## CS473-4

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## 1

Assume that before QUERY,  $PRE\_PROCESS$  is called and indexed array will be used for query. An example of code piece is as below.

```
C[0..k] \leftarrow PRE\_PROCESS(A)
number \leftarrow QUERY(C, a, b)
 PRE\_PROCESS(A)
C[0..k] \leftarrow \text{new array}
\mathbf{for}\;i=0\;\mathrm{to}\;k\;\mathbf{do}
  C[i] \leftarrow 0
end for
for j = 1 to A.length do
  C[A[j]] \leftarrow C[A[j]] + 1
end for
\mathbf{for}\;i=1\;\mathrm{to}\;k\;\mathbf{do}
  C[i] \leftarrow C[i] + C[i-1]
end for
return C
 QUERY(C, a, b)
// Assume a > 0 and b < k.
if a == 0 then
  return C[b]
else
  return C[b] - C[a - 1]
end if
```

```
SPARSE-TRANSPOSE(R, C, V, m, n, k)
   R'[0..n+1] \leftarrow \text{new array}
   C'[0..k] \leftarrow \text{new array}
   V'[0..k] \leftarrow \text{new array}
   fill R' with 0s
   for i = 0 to k do
       R'[C[i]] \leftarrow R'[C[i]] + 1
   end for
   cumsum \leftarrow 0
   \mathbf{for}\ \mathrm{col} = 0\ \mathrm{to}\ \mathrm{n}\ \mathbf{do}
       tmp \leftarrow R'[col]
       R'[col] \leftarrow \text{cumsum}
       cumsum \leftarrow cumsum + tmp
   end for
   \mathbf{for}\ i = to\ m\ \mathbf{do}
       \label{eq:for j and kinetic model} \ensuremath{\mathbf{for}}\ j = R[i]\ \mathrm{to}\ R[i+1]\ \mathrm{\mathbf{do}}
           \mathrm{col} \leftarrow C[j]
           \text{dest} \leftarrow R'[col]
           C'[dest] \leftarrow \text{row}
           V'[dest] \leftarrow V[j]
           R'[col] \leftarrow R'[col] + 1
       end for
   end for
   \mathsf{last} \leftarrow 0
   for i = 0 to n + 1 do
       tmp \leftarrow R'[i]
       R'[i] \leftarrow \text{last}
       last \leftarrow tmp
   end for
   return R', C', V'
```

```
Pre:
  struct element {
     int value;
     element * next;
  int \max = -\infty;
  for i = 0 to n do
     if A[i] > \max then
        \max \leftarrow A[i]
     end if
  end for
  element hash[max];
  \mathbf{for}\;i=0\;\mathrm{to}\;n\;\mathbf{do}
     element e \leftarrow new element;
     e.value \leftarrow i
     e.next \leftarrow hash[A[i]]
     hash[A[i]] \leftarrow e
  end for
```

The Greatest element is found and indexing is done by it. Thus, collisions of different values in hashing are prevented. In a row there is only items whose values are same as their index(i). Insertion is done to head (last one in unsorted array is the head in the list).

```
\begin{aligned} Post: & & \textbf{for i} = n - 1 \text{ to } 0 \text{ do} \\ & & A[i] \leftarrow hash[A[i]].value \\ & & hash[A[i]] \leftarrow hash[A[i]].next \\ & \textbf{end for} \end{aligned}
```

All elements from the sorted array replaced.