# The System of Automatic Searching for Vulnerabilities or how to use Taint Analysis to find vulnerabilities

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## Who is Alex Bazhanyuk

Security Researcher

Organizer of Defcon Ukraine Group

Working in UC Berkley in BitBtlaze project

Solves problems of automation of RE

#### Who is Nikita Tarakanov

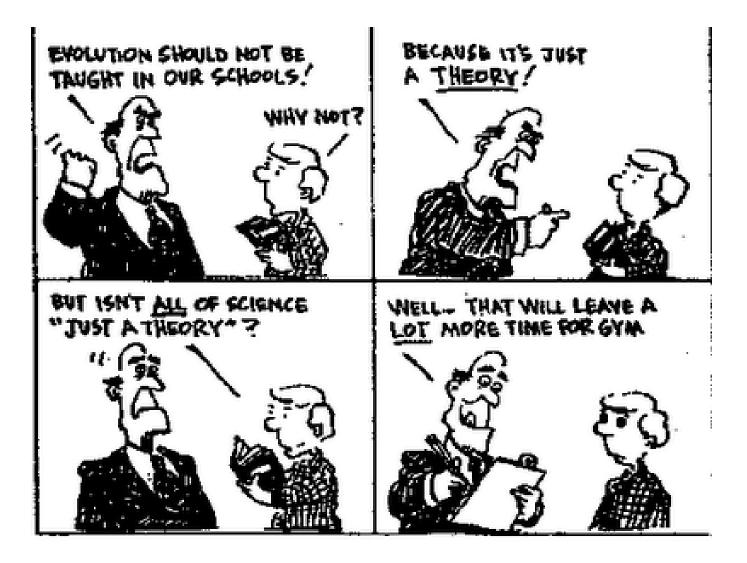
- Independent Security Researcher
- Author of some articles in ][akep magazine
- Likes to reverse engineer r0 parts
- Discovered a lot of LPE vulnerabilities
- Solves problems of automation of RE

## SASV main parts

IDA Pro plugins

BitBlaze: Vine+utils, TEMU + plugins

## Theory



## **Tainting**

Taint sources:

Network, Keyboard, Memory, Disk, Function outputs etc.

Taint propagation: a data flow technique

Memory

Whole-system

Across registers/memory/disk/swapping

## Fundamentals of taint analysis



#### **Taint propagation**

- •If an operation uses the value of some **tainted** object, say X, as assignes value to another, say Y, then object Y becomes **tainted**. Object X taints the object Y
- Taint operator t
- $\bullet X \rightarrow t(Y)$
- Taint operator is transitive
- $X \rightarrow t(Y)$  and  $Y \rightarrow t(Z)$ , then  $X \rightarrow t(Z)$

#### BitBlaze: Binary Analysis Infrastructure



- Automatically extracting security-related properties from
- binary code
- Build a unified binary analysis platform for security
- Static analysis + Dynamic analysis + Symbolic Analysis
- Leverages recent advances in program analysis, formal methods, binary instrumentation...

Solves security problems via binary analysis

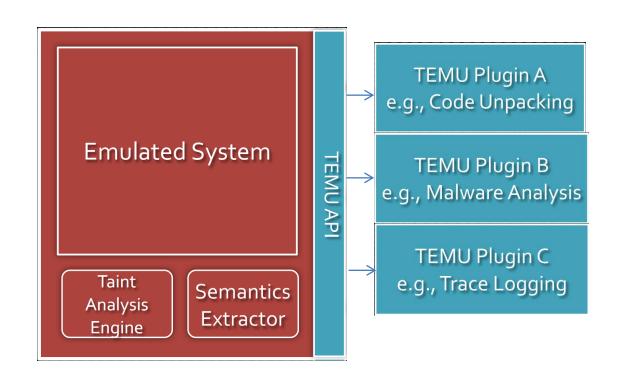
- More than a dozen different security applications
- Over 25 research publications

#### BitBlaze

- http://bitblaze.cs.berkeley.edu/
- TEMU,VINE
- Rudder, Panorama, Renovo

Static Analysis	Dynamic Analysis	Symbolic Exploration
Component	Component	Components
VINE	TEMU	Rudder/ BitFuzz/FuzzBall

#### **TEMU**



#### Limitations of TEMU

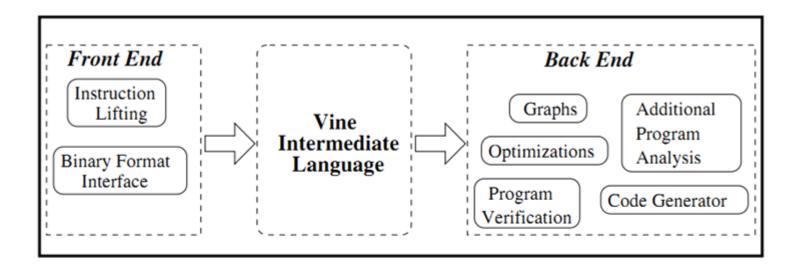
Qemu 0.9.1 - TEMU

Qemu 0.10 - TCG(Tiny Code Generator)-TODO

Qemu 0.10 ⇔ Qemu 1.01

#### **VINE**





## The Vine Intermediate Language

```
program ::= decl* instr*
            ::= var = exp \mid jmp \ exp \mid cjmp \ exp, exp, exp \mid halt \ exp \mid assert \ exp
instr
                 | label integer | special id_s
            ::= load(exp, exp, \tau_{reg}) | store(exp, exp, exp, \tau_{reg}) | exp \diamondsuit_b exp | \diamondsuit_u exp
exp
                 | const | var | let var = exp in exp | cast(cast\_kind, \tau_{reg}, exp)
cast_kind ::= unsigned | signed | high | low
decl
            ::= var var
          ::= (string, id_v, \tau)
var
\Diamond_b
           := +, -, *, /, /_s, \mod, \mod_s, \ll, \gg, \gg_a, \&, |, \oplus, ==, \neq, <, \leq, <_s, \leq_s
\Diamond_n
           ::= - (unary minus), ! (bit-wise not)
           ::= const \mid \{ n_{a1} \rightarrow n_{v1}, n_{a2} \rightarrow n_{v2}, \dots \} : \tau_{mem} \mid \bot
value
           := n : \tau_{\text{reg}}
const
            := \tau_{reg} \mid \tau_{mem} \mid Bot \mid Unit
	au
            ::= reg1_t | reg8_t | reg16_t | reg32_t | reg64_t
	au_{reg}
           ::= mem_t(\tau_{endian}, \tau_{reg})
	au_{mem}
           ::= little | big | norm
	au_{endian}
```

#### Example of disasm:

```
rep stos \%eax,\%es:(\%edi) R@eax[0x00000000][4](R) T0
fc32dcec:
R@ecx[0x00000002][4](RCW) T0 M@0xfb7bfff8[0x00000000][4](CW) T1 {15
(1231, 69624) (1231, 69625) (1231, 69626) (1231, 69627) }
            rep stos %eax,%es:(%edi) R@eax[0x00000000][4](R) T0
fc32dcec:
R@ecx[0x00000001][4](RCW) T0 M@0xfb7bfffc[0x00000000][4](CW) T1 {15
(1231, 69628) (1231, 69629) (1231, 69630) (1231, 69631) }
            mov %edx,%ecx R@edx[0x0000015c][4](R) T0
fc32dcee:
R@ecx[0x00000000][4](W) T0
fc32dcf0:
           and $0×3,%ecx I@0×0000000[0x00000003][1](R) T0
R@ecx[0x0000015c][4](RW)
                             T0
fc32dcf5: and $0\times0,-0\times4(\%ebp) [@0\times00000000[0x00000000][1](R) T0
M@0xfb5ae738[0x00000002][4](RW) T0
fc32dcf9:
           jmp 0x0000000fc32c726
                                       J@0×0000000[0xffffea2d][4](R) T0
            cmpl $0 \times 0, -0 \times 58 (%ebp) 1@0 \times 00000000[0x000000000][1] (R) TO
fc32c726:
M@0xfb5ae6e4[0x00000000][4](R) T0
```

#### Taint info

- T0 means that the statement is not tainted.
- T1 means that the statement is tainted.
- Here's an example of:
- fc32dcec: rep stos% eax,% es: (% edi) R @ eax [0x00000000]
   [4] (R) TO R @ ecx [0x00000001] [4] (RCW) TO M @ 0xfb7bfffc
   [0x00000000] [4] (CW) T1 {15 (1231, 628) (1231, 629) (1231, 630) (1231, 631)}
- 4 bits of information tainted and they depend on the offset: 628, 629, 630, 631. 1231 this number is origin(kind of ID that TEMU plugin sets), and 15 this number of the source type.

#### appreplay

• ./vine-1.0/trace\_utils/appreplay -trace font.trace -ir-out font.trace.il -assertion-on-var false-use-post-var false

#### where:

- appreplay ocaml script that we run;
- -trace the way to the trace;
- -ir-out the path to which we write IL code.
- -assertion-on-var false-use-post-var false flags that show the format of IL code for this to false makes it more readable text.

#### Example of IL code:

- Begins with the declaration of variables:
- INPUT it's free memory cells, those that are tested in the very beginning (back in TEMU), input into the program from an external source.

```
var cond_000017_0x4010ce_00_162:reg1_t;

var cond_000013_0x4010c3_00_161:reg1_t;

var cond_000012_0x4010c0_00_160:reg1_t;

var cond_000007_0x4010b6_00_159:reg1_t;

var INPUT_10000_0000_62:reg8_t;

var INPUT_10000_0001_63:reg8_t;

var INPUT_10000_0002_64:reg8_t;

var INPUT_10000_0003_65:reg8_t;

var mem_arr_57:reg8_t[4294967296]; — memory as an array

var mem_35:mem32l_t;
```

```
R_EAX_5:reg32_t =
0×73657930:reg32 t;
var idx_144:reg32_t;
var val_143:reg8_t;
idx_144:reg32_t =
0x12fef0:reg32 t;
val 143:reg8 t =
INPUT 10000 0000 62:reg
8 t;
mem_arr_57[idx_144:reg32
_{t} + 0:reg32_{t}:reg8_{t} =
cast((val_143:reg8_t &
Oxff:reg8_t) >>
0:reg8_t)L:reg8_t;
```

```
T 32t2 \ 60:reg32 \ t = R \ ESP \ 1:reg32 \ t;
T 32t1 59:reg32 t = T 32t2 60:reg32 t
    + 0x1c8:reg32 t;
T_32t3_61:reg32_t = ((
cast(mem_arr_57[T_32t1_59:reg32_t +
    0:reg32 t]:reg8 t)U:reg32 t
<< 0:reg32 t
cast(mem arr 57[T 32t1 59:reg32 t +
    1:reg32 t]:reg8 t)U:reg32 t
<< 8:reg32 t)
cast(mem_arr_57[T_32t1_59:reg32_t +
    2:reg32 t]:reg8 t)U:reg32 t
<< 0×10:reg32 t)
cast(mem_arr_57[T_32t1_59:reg32_t +
    3:reg32 t]:reg8 t)U:reg32 t
<< 0×18:reg32 t
R EAX 5:reg32 t = T 32t3 61:reg32 t;
```

#### What is STP and what it does?

- STP constraint solver for bit-vector expressions.
- separate project independent of the BitBlaze
- To produce STP code from IL code:
- ./vine-1.0/utils/wputil trace.il -stpout stp.code
- where the input is IL code, and the output is STP code

## STP program example

```
mem arr 57 8: ARRAY BITVECTOR(64) OF BITVECTOR(8);
INPUT 10000 0000 62 4:BITVECTOR(8);
ASSERT( 0bin1 =
(LET R EAX 5 232 =
0hex73657930
IN
(LET idx 144 233 =
0hex0012fef0
IN
(LET val 143 234=
INPUT 10000 0000 62 4
IN
(LET mem arr 57 393 =
idx 144 233,0hex00000000))]:=(val 143 234;0hexff)[7:0])
IN
(cond_000017_0x4010ce_00_162_392;0bin1))))));
Is this expression false?
QUERY (FALSE);
And give a counter example:
COUNTEREXAMPLE:
```

#### STP output example

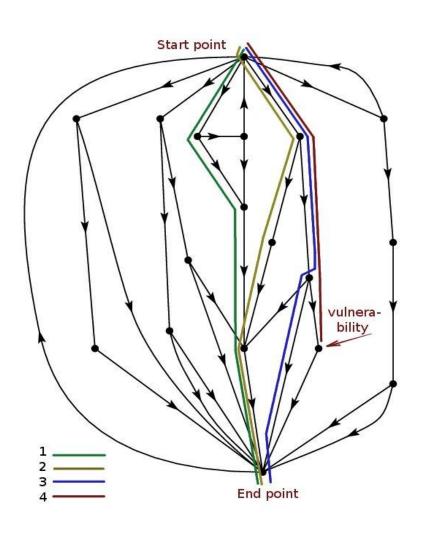
- ./stp stp.code
- Example of STP output:

```
ASSERT(INPUT_10000_0001_63_5 = 0 \times 00);
ASSERT(INPUT_10000_0002_64_6 = 0 \times 00);
ASSERT(INPUT_10000_0000_62_4 = 0 \times 61);
ASSERT(INPUT_10000_0003_65_7 = 0 \times 00);
Invalid.
```

#### **SASV Components:**

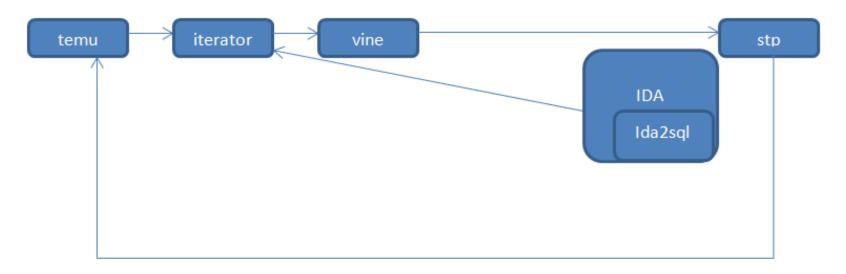
- **Temu** (tracecap: start/stop tracing. Various additions to tracecap(hooks etc.))
- Vine (appreplay, wputil)
- STP
- IDA plugins:
- DangerousFunctions finds calls to malloc,strcpy,memcpy etc.
- IndirectCalls indirect jumps, indirect calls.
- ida2sql (zynamics) idb in the mysql db. (<a href="http://blog.zynamics.com/2010/06/29/ida2sql-exporting-ida-databases-to-mysql/">http://blog.zynamics.com/2010/06/29/ida2sql-exporting-ida-databases-to-mysql/</a>)
- **Iterators** wrapper for temu, vine, stp.
- Various publishers for DeviceloControl etc.

#### How does SASV work?



#### **SASV**

• Scheme:

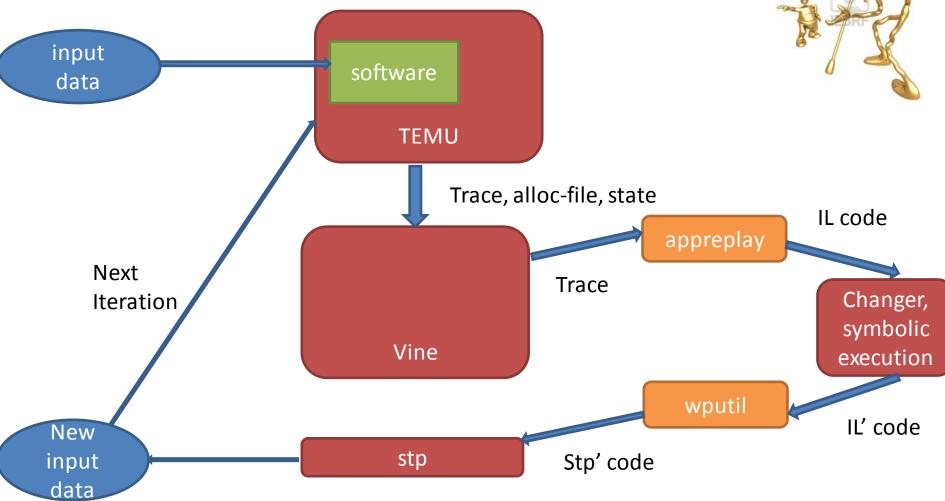


- Min Goal: max coverage of the dangerous code
- Max Goal: max coverage of the all code

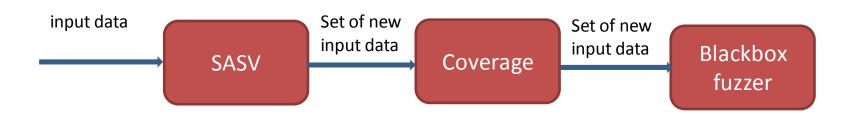
#### SASV basic algorithm

- 1. Work of IDA plugins -> dangerous places
- 2. Publisher(s) -> invoke targeted code
- 3. TEMU -> trace
- 4. Trace -> appreplay -> IL
- 5. IL -> change path algo -> IL'
- 6. IL' -> wputil -> STP\_prorgam'
- 7. STP\_prorgam' -> STP -> data for **n+1** iteration
- 8. Goto #2

## Diagram for new path in graph



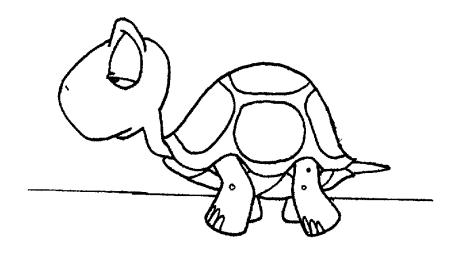
## Combo system: Dumb+Smart



#### Disadvantages

- Definition of the vulnerability is difficult task.
- Performance speed of tracing in TEMU is

## **AWFUL**





#### Get rid of that damned QEMU!

Move taint propagation to Hypervisor!

Damn good idea!

But not implemented yet ☺



#### Vulnerabilities in drivers

- Overflows: stack, pool, integer
- Pointer overwrite
- Null pointer dereference
- Race condition
- Various logical vulnerabilities

## Attack vectors(r3->r0)

IOCTL

SSDT hooks(Native & Shadow)

various notification routines

#### DeviceloControl

#### Parameters:

- hDevice
- dwIoControlCode
- IpInBuffer
- nInBufferSize
- IpOutBuffer
- nOutBufferSize
- IpBytesReturned
- IpOverlapped

#### Concept

#### **IOCTL**:

#### Data to taint:

- dwloControlCode to get list of supported ioctl codes
- IpInBuffer pointer(METHOD\_NEITHER) and data (METHOD\_BUFFERED)
- nInBufferSize size ranges
- IpOutBuffer pointer(METHOD\_NEITHER) and data (METHOD\_BUFFERED)
- nOutBufferSize-size ranges

#### Tracing only targeted driver code

# •Odayz Time!

- loctl code 0x220044 (METHOD\_BUFFERED)
- No range check for size
- Just check for correct address NPD check (MmIsAddressValid)
- Pool corruption in cycle
- No control of overflowing data

```
.text:0001D881
```

- .text:0001D884
- .text:0001D885
- .text:0001D887
- .text:0001D889
- .text:0001D88F
- .text:0001D892
- .text:0001D893
- .text:0001D898
- [..]

```
mov edi, [ebx+0Ch]
push edi
               ; our buffer
call esi; MmIsAddressValid
test al, al
    loc 1DDAB
İΖ
push [ebp+output_buff_size]
push edi
push offset rules_list
call
     ioctl 0x220044 vuln
```

```
ebx, [ebp+our_buffer_size_controlled]
.text:000156EA
                       mov
.text:000156ED
                              [ebp+NewIrql], al
                       mov
.text:000156F0
                             eax, dword 22CA0
                       mov
                              edx, offset dword 22CA0
.text:000156F5
                       mov
.text:000156FA
                              eax, edx
                       cmp
.text:000156FC
                       įΖ
                            short loc 15748
[..]
.text:00015700
                              ecx, [eax+0Ch]
                       mov
.text:00015703
                              [ebx], ecx
                       mov
.text:00015705
                              ecx, [eax+10h]
                       mov
                             [ebx+4], ecx
.text:00015708
                       mov
                              ecx, [eax+14h]
.text:0001570B
                       mov
.text:0001570E
                              [ebx+8], ecx \leftarrow write outside of the pool chunk
                       mov
.text:00015711
                              ecx, [eax+18h]
                       mov
.text:00015714
                              [ebx+0Ch], ecx
                       mov
```

- loctl code 0x220030
- Range check for inbuff\_size >= 0x2AA
- Range check for outbuff\_size >= 0x4D0
- Allocs pool memory for const size 0x4D0
- And...
- Zeroing it with outbuff\_size length! LOL

```
.text:0001D704
                            [ebp+inbuff size], 2AAh
                      cmp
.text:0001D70B
                          loc_1DDAB
                      jb
                      mov esi, 4D0h
.text:0001D711
                      cmp [ebp+output_buff_size], esi
.text:0001D716
                          loc 1DDAB
.text:0001D719
                      jb
                      push 746D74h
.text:0001D71F
                                        ; Tag
.text:0001D724
                      push esi
                                    ; NumberOfBytes
                      push 0
.text:0001D725
                                    ; PoolType
                      call ds:ExAllocatePoolWithTag
.text:0001D727
```

[..]

```
push edi; pool_mem_const_size
.text:0001D74B
                            eax, [ebp+output_buff_size]
 .text:0001D74C
                       lea
.text:0001D74F
                       push eax ; output_buff_size
                       push [ebp+NewIrql] ; inbuff
.text:0001D750
                       push 220030h ; ioctl code
.text:0001D753
                            ioctl several ioctl codes
.text:0001D758
                       call
[..]
                              esi, [ebp+output_buff_size]
.text:00014918
                       mov
[..]
                             dword ptr [esi];
.text:00014977
                       push
.text:00014979
                       push 0
                       push [ebp+pool_mem_const_size];
.text:0001497B
 .text:0001497E
                       call
                            memset
```

## Pitfalls of tainting r0

Taint info lost

Check of system environment variables

System defense mechanism(s) (win32k.sys

WATCHDOG BugCheck)

## Pitfalls of tainting r0(IOCTL)

KeGetPreviousMode

IoGetCurrentProcess

Even hooking NtDeviceIoControlFile!

#### How some AVs kill LPE Odayz

Check for previous mode:

• .text:0001DC32 cmp byte ptr [ebx+**20**h], 0

.text:0001DC36 jnz loc\_1DDAB

.text:0001DC3C mov eax, [edi]

The vuln is here, dword\_22934 is function ptr

.text:0001DC3E mov dword\_22934, eax

## Thanks, ©

## •Questions?

http://twitter.com/#!/ABazhanyuk
http://twitter.com/#!/NTarakanov