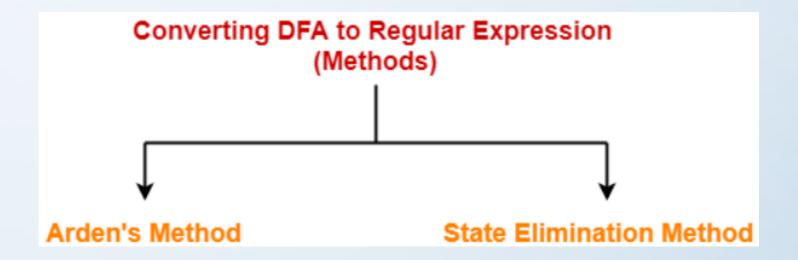


CSC-257 Theory Of Computation (BSc CSIT, TU)

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Conversion of DFA to Regular Expression

A DFA can be converted into RE using two methods.



Conversion of DFA to Regular Expression

- Arden's Theorem: Let p and q be two regular expressions over the alphabet Σ , if p does not contain any empty string then $\mathbf{r} = \mathbf{q} + \mathbf{rp}$ has a unique solution $\mathbf{r} = \mathbf{qp}^*$.
- **Proof**: Here, we have, r = q + rp(i)
- putting the value of r = q + rp on the right hand side of the relation (i),
- We have, r = q + (q + rp)p
- $r = q + qp + rp^2$(ii)
- Again putting value of r = q + rp in relation (ii),
- we get, $r = q + qp + (q + rp)p^2$
- $r = q + qp + qp^2 + rp^3$

Conversion of DFA to Regular Expression

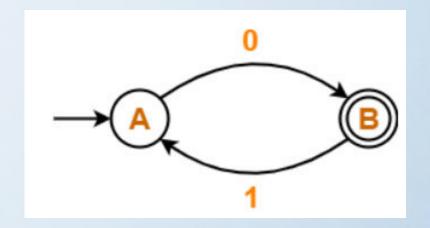
- Continuing the process in same way, we will get as;
- $r = q + qp + qp^2 + qp^3 + \dots$
- $r = q(E + p + p^2 + p^3 + \dots$
- Thus, $r = qp^*$ Proved.

- To convert the given DFA into a regular expression, here are some of the assumptions regarding the transition system:
- 1. The transition diagram should not have the E-transitions
- 2. There must be only one initial state
- 3. Form an equation for each state considering the transitions which comes towards that state. The vertices or the states in the DFA are as; q_1,q_2,\ldots,q_n (Any q_i is final state).
- 4. Note: If there exists multiple final states, then:
 - Write a regular expression for each final state separately.
 - Add all the regular expressions to get the final regular expression.

5. w_{ij} denotes the regular expressions representing the set of labels of edges from q_i to q_i , We can get the following conditions as:

- 6. Add '∈' in the equation of initial state
- 7. Bring final state in the form r = q + rp to get the required regular expression
- Hence, solving these equations for q_i in terms of w_{ij} gives RE

- Convert the following DFA to RE
- Solution :
- Form equation for each state as:
 - $A = \in + B1$ (1)
 - B = A0 (2)
- Bring final state in the form r = q + rp.
- Using (1) in (2), we get :
 - $B = (\in + B1)0$
 - $B = \in 0 + B10$
 - -B = 0 + B(10) (3)
- Using Arden's Theorem in (3), we get:
 - $-B = 0(10)^*$, where r = B, q = 0 and p = (10)
- Thus, Regular Expression for the given DFA = $0(10)^*$ [since, r = qp*]



- Convert the following DFA to RE
- Solution :
- Forming equation for each state as :

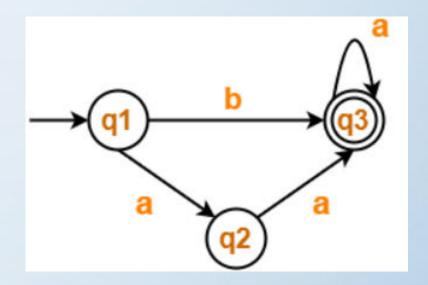
$$- q_1 = \in$$
 (1)
 $- q_2 = q_1 a$ (2)

$$- q3 = q_1b + q_2a + q_3a$$
(3)

- Bring final state in the form r = q + rp as :
- Using (1) in (2), we get:

$$- q_2 = \epsilon a$$

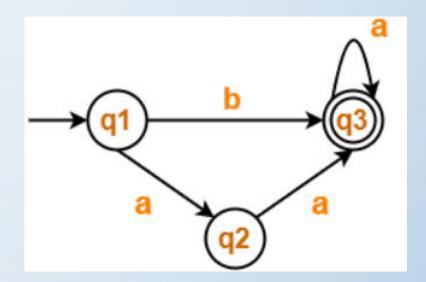
$$- q2 = a$$
 (4)



- Using (1) and (4) in (3), we get:
 - $q_3 = q_1b + q_2a + q_3a$
 - $q_3 = \in b + aa + q_3a$
 - $q_3 = (b + aa) + q_3 a$ (5)
- Using Arden's Theorem in (5), we get:

$$- q_3 = (b + aa)a^*$$





Convert the following DFA to RE

Solution :

Here, the equations are

•
$$q_1 = q_2 1 + q_3 0 + \epsilon$$
....(i)

•
$$q_2 = q_1 0$$
.....(ii)

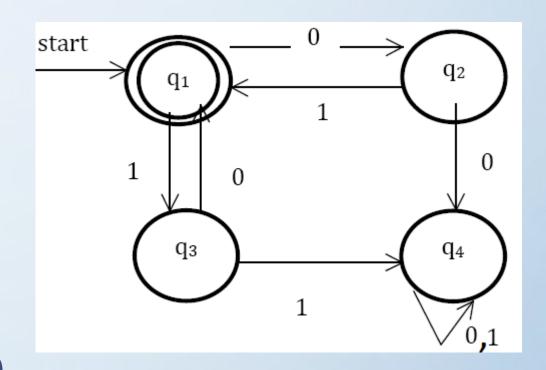
•
$$q_3 = q_1 1$$
.....(iii)

•
$$q_4 = q_2 0 + q_3 1 + q_4 0 + q_4 1 \dots (iv)$$

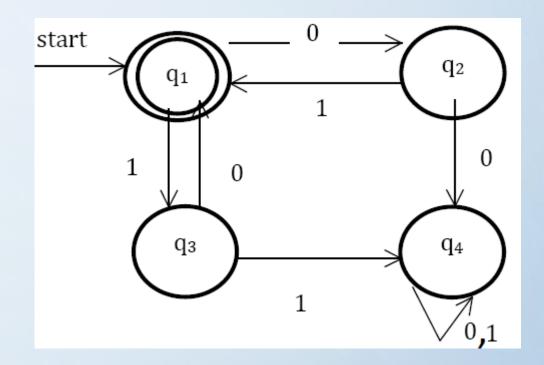


•
$$q_1 = q_101 + q_110 + \epsilon$$

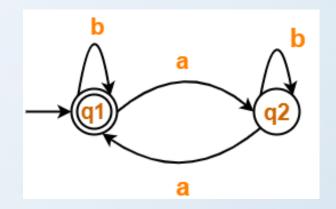
•
$$q_1 = \varepsilon + q_1(01+10)$$



- Let, $- q = \varepsilon$, $- r = q_1$, and - p = (01 + 10)
- Therefore, according to Arden's rule,
- $q1=\epsilon (01 + 10)* [since, r = qp*]$
- since, q1 is the final state,
- so, RE= ε (01+ 10)* = (01+ 10)* is the required RE for the given DFA



Convert the following DFA to RE



Convert the following DFA to RE

