COL 351: Analysis and Design of Algorithms

Tutorial Sheet - 4

Question 1 Complete the pseudo-code below to obtain an O(|V| + |E|) time algorithm for computing cut-vertices of an n-vertex undirected (possibly unconnected) graph G = (V, E).

```
1 VISITED[v] \leftarrow True;
2 HIGH-POINT[v] \leftarrow LEVEL[v];
3 foreach w \in N(v) do
       if (VISITED[w] = False) then
           Set PARENT[w] \leftarrow v and LEVEL[w] = 1 + LEVEL[v];
6
7
           if
8
           end
10
       else if (PARENT[w] \neq v) then
11
           \texttt{HIGH-POINT}[v] \leftarrow \min\{\texttt{LEVEL}[w], \ \texttt{HIGH-POINT}[v]\};
12
       end
14 end
```

Procedure DFS(v)

Procedure Compute-Cut-vertices(G)

Question 2 A topological order of a DAG is a linear ordering of its vertices such that for every directed edge (x, y), x comes before y in the ordering. Let G = (V, E) be a DAG on n vertices, and $L = (v_1, \ldots, v_n)$ be a list of vertices of G satisfying

$$FINISH-TIME(v_1) > FINISH-TIME(v_2) > \cdots > FINISH-TIME(v_n)$$

with respect to some DFS traversal of G. Then prove that L is a topological ordering of G.

Question 3 Let G = (V, E) be a DAG and s be a source vertex such that all vertices of G are reachable from s. Let y and z be any two vertices in G satisfying that there is a path from y to z. Prove that the following two properties hold true for each DFS traversal carried out from s in G.

- **1** : FINISH-TIME(z) < FINISH-TIME(y).
- 2 : If START-TIME(y) < START-TIME(z), then z must be a descendant of y in the DFS tree.

Show that if G was a general directed graph, and vertices y,z satisfies that there is a path from y to z in G. Then, START-TIME(y) < START-TIME(z) does not imply that z is a descendant of y in the DFS tree.

Question 4 Prove that any n vertex undirected graph contains at most n-1 bridge-edges and at most n-2 cut-vertices.

Question 5 Let $A = (a_1, \ldots, a_n)$ and $B = (b_1, \ldots, b_m)$ be two sequences, where a_i 's and b_j 's are positive integers. Devise an O(mn) time algorithm to compute a **longest common increasing subsequence** of A and B.

Question 6 Let $A = (a_1 \dots a_n)$ and $B = (b_1 \dots b_m)$ be two strings, where a_i 's and b_j 's are English alphabets. Devise an O(mn) time algorithm to compute a **longest common substring** of A and B.