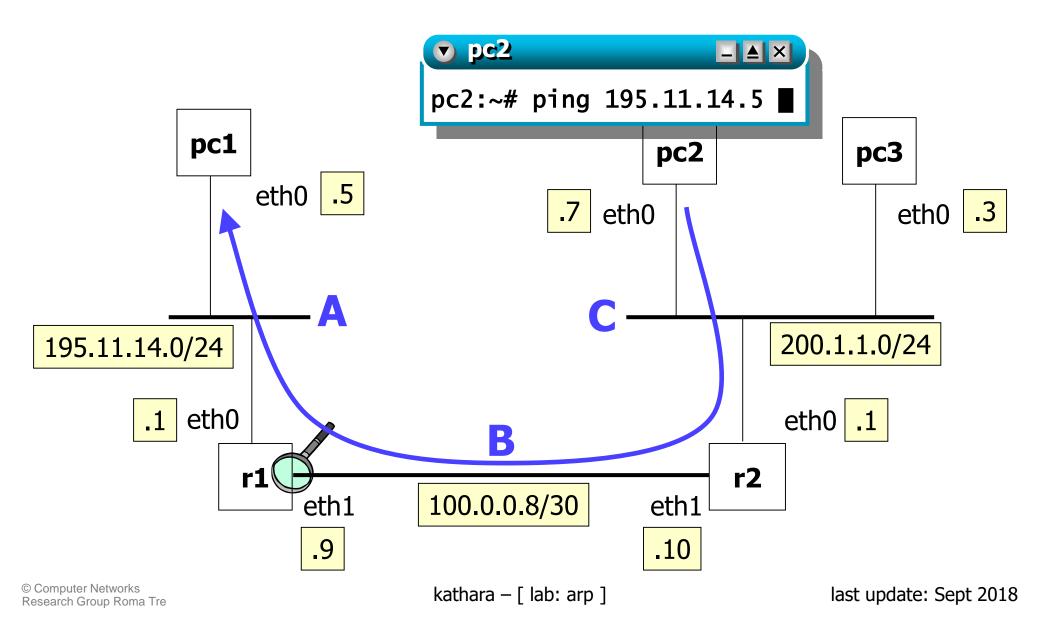
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on collision domain C

```
r2:~# tcpdump -e -t -i eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 96 bytes
fe:fd:c8:01:01:07 > Broadcast, ethertype ARP (0x0806), length 42: arp who-has
200.1.1.1 tell 200.1.1.7
fe:fd:c8:01:01:01 > fe:fd:c8:01:01:07, ethertype ARP (0x0806), length 42: arp reply
200.1.1.1 is-at fe:fd:c8:01:01:01
fe:fd:c8:01:01:07 > fe:fd:c8:01:01:01, ethertype IPv4 (0x0800), length 98: IP
200.1.1.7 > 195.11.14.5: icmp 64: echo request seq 1
fe:fd:c8:01:01:01 > fe:fd:c8:01:01:07, ethertype IPv4 (0x0800), length 98: IP
195.11.14.5 > 200.1.1.7: icmp 64: echo reply seq 1
```

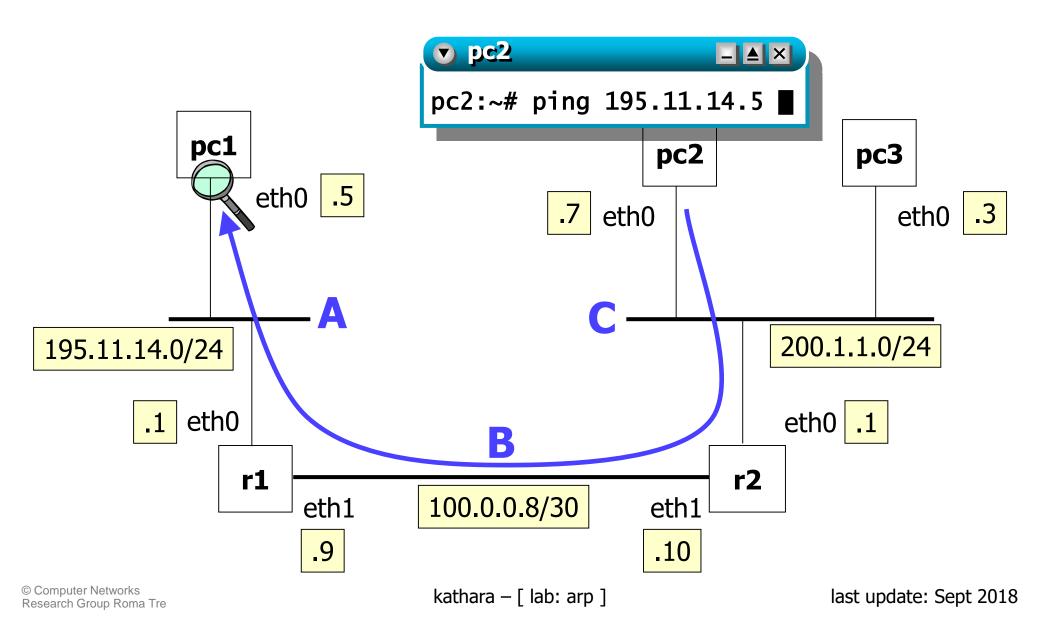
- 1. pc2 asks all the stations on collision domain C: "who has 200.1.1.1?" (200.1.1.1 is pc2's default gateway)
- 2. r2 replies  $\Rightarrow$  both pc2 and r2 update their arp cache
- 3. pc2 sends to r2 the ip packet (icmp echo request) for pc1
- 4. r2 sends to pc2 the corresponding echo reply (generated by pc1)



on collision domain B

```
r1:~# tcpdump -e -t -i eth1
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth1, link-type EN10MB (Ethernet), capture size 96 bytes
fe:fd:64:00:00:0a > Broadcast, ethertype ARP (0x0806), length 42: arp who-has
100.0.0.9 tell 100.0.0.10
fe:fd:64:00:00:09 > fe:fd:64:00:00:0a, ethertype ARP (0x0806), length 42: arp reply
100.0.0.9 is-at fe:fd:64:00:00:09
fe:fd:64:00:00:0a > fe:fd:64:00:00:09, ethertype IPv4 (0x0800), length 98: IP
200.1.1.7 > 195.11.14.5: icmp 64: echo request seq 1
fe:fd:64:00:00:09 > fe:fd:64:00:00:0a, ethertype IPv4 (0x0800), length 98: IP
195.11.14.5 > 200.1.1.7: icmp 64: echo reply seq 1
```

- 1. r2 asks all the stations on collision domain B: "who has 100.0.0.9?" (100.0.0.9 is the next hop obtained from the routing table)
- 2. r1 replies  $\Rightarrow$  both r1 and r2 update their arp cache
- 3. r2 sends to r1 the echo request generated by pc2 for pc1
- 4. r1 sends to r2 the echo reply generated by pc1 for pc2

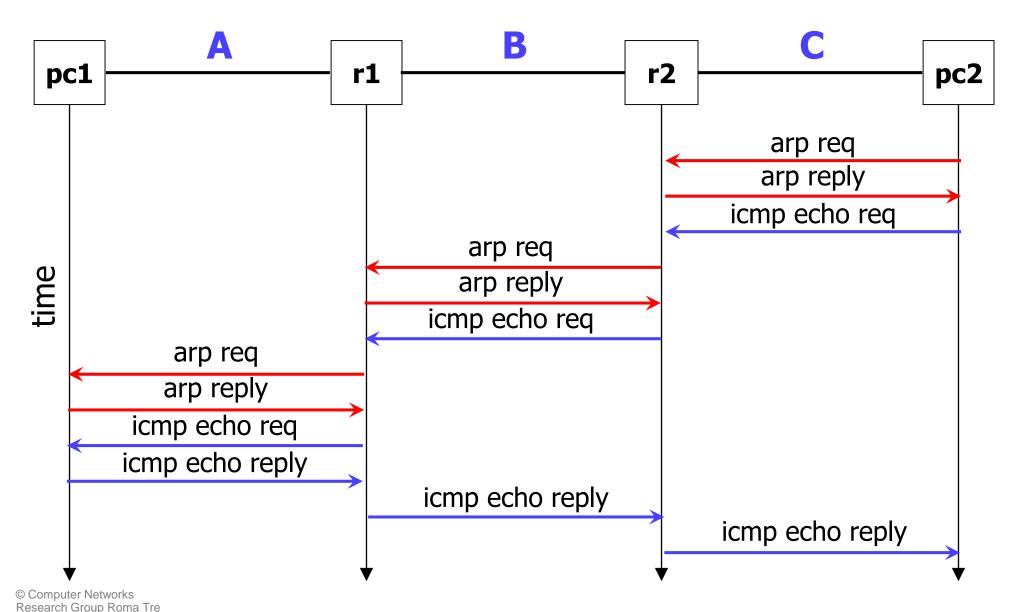


on collision domain A

```
pc1:~# tcpdump -e -t -i eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 96 bytes
fe:fd:c3:0b:0e:01 > Broadcast, ethertype ARP (0x0806), length 42: arp who-has
195.11.14.5 tell 195.11.14.1
fe:fd:c3:0b:0e:05 > fe:fd:c3:0b:0e:01, ethertype ARP (0x0806), length 42: arp reply
195.11.14.5 is-at fe:fd:c3:0b:0e:05
fe:fd:c3:0b:0e:01 > fe:fd:c3:0b:0e:05, ethertype IPv4 (0x0800), length 98: IP
200.1.1.7 > 195.11.14.5: icmp 64: echo request seq 1
fe:fd:c3:0b:0e:05 > fe:fd:c3:0b:0e:01, ethertype IPv4 (0x0800), length 98: IP
195.11.14.5 > 200.1.1.7: icmp 64: echo reply seq 1 ■
```

- 1. r1 asks all the stations on collision domain A: "who has 195.11.14.5?" (195.11.14.5 is the destination address of the icmp request obtained from the ip header)
- 2. pc1 replies  $\Rightarrow$  both pc1 and r1 update their arp cache
- 3. r1 sends the ip packet (echo request) to pc1
- 4. pc1 generates the corresponding echo reply for pc2 and sends it to r1

#### step 6 – understanding the whole picture



#### step 7 – arp implementation details

```
r2:~# tcpdump -e -t -i eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 96 bytes
.....

fe:fd:c8:01:01:01 > fe:fd:c8:01:01:07, ethertype ARP (0x0806), length 42: arp who-has
200.1.1.7 tell 200.1.1.1
fe:fd:c8:01:01:07 > fe:fd:c8:01:01:01, ethertype ARP (0x0806), length 42: arp reply
200.1.1.7 is-at fe:fd:c8:01:01:07
```

- arp requests are usually in broadcast
- it may also happen that a station (router/pc) sends a unicast arp request to check if an entry of the arp cache is still valid (suggested by the standard, see rfc 826)
- unicast arp requests may be performed periodically on each entry of the arp cache, depending on the implementation

#### proposed exercises

what packets can we observe over all the three collision domains as the ping from pc2 to pc1 continues (ignore any implementation dependent arp behavior)?

### proposed exercises

- check the different error messages obtained by trying to ping an unreachable destination in the case of
  - local destination
  - non local destination
- which packets are exchanged in the local collision domain in the two cases?