

OSY.SSI[2019][6]

In the last episode...

Internet vs the Classical Theory

In the news...

- ▶ US Attorney General William Barr will present an open letter to Facebook and its CEO, Mark Zuckerberg, cosigned by British and Australian officials, asking the company not to implement end-to-end encryption protections [src]
- ▶ HTTPS-snooping malware 'patches' random number generator, decodes traffic and installs rogue certificates [src]
- ▶ CVE-2019-2215: Local privilege escalation vulnerability affecting Samsung, Xiaomi, Huawei, Google Pixel amongst others. Exploited in the wild. [src]
- ▶ CVE-2019-11932: WhatsApp RCE triggered by sending a GIF [src]
- ▶ Former Yahoo engineer pleads guilty to hacking thousands of accounts in hunt for nudes [src]
- ▶ ANU Vice Chancellor releases detailed report of cyberattack [src]
- ▶ Egyptian Hackers Blamed For Cyberattack On Android Devices—But It's Complicated [src]
- ▶ DHS and FDA warn about much broader impact of Urgent/11 vulnerabilities [src]
- ▶ A recent attack aimed at a U.S.-based oil, gas and chemical supplier leverages the company's use of the enterprise-class Asterisk open-source PBX software, used for VoIP services. [src]
- ▶ EA Games Leaks Personal Data of 1600 FIFA 20 Competitors [src]

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Do programmers know this and its consequences when they design and implement programs?

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- ▶ USS Yorktown (international waters, 1997): x/0.
 - ▶ Computer system down. Propulsion system down. In the middle of nowhere.
- ▶ And various major vulnerabilities based on number play
 - ▶ WebKit/Chrome (CVE-2009-2195)
 - ▶ Ruby (CVE-2013-4164)
 - ▶ Adobe Flash (CVE-2015-3077, CVE-2014-0502, etc.)
 - ▶ Mozilla Thunderbird (CVE-2017-5407): Pixel and history stealing via floating-point timing side-channel!

The overarching principle

It is obvious, but precisely for this reason it requires insisting

Mathematical model \neq Program \neq What the computer should do \neq What it does

While it is relatively rare, it *does* happen that the CPU returns the *wrong* answer to an arithmetic question (as was famously the case for the first Intel Pentium with FDIV and F00F).

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Many programmers forget that. Please, don't.

Example: Intel Pentium F00F, Reloaded

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Can wipe the UEFI, re-infect the OS after a clean install, bypass “Secure Boot” sequences, etc.

Part II

Physics, Caches, Injection, Beyond the Classical Theory

Cache-cache

A **cache** is some fast-and-expensive memory, that is used to avoid fetching slow-and-cheap memory. It's a form of 'look-up table'.

HDD	DRAM	L3	L2	L1	registers
1 ms	100 ns	80 cy	10 cy	4 cy	1 cy
1 TB	16 GB	8 MB	256 kB	32 kB	8 B

- ▶ If we've been clever and pre-loaded the right stuff, data is ready and we can proceed: it's a **cache hit**
- ▶ Otherwise, we have to ask the next cache in line, it's a **cache miss**

The difference is small, but measurable: at the very least we found a **timing channel**.

Cache-cache

More specifically, we can have

- ▶ A **cold-start miss**, if the requested data is missing from the cache,
- ▶ A **capacity miss**, if the requested data is larger than the cache,
- ▶ A **conflict miss**, if the requested data could fit, but won't due to e.g. alignment constraints

Since most architectures have only few registers, the vast majority of programs use caches at some point.

Reset, initialisation, microarch

- ▶ **Reset attack:** Cache memory is volatile, and often easy to erase (cold boot). It is also possible to fill it with crap.
 - ▶ **Initialisation attack:** if we can run a program on the same system, it will fill the caches in a controlled way.
 - ▶ **Microarchitectural attack:** the status of the cache is known and contains everything needed for the targeted program.
-
- ▶ Flush + Reload / Evict + Reload
 - ▶ Flush + Flush (why ?)
 - ▶ Prime + Probe

What do I need to run a cache attack?

Does it work in practice?

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It depends.

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It depends.

Yes. Maybe. Definitely!

First demonstrated in 2005

- ▶ OpenSSL, full secret AES key extracted in 13 ms
- ▶ Android dm-crypt, full secret AES key in 65 ms

But wait, there's more!

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Building on this observation since 2017 about two dozen attacks have been described targeting modern CPUs:

- ▶ Spectre (1, 1.2, 2, NG, NG1.1, NG3a, NG4, RSB)
- ▶ Meltdown
- ▶ SWAPGS
- ▶ MDS: Fallout, RIDL, ZombieLoad (MSBDS, MLPDS, MFBDS, MDSUM)
- ▶ Foreshadow/L1TF

How do you detect/mitigate such an attack?

Virtual Cache-cache

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On Amazon EC2:

- ▶ Keystroke recovery (SWT01)
- ▶ Stealing secret keys (ZJRR12, IGIES15, IIIES14)

No amount of sandbox or virtual machine would change the hardware.

Bonus: Clémentine Maurice's [Intro to cache attacks]

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Ok so we just discussed caches and how they give us time. But what is time anyway?

Computers **are not theoretical devices**.

They are physical devices, and the usual (software-oriented) engineer makes the confusion between model and reality.

All models are lies, some are useful

Motivating, textbook example

Exercise: write a program that takes as input k and outputs g^k . The program must not reveal k .

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Exercise: write a program that takes as input k and outputs g^k . The program must not reveal k .

Ok so now you already realise you can't just compute $g \times \cdots \times g$. (tell me why)

But you also realise that the operation must somehow depend on k .

The general idea

- ▶ The system depends on some unknown k
- ▶ Predict what the system would do if $k = x$
- ▶ Compute correlation between prediction and reality

If it doesn't work: wrong hypothesis, try another x . Otherwise...

Example (demo?): password finding with timing information.

A better algorithm for g^x

The square-and-multiply algorithm (x_bin is the binary representation of x)

```
def power(g,x_bin):
    result = 1
    while i = x_bin.pop() and i != None:
        if i == 1:
            result = result * g
        g = g*g
    return result
```

Side note: this operation is essential to many cryptographic algorithms.

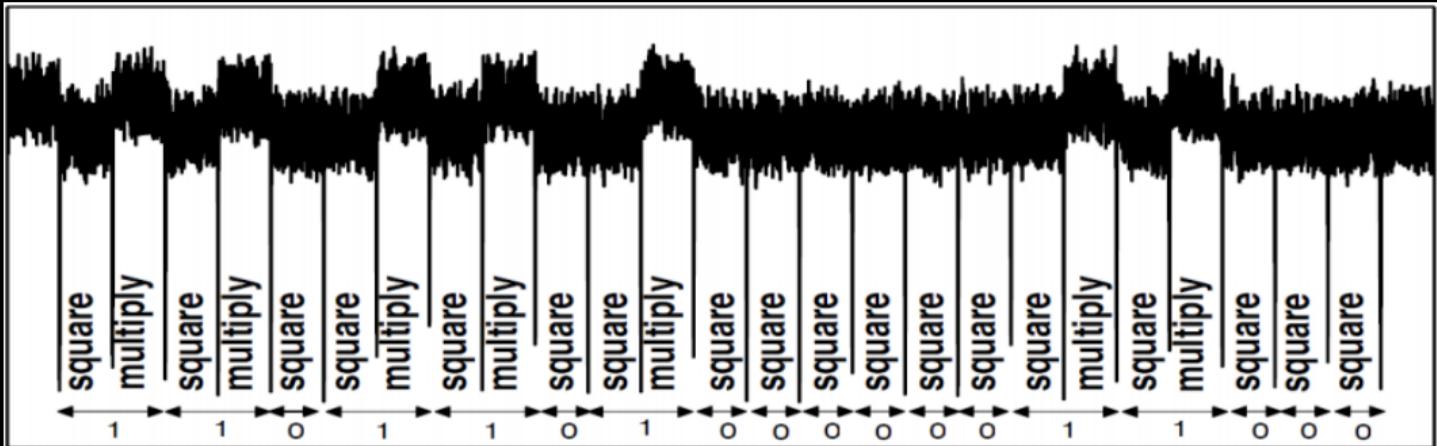
Look ma'! No timing channel!

At first glance this is a better algorithm than the naive one:

- ▶ Much faster: complexity $O(\log x)$ vs $O(x)$
- ▶ No timing channel: if x has a fixed size we are fine

So why am I talking about it?

Simple power analysis: square and multiply



We literally read the secret bits from the power consumption trace!

SPA: A few comments

- ▶ Here 'correlation' really is obvious (sometimes, real math needed)
- ▶ Especially effective against embedded devices
- ▶ How do you prevent this?

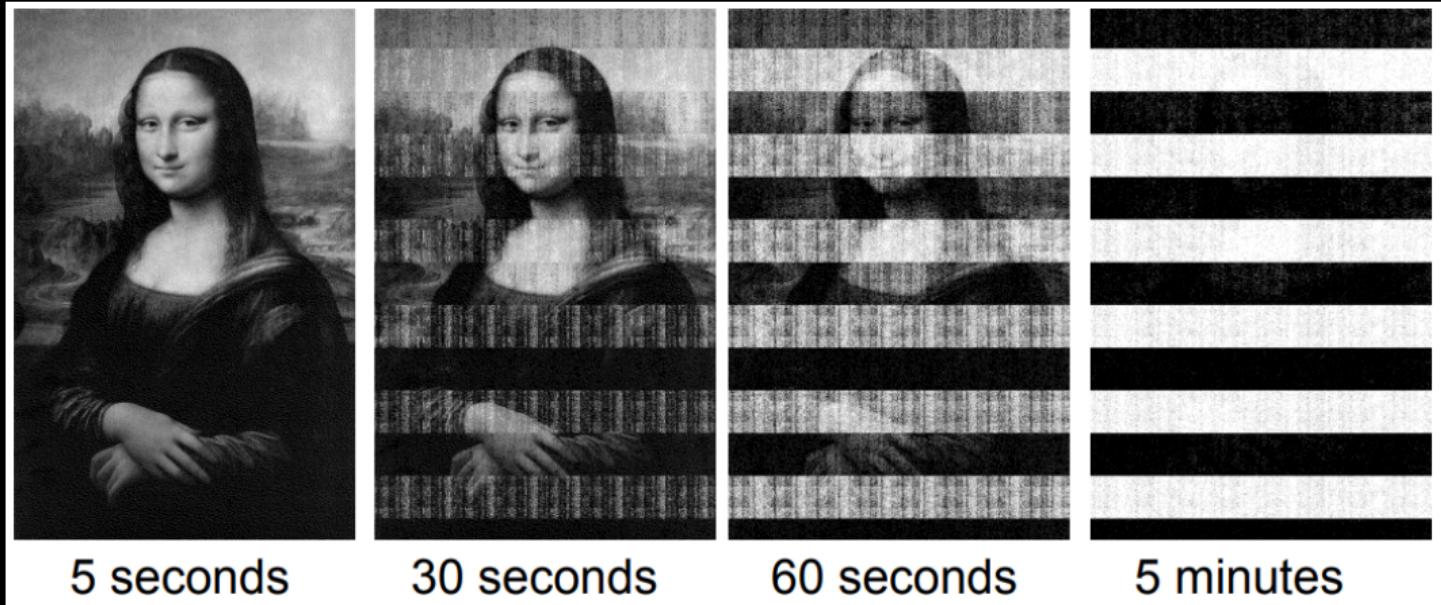
Photons?

Transistors emit photons when they switch

- ▶ 10^2 to 10^4 photons per switch with peak in NIR region (900–1200 nm)
- ▶ Can be detected with photomultipliers and CCD cameras
- ▶ Comes from area close to the transistor drain (mostly NMOS side)

So we can literally see transistors switching.

Data remanence



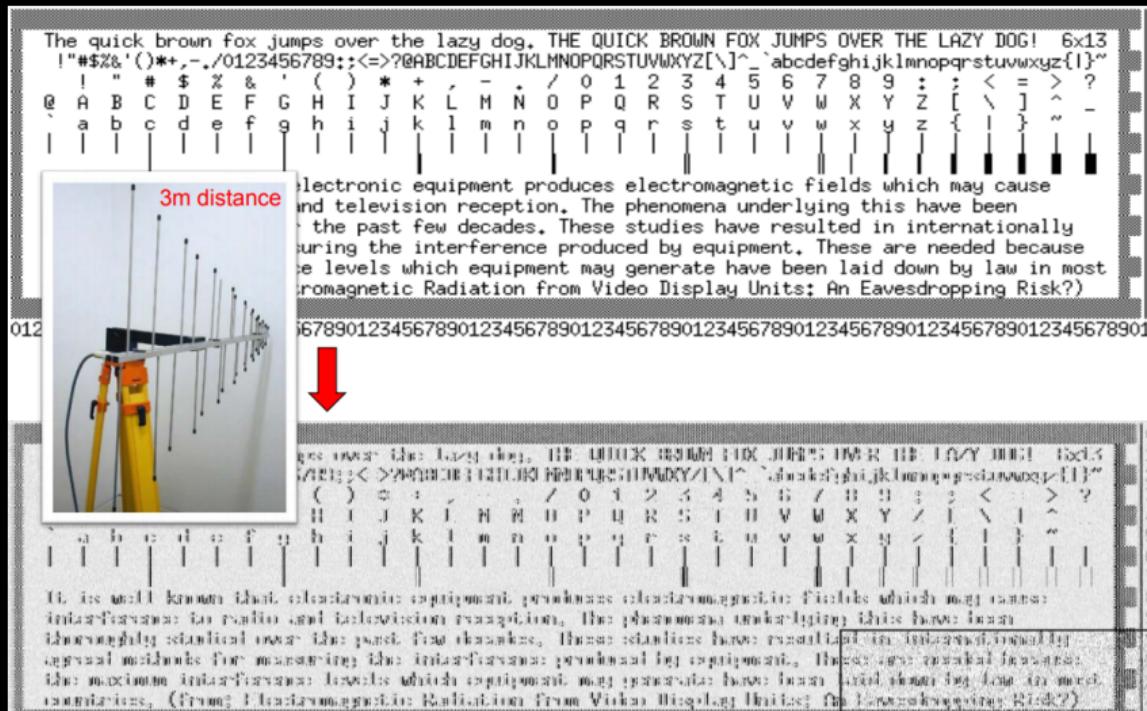
At -50 degrees C, memory remains unchanged for about 10 hours ⇒ Cold boot attacks

Reading screens from 8 km away

How do you call a wire with varying current?

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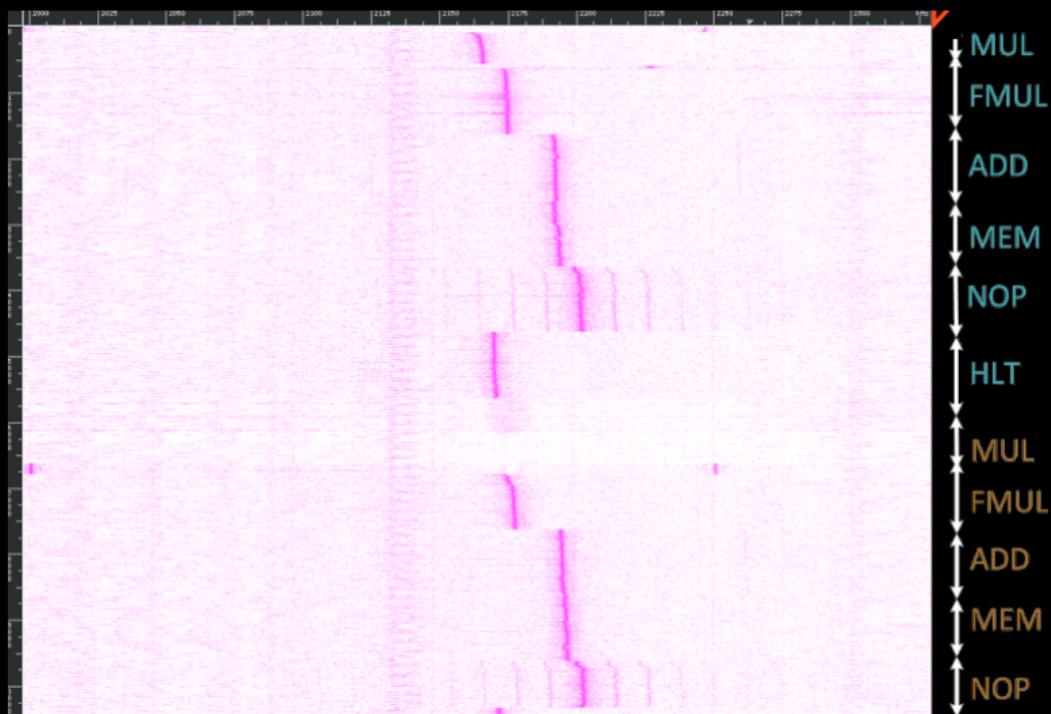
Video?

Reading keystrokes from 25 m away

How do you call a free USB charger/HDMI cable/Power outlet?

Reading keystrokes from 25 m away

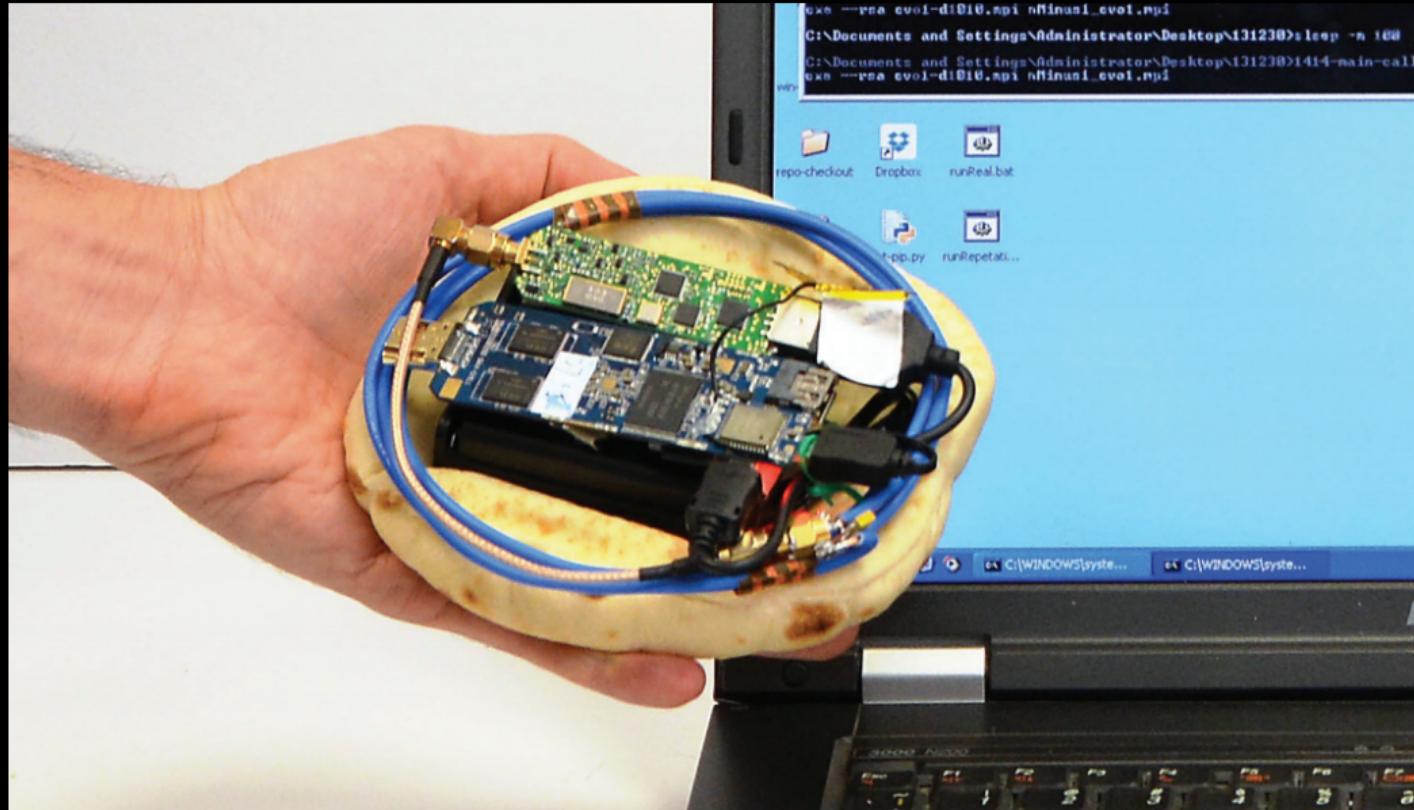
How do you call a free USB charger/HDMI cable/Power outlet?



The “ground” leaks information up to several meters! Even through contact.

Have some hummus?

Portable Instrument for Trace Acquisition – PITA



Acoustic side-channels

Have you ever noticed that your computer makes noise?

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RSA 4096-bit key extraction from 10 meters away using a parabolic microphone.

Hardware trojan

Sometimes there isn't a readily usable side or covert channel.

So maybe we can create one?

This is easier if we are / have access to the manufacturer.

The Clipper Chip 1993



Hardware trojans (cont'd) / DAPINO GAMMA 2010

February 19 2015, 8:25 p.m.

A

MERICAN AND BRITISH spies hacked into the internal computer network of the largest manufacturer of SIM cards in the world, stealing encryption keys used to protect the privacy of cellphone communications across the globe, according to top-secret documents provided to *The Intercept* by National Security Agency whistleblower Edward Snowden.

The hack was perpetrated by a joint unit consisting of operatives from the NSA and its British counterpart Government Communications Headquarters, or GCHQ. The breach, detailed in a secret 2010 GCHQ [document](#), gave the surveillance agencies the potential to secretly monitor a large portion of the world's cellular communications, including both voice and data.

SECRET STRAP 1

CNE access to core mobile networks

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 - Billing servers to suppress SMS billing
 - Authentication servers to obtain K's, Ki's and OTA keys
 - Sales staff machines for customer information and network engineers machines for network maps
 - GEMALTO – successfully implanted several machines and believe we have their entire network – TDSD are working the data

This information is exempt under the Freedom of Information Act 2000 (FOIA) and may be exempt under other UK information legislation. Refer any FOIA queries to GCHQ on [REDACTED] or [REDACTED]. © Crown Copyright. All rights reserved.

SECRET STRAP 1



Hardware trojans (cont'd)

- ▶ 3500 counterfeit Cisco network components were discovered in the US (FBI)
- ▶ Backdoor in a military grade FPGA device (Skorobogatov 2012)
- ▶ Mobile phones developed by Chinese device manufacturer ZTE have been found to carry a backdoor to instantly gain root access (Alperovitch 2012)

China anti-terrorism bill, passed at the end of 2015, took effect on the first day of 2016. Amongst the provisions, is the requirement that every manufacturer provides the Chinese government with “backdoor” access to their products.

About 80% of all electronic devices in circulation are manufactured in China.

TOP SECRET//COMINT//REL TO USA, FVEY

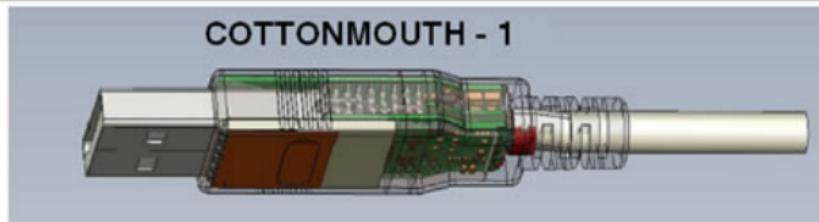


COTTONMOUTH-I

ANT Product Data

08/05/08

(TS//SI//REL) COTTONMOUTH-I (CM-I) is a Universal Serial Bus (USB) hardware implant which will provide a wireless bridge into a target network as well as the ability to load exploit software onto target PCs.



(TS//SI//REL) CM-I will provide air-gap bridging, software persistence capability, "in-field" re-programmability, and covert communications with a host software implant over the USB. The RF link will enable command and data infiltration and exfiltration. CM-I will also communicate with Data Network Technologies (DNT) software (STRAITBIZARRE) through a covert channel implemented on the USB, using this communication channel to pass commands and data between hardware and software implants. CM-I will be a GENIE-compliant implant based on CHIMNEYPOOL.



Così fan tutte

To what extent can we trust hardware components?

What techniques and countermeasures can be used to prevent, detect, or circumvent side-channels/covert channels/hardware trojans?

While you think about it

Let's have a break