THE PRINT SHOP COMPANION



<u> 2017-06-17</u>



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2017-06-17 A 4am crack

-----The Print Shop Companion-----

Name: The Print Shop Companion Version: 1.2 Genre: graphics

Year: 1985

OS: custom

Credits: Roland Gustafsson Publisher: Broderbund Software

Platform: Apple **JC**+ or later (64K)

Media: single-sided 5.25-inch floppy

Broderbund Software PRESENTS...

Previous cracks: none (of this version)

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In Which	Chapter 0 Various Automated Tools Fail In Interesting Ways

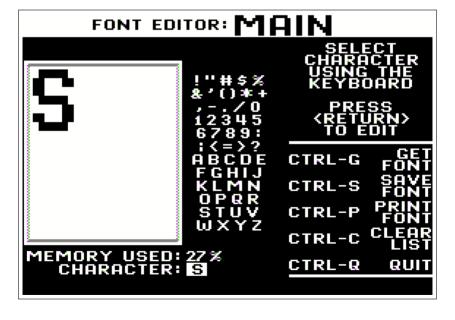
Locksmith Fast Disk Backup unable to read track \$22 copy displays title screen then hangs on main menu

read error on last pass

COPYA

read error on track \$22 copy displays title screen then hangs on main menu

EDD 4 bit copy (no sync, no count)



```
Copy JC+ nibble editor
  track $22 is quite unusual, with
  repeated sequences of $D4 $D5 $DE $D4
  and other nibbles, with and without
  timing bits
                    --0--
COPY JE PLUS BIT COPY PROGRAM 8.4
(C) 1982-9 CENTRAL POINT SOFTWARE, INC.
TRACK: 22 START: 1800 LENGTH: 3DFF
1828: 9A 9A BB DA 97 BE BD D5+ VIEW
1830: FF+CB+A6+AE F6 A5 D7 DF
1838: AD+FF+A5+D4+D5 DE D4 BB
1840: FA F5 AA+FF+BD A9 AB BD
1848: A9 EF EF D6 FF+FF+FF+D4
1850: D5 DE D4 BA FE F5 AA AA+
1858: FF+D4 D5 DE D4 BA FF F5
1860: AA AA AA AA+D4+D5 DE D4
1868: BB FE F5 AA AA AA FF+A9
  A TO ANALYZE DATA ESC TO QUIT
  ? FOR HELP SCREEN / CHANGE PARMS
  Q FOR NEXT TRACK SPACE TO RE-READ
```

catalog that lists "FONT." and "BORD." files (presumably fonts and borders loadable from the program) no way to read track \$22 (no sectors) Why didn't COPYA work? non-standard structure on track \$22 Why didn't Locksmith FDB / EDD work? presumably there is some runtime protection check that triggers after displaying the main menu, which checks the structure of track \$22 Next steps: 1. find runtime protection check

track 0 has a custom bootloader track \$11 has a DOS 3.3-style disk

Disk Fixer

2. disable it 3. declare victory (*)

Chapter 1 In Which We Will Not Be

Going To The Gym Any Time Soon

they need to turn on the drive motor by accessing a specific address in the \$C0xx range. For slot 6, it's \$C0E9, but to allow disks to boot from any slot, developers usually use code like this:

LDX <slot number x 16>
LDA \$C089,X

There's nothing that says where the slot number has to be, although the disk controller ROM routine uses zero page \$2B and lots of disks just reuse that. There's also nothing that says you have to use the X-register as the

Since my copy goes down a different code path than the original, I'm

guessing there is a runtime protection check somewhere. One thing that all protection checks have in common is

Also, since developers don't actually want people finding their protection-related code, they may try to encrypt it or obfuscate it to prevent people from finding it. But eventually, the code must run, and it must run on my machine, and I have the final say on what my machine

accumulator as the load register. But most RWTS code does, out of convention I suppose (or possibly fear of messing up such low-level code in subtle ways).

But sometimes you get lucky.

does or does not do.

index, or that you must use the

```
Turning to my trusty Disk Fixer sector
editor, I search the non-working copy
for "BD 89 CO", which is the opcode
sequence for "LDA $C089,X".
□Disk Fixer]
 Ē"F"ind∃
   ["H"ex]
     E"BD 89 C0"]
                --u--
        --- DISK SEARCH ---
$0D/$0B-$D8
            PRESS ERETURNI
The first match on track 3 might be my
jackpot. The protection routine appears
to start at offset $19:
               ----
T03,S06
       --- DISASSEMBLY MODE --
; clear page of memory at $BB00
0019:A0 00
                  LDY #$00
001B:A9 FF
                 LDA #≸FF
                 STA $BB00,Y
001D:99 00 BB
                 ĪNY
BNE $001D
0020:C8
0021:D0 FA
```

```
; turn on drive motor manually (always:
; suspicious)
0023:AE F8 05
0026:BD 89 C0
                    LDX $05F8
                    LDA $C089,X
; look for that nibble sequence I saw
; in the nibble editor: $D4 $D5 $DE $D4
0029:BD 8C C0
002C:10 FB
002E:C9 D4
0030:D0 F1
                    LDA $C08C,X
                    BPL $0029
                    CMP #$D4
BNE $0023
; This subroutine just does the LDA/BPL
; loop to get the next nibble. But it
; tells me that this code is probably
; executing from $B6xx in memory.
0032:20 E5 B6
                    JSR $B6E5
0035:C9 D5
                    CMP #$D5
0037:D0 F5
                    BNE $002E
                    JSR $B6E
CMP #$DE
0039:20 E5 B6
                          $B6E5
003C:C9 DE
003E:D0 F5
                   BNE $0035
0040:20 E5 B6
                   JSR $B6E5
                    CMP #$D4
BNE $003C
0043:C9 D4
                          $003C
0045:D0 F5
0047:EA
                    NOP
; get a 4-and-4 encoded value from the
; next two nibbles
0048 BD 8C C0
                    LDA
                          $C08C,X
                    BPL
004B:10 FB
                           $0048
004D:2A
                    ROL
                 STA
                         $26
004E:85 26
0050:BD 8C
          CØ.
                   LDA $C08C,X
0053:10 FB
                    BPL
                          $0050
0055:25 26
                    AND
                           $26
```

```
; continue below
0058:4C B5 B6
                    JMP $B6B5
; now check for an epilogue of sorts,
; the two-nibble sequence $F5 $AA
00B5:20 E5 B6
                    JSR $B6E5
                   CMP #$F5
00B8:C9 F5
                   BNE $00DC
00BA:D0 20
                   JSR $B6E5
CMP #$AA
BNE $00DC
00BC:20 E5 B6
00BF:C9 AA
00C1:D0 19
; Using the 4-and-4 encoded value as
; a lookup into the page we initialized
; earlier, see if we've seen this value
; already. (We initialized all 256
; addresses with #$FF.)
00C3:B9 00 BB LDA $BB00,Y
; if we've found this value already,
; skip ahead
00C6:10 14
                    BPL $00DC
; otherwise mark this value as "found"
00C8:A9 00 LDA #$00
00CA:99 00 BB STA $BB00,Y
```

```
; loop through the array of "found"
; values and count how many unique
; values we've seen
00CD:AA
                         TAX
00CE:A8
                         TAY
00CF:B9 00 BB
                        LDA $BB00,Y
                        вмі ≴аа́∩5
00D2:30 01
00D4:E8
                        INX
00D5:C8
                         INY
00D6:D0 F7
                         BNE $00CF
; have we seen #$A0 unique values yet?
00D8:E0 A0
                         CPX #$A0
; yes, we're done
00ĎA:B0 03
                         BCS $00DF
; no, loop back and read some more
00DC:4C 23 B6
                        JMP $B623
; turn off drive motor and exit
00DF:AE F8 05 LDX $05F8
                        LDX $05F8
                       LDA $C088,X
LDA $C08C,X
BPL $00E5
00E2:BD 88 C0
00E5:BD 8C C0
00E8:10 FB
00EA:60
                         RTS
Well that certainly explains why my
Locksmith Fast Disk Backup copy hung
it's looking for a nibble sequence
($D4 $D5 $DĒ $D4) that doesn't exist
because it didn't read anything from
track $22. But why can't EDD copy this?
What's so special about track $22?
```

Chapter 2 In Which We See Differences But No Clarity

```
On my non-working bit copy (created by 
EDD 4), I can edit T03,S06 to insert
some useful code before we start the
protection check. That sector is loaded
at $B600, and the code I want looks
like this:
[S6,D1=non-working copy]
       20 EB B6
B619-
                   JSR
                          $B6EB
       EΑ
B61C-
                   NOP
B61D- 99 00 BB
                   STA
                          $BB00,Y
B620- C8
                   INY
B621- D0 FA
                    BNE
                          $B61D
; all new code here --
; set reset vector to jump to monitor.
B6EB- A9 69
                    LDA #$69
B6ED- 8D F2 03
B6F0- A9 FF
                    STA
LDA
                         $03F2
                         #$FF
                    STA $03F3
B6F2- 8D F3 03
B6F5- 49 A5
                   EOR #$A5
B6F7- 8D F4
              ΩЗ.
                    STA $03F4
; reproduce original code from $B619
B6FA- A0 00
                  LDY #$00
B6FC- A9 FF
                    LDA #$FF
B6FE- 60
                    RTS
Now I can press (Ctrl-Reset) and jump
to the monitor while the protection
check is runnina.
*C600G
...protection check runs and hangs...
<Ctrl-Reset>
```

Now I'm in the monitor and have unfettered access to all of memoru. Let's look at the buffer at \$BB00 and see which encoded values are being which aren't. found and *BB00.BBFF BB00-ЙΘ 00 ЙΘ ЙΘ FF FF ЙΘ FF BB08-FF FF ЙΘ FF ЙΘ ЙΘ ЙΘ ЙΘ BB10-ЙΘ ЙΘ FF FF 00 99 ЙΘ ЙΘ BB18-FF FF FF FF 00 FF 00 00 BB20-00 00 FF FF 00 00 00 00 BB28-FF FF 00 00 FF FF 00 00 BB30-FF 00 00 00 00 FF 00 FF BB38-FF FF 00 00 00 00 00 00 BB40-FF 00 00 00 00 FF 00 00 BB48-FF FF FF 00 00 00 FF 00 BB50-FF 00 FF 00 FF FF 00 FF BB58-00 00 00 FF 00 FF 00 00 BB60-FF FF 99 00 00 FF FF FF BB68-FF FF FF FF FF FF 00 00 BB70-FF 00 00 00 FF FF 00 00 BB78-FF FF FF FF FF FF 00 00 BB80-FF FF FF FF FF ЙΘ ЙΘ FF BB88-FF FF FF FF FF 00 00 00 BB90-FF 00 FF 00 FF FF FF FF BB98-00 FF 00 FF FF 00 00 FF BBA0-FF 00 00 00 FF FF ЙΘ FF BBA8-FF FF 00 FF FF 00 00 00 BBB0-FF 00 FF FF 00 FF FF 00 BBB8-FF 00 00 FF FF FF 00 00 BBC0-00 00 FF 00 FF FF 00 FF BBC8-FF FF FF 00 FF 00 00 00 BBD0-FF FF 00 00 FF 00 00 FF BBD8-FF 00 FF FF FF FF FF 00 BBE0-FF FF FF FF 00 FF 00 00 BBE8-FF FF 00 FF 00 00 00 FF BBF0-FF 00 FF FF FF 00 FF 00 BBF8-00 FF 00 FF FF 00 00 00

I can also reuse the counting routine (at \$B6CA) to find out how many values we're findina. ***B6CAL** B6CA-A9 00 LDA #\$00 B6CC-EΑ NOP: B6CD-AΑ TAX B6CE-TAY **A8** B6CF-B9 00 BB ≴BB00,Y LDA. B6D2-30 01 BMI **\$**R605 B6D4-E8 INX C8 B6D5-INY B6D6-D0 F7 BNE \$B6CF ; count of zero values is in X register B6D8- 8A TXA

; print it (as a hex value) and exit

B6D9− 4C DA FD JMP \$FDDA

*B6CAG 7E

The minimum number of unique 4-and-4 encoded values required to pass the protection is #\$A0 (160), but we're only seeing #\$7E (126). So we're way

short.

```
To find out why, I want to know exactly
which values I'm missing. I would
really like the original disk to break
into the monitor after it successfully
completes this protection check, so I
can see what the array at $BB00 looks
like. Of course I'm not going to modify
my original disk (NEVER EVER DO THAT).
But I can add some code on my non-
working copy to wait for a keypress
before starting the protection check,
then swap in the original disk before
continuina.
That code, still on T03,S06, looks like
this:
; wait for keystroke
B6EB- 2C 10 C0
                   BIT $C010
B6EE- AD 00 CO
                   LDA $C000
B6F1- 10 FB
                  BPL $B6EE
B6F3- 2C 10 C0
                   BIT
                        $0010
; original code (from $B619)
                   LDY #$00
B6F6- A0 00
B6F8- A9 FF
                   LDA #$FF
B6FA- 60
                   RTS
```

Why are we missing so many values?

4C 23 B6 B6DC-JMP **\$**B623 B6DF-4C 59 FF \$FF59 JMP Rebooting my non-working copy, it gets as far as the main menu and pauses to wait for a key (at \$B6EB). Ejecting my copy and inserting the original, then pressing a key, it runs the protection check on the original disk and breaks to the monitor with a very satisfying beep. (beep)

And after succeeding, I want to jump to

CPX

BCS

#\$A0

\$B6DF

the monitor:

EO AO

B0 03

B6D8-

B6DA-



Now I can see what the \$BB00 buffer like after a successful looks run on an oriqinal disk: *BB00.BBFF BB00-FF 00 00 00 00 FF 00 00 BB08-00 FF 00 00 00 00 ЙΘ 00 BB10-ЙΘ 00 00 FF ЙΘ ЙΘ ЙΘ ЙΘ BB18-FF FF FF 00 00 ЙΘ FF ЙΘ BB20-99 00 FF 00 00 αа 00 00 BB28-FF 00 00 FF FF 00 00 00 BB30-FF 00 00 00 FF 00 FF 00 BB38-00 FF 00 FF 00 00 00 00 BB40-00 FF FF 00 FF 00 00 00 00 BB48-FF 00 00 FF 00 00 00 BB50-FF FF FF FF 00 00 00 00 BB58-00 00 FF FF 00 00 00 00 BB60-FF FF FF 00 FF 00 00 FF BB68-00 00 FF FF 00 00 00 00 BB70-00 FF 00 00 00 00 00 00 BB78-99 FF FF 00 00 00 00 00 BB80-FF FF 00 00 FF FF 00 00 BB88-FF FF FF FF 00 00 00 00 BB90-00 00 FF 00 00 FF 00 FF BB98-FF 00 00 00 FF 00 00 FF 00 00 BBA0-FF 00 FF FF 00 00 BBA8-FF FF 00 FF FF 00 00 00 FF 00 FF BBB0-00 00 FF FF FF BBB8-FF 00 00 00 FF 00 00 00 BBC0-FF 00 FF FF 00 00 FF 00 BBC8-FF FF 00 FF FF 00 00 00 BBD0-FF 00 FF 00 00 FF FF 00 00 FF FF 00 BBD8-00 00 00 FF BBE0-00 FF 00 FF FF 00 00 FF BBE8-FF FF 00 00 00 FF 00 FF BBF0-FF FF FF 00 00 00 00 FF BBF8-00 FF FF FF 00 00 00 00 Superimposing the two arrays, I can see which 4-and-4 encoded values are being

		ne orig: Jorking			ver	found
BB00- BB08- BB10- BB18- BB20- BB28- BB30- BB38-		9				
And s	o on.	No obv:	ious	patter	n em	erges.
Hmm.						
at tĒ	ne firs	o swito st valuo isk but	e tha	t appe	ars	on the
myste them secto table	eries. in the or. The es, eve	coded value of the code value of the code value of the code of the	nprot: ss fi: e pub he 80:	ected eld of lished s. Her	disk eve loo e's	s use ry
Dec	Hex	Binary	<u> </u>	4-and	-4	
0 1 2 3 5 6 7 8	\$01 \$02 \$03 \$04 \$05 \$06 \$08	000000 000000 000000 000000 000000 00000	301 310 311 100 101 110 111	AA AB AB AB AA AE AB AE AB AF AE AA		

```
The first value missing from my bit
copy was 5 (in location $BB05),
let's search the original disk for the
nibble sequence "AA AF F5". That's the
4-and-4 encoded value ("AA AF"), plus
the first nibble of the expected
epiloque ("F5 AA").
                  --0--
TRACK: 22 START: 1800 LENGTH: 3DFF
2AC8:
         AΒ
            AA F5
                                 UTFW
      D4
                  AA AA
                         AΑ
                            AΑ
               D4 AB AB F5
2AD0: D4
         D5 DE
                            AΑ
2AD8: FF FF 9A 9A BB DA 95 AB
2AE0: BD D5 FF 9B D4 D5 DE
                            D4
                  ~~~~~~~~
2AE8:
      AA AF F5 AA FF FF A6 AE
                                <-2AE8
           ~~~~
      ~~~~
2AF0: F6 A5 BA EF B5 FF FF B5
2AF8: D5 DE D4 AB AF F5 AA FF
           D4 AB AF F5 AA FF
2800: FF 9A BB DA 95 D5 BD D5 FIND:
2B08: FF EC
           D4 D5
                  DE
                      D4
                         ΑE
                            AB
                                AA AF F5
I highlighted the three relevant
sequences, starting at offset $2AE4:
  - D4 D5 DE D4 (prologue)
               (4-and-4 encoded value)
  – на АF
– F5 АА
                (epiloque)
In that same screenshot, we can see
evidence of several other values:
    "AB AA" (2) at offset $2AC9
  - "AB AB" (3) at offset $2AD4
    "AE AB" (9) at offset $2B0E
```

Now let's see what copy looks like.	My	non-working	bit

So this confirms what running the protection code already told me: the original disk contains the group of (prologue)(4-4 encoded value)(epilogue)

where the encoded value is 5.

ES6,D1=non-working copy] --v- TRACK: 22 START: 1800 LENGTH: 3DFF 2430: AA FF D4 D5 DE D4 AA AB VIEW 2438: F5 AA AA AA D4 D5 DE D4 2440: AB AA F5 AA AA AA D4 D5

~~~~

2448: DE D4 AB AB F5 AA FF B3

^^^^
2450: A6 AE F6 A5 AA EF B5 FF <-2450
2458: FF A9 AB BD A9 AA BF D6
2460: FF FF D4 D5 DE D4 AB AE

^^^^^^^^^^^^
2468: F5 AA FF FF A6 AE F6 A5 FIND:
2470: BA FF B5 96 FF D4 D5 DE D4 D5 DE

 $--\wedge--$ 

Highlighting prologues ("D4 D5 DE D4"), we can see the track contains lots of groups with 4-and-4 encoded values: - "AA AB" (1) at offset \$2436 - "AB AA" (2) at offset \$2440

- "AB AA" (2) at offset \$2440 - "AB AB" (3) at offset \$244A - "AB AE" (6) at offset \$2466 But no group contains "AA AF" (5),

anywhere in the track. I am no closer to understanding why.



Chapter 3 Now You See It, Now You Don't Let's return to the original disk and examine the protection track again.

[S6,D1=original disk]

--v--

TRACK: 22 START: 1800 LENGTH: 3DFF

D4 D5 DE D4 AB AB

AB.

2F88: 9A 9A BB DA 95 EB FD

2F90: FF FF 9A 9A BB DA 95 D5 2F98: BD D5 FF FF 9A 9A BB DA

2F68: AA FF 9A 9A BB DA 95 AB

2F58: DE D4

AA.

2F60:

2F70: D5 FF FF B5 D5 DE BD D4 ??^^^^ 2F78: F5 AA AA FF A6 <-2F78 AA AF AΑ ~~~~ 2F80: AE F6 A5 BA EF B5 FF FF

AA F5 AA AA

AA.

F5

D5

UIEW

FIND:

AA AF F5

--^-Waaaaait a minute. That can't be right.
There's the 4-and-4 encoded value for 5
("AA AF", at offset \$2F78), and the
epilogue after it ("F5 AA", at offset \$2F7A). But look at the prologue before it. It should be "D4 D5 DE D4" starting

is missing. Maybe there are two copies of this group and I just found the other one this time? Nope, this is only instance of "AA AF F5" on the track.

at offset \$2F74, but the first nibble

this is always a reasonable question.) Nope, this is the same original disk I read last time. Side note: I realize the offsets are different, but they're not relevant to this mystery. They depend on where the disk was spinning when Copy II Plus started reading the track this time. That could be anywhere; I'd expect the offsets to change every time I re-read a track. But the nibbles themselves should be the same. That's the data on the disk. Data on the disk shouldn't change every time you read it. That's kind of the point of a storage device. Unless... Oh no. Oh God. Oh God no. There's only one thing you can put on a disk that will change every time you read it: nothing. And by "nothing," I mean "a long sequence of zero bits." And that's what is on the original disk between each of these groups: nothing.

Did I insert the wrong disk? (I do have an overabundance of floppy disks, so a magnetic state change." The Disk II drivē isn't digital; īt's analog. If it doesn't see a state change in a certain period of time, it calls that a "0". If it does see a change, it calls that a "1". But the drive can only tolerate a lack of state changes for so long about as long as it takes for two bits to go by. Fun fact(\*): this is why you need to use nibbles as an intermediate on-disk format in the first place. No valid nibble contains more than two zero bits consecutively, when written from mostsignificant to least-significant bit. So what happens when a drive doesn't see a state change after the equivalent of two consecutive zero bits? The drive thinks the disk is weak, and it starts increasing the amplification to try to compensate, looking for a valid signal. But there is no signal. There is no data. There is just a yawning abyss of nothingness. Eventually, the drive gets desperate and amplifies so much that it starts returning random bits based on . ambient noise from the disk motor and the magnetism of the Earth. Seriouslu. (\*) not guaranteed, actual fun may vary

A bit of background. When we say a

"zero bit," we really mean "the lack of

8 zero bits -- like any other 8-bit nibble. You can write whatever you want to a disk; it doesn't need to be what DOS would consider a "valid" nibble. But when you read that nibble back, the drive can<sup>i</sup>t handle 8 zero bits in a row, so it will actually return some random bits. Which is why no one does that. Returning random bits doesn't sound very useful for a storage device, but it's exactly what the developer wanted, and that's exactly what this copy protection scheme depends on. Here's whu: Bit copiers can't duplicate a long sequence of zero bits. Why? Because that's not what they see. What they see is some random bits -the real zero bits interspersed with phantom "1" bits. So that's what they write to the target disk. Whatever randomness they get when they read the original disk will essentially get "frozen" onto the copy.

It's trivial to write zero bits to a disk; just write a #\$00 nibble to write

at the protection code again. It looks for the custom prologue "D4 D5 DE D4", then checks the 4-and-4 encoded value that follows and marks that value as "found" in the buffer at \$BB00. But each prologue is preceded by a bitstream that changes every time it's read. Sometimes those random bits will align in such a way that they form two full, valid nibbles, and the next prologue will be read correctly. Other times, they will only form a partial nibble that is completed by the first

few bits of the prologue, so the

As far as I can tell, the sequence of zero bits is 18 bits long. I'm not sure the exact length matters. The important

proloque will be missed.

part is the randomness.

Now, why does this matter? Let's look

Here's what's on the disk (the \$D4 and \$D5 are the start of the proloque): /--00--\/--00--\/--D4--\/--D5--\ 00000000000000000000110101010110101 But remember, more than two consecutive zero bits will form an "abyss" that the floppy drive will fill with randomness. Some of those zero bits before the prologue will randomly transform into "1" bits, and it'll be a different set every time you read the disk. How does that affect the protection check?

Here's one of many possible bitstreams

that might come out of the abuss:

/--FF--\/--9B--\/--D4--\/--D5--\

00111111111001101111101010011010101

That's what I saw the first time I read

the original disk with the Copy II Plus

nibble editor: the abyss coalesces into

two full nibbles (\$FF<sup>-</sup>\$9B), then I read

the proloque (\$D4 \$D5 and so on).

Here's another possible bitstream that might come out of the abyss: That's what I saw (on the same disk!) the second time I read it, as shown at the beginning of this chapter. This time, the abyss coalesces into two and a half nibbles -- \$FF, \$FF, and a few extra bits. Those extra bits combine with the bits that are supposed to be part of the \$D4 nibble, but because . we're in the middle of a nibble when the \$D4 bits start, we end up finishing that nibble (\$B5) instead, and the \$D4 nibble disappears. The first nibble of the proloque, \$D4, gets consumed by the abyss. The second nibble of the prologue, \$D5, survives unscathed, but by then it's too late. Without the full prologue, we'll skip this entire group and the 4-and-4 encoded value within it. The protection code requires finding 160 unique 4-and-4 encoded values after a full prologue (\$D4 \$D5 \$DE \$D4). The loop at \$B6CD counts the number of values it's found by incrementing the X register for every value marked "found" in the buffer at \$BB00. At \$B6D8, it compares X to #\$A0 and branches to the success path when it's found enough unique values. It doesn't care which values it finds, or in which order; it only cares about the total count.

you read the track, but you'll (eventually) find them all as random bits fluctuate. So the check only passes after some number of reads of a nondeterministic bitstream that sometimes (but not always) corrupts the data after it. But bit copiers don't preserve long streams of zero bits. Instead, they write out whatever phantom "1" they find. On a copy, you'll also miss some percentage of prologues -- and thus skip over some percentage of the 4-and-4 encoded values -- each time you read the track. But that percentage will never increase no matter how many times you read it. No more randomness. No morė "eventually."

God, I hate physical objects.

If you re-read the original disk enough times, each stream of random bits will (eventually) align so the next prologue

(eventually) marked as "found." You'll miss some of the prologues each time

is (eventually) read correctly and enough different encoded values are

P

Chapter 4 In Which The Solution Is Trivial Except When It Isn't protection routine at \$B619 so it returns gracefully and unconditionally. Making a fresh copy (no more hacks), I can turn to my trusty Disk Fixer sector editor and put an "RTS" at the start of the protection code: T03,S06,\$19: A0 -> 60 JPR#6 ...boots to main menu, at last... Except... In the course of testing, I discovered one final small slight really hardly noticeable little problem... It doesn't print. I mean, it prints, but it doesn't print what I want it to print. It's all just garbage. At first, I assumed this was a problem with my emulator. (I no longer have a real printer to test with.) But no. Through some emulator trickery, I made an unpatched copy "pass" the protection check by altering the CPU flow without altering the code on disk or in memory, and that copy printed iust fine. Something is detecting that I modified the protection code.

None of this affects how I can make my copy work. I just need to patch the

realistically it's probably just XORing the bytes together. Another possibility is ADDing the bytes together. But throw away your modern conceptions of some sort of cryptographic hash function. In the 80s, any "encryption" was probably just XOR, and any "checksum" was probably one byte long and was the cumulative XOR of all the other butes. (Story time! Sierra was notorious for their anti-tamper checks. Games would have a self-decrypting protection check and a separate anti-tamper check that XOR'd the encrypted protection code. But the first two bytes of the code were always the same, and word got around of how you could defeat the protection AND the anti-tamper check in all their games by searching for the first two bytes and changing them to known values. In response, Šierra added a second anti-tamper check that checksummed only the odd bytes.) What I'm saying is that I can probably fool this tamper check without even finding it, by applying the principle of transitivity. I changed an \$A0 byte to \$60 to defeat the protection. Now I need to find a \$60 byte and change it to \$A0 to defeat the anti-tamper check.

Now, there are any number of ways to checksum a region of memory, but the Apple II is not a fast machine, and

entire thing). T03,S06,\$EA: 60 -> A0 ]PR#6 ...works, and prints correctly... Sometimes it's better to be lucky than

After extensive testing, there doesn't appear to be any further protection.

And there is one, in the same sector, at offset \$EA: the "RTS" instruction at the end of the subroutine that starts at \$B6E5 (but is no longer called since I put an "RTS" at the very beginning of the protection code to bypass the

Quod erat liberandum.

good.