# COL759: CRYPTOGRAPHY AND COMPUTER SECURITY

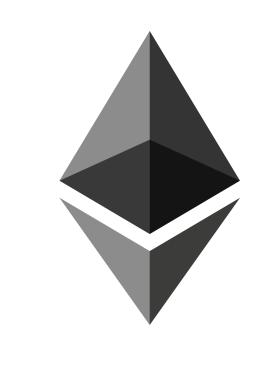
2022-23 (SEMESTER 1)

LECTURE 1: INTRO

#### COL759 IS NOT ABOUT ...

Computer Security (SIL 765)





Symmetric Key Encryption

Message Authentication Codes

Collision resistant hash functions

**Authenticated Encryption** 

**Public Key Encryption** 

Signatures

Zero Knowledge Proofs

COL759



Formal security definitions, and security proofs

#### **COURSE OBJECTIVES**

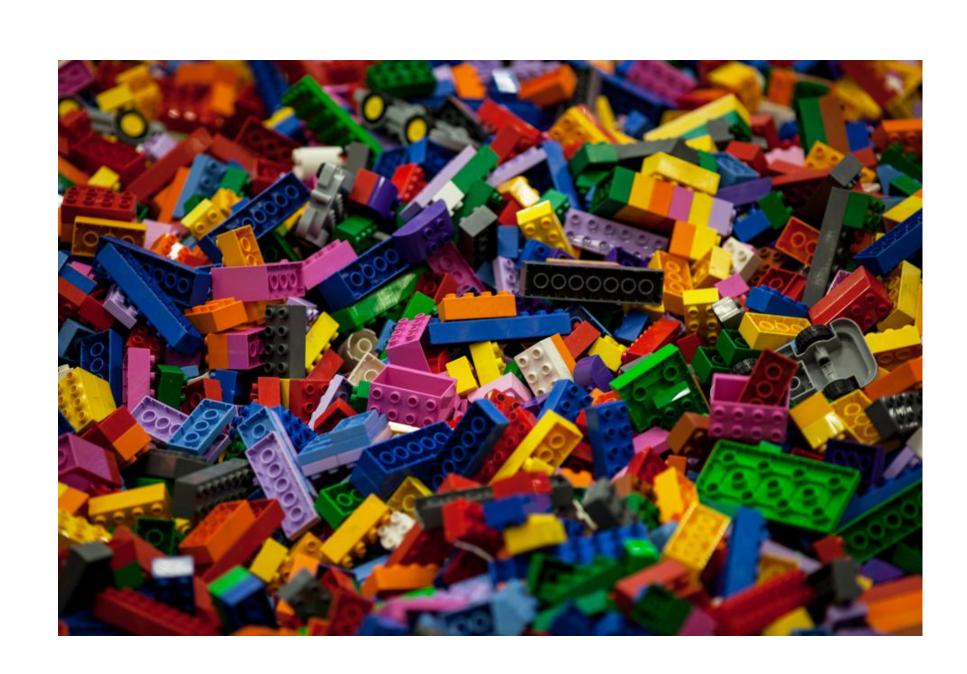
- Develop 'crypto mindset'
- Modelling threats via security definitions
- Learn the 'atomic' building blocks of crypto, and how they're used for building more complex primitives
- Understand how to prove security

#### 4-STEP RECIPE FOR A CRYPTO PRIMITIVE

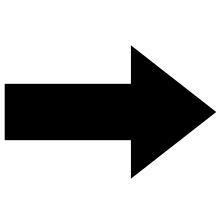
- 1. Define security for the primitive
- 2. Figure out the required building blocks
- 3. Propose construction
- 4. Prove construction satisfies security definition

ANALOGY: LEGO BLOCKS

#### THE LEGO ANALOGY









Our desired crypto primitive

#### THE LEGO ANALOGY

**LEGO** 

**CRYPTO** 

Lego blocks are immutable.

Can't manufacture your own Lego blocks

DO NOT modify the building blocks

DO NOT implement the building blocks (use existing libraries)

# WHY FORMAL DEFINITIONS AND PROOFS OF SECURITY?

#### SYMMETRIC KEY ENCRYPTION

#### Key Space $\mathcal{K}$

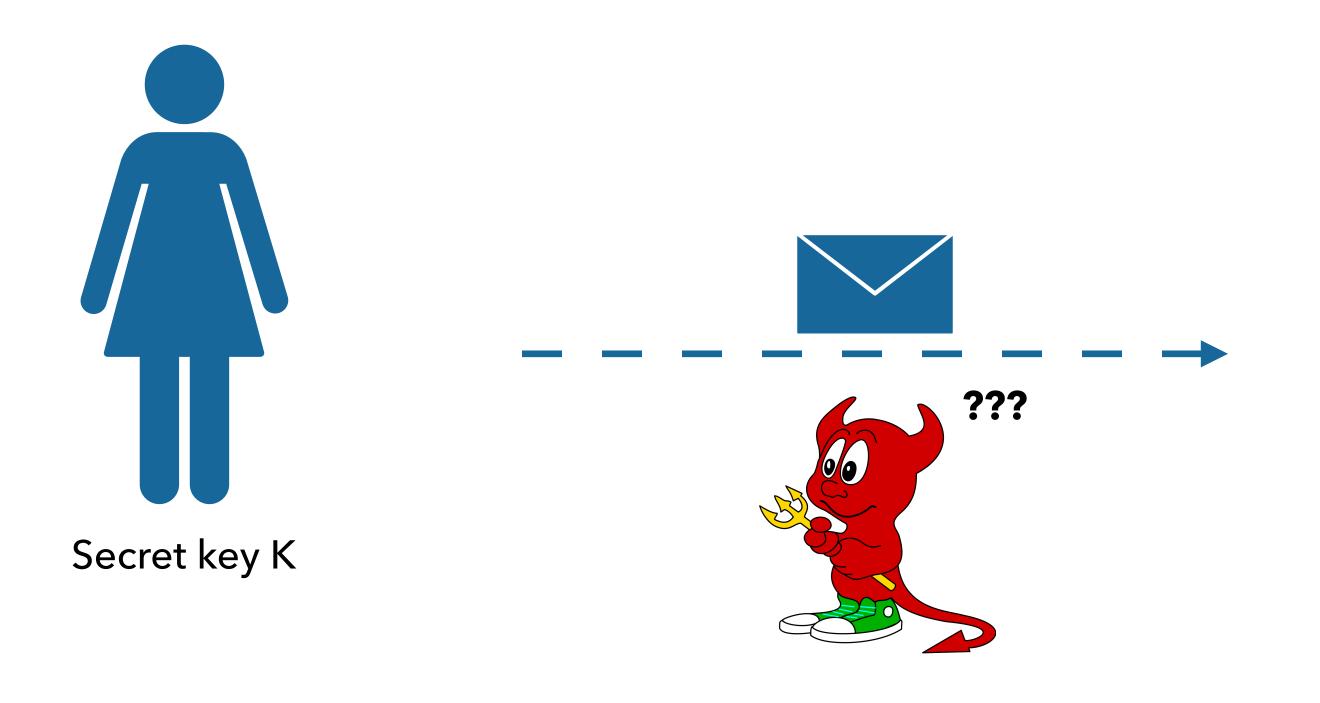
Encrypt(mesage, secret key) → ciphertext

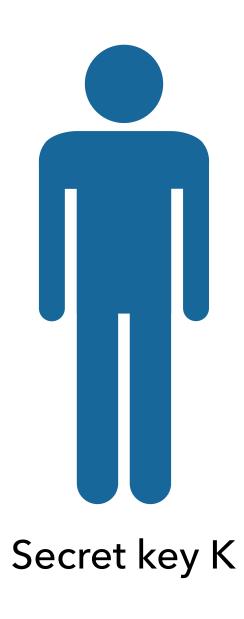
Decrypt(ciphertext, secret key) → message/ ⊥

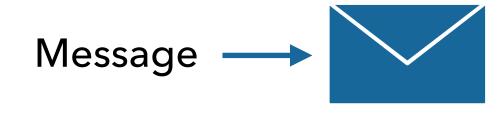
#### CORRECTNESS

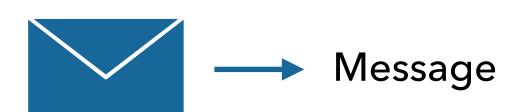
Decrypt(Encrypt(m, k), k) = m

#### SYMMETRIC KEY ENCRYPTION









### 1. Caesar's cipher

- Secret key:
  Integers
- Encrypt:
   shift each character forward by s positions

#### **BROKEN CIPHER!!**

Attack via brute force

### 2. Substitution cipher

- Secret key:
   Permutation over the alphabets
- Encrypt:
   substitute each character according to
   permutation in secret key

#### **BROKEN CIPHER!!**

Attack via frequency analysis

# 3. Vigenère's cipher

- Secret key: n different permutations
- Encrypt:

   For character at position i, use substitution
   with (i mod n)<sup>th</sup> permutation

le chiffrage indéchiffrable (the indecipherable cipherable cipherable cipher)

# 3. Vigenère's cipher

- Secret key:
   n different permutations
- Encrypt:

   For character at position i, use substitution
   with (i mod n)<sup>th</sup> permutation

Efficient algo. for breaking substitution cipher



Efficient algo. for breaking Vigenère cipher



Is 'Double-Encrypt-Vigenere' secure?
What if n = length of message to be encrypted?

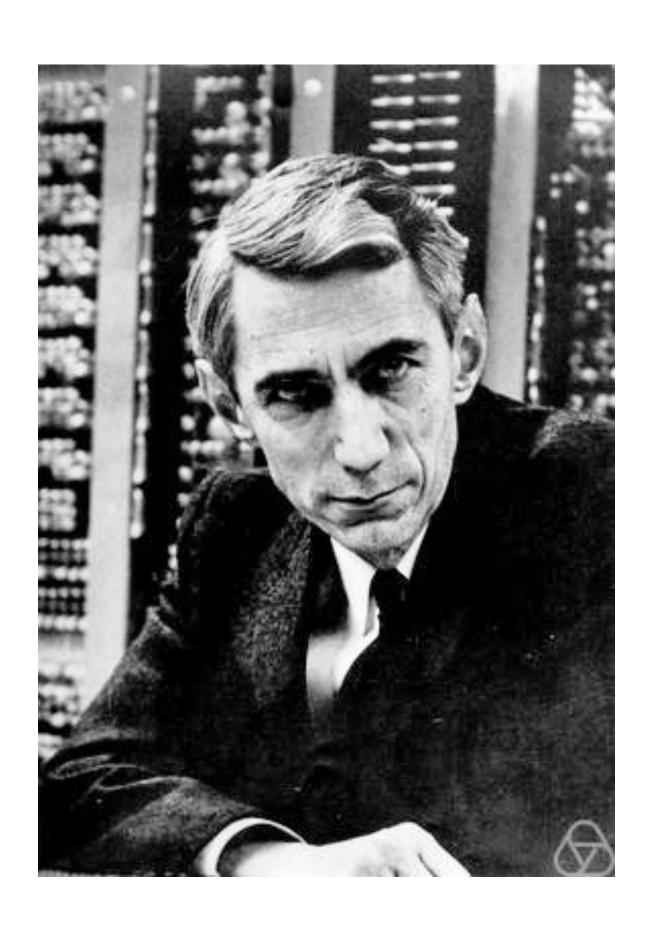
# 4. Rotor machines (Enigma)



#### **BROKEN DURING WW-II**



#### 5. Shannon's One-Time Pad



"Communication Theory of Secrecy Systems" (1945)

Perfectly secure encryption scheme

Key cannot be reused

Must be as large as the message

# 6. Data Encryption Standard (DES) (1970s)

Key Space :  $\{0,1\}^{56}$ 

DES:  $\{0,1\}^{64} \times \{0,1\}^{56} \rightarrow \{0,1\}^{64}$ 

 $\mathsf{DES}^{-1}: \{0,1\}^{64} \times \{0,1\}^{56} \to \{0,1\}^{64}$ 

DES can 'encrypt' 64 bit messages. How to encrypt longer messages?

#### **ATTACKS**

Exhaustive search for key: ~ 2<sup>56</sup> steps Best known attack: ~ 2<sup>44</sup> steps

Feasible using modern supercomputers

Can we extend key space appropriately?

### 6.1. Double DES (2DES)

Key Space :  $\{0,1\}^{112}$ 

$$2DES(m, (k_1, k_2)) = DES(DES(m, k_1), k_2)$$
$$2DES^{-1}(ct, (k_1, k_2)) = DES^{-1}(DES^{-1}(ct, k_2), k_1)$$

#### **ATTACKS**

Previous attacks infeasible.

But a different (very simple) attack breaks 2DES

Triple DES (3DES)?

# 7. Advanced Encryption Standard (AES) (1990s)

Key Space :  $\{0,1\}^{128}$ 

AES:  $\{0,1\}^{128} \times \{0,1\}^{128} \rightarrow \{0,1\}^{128}$ 

 $AES^{-1}: \{0,1\}^{128} \times \{0,1\}^{128} \rightarrow \{0,1\}^{128}$ 

AES can 'encrypt' 128 bit messages. How to encrypt longer messages?

#### **ATTACKS**

Exhaustive search for key:  $\sim 2^{128}$  steps Best known attack:  $\sim 2^{126}$  steps

Most widely used crypto algorithm

#### HARNESSING COMPUTATIONAL HARDNESS FOR CRYPTOGRAPHY

No efficient algo. for breaking AES



Secure encryption scheme

#### CENTRAL THEME IN MODERN CRYPTO

Building blocks: hard computational problems

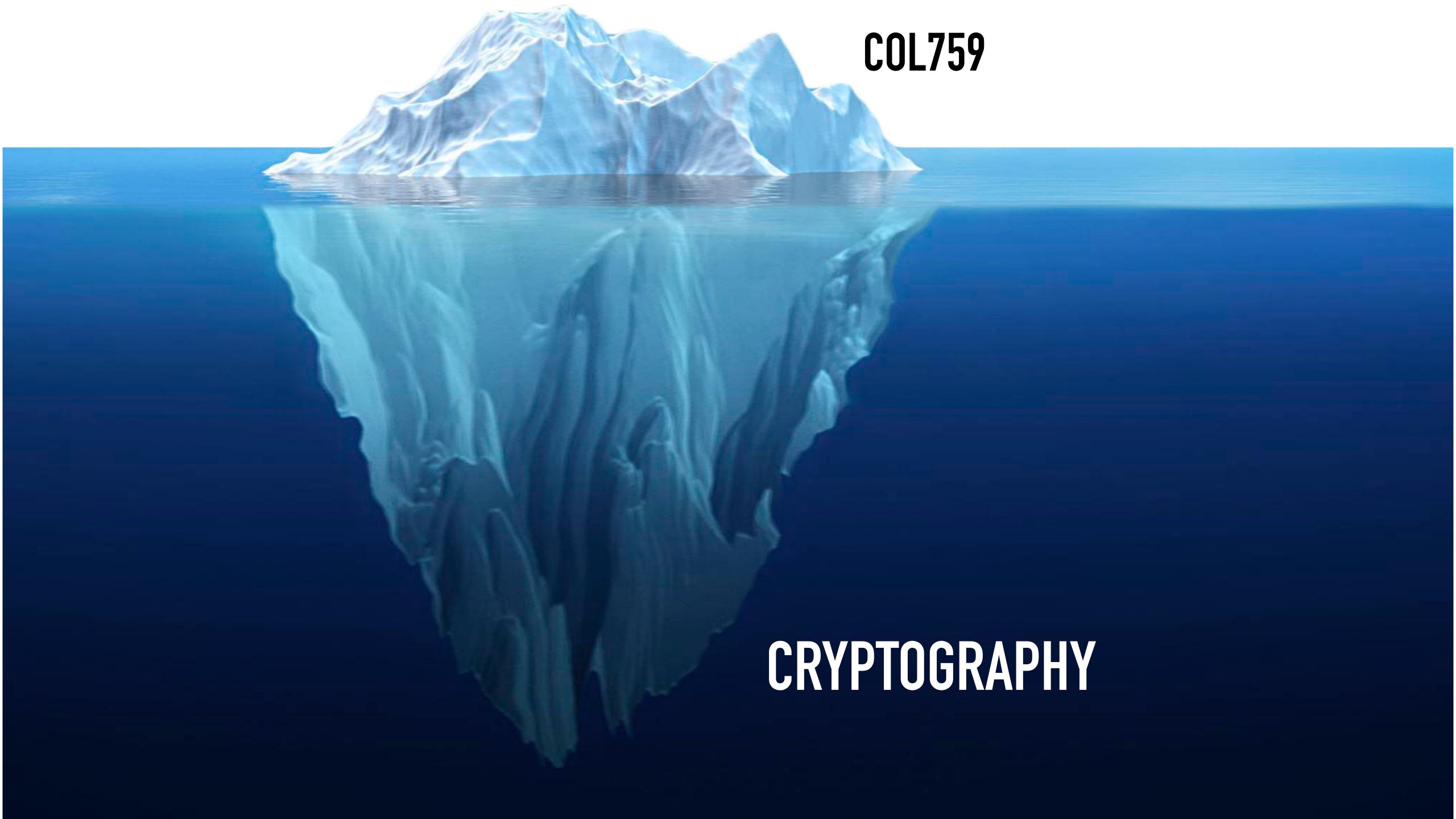
Proof of security: No efficient algo for problem



Cryptosystem is secure

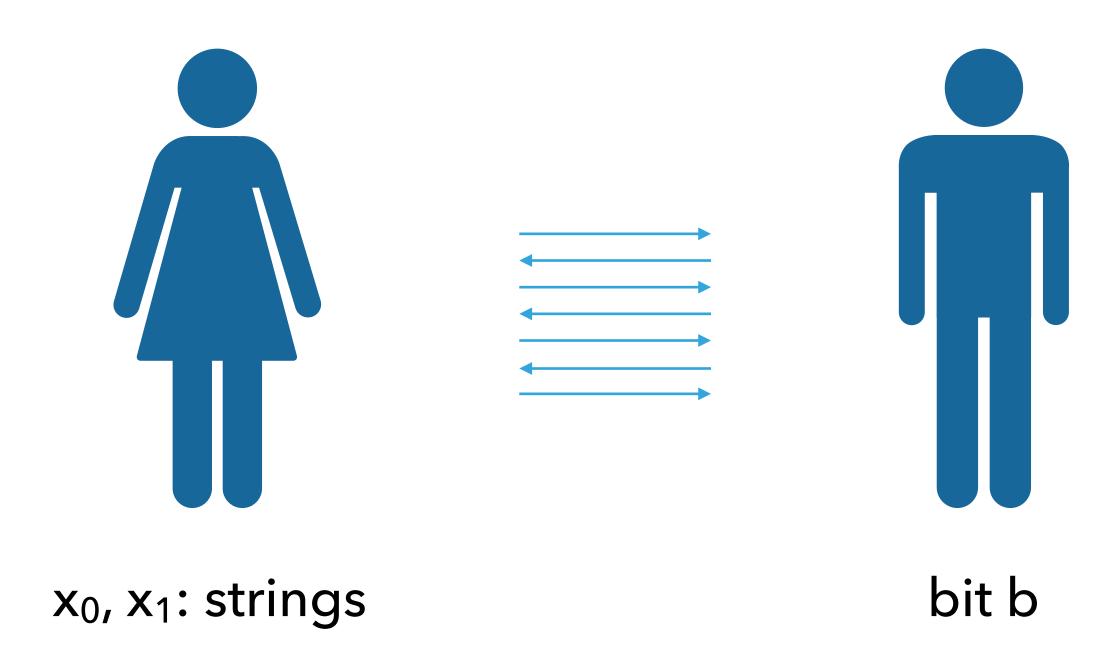
#### SOURCES OF HARD COMPUTATIONAL PROBLEMS

- Cryptographic standards: AES, SHA etc
- Number Theory
- Geometry
- Combinatorics ...



#### CRYPTO IS MAGIC!

# Magic #1



#### **DESIDERATA**

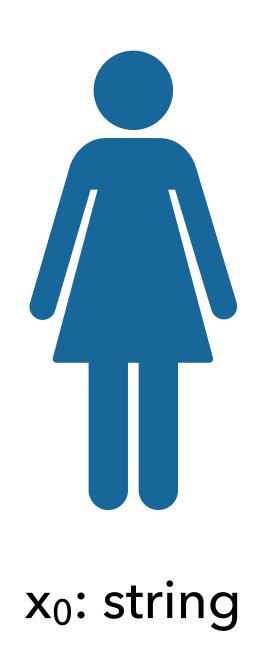
Bob should learn x<sub>b</sub>

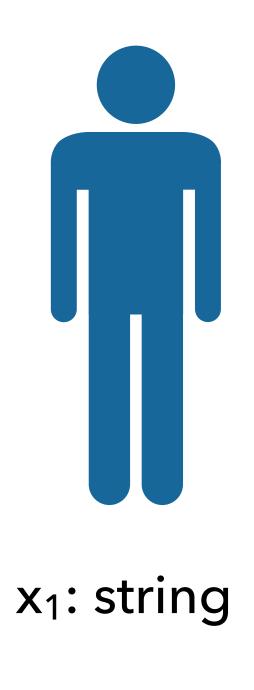
- Alice should not learn b Bob should not learn  $x_{1-b}$ 

'Oblivious Transfer'

#### CRYPTO IS MAGIC!

# Magic #2





#### **DESIDERATA**

Alice and Bob should learn f(x<sub>0</sub>, x<sub>1</sub>)

Alice should not learn x<sub>1</sub>
 Bob should not learn x<sub>0</sub>

'Multiparty Computation'

#### CRYPTO IS MAGIC!

# Magic #3

			7					
		2				3		
4				5	9		1	
			4				7	
	3			7	6	1		
6			8					
9				6	1		5	
					8			
	4							9

#### **DESIDERATA**

- Convince you that puzzle is solvable
- Without revealing any hint about solution

'Zero Knowledge Proofs'

#### **COURSE INFO**

Webpage: <a href="https://www.cse.iitd.ac.in/~koppula/courses/COL202\_2102.html">https://www.cse.iitd.ac.in/~koppula/courses/COL202\_2102.html</a>

Lecture notes on OneNote (\_Content Library -> Lectures -> Lecture #)

#### **COURSE POLICY**

#### NON NEGOTIABLES

- Theoretical course: assignments and exams will involve writing formal security proofs
- Assignments: groups of size 1 or 2. Must be typed in Latex
- Strict plagiarism policy
- Please avoid using laptops/phone during the lecture

#### **COURSE POLICY**

#### **NEGOTIABLES**

#### - Grading: relative

- Minor: 25%

Open-notes (handwritten notes only)
Longer duration for exams?

- Major: 30%

- Assignments (4): 25%

- Quizzes (best 5 of 6): 20%

#### - Late Policy for assignments:

- Three late days (cumulative)

- Please participate actively in class. Ask lots of questions
- Assignments: start early

THANK YOU!

(see you tomorrow)