COD310 Notes

1 Stuff to Read

- 1. 3D memory systems 3D DRAM
- 2. Memory leakage power
- 3. SPEC CPU2006 benchmarks
- 4. Hybrid Memory Cube (HMC), High Bandwidth Memory (HBM) and Wide IO (WIO)
- 5. CACTI-3DD

2 Doubts

1. What mechanism controls the memory management in memory devices

3 Leakage-Aware Dynamic Thermal Management of 3D Memories

3.1 Overview (Abstract)

- 1. Controlling leakage by monitoring temperature
- 2. Turn off specific memory channels to control temperature (before turning off, migrate data to 2D memory) FastCool
- 3. Energy-Efficient FastCool (EEFC) decides which channels to be closed

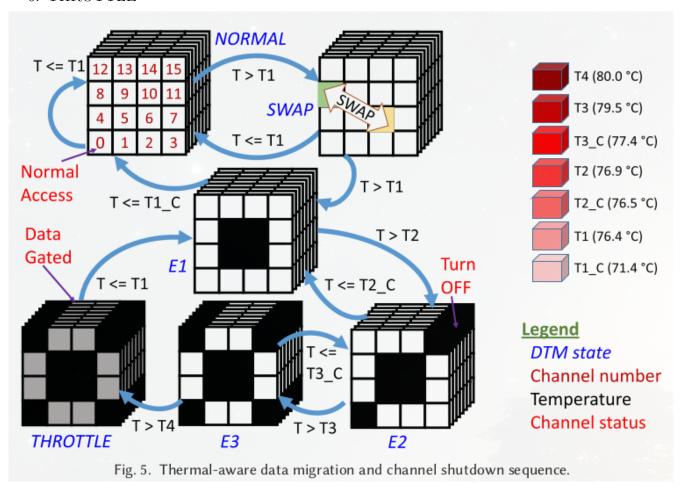
3.2 Introduction

- 1. 3D memory is stacked 2D DRAM, thus has higher power density
- 2. Power consumption involves dynamic (48%) and static/leakage power (52%)
- 3. Static power increases exponentially with temperature, thus a positive feedback between temperature and leakage

3.3 Proposed DTM Strategies

3.3.1 TAM (Thermal-Aware Migration) States

- 1. NORMAL
- 2. SWAP
- 3. E1 (thermal Emergency 1)
- 4. E2
- 5. E3
- 6. THROTTLE



3.3.2 Memory Delay Models

2D memory request time = Data Migration Delay (DMD) + Data Access Delay (DAD)

Symbol	Description	Typical value
D	Time duration for an epoch	1 ms
С	Cache Line Size	64 Bytes
3D Memory Parameters		
$b_{ m 3D}$	Per channel 3D memory bandwidth	8 GBps
L_{3D}	3D memory access latency	29 ns
s_{3D}	Size of data migrated, $3D \rightarrow 2D$	Runtime ¹
A_{3D}	Total accesses made to s_{3D} data (in last epoch)	Runtime ¹
n_{3D}	Number of 3D channels (migrating 3D \rightarrow 2D)	Runtime ¹
2D DRAM Parameters		
$b_{ m 2D}$	Per channel bandwidth	12.8 GBps
L_{2D}	2D memory access Latency	45 ns
s_{2D}	Size of data migrated, $2D \rightarrow 3D$	Runtime ¹
A_{2D}	Accesses made to 2D memory (in last epoch)	Runtime ¹
$N_{ m 2D}$	Number of 2D memory channels	4
R_{2D}	Number of ranks per channel	2
BA_{2D}	Number of banks per rank	8
	Page Policy	Closed

¹Runtime - Values are determined by DTM policy with the help of hardware counters.

3.3.2.1 DMD

$$DMD = (s_{3D} + s_{2D}) \times \max\left(\frac{1}{B_{3D}}, \frac{1}{B_{2D}}\right)$$

$$B_{3D} = b_{3D} \times n_{3D} \times (1 - 3D \text{ Memory Refresh Overhead})$$

$$B_{2D} = b_{2D} \times N_{2D} \times (1 - 2D \text{ Memory Refresh Overhead})$$

$$Refresh Overhead = \frac{\text{Time required to refresh a row} \times \text{Number of rows}}{\text{Refresh interval}}$$

3.3.2.2 DAD

$$DAD = DAD_B + DAD_L$$

$$A = A_{2D} + A_{3D}, \text{ total 2D memory accesses}$$

$$DAD_B = \frac{A \times C}{B_{2D}}$$

$$DAD_L = QD + LD$$
, queuing delay + latency delay

3.3.2.2.1 Latency Delay

$$LD = \frac{A \times L_{2D}}{BA_{2D} \times R_{2D} \times N_{2D}}$$

3.3.2.2.2 Queuing Delay

Uses Queuing Theory to model the waiting time. M/M/1 model is used, having a single queue for each server, and arrival and service rates are Poisson and exponential respectively.

$$QD = \frac{A \times C}{T \times B_{2D}} \times \frac{N_{2D}}{B_{2D} - (A \times C/T)}$$

$$\left(\lambda = \frac{A \times C}{T \times N_{2D}}, \mu = \frac{B_{2D}}{N_{2D}}, \text{expected time} = \frac{\lambda}{\mu \times (\mu - \lambda)}\right)$$

3.3.3 FastCool

1. Transition to E1 happens only if total access count of channels $\{5, 6, 9, 10\}$ exceeds A_{MIN} $(A_{MIN} = 160, 000)$

$$A_{3D} > A_{MIN}$$

2. Queuing stability needs to be ensured before migrating $(\lambda < \mu)$

$$A < \frac{B_{2D} \times T}{C}$$

3. Ensure transfer to 2D memory happens only if 2D delay is below a certain threshold to prevent slow down of operations

$$Delay < D_{MAX} (= 8.415ms)$$

3.3.4 FC Policy Improvements

later

3.3.5 EEFC

later

3.3.6 Leakage Current Estimation

later

3.3.7 DTM Policy Implementation

later

4 PredictNcool: Leakage Aware Thermal Management for 3D Memories Using a Lightweight Temperature Predictor

4.1 Overview

- 1. Instead of reacting to temperature changes, this model attempts to utilise predicted temperature changes to reduce application runtime and memory energy
- 2. Symmetries in floor-plan and other design insights are used to reduct the predictor model parameters

5 CoreMemDTM: Integrated Processor Core and 3D Memory Dynamic Thermal Management for Improved Performance

5.1 Overview

Independent thermal management of core and memory leads to inefficient management since both cores and memories slow down