

# CSC-257 Theory Of Computation (BSc CSIT, TU)

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# Chomsky Normal Form (CNF)

• A context free grammar G = (V, T, P, S) is said to be in Chomsky Normal Form (CNF) if every production in G are in one of the two forms :

```
A \rightarrow BC and
```

 $A \rightarrow a$ 

Where A, B, C  $\in$  V and a  $\in$  T

- Thus, a grammar in CNF is one which should not have:
  - E-production
  - Unit productions
  - Useless symbols

• **Theorem**: Every context free language (CFL) without E-production can be generated by grammar in CNF.

#### Proof:

- 1. Simplifications: Eliminate E-productions if any, Eliminate unit productions if any, and Eliminate useless symbols if any
- 2. If all the productions are of the form A  $\rightarrow$  a and A  $\rightarrow$  BC with A, B, C  $\in$  V and a  $\in$  T, we have done. Otherwise,
- 3. If the right side is in the form  $A \rightarrow aB$  or  $A \rightarrow Ba$  where  $a \in T$  and  $A,B \in V$ , then replace the production by  $A \rightarrow XB$  and  $X \rightarrow a$ . Repeat this process for every production of the form  $A \rightarrow aB$  or  $A \rightarrow Ba$
- 4. Replace each production of the form  $A \rightarrow B_1B_2 \dots B_n$  where n > 2 with  $A \rightarrow B_1C$  where  $C \rightarrow B_2 \dots B_n$ . And repeat this step for all the productions having two or more symbols on the right side.  $[A, B_1, B_2, \dots B_n \in V]$
- 5. Finally, all of the productions will be in the form as : A  $\rightarrow$  BC or A  $\rightarrow$  a. This is certainly a grammar in CNF.

Example: Convert the following grammar into CNF.

```
S \rightarrow AAC
A \rightarrow aAb \mid \in
C \rightarrow aC \mid a
```

- Solution:
- 1. Similfication: grammar after removal of E-productions is:

```
S \rightarrow AAC \mid AC \mid C
A \rightarrow aAb \mid ab
C \rightarrow aC \mid a
```

grammar after removal of unit productions is:

```
S \rightarrow AAC \mid AC \mid aC \mid a

A \rightarrow aAb \mid ab

C \rightarrow aC \mid a
```

- There are no any useless symbols
- 2. Resulting grammar is not in CNF
- 3. Here, suppose A'  $\rightarrow$  a and B'  $\rightarrow$  b, then our grammar will look like

```
S \rightarrow AAC \mid AC \mid A'C \mid a
A \rightarrow A'AB' \mid A'B'
C \rightarrow A'C \mid a
A' \rightarrow a
B' \rightarrow b
```

• 4. Now suppose,  $X \rightarrow AC$  and  $Y \rightarrow AB'$ , then grammar will be :

```
S \rightarrow AX \mid AC \mid A'C \mid a
A \rightarrow A'Y \mid A'B'
C \rightarrow A'C \mid a
X \rightarrow AC
Y \rightarrow AB'
A' \rightarrow a
B' \rightarrow b
```

Now, above grammar is in CNF.

• Exercises: Convert the following grammars in CNF.

1.

```
S \rightarrow ASA \mid aB

A \rightarrow B \mid S

B \rightarrow b \mid \epsilon
```

[ **Solution**: https://www.youtube.com/watch?v=FNPSlnj3Vt0]

2.

$$S \rightarrow a \mid aA \mid B$$
  
 $A \rightarrow aBB \mid C$   
 $B \rightarrow Aa \mid b$ 

[ **Solution**: https://www.youtube.com/watch?v=7G0PwGrdlH8 ]