## CS476: Automata Theory and Formal Languages Homework 3

Assigned: 23/04/2014 Due: 12/05/2014, 17:00

## Questions

- 1. (20pts) State whether the following statements are true or not. You must give a BRIEF explanation or show a counter example to receive full credit.
  - (a) If a language L is decidable then the language  $\bar{L}$  is also decidable.
  - (b) Every deterministic Turing machine (DTM) has an equivalent nondeterministic Turing machine.
  - (c) There exists a polynomial-time reduction from Boolean Satisfiability Problem (SAT) to Subset Sum Problem (SS).
  - (d) If a problem P is NP-complete then the language induced by problem P is undecidable.
- 2. (20pts) Give a Turing Machine that decides language  $L = \{w1^n : |w| = n, w \in \{0, 1\}^*\}$ .
- 3. (30pts) Disprove (by reduction) or prove that the following languages are decidable.
  - (a)  $L = \{\langle A \rangle : A \text{ is a DFA and } L(A) = L'(A)\} \text{ where } L'(A) = \{\bar{w} : w \in L(A)\}$
  - (b)  $L = \{ \langle A, S \rangle : A \text{ is a TM and } L(A) \subseteq S \}$
  - (c)  $L = \{\langle M_1, M_2 \rangle : M_1 \text{ and } M_2 \text{ are TMs such that } M_1 \text{ accepts } \langle M_2 \rangle \text{ and } M_2 \text{ accepts } \langle M_1 \rangle \}$
- 4. (30pts) Number 3 is a lucky number (this is the reason why you have 3 assignments). So, we have two problems (both are related with number "3"): 3-SET-PARTITION and 3-GRAPH-PARTITION. Prove that they are NP-Complete.
  - (a) 3-SET-PARTITION: Let A, B, C be three finite, disjoint sets and let T be a subset of  $A \times B \times C$ . That is, T consists of triples (a, b, c) such that  $a \in A, b \in B$  and  $c \in C$ . Given A, B, C, T and an integer k, the 3-SET-PARTITION problem decides whether there is a subset M of T (i.e.,  $M \subseteq T$ ), such that  $|M| \ge k$  and for any two distinct triples  $(a_1, b_1, c_1) \in M$  and  $(a_2, b_2, c_2) \in M$ , we have  $a_1 \ne a_2, b_1 \ne b_2$  and  $c_1 \ne c_2$ .
  - (b) 3-GRAPH-PARTITION: Given an undirected graph G=(V,E), the 3-GRAPH-PARTITION problem decides whether the vertex set V can be partitioned into three disjoint subsets  $V_1$ ,  $V_2$  and  $V_3$ , such that for any two vertices  $v_a \in V$  and  $v_b \in V$ , if  $(v_a, v_b) \in E$ , then  $v_a$  and  $v_b$  cannot be placed into the same partition  $V_i$ ,  $i \in \{1, 2, 3\}$ .