

CSC-257 Theory Of Computation (BSc CSIT, TU)

Ganesh Khatri kh6ganesh@gmail.com

- Step 1: Convert the given productions of CFG into GNF.
- **Step 2**: The PDA will only have one state {q}.
- Step 3: The initial symbol of CFG will be the initial symbol in the PDA.
- **Step 4:** For non-terminal symbol A, add the following rule : $\delta(q, \epsilon, A) = (q, a)$, where the production rule is $A \rightarrow a$
- Step 5 : For each terminal symbols, add the following rule : $\delta(q, a, a) = (q, \epsilon)$ for every terminal symbol a

• Example : Convert the following CFG into PDA

$$S \rightarrow 0SX \mid 1SY \mid \epsilon$$

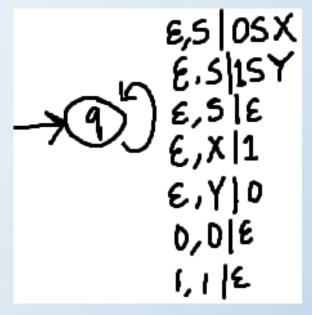
$$X \rightarrow 1$$

$$Y \rightarrow 0$$

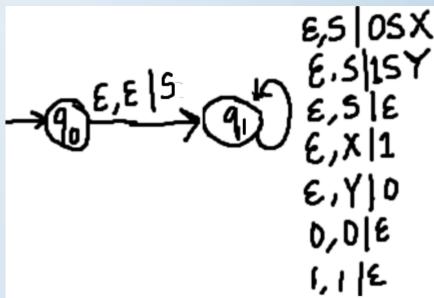
R5: $\delta(q, 1, 1) = \{(q, \epsilon)\}$

- Here, given grammar is already in GNF
- The PDA can be given : $P = \{(q), (0, 1), (S, X, Y, 0, 1), \delta, q, S, q\}$
- The production rule δ can be :

R1:
$$\delta(q, \epsilon, S) = \{(q, 0SX) \mid (q, 1SY) \mid (q, \epsilon)\}$$
R2: $\delta(q, \epsilon, X) = \{(q, 1)\}$
R3: $\delta(q, \epsilon, Y) = \{(q, 0)\}$
R4: $\delta(q, 0, 0) = \{(q, \epsilon)\}$



or



• Testing 1010 against PDA:

•	$\delta(q,$	1010, S)	$\vdash \delta(q,$	1010,	1SY)	R1
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 \vdash δ(q, 010, SY) **R5**

 \vdash δ(q, 010, 0SXY) **R1**

 $\vdash \delta(q, 10, SXY)$ **R4**

 $\vdash \delta(q, 10, XY)$ **R1**

 \vdash δ(q, 10, 1Y) **R2**

 $\vdash \delta(q, 0, Y)$ **R5**

 $\vdash \delta(q, 0, 0)$ **R3**

 $\vdash \delta(q, \epsilon, \epsilon)$ R4

String accepted by empty stack method.

Rules:

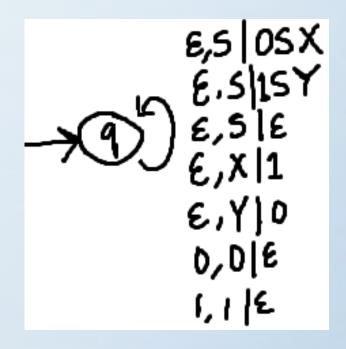
R1: $\delta(q, \epsilon, S) = \{(q, 0SX) \mid (q, 1SY) \mid (q, \epsilon)\}$

R2: $\delta(q, \epsilon, X) = \{(q, 1)\}$

R3: $\delta(q, \epsilon, Y) = \{(q, 0)\}$

R4: $\delta(q, 0, 0) = \{(q, \epsilon)\}$

R5: $\delta(q, 1, 1) = \{(q, \epsilon)\}$



Exercise :

 Construct PDA for the given CFG, and test whether 010000 is acceptable by this PDA.

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S \rightarrow 0BB
B \rightarrow 0S | 1S | 0
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Construct PDA for the given CFG, and trace the acceptance of a + (a*a)

$$E \rightarrow T \mid E + T$$

 $T \rightarrow F \mid T * F$
 $F \rightarrow a \mid (E)$