

**CROWDY**  
**GENERAL-PURPOSE, EXTENSIBLE**  
**CROWDSOURCING FRAMEWORK**

A THESIS

SUBMITTED TO THE DEPARTMENT OF COMPUTER ENGINEERING  
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MASTER OF SCIENCE

By  
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April, 2014

I certify that I have read this thesis and that in my opinion it is fully adequate,  
in scope and in quality, as a thesis for the degree of Master of Science.

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# ABSTRACT

## CROWDY GENERAL-PURPOSE, EXTENSIBLE CROWDSOURCING FRAMEWORK

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M.S. in Computer Engineering

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April, 2014

Write the abstract here. Please note that an abstract is not the narrative form of contents page. Instead, it should give the same content as in the thesis, and must not be just a pointer to the content of the thesis. Avoid unnecessary words like “In this thesis,”; it is clear that this is the abstract of this thesis, anyway. A good abstract should contain balanced material from each and every chapter of the thesis, and tell essentially the same thing. A common mistake is to write too much from introduction and conclusion chapters, but too little from anything in between.

*Keywords:* Key one, key two.

ÖZET

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# Chapter 1

## Introduction

The asynchrony, heterogeneity and inherent loose coupling that characterize applications promote system of systems as a natural design abstraction for growing class of software systems. This new concept goes beyond the size of current system definition by several measures such as number of people the system employed for different purposes; number of connections and interdependencies among components; number of hardware elements; amount of data stored, accessed, manipulated, and refined and number of lines of code. Such systems are called Ultra-Large-Scale (ULS) systems and perceived as socio-technical ecosystems [1].

The socio-technical aspect of a system is a result of decentralized and dynamic structure formed by people and software components interacting in complex ways. People become not only users, but also an integral part of the system providing content and computation, and the overall behavior [1]. The difference between the roles concerning system components and humans (user, developer) becomes less distinct. The homogeneity of components ceases together with the increasing scale and variety of people and software components involved within the system. Thus, system gains a social and technical characteristic with the ability to solve numerous problems, even the ones requiring human intelligence.

Social aspect of systems has expanded the scale of collaboration from small

or medium sized to internet-scale [2] leading to new era of computation. Collaboration of creative and cognitive people with number-crunching computer systems have appeared under many names such as crowdsourcing, human computation, collective intelligence, social computing, global brain etc, for which you can find detailed studies on classification of systems and ideas in [3, 4] collected under distributed human computation term.

The term *crowdsourcing*, which is the main consideration in this body of work, is first coined by Jeff Howe in the June 2006 issue of Wired magazine [5] as an alternative to the traditional, in-house approaches focusing on assigning tasks to employees in the company for solving problems. Crowdsourcing describes a new, mainly web-based business model that exploits collaboration of individuals in a distributed network through an open call for proposals. The term is described by Howe as:

Simply defined, crowdsourcing represents the act of a company or institution taking a function once performed by employees and outsourcing it to an undefined (and generally large) network of people in the form of an open call. This can take the form of peer-production (when the job is performed collaboratively), but is also often undertaken by sole individuals. The crucial prerequisite is the use of the open call format and the large network of potential laborers. [6]

Crowdsourcing as the new and powerful mechanism of computation has become compelling to accomplish work online [7]. Over the past decade, numerous crowdsourcing systems have appeared on the Web (Threadless, iStockphoto, Innocentive etc). Such systems enabling excessive collaboration of people have provided solutions to the problems and tasks that are trivial for humans, but cannot be easily completed by computers or computerized. Hundreds of thousands of people have worked on various tasks including deciphering scanned text (recaptcha.net), discovering new galaxies (galaxyzoo.org), seeking missing people (helpfindjim.com), solving research problems (Innocentive), designing t-shirts (Threadless). Even Wikipedia and Linux can be viewed as crowdsourcing systems from a point of view that conceives crowdsourcing as explicit collaboration

of users to build a long-lasting and beneficial artefact [8].

Indeed, one of the most interesting developments in terms of crowdsourcing is Amazon’s Mechanical Turk (MTurk), which is a general purpose crowdsourcing platform recruiting large numbers of people to complete diverse jobs. Assignments on MTurk range from labeling images with keywords, transcribing an audio snippet, finding some piece of information on the Web. Requesters submit jobs, which are called Human Intelligent Tasks or HITs in MTurk parlance, as HTML forms. Respectively workers, who are the crowd of users, (called Turkers on MTurk) perform or complete these jobs by inputting their answers and receiving a small payment in return. This actual platform and other example systems listed above present the potential to accomplish work in different areas within less time and money required by traditional methods [9, 10].

## 1.1 Problem

Although current crowdsourcing systems led by MTurk allow a variety of tasks to be completed by people, the tasks requested for completion are typically simple. Tasks, often described as micro-tasks, have the following two fundamental characteristics:

**Difficulty.** Tasks are narrowly focused, low-complex and require little expertise and cognitive effort to complete (taking a couple of seconds to a few minutes).

**Dependency.** Tasks assigned to humans are independent of each other. The current state of one job has no effect on the other. The result of one job cannot be input to the other to create some information flow.

In that sense, simplicity makes the division and distribution of tasks among individuals easy [11], and independency enables parallelizing and bulk-processing tasks. However, solving more complex and sophisticated problems requires effective and efficient coordination of computation sources (human or software) rather

than creating and listing a series of micro-tasks to-be-completed.

Recently detailed analyses on current mechanisms based on foundations of crowdsourcing reveal the necessity of a more sophisticated problem-solving paradigm. Researchers explicitly state the need for a new generic platform with ability to tackle advanced problems. Kittur et al. [12] suggest researchers to form new concepts of crowd work beyond the simple, independent and deskilled tasks. Based on the fact that complex work cannot be accomplished via existing simple and parallel approaches, the authors state the requirement for a platform to design multi-stage workflows to complete complex tasks, which can be decomposed into smaller subtasks, by appropriate groups of workers selected through a set of constraints.

In another piece of work, Bernstein et al. [13] regard all the people and computers as constituting a *global brain*. Authors indicate the need for new powerful programming metaphors that can more accurately demonstrate the way people and computer work in collaboration. These metaphors are expected to solve dependent sections of more complex problems by decompositions and management of interdependencies. Further, the specification of task sequence and information flow are expected to enable deliberate collaboration over solutions.

However, recent research only partially addresses these challenges by providing programming frameworks and models [7, 14, 15, 16, 17, 18, 19, 20] for massively parallel human computation, limited-scope and ability user interfaces [10, 21, 22, 23], concepts for planning [24], analysis of collaboration [25]. These works fail to tackle challenges of crowdsourcing due to various reasons: having rigid structure and requirements, being only applicable to a small and bounded problem-set, focusing on a specific aspect of crowdsourcing, being developed in an ad hoc manner, requiring a significant amount of work in order to implement and integrate.

Further, human workers are often regarded as homogeneous and interchangeable due to the issues of scalability and availability in existing mechanisms [14]. However, people in a crowd have different skills, and can perform different roles based on their interests and expertise [11]. Current services are created without

considering the availability and preferences of people, constraints and relationships, and the support of dynamic collaborations [26]. Thus, human involvement in current mechanisms should be rethought due to limited support for collaboration and ignorance of collaboration patterns in problem-solving [27] over general-purpose infrastructures that can more accurately reflect the collaboration of people and computers [9, 13].

Nevertheless, the development of more generic crowdsourcing platforms along with new applications and structures are expected by the research community [8].

## 1.2 Purpose

To address these issues, a new general-purpose and extensible platform for crowdsourcing is proposed in this work. The new platform, *Crowdy*, is a computing stack to design and develop crowdsourcing applications for effective and efficient collaboration of human and software components. The platform consists of an application editor, a runtime environment and computation resources. Users design applications by simply creating and connecting operators together. These applications are submitted to runtime environment. The runtime environment executes applications by creating processes. A process is performed via corresponding computation resources, which can be either people or software. In the case of human computation existing crowdsourcing services are used. Otherwise, computers are utilized.

This platform is concentrated on providing mechanisms that can be used to decompose the implementation of an application into a set of components as a close representation of the real-world problem. The main characteristic of this work is to show how sophisticated problems can be accomplished cleanly and easily relying on component-based model.

## 1.3 Thesis Plan

The remainder of the work continues as follows. In Chapter 2, background information related to this study is given and previous works are examined. Chapter 3 provides the details of the proposed platform. Motivating examples are described and developed using the proposed platform in Chapter 4. In Chapter 5, a brief discussion of future work is given. A final chapter concludes this study.

# Chapter 2

## Background

In the following, first the preliminary information about crowdsourcing concepts is summarized, since crowdsourcing is a relatively area of computation. Later, the building blocks of a crowdsourcing platform are described, and related studies are examined thoroughly over this description.

### 2.1 Preliminaries

Crowdsourcing term is first proposed by Jeff Howe [5] in 2006. Although there is a long list of terms (collective intelligence, social computing, human computation, global brain etc.), in which some of them are new and some others are old, the use of term "crowdsourcing" in the academia (demonstrated in Figure 2.1) indicates that research on this domain is promisingly increasing.

Researchers have been exploring different approaches to employ crowd of people in solving various problems. From task creation to quality control there has been a lot of research in crowdsourcing. In most of these studies, Amazon's Mechanical Turk (MTurk) is the main consideration on crowdsourcing where new proposals are implemented and executed on this platform. MTurk is an intermediary platform between employers (known as requesters) and employees (workers



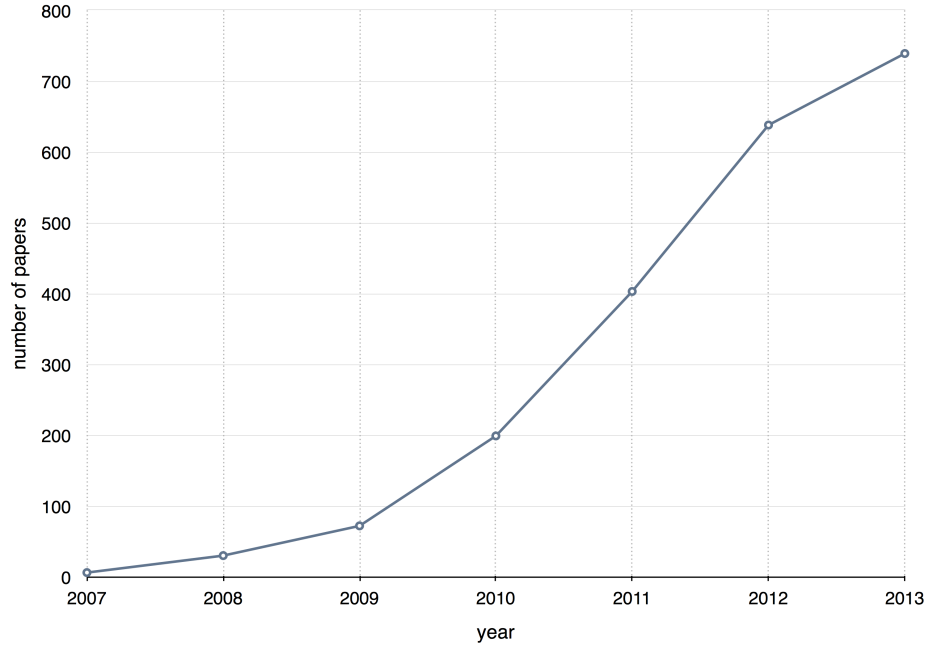


Figure 2.1: The frequency of crowdsourcing keyword in academic papers.<sup>1</sup>

or turkers) for various-sized and difficulty assignments. The platform has become successful for tackling the problems that cannot be solved by computers and subject to numerous research studies.

Similar to other academic studies, MTurk is considered in this work to discuss the common terms and form a preliminary list of concepts. In the following, the fundamental concepts of crowdsourcing is explained.

### 2.1.1 Tasks

Task is a piece of work to be done by the crowd. Terms task, micro-task and job are interchangeably used. The term Human Intelligent Task or HIT is commonly used as well. A task is often expressed over an HTML form in which there are three different input types: single selection (radio buttons), multiple selection (checkboxes) and free text (text areas). For single and multiple selection one or more items can be selected from a list of possible options. Considering free text

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<sup>1</sup>The statistics are gathered through ACM Digital Library.

types workers are supposed to enter a textual response that can be paragraph(s) or sentence(s) or number(s).

In addition to the labels attached to input forms, task has a short description that provides the instructions and keywords, which will help workers search for specific tasks. Further, number of assignments (copies) that requester wants completed per task can be set. Additional copies of the same task allows parallel processing. In that case, the system is supposed to ensure that the parallel workers are unique (i.e., no single worker complete the same task more than once).

A task is generally simple requiring small amount of time and expertise to complete. In the following, sample tasks from MTurk are presented.

[Sample Tasks Figure Will Come Here]

Tasks can be grouped into task groups. Task groups are for the tasks that share similar qualities such as tasks to tag images of nature and people or tasks asking for translation of a text snippet from a language to another.

#### **2.1.1.1 Time**

Often time scale requiring to complete a task ranges from seconds to minutes. Nearly 20% of tasks takes less than 1-hour, and more than half of tasks does not take more than 16-hours [28]. Maximum time allotted per assignment can be set by requester.

#### **2.1.1.2 Payment**

Compensation or reward for completing tasks range from a single penny to dollars. An analysis of MTurk showed that 90% of tasks pays \$0.10 or less [28].

### **2.1.1.3 Acceptance**

Once a worker completes and submits an assignment, the requester is notified. The requester has an option to either accept or reject the completed assignment. Acceptance of an assignment, which indicates the work done is satisfactory, makes the worker who completed it get paid, on the other hand rejection withholds payment for which the reasoning may or may not be provided by the requester. Another option is automatic approval, which is the case when requester does not review work after a some time that can be set by the requester.

### **2.1.1.4 Expiration**

Tasks have a lifetime limit that decides the amount of time that a particular task remains in the listings. Lifetime can be set by requester. After a task reaches end of its lifetime, it is automatically pulled from the listings.

Another type of expiration can happen while a worker is operating on an assignment. When workers accept an assignment, the assignment is reserved for them making no other worker to accept it. The reservation is for particular piece of work (in this case assignment) and time-limited. If worker does not submit the completed assignment in allotted time, then reservation is cancelled and assignment is made available for others again.

## **2.1.2 Requesters**

Requesters are the employers who post tasks, obtain results and pay workers. Requesters are expected to design tasks with all the details (description, instructions, question types, payment, expiration settings etc.) After completed assignments are submitted, requesters can review them by accepting or rejecting, and collect the completed work.

These operations can be done via user interface or application programming

interface (API) if there is one.

### **2.1.3 Workers**

Workers are the online users or someone from the crowd who work on and complete assignments in exchange for a small payment. On some platforms (currently not available on MTurk), detailed information about workers are gathered and kept in a database. This information can be later utilized by requesters while associating specific constraints to the assignments such as limiting age to some range for a specific task.

## **2.2 Building Blocks of a Platform**

Kittur et al. provides a detailed description for the future crowdsourcing platform in [12]. Based on this description building blocks of a crowdsourcing platform is defined in the following.

A crowdsourcing platform is a platform to manage tasks and workers in the process of solving problems through multi-stage workflows. Platform should enable decomposition of complex tasks into subtasks and assign them appropriate group of workers. Workers are formed by people with different skills and expertise. The motivation of workers is guaranteed through various approaches such as reputation or payment. Quality assurance is required to ensure the output of a task is high quality and contributes to solution.

Regarding this definition, a crowdsourcing platform consists of three basic elements:

### 2.2.1 Workflow Design

Complex and sophisticated problems cannot be solved by current simple parallel approach. The reason is that these problems have interdependent portions with changing space, time and skill-set requirements. On the other hand, current approaches undertake problems through bulk processing of independent and simple tasks that require no specific expertise. Therefore, a more appropriate problem solving approach is required.

The need for a more advanced approach is mentioned in [12]. Authors indicate *workflows* in the sense of decomposing problems into smaller chunks, coordinating dependencies among these chunks and combining results from them. In that sense, workflow design refers to the development of applications to solve complex problems. It considers design and management of computation resources to solve the problem.

Researchers often apply existing programming paradigms to problem specification and workflow design. MapReduce [7, 14], Divide-and-Conquer [19], Iterative [16], Workflow [15] are some of these approaches taken for organizing and managing complex workflows. Although these paradigms perfectly fit to some problems, there are other problems not suited well to those.

In terms of crowdsourcing, workflow could involve different scale of operations that can only be solved by diverse set of actors. A workflow may tackle highly dependent tasks (e.g., translating a poem and assuring its quality) or massively parallel tasks (e.g., finding some piece of information on the Web). Therefore, the workflow design has diverse space requirements.

### 2.2.2 Task Assignment

In the context of crowd work, coordination of limited resources is prominent, because problems has various types of tasks that require coordination and collaboration of different computation resources. Crowdsourcing problems have pieces

that cannot be computerized, but require human intelligence. In addition, there may be other problems that require both human and software resources.

Although availability and assignment of software resources can be managed by certain algorithms, the case for human resources is rather unpredictable and hard to solve. Human resources are formed by workers and the availability of a worker cannot be predicted or known at any time. In fact, human resources are not homogenous like computers. They are rather heterogeneous and each worker in the crowd has a specific set of skills. Therefore, there may be a situation in which a task couldn't be assigned to any available workers due to mismatch between task's requirements and skill-set of available workers. Ideally workers are employed with tasks that fit their area of expertise. In practice, requesters may disregard this constraint, but results would not be useful at all.

One important point is that existing systems are misleading by assuming human crowd is homogeneous. Task assignment needs to be redefined by demonstrating aspects of human resources.

### **2.2.3 Quality Control**

Quality control is a big challenge for crowdsourcing platforms, since low quality work is common. It is shown that low quality submissions can compromise up to one third of all submissions [21]. As a result, researchers have started to investigate several ways of detecting and correcting for low quality work.

Visualization of workflow is one of the methods employed for which directed graphs are used to show the organization of crowdsourcing tasks, allowing end-users to better understand the problem and proposed solution design [19, 20].

Inserting 'gold standard' questions into an assignment by which workers who answer them incorrectly can be filtered out or given feedback [29]. However, writing validation questions create extra burden to requesters and may not be applied to all types of tasks.

Majority voting to identify good submissions is proposed as another option [29, 21], but this technique can be affected by majority (especially when the possible options are constrained) or failed in situations where there are no answers in common such as creative or generative work [23].

Systems (including MTurk) often do not apply any of these quality control approaches, but provide other ways to achieve good workers and discourage bad ones. Currently each worker on MTurk has an acceptance rate that is updated after requester’s review on completed assignments. However, this feature does not differentiate one type of task from the other in terms of effect on acceptance rates and that makes it a limited utility. A worker who is skilled in audio transcription would probably have high accuracy rating in related tasks. However, there is no way to reason that this worker can also perform English-to-Turkish translation tasks. Even worse is that workers who pick and complete easy tasks would probably have higher accuracy rates than the ones who choose to perform tasks that require time and expertise [18].

Requesters can utilize acceptance rates by assigning a low limit of them to assignments, which allow only workers with acceptance rates above-limit to accept assignments.

## **2.3 Related Work**

In response to the challenges of crowdsourcing a number of attempts have been made by the researchers. In the following, the studies related to this work are listed, described and examined in detail. The studies, which try to address the challenges of crowdsourcing and proposing new way of creating crowdsourcing workflows, are considered as related only. Otherwise, the related work list would be too long, since many views Wikipedia and Linux as crowdsourcing systems.

### 2.3.1 Crowdsourcing Platforms

The following is a list of crowdsourcing studies that propose a new platform to create crowdsourcing services. Each study is discussed extensively and compared to other studies in Table 2.1. The comparison of related works is performed over the building blocks of a crowdsourcing platform that is given in the previous section. The following set of features is chosen to assess related works:

- concept: The concept or paradigm that platform is based on. This basically defines the process of workflow design.
- programming: Requirement of programming for requester to design a workflow.
- user interface: Availability of a user interface to requester to design a workflow.
- human resources: Availability of human resources in task assignment.
- software resources: Availability of software resources in task assignment.
- heterogeneity: Whether heterogeneity of crowd is available to requester or not in task assignments.
- quality control: Existence of a quality control mechanism.
- progress: Availability of a view that displays the progress of a workflow (useful for quality control).

#### 2.3.1.1 Jabberwocky

Jabberwocky [14] is a social computing stack consists of three main components: Dormouse, ManReduce, and Dog. Dormouse is created to enable cross-platform human and machine resource management. It acts like a "virtual machine" layer in the computing stack, consisting of low-level software libraries that interact with people and computers. ManReduce is a parallel programming framework



for human and machine computation working on top of Dormouse. ManReduce is inspired by MapReduce [30] that is basically mapping problem into a set of small chunks of work, and then reducing them on an output that aggregate the responses or solutions. Dog is a high-level programming language on top of ManReduce. The language is formed by a small set of primitives for requesting computational work from people of machines.

Jabberwocky is limited in several ways. The computing stack is a command-line tool supported by restricted built-in libraries and can only be run on Dormouse server. The high-level language, Dog, seems simple, clear and expressive, but still there is still a learning curve based on the fact that not all crowdsourcing requesters are developers. This is the issue for the command-line tool as well.

Another important limitation is lack of progress idiom within the system. Jabberwocky receives script definitions (deployment to Dormouse server), a process starts. Once all of them are completed, the output is written to a destination file and process terminates. In that respect, there is no way that end user can observe the current state of problem-solving.

Even if this work aims at solving general-purpose and real-world problems, it is mentioned that real-time and single-worker sequential applications are not well-suited. Despite MapReduce paradigm is simple and many social computing problems fit naturally to this paradigm, it is obvious that only some set of problems are appropriate to be solved using such a paradigm or system, which is another limitation.

However, Jabberwocky’s notion of real identity and social structure, which would allow end users to define person-level constraints, is noteworthy.

### **2.3.1.2 WeFlow**

WeFlow [15] is a collaborative application specification, application generator and execution engine proposed in a master thesis [31]. A framework is introduced to create and run collaboration-based applications. That framework allows end users

to decompose problem into tasks, describe the computation resources, define the control and data flow.

A task consists of input(s), instruction describing the expected action from user and output(s). Task can be a type of basic (the most simple, atomic work definition), conditional (execution based on a condition), repetition (recurring execution based on a conditional), doall (groups of tasks executed in parallel) or collective (multiple times execution). The framework is definitely restricted by these predefined task definitions.

Besides task specification, end user should designate the control flow and data flow. Control flow determines the order of tasks to-be-executed. On the other hand, data flow is defined to map data between and/or within tasks. All these specifications are done via XML. Despite the fact that the framework is mentioned to require no programming skills and XML is widely used format for representing arbitrary data structures, there is a learning process of WeFlow specifications not to mention the lack of development and debugging environments.

WeFlow differentiates human workers from each other depending on a role definition. Participants are linked to some role. Likewise a task is associated with a role. Although the whole role definition would allow end user handle access control by describing permission-like roles, it does not discriminate human performers based on their characteristics such as age, gender, interest, expertise etc.

### **2.3.1.3 TurKit**

TurKit [16] is a toolkit for deploying iterative tasks to MTurk. Toolkit is based on a model that concentrate on iterative work in which series of individuals work on tasks that are previously completed by others. Although creators of TurKit apply several nontrivial problems (image description, brainstorming, writing tasks, sorting etc.) to the iterative model, this work majorly restricted by the iterative paradigm that toolkit operates on. The complex and sophisticated problems that are expected to be addressed by crowdsourcing systems do not often correspond

to iterative model.

TurKit API is defined to help writing iterative MTurk tasks. However, TurKit expects end user to be a programmer and create HTML pages for tasks, and write Javascript files using API. In fact, end user is responsible for testing and making sure that the pages with TurKit functionality interact properly with MTurk. This just reveals another major limitation of TurKit on its usefulness.

However, the toolkit have some notion of fault tolerance preventing wasted money or time due to bugs and system crashes. Toolkit stores information about the trace of a program's execution. This trace is used whenever program crashes and to put the program back into its previous state. Thus, toolkit does not re-execute the whole program, but the sections where they are left unfinished.

#### **2.3.1.4 CrowdLang**

CrowdLang [17] is a general-purpose framework and a concept of an executable, model-based programming language for workflow definition. The framework is developed based on the assumption that a complex problem is characterized by defining the problem, decomposing the problem into subproblems, planning subproblems, executing the plan and aggregating the solutions of subproblems.

CrowdLang programming language is based on a small set of operators: Divide-and-Conquer, Aggregate, Multiply, Merge, Router and Reduce (functionality of operators can be understood by their names). These operators are combined together to solve complex problems by routing, distributing and task decomposition. In addition, different types of genes are defined to address various participation patterns.

This work introduces a good and novel concept for general-purpose crowdsourcing that is suitable to most problems. The framework is not bounded to some other programming paradigm or limited to only one aspect of crowdsourcing. It rather supports complex coordination mechanisms. In this respect, translation, which is a sophisticated and difficult problem for crowdsourcing, is attempted

in [9] and the results (translation of 15 different articles in less an hour) are promising.

#### **2.3.1.5 AutoMan**

AutoMan [18] is described to be a fully automatic crowd programming system. It is a programming system integrating human and computer computation. On top of AutoMan system, a domain specific programming language is defined. It is implemented as embedded domain-specific language for Scala.

The system’s whole crowdsourcing concept is based on Question objects. Question can be type of radio-button, checkbox and restricted free-text. The human computation aspect of system only corresponds to the answers given by individuals. Further, system provides no mechanism to design complex problems through a set of subproblems or tasks. This makes system no advantageous to the existing systems depending on simple tasks.

In fact, end users are supposed to write programs in AutoMan DSL and provide them to the system. Considering the examples in the paper, the learning curve for this language can be expected to be high, because it requires knowledge of Scala language (comparing with Dog introduced in [14]).

However, scheduling, pricing and quality control mechanisms are significant components of AutoMan. The runtime system has a scheduling component that periodically checks the current situation of the results, and reprices and restarts human tasks as necessary. By making this, system tries to achieve the predefined confidence level (by end user) while staying under budget.

#### **2.3.1.6 Turkomatic**

Turkomatic [19] is a tool that recruits workers for planning and solving complex tasks. The system executes price-divide-solve approach by asking workers to divide complex steps into simpler ones until they are at a simple level, then

to solve them. The approach simply uses divide-and-conquer strategy, but the division is done by the crowd different from other crowdsourcing systems.

The system has a set of pre-structured task templates. End users can produce workflows by combining templates together without implementing any software or designing intermediate tasks. The price-divide-solve approach is expected to produce a directed, acyclic task graph in which nodes represent subtasks and links describe task dependencies. The user interface visualizes the task graph and enables endusers delete or modify them in real-time.

The system is significant by demonstrating complex problems through acyclic task graphs. The graphs are not just shown to the endusers, but also implemented in a way that enables real-time modification. However, initial workflow design is left to workers by a single instruction (a few sentences) telling what the enduser achieve in the end. This approach is limited by the efficiency of instruction and understanding of workers, and reportedly not really successful. It is mentioned that for complex work manual intervention and editing of workflow is effective.

Currently the system's crowd planned workflows are not guaranteed to be context free. This restricts the set of problems that the system can tackle, but still this study presents promising results by demonstrating and managing complex problems through acyclic graphs.

#### **2.3.1.7 CrowdForge**

CrowdForge [7] is a general-purpose framework and toolkit for accomplishing complex and interdependent tasks using micro-task markets. The framework is built on MapReduce [30] approach, which first breaks down a complex problem into a sequence of subproblems, and combines the results later. A similar approach is taken for the design of ManReduce in Jabberwocky stack [14].

Although CrowdForge approach is designed to fulfill complex problems, it is still limited to the capabilities of MapReduce approach. Some problems may not be addressed by this approach such as the case for tasks that cannot be really

decomposed. Another case would be when subtasks are not independent, but the state or result of one is important to complete the other. All these exemplifies the limitations of the approach taken from distributed computing field and expected to fit well in crowdsourcing.

Additionally system has no support for iteration or recursion. It requires the end user (task designer) to specify each stage (partition, map, reduce) in the task flow.

#### **2.3.1.8 CrowdWeaver**

CrowdWeaver [20] is a system to visually manage crowd work. This system comes forward among various other works with it's visual representation abilities. High-level representation of a workflow makes it easier for end user to grasp, design and develop crowdsourcing programs.

The system has a predefined set of task templates. Workflows are created by creating tasks and linking them with each other in various ways. Branching and combining multiple data flows are supported and that makes design of complex problems possible.

Besides visualization, CrowdWeaver has a notion of tracking and notification. The task progress component monitors the current state and depicts the current state via graphs. Along with notification component, users can be notified about the progress based on the predefined conditions (specified by users). Further, it is possible to stop a task and make changes on existing branch in real-time.

Despite CrowdWeaver demonstrates a system that can be easily utilized by users with no programming background, the system is currently limited by predefined set of templates. In fact, it's visual abilities are restricted by only showing the workflow, but not enabling users make changes on it directly.

Feature	Jabberwocky	WeFlow	TurKit	CrowdLang	AutoMan	Turkomatic	CrowdForge	CrowdWeaver
Concept	MapReduce	Workflow	Iterative	N/A	N/A	Divide&Conquer	MapReduce	Data Flow
Programming	●	○	●	●	●	○	○	○
User Interface	○	○	○	○	○	●	●	●
Human Resources	●	●	●	●	●	●	●	●
Software Resources	●	●	○	●	●	○	○	●
Heterogeneity	●	●	○	○	○	○	○	○
Quality Control	○	○	●	○	●	○	○	●
Progress	○	○	○	○	○	○	○	●

●: Yes

◐: Partially

○: No

Table 2.1: Comparison of existing crowdsourcing platforms.

## **2.3.2 Other Studies**

### **2.3.2.1 Soylent**

Soylent [21] is a word processing interface enabling Microsoft Word users to employ MTurk workers to shorten, proofread and edit parts of their documents on demand. The system is developed with respect to the Find-Fix-Verify pattern that is introduced in the same study.

The pattern approaches complex tasks (focused on text editing) by splitting them into a series of generation and review stages. Rather than asking a single worker to read and edit an entire section, the pattern first recruits a set of workers to find areas of improvements. Then, another set of workers review the candidate areas and filter out incorrect ones. Finally, in the verify stage workers perform quality control on previous submissions. Throughout the process the pattern utilizes independent agreement and voting.

Both the system and the pattern concentrates on a small set of use cases of crowdsourcing. The system can only be employed by Microsoft Word users for editing text. The pattern is limited to the problems where decomposition into subproblems is possible.

### **2.3.2.2 Qurk**

Qurk [10, 22] is query processing system that automatically translates queries into tasks to-be-completed by humans. The system has domain-specific language to express tasks. A UDF-like approach is taken by integrating SQL with MTurk expressions.

The approach is concentrated on a MTurk-aware database system rather than a general-purpose crowdsourcing mechanism. Thus, the set of tasks that can be introduced via the system is limited by associated database concepts such as joining, sorting etc. Although generative tasks for which workers give unconstrained input are made available, still processing is done on input data obtained from the



underlying database system. However, this study provides an important example by employing people in a database system for managing various queries.

#### **2.3.2.3 Mobi**

Mobi [24] is a system to crowdsource itinerary plans. The system illustrates a use case of crowdware paradigm. The paradigm focuses on tasks with global requirements and provides a single workspace in which a crowd of individuals contribute opportunistically to the global solution based on their knowledge and expertise. Itinerary planning is taken as a case study and implemented in Mobi system through a single interface.

In the system and paradigm, the problem definition is limited to the tasks with global constraints. This is a clear example of a study that concentrates on only one aspect of crowdsourcing. However, the study provides insights on the effectiveness of using unified solution context for workers or directing crowd's submissions through a structured guide or benefits of iterative refinements.

#### **2.3.2.4 CrowdSpace**

CrowdSpace [23] is a system that supports the evaluation of complex crowd work by combining information about worker results through visualization. This system focuses on exploration and assessment of worker performance and behavior rather than providing ways to manage complex workflows.

The system presents a unified user interface with four components: a scatter plot of aggregate behavioral features (time spent on the task, number of keys pressed while processing the task etc.), distribution of each behavioral features, traces of worker/output pairs and overall worker behavior based on their answers.

Low quality work is common in crowdsourcing. However, this system provides great insights by combining worker behavior with their submissions, and enabling end users to identify the behavior of workers who have good or bad outputs.

Although this approach is limited by several aspects such as being only applicable to pages in which Javascript can be inserted or assuming that worker does all the processing on the page etc, the quality control approach taken in the study can be used to better understand and address the nature of the crowd.

#### **2.3.2.5 Human Architecture**

Human Architecture [25] is an adaptation mechanism to the changing requirements of the various type collaborations. Unlike other studies, this work focuses on collaboration problem from the architectural perspective. An architecture description language is proposed to describe collaboration dynamics. Considering software architecture human components and collaboration connectors are introduced to demonstrate coordination dependencies.

However, the proposed human architecture approach is too architecture focused and highly complex. The collaboration of individuals is narrowed down to component and connectors, but in the context of complex problems collaboration is dynamically changing and highly interactive due to the data items that are output from one and input to the other.

# Chapter 3

## Platform

*Crowdy* is an extensible, general-purpose crowdsourcing platform to solve complex problems. The platform is developed over the fundamentals of stream processing paradigm for which a series of operations are applied to the continuous stream of data elements via operators.

*Crowdy* is an operator-centric platform. Using this platform, a requester with no requirement of a programming background can quickly translate a complex problem into a crowdsourcing application by simply selecting operators and connecting these operators together. As a result of *Crowdy*'s focus on operators, requesters can design applications by selecting right set of building blocks that are necessary to solve their problem, and customizing these blocks particular to the computation to-be-conducted.

*Crowdy* embodies several features:

- A standard toolkit of operators that can both human and software resources (human or software) to accomplish various tasks
- Configuration support to control and coordinate resource utilization
- Customizable collaborations over parameterization
- Application runtime interface

### 3.1 Architecture

*Crowdy* platform is implemented as a REST [32] architecture shown in Figure 3.1. Applications are developed, configured and validated on the client-side. These applications are submitted to server-side via an application programming interface (API). Applications are executed, and results are saved on the server-side.

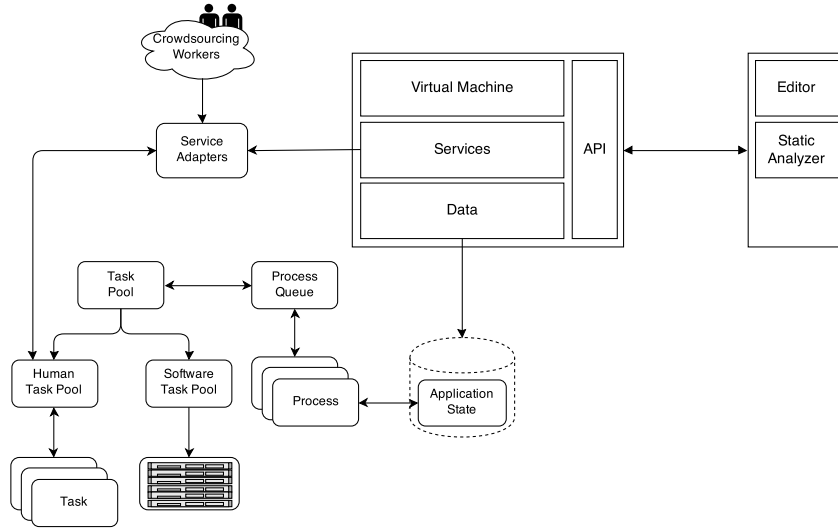


Figure 3.1: Architecture of the platform.

The client-side provides an editor to easily create *Crowdy* applications. Requesters can drag and drop various types of operators and connect them together to create a data flow, which conforms to a *Crowdy* application. These applications are further analyzed and validated before submission. A valid application is submitted to the virtual machine on the server-side.

The server-side consists of three layers: virtual machine, services and data. In addition, an API is provided to submit new applications to the virtual machine and retrieve details of existing ones. Virtual machine sits on top of existing crowdsourcing services and provides service-independent environment for human computations. Once a new application is submitted, virtual machine first allocates required computation resources, then runs the application and finally saves the results. Services layer is formed by different crowdsourcing services such as Amazon’s Mechanical Turk or microworkers. The data later is basically where

application state is kept.

In the remainder of this chapter, the fundamental concepts of *Crowdy* are explained in more detail and features are explored as we look into various aspects of application development using *Crowdy* platform.

## 3.2 Concepts

*Crowdy* applications are developed to solve complex and sophisticated problems that require both human intelligence and computing power. A typical application contains three main high-level components: data ingest, processing and data egress.

Let's consider a minimal "Hello World" application that has these three components in total, connected in a simple pipeline topology. Figure 3.2 demonstrates the application in the form of a flow graph. On the ingest side, there is a *source operator*, which acts like a data generator. Source operator produces data tuples that are processed down the flow by the *processing operator*. Finally, the *sink operator* simply converts the tuples in such a form (text file, email etc) that can be easily interpreted by requesters.

The application flow graph is specified as a series of operator instances and connections (data flows) that link them. A data flow basically transfers data tuples produced by an operator to another. One or more data segments can be assembled in a data tuple via output specification of an operator instance (see Section 3.3). In addition, several options can be specified to configure an operator instance. These include parameters, operator-specific rules, which are studied in the rest of this chapter.



Figure 3.2: A sample, minimal *Crowdy* application.

In a more realistic application, one or more source operators can be employed to produce various data tuples that differ in both size and specification. Similarly, the application would have one or more processing operators along with other types of operators organized in a way that is significantly more complex than this example.

### 3.2.1 Operator

An operator is the basic building block of an application. Operator has a type that is specified at the time of creation. This type determines configuration respectively. Also a unique ID is assigned to an operator. In addition, operator has optional name and description fields that can be used for bookkeeping purposes.

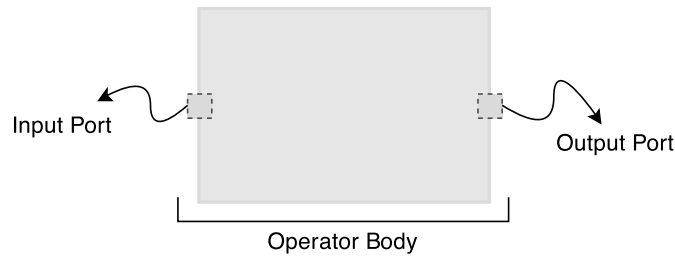


Figure 3.3: Base operator representation.

Operator may have an input port or output port or both corresponding to its type definition. Figure 3.3 demonstrates a base operator, which consists of a body and ports. Although it is not shown here, each operator presentation has a specific icon on their body associated with its type.

An operator may output tuples to any number of operators, but it can only receive tuples from one operator unless it is a union operator (see Section 3.2.1.5). Union operator can receive tuples from multiple operators and aggregate them, if these tuples have the same specification. Therefore, consistency of incoming flow specification for each operator type is ensured. This is a significant feature to guarantee operators functionality, because an operator (excluding source operators) uses and operates specifically on the information from incoming flow.

*Crowdy* provides a set of built-in operators that can be used to build applications. In general, these operators perform common tasks associated with data generation, processing and outputting.

Operators are generally cross-domain to allow general-purpose computation possible. They are grouped under six main categories: *source*, *sink*, *processing*, *relational*, *utility*, *adapter*.

### 3.2.1.1 Source operators

The set of source operators generates data tuples. These operators do not have an input port, but have an output port, which produces data tuples. Figure 3.4 represents a source operator.

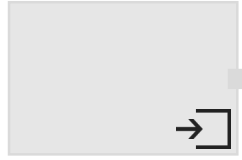


Figure 3.4: Source operator representation.

Source operators together with processing operators are the ones that can be used to specify data flow coming out of an operator. Output specification is an action to identify the data tuple with a series of segments. Other operator types cannot make changes on output specification, but can manipulate the flow by dropping or copying data tuples.

**human.** The *human source operator* is a stateless operator used to produce new data tuples via human workers. Existing crowdsourcing services such as MTurk is used to produce new tuples. A new data tuple is produced per successfully completed human intelligent task. These tasks are automatically created and posted with respect to the specified parameters of the operator.

Human source operator has the following parameters:

- **number of copies:** The maximum number of data tuples can be generated

by the operator. It's value ranges from 1 to 1000.

- **max allotted time:** The maximum time in seconds given to a human worker to solve and submit the task. It's value ranges from 10 to 300.
- **payment:** The payment in cents to be given to the human worker in case of successful task completion. It's value ranges from 5 to 500.
- **instructions:** The detailed information for human workers on how to complete the task.
- **question:** The set of sentences asking for specific information from human workers.
- **input list:** The list of **inputs** that will be shown to human worker to fill in. An **input** can be a type of *text*, *number*, *single choice* or *multiple choice*. Each of these types corresponds to an HTML element. Table 3.1 lists types and their details.

A text-input presents an input field where the human worker can enter data. The maximum number of characters that can be entered to the field can be set by the requester. Similarly number-input corresponds to a input field where only numbers can be fed in, and the maximum and minimum value for the field are set by the requester. The other two input types conform to input fields where the options given by the requester are presented to the human worker as a list. Human worker is expected to select only one and one or more options for single choice and multiple choice types respectively.

In fact, each **input** conforms to a segment in the data tuple.

Table 3.1: List of inputs and options

type	parameters	HTML element
text	max number of characters	input [type=text]
number	min value, max value	input [text=number]
single choice	options	input [type=radio]
multiple choice	options	input [type=checkbox]

**manual.** The *manual source operator* is a stateless operator to produce new data tuples.



Manual source operator has the following parameters:

- **manual entry**: The manual text to be parsed and used to produce new tuples.
- **delimiter**: Delimiter to determine segments in a tuple. This can have one of the following values: none (``'``), white space (``' ``), tab (``'\t'``), comma (``', '``), column (``': '``).

Manual source operator uses manually entered text to create new tuples. Operator retrieves the manual text, parses it line by line and then applies the delimiter. Therefore, each line constitutes a data tuple, and delimiter is used to create segments in a tuple.

For example, if **delimiter** is chosen to be white space and the following is entered to **manual entry**,

```

Lorem ipsum
Consectetur adipiscing
Phasellus vehicula
```

the following data tuples will be generated:

```
[
{"segment_1": "Lorem", "segment_2": "ipsum"},
{"segment_1": "Consectetur", "segment_2": "adipiscing"},
{"segment_1": "Phasellus", "segment_2": "vehicula"}
]
```

It is possible to have such a manual entry that ends up in different number of segments for different lines. To prevent this happening manual source operator uses the first line to generate output specification. If more segments are generated in the following lines, they are discarded. If there is not enough segments in another line, then the corresponding segments are emptied and then outputted.

### 3.2.1.2 Sink operators

The set of sink operators is where data tuples are serialized and converted into the formats that can be used by requesters with ease. These operators have one input port, but no output port. Figure 3.5 demonstrates a sink operator.



Figure 3.5: Sink operator representation.

**email.** The *email sink operator* is a stateless operator to convert data tuples into a text format and email them to requesters. **email** parameter specifies the requester's email address.

**file.** The *file sink operator* is a stateless operator to serialize the data tuples into a file. Operator has one parameter **filename** that is used to specify the name of file in which tuples will be written.

### 3.2.1.3 Processing operators

The set of processing operators provides data tuple processing via human workers. These operators have both input and output ports. Figure 3.6 shows a processing operator.

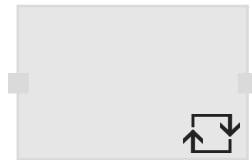


Figure 3.6: Processing operator representation.

As mentioned before, processing operators can manipulate the data flow specification in addition to source operators. These operators can change the existing flow specification by adding, deleting or editing data segments.

**human.** The **human processing operator** is almost same as human source operator. The difference is human processing operator has an input port. That means there is a flow of data tuples coming to operator. These incoming tuples are made available to requesters via their specification.

The parameters of human processing operator is no different than the parameters of human source operator. Additionally processing operator has available segment list, which provides placeholders for the segments of an incoming data tuple. Requesters can place these placeholders in **instructions**, **question** and **input list** (applicable to single choice and multiple choice inputs).

At runtime when a new tuple arrives, each placeholder in parameters is replaced with the corresponding value of incoming tuple's segment. This enables dynamically created human tasks. Therefore, requesters can create an information flow from one operator to another.

#### 3.2.1.4 Relational operators

The set of relational operators enables fundamental manipulation operations on the flow of data tuples. Each relational operator implements a specific functionality providing continuous and non-blocking processing on tuples. Therefore, these operators have both input and output port.

**selection.** The *selection operator* is a stateless operator used to filter tuples. A typical selection operator is shown in Figure 3.7.



Figure 3.7: Selection operator representation.

On a per-tuple basis a boolean predicate is evaluated and a decision is made as to whether to filter the corresponding tuple or not. Boolean predicates are specified by requesters as part of operator parameterization. These predicates,

which are identified in **rules**, are the only members of parameters.

A selection operator has zero or more rules to filter data tuples. When there is no rule specified, then no filtering will be done, and all data tuple will be passed down to data flow. Otherwise, each rule is evaluated on an incoming data tuple. Whenever a rule is evaluated to be true, then corresponding action is carried out that is either filter in or out the tuple. If no rule is evaluated to be true on a data tuple, then tuple is still passed to the next operator. Rules share the following predefined format:

*Filter (in/out) when boolean-predicate*

Similarly **boolean-predicate** has the following format

**segment-name** (*equals/not equals/contains*) **query**

where **segment-name** is one of the segments of incoming tuple, and **query** is a free-text to be filled by requester.

**sort.** The *sort operator* is a stateful and windowed operator used to first group tuples and then sort them based on the specified data segment and order. Figure 3.8 illustrates the operator.



Figure 3.8: Sort operator representation.

Sorting is performed and results are produced (outputted one by one) every time window trigger policy fires. The policy is basically triggered when the number of data tuples reaches **window size**, which is specified as a parameter. **window size** can have a value between 1 and 100.

Similar to selection operator, sort operator implements a set of rules that simply identifies the segment and the order to be used for sorting. If no rule is specified, then tuples are sorted in the ascending order with respect to the first

segment in the incoming data tuples. If there are more than one rule is given, then these rules are applied in the order they are specified by requester.

Sorting rules share the following predefined format:

*Sort using **segment-name** in (ascending/descending) order*

where **segment-name** is one of the segments of incoming tuple.

### 3.2.1.5 Utility operators

The set of utility operators provides flow management functions. These operators handle operations such as separating a flow into multiple flows or joining multiple flows into a single one.

**enrich.** The *enrich operator* is a stateless operator used to enrich data flow by replaying incoming data tuples. Figure 3.9 shows an example view of the operator.



Figure 3.9: Enrich operator representation.

Enrich operator has one parameter called **number of copies**. This parameter determines how many copies will be produced for each incoming tuple. It's value ranges from 1 to 10.

**split.** The *split operator* is a stateless operator used to divide an inbound flow into multiple flows. This operator has one incoming flow and can have one or more outgoing flow. Figure 3.10 provides the presentation of a split operator.

The boolean predicates specified in rules are evaluated whenever a new tuple arrives. Then, a decision is made to where to send the tuple. Similar to relational operators, rules are specified by requesters as part of operator parameterization.

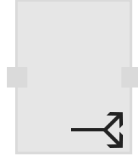


Figure 3.10: Split operator representation.

A split operator has zero or more rules. When there is no rule specified, then tuples are passed to every operator connected down the flow. Otherwise, each and every rule is evaluated on an incoming data tuple. Whenever rule's predicate is evaluated to be true, then tuple is passed to the corresponding operator. If no predicate turns out to be true for a tuple, then it is dropped.

Rules share the following predefined format:

*Send to next-operator when boolean-predicate*

in which **next-operator** is the connected operators down the flow, and **boolean-predicate** has the same definition given for selection operator.

**union.** The *union operator* is a stateless operator used to join two or more data flows into one. Different than other operators, this operator can receive more than one flow. These flows are combined by the operator and outputted as if they are a single flow. Figure 3.11 demonstrates the operator.

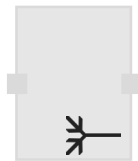


Figure 3.11: Union operator representation.

Union operator requires incoming flows to have the same specification. Otherwise, union operation will basically fail.

### 3.2.2 Flow

Flow is the connection between operators used to move information from one operator to another. A flow is formed by one or more data segments. A segment has an identifier and corresponding value. In that sense, flow is like a hash table where segment identifiers are keys, and values stored in segments are values as demonstrated in the following:

```
{  
"segment-1": "value-1",  
"segment-2": "value-2",  
...  
"segment-N": "value-N",  
}
```

The details of a flow can be specified by either source or processing operators. It is immutable in other operators. Specification can be achieved via output specification section in operator configuration. Requesters can create new segments, edit or delete existing ones. Source operators are useful to specify flow specification in the early stages of workflow design and generate tuples using that specification. The flow specification is crucial, since it determines the specifics of the information that is carried over in the application. Still this specification can be redefined in processing operators.

## 3.3 Application Development

The application development using *Crowdy* can be achieved by simply creating new operators, configuring them and connecting them together. However, there are certain predefined rules that should be satisfied to develop a valid application and execute it.

An application is formed by a workflow. Although it is not suggested, application can have multiple workflows. A workflow is basically a set of operators

connected together. An operator can be connected to another operator by attaching a flow from the output port of an operator to the input port of the other operator. Workflow must contain at least one source and one sink operator. There may be operators that are not connected to any other operator within the application. Operators with no connections are discarded and they do not cause validation failures.

A valid workflow means that each operator employed in that workflow is valid. Validity of an operator depends on its type and therefore its configuration. In the following each operator group is examined in terms of validation.

#### 3.3.0.1 Source operators

A source operator has one output port, which outputs the data tuples with given specification. Therefore, source operator is required to have an output specification that has at least one data segment in it.

**human.** Parameters of a human operator should be valid. That means `number of copies` is an integer between 1 and 1000, `max allotted time` is an integer between 10 and 300, `payment` is an integer between 5 and 500. Although `instructions` can be left empty by the requester, `question` cannot be empty. There must be at least one `input` defined in the `input list` to satisfy that source operators have an output specification with at least one segment defined.

**manual.** `manual entry` and `delimiter` should not be empty. Nonempty `manual entry` and `delimiter` correspond to an output specification with at least one segment in it.

#### 3.3.0.2 Sink operators

A sink operator has one input port, which receives data tuples.

**email.** Operator must have a valid email address given in `email` parameter.



**file.** Operator must have a valid file name given in `filename` parameter.

### 3.3.0.3 Processing operators

A processing operator has one input and one output port. It receives data tuples, processes them and outputs them to next operator down to flow. Similar to source operators, processing operators must have an output specification that has at least one data segment in it. Since human operator is the only processing operator, the validity rules for this operator is same as human operator of type source operator.

### 3.3.0.4 Relational operators

A relational operator has one input and one output port. It is useful to manipulate flow. Operators of this type has rules to determine manipulation.

**selection.** Each rule in the rule list must be valid, which means a rule with valid `segment-name`.

**sort.** Operator must have a valid `window size`. This value can be an integer between 1 and 100. Additionally each rule in the rule list must be valid by being formed by a valid `segment-name`.

### 3.3.0.5 Utility operators

A relational operator has one input and one output port. This type of operators is useful for flow management.

**enrich.** Operator must have a valid `number of copies`. This value can be an integer between 1 and 10.

**split.** Split operator must have an incoming flow to-be-split and at least one outgoing flow. Split operation is defined by rules. Each rule in the rule list

must be valid by specifying a valid **next-operator**.

**union.** Union operator must have at least one incoming flow to-be-merged and one outgoing flow. Each incoming flow must have the same specification, which is ensured by application editor by preventing connecting flows with different specification to the same union operator.

# Chapter 4

## Motivating Examples

- Translation

TODO

# Chapter 5

## Discussion

- Extending operator set - More quality control - Improvements in human task assignment

TODO

# Chapter 6

## Conclusion

TODO

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