5 wires, 5 ways to get root

Speaker Bio

- Electrical Engineering education with focus on CS and Infosec
- 10 years of fun with hardware
 - o silicon debug
 - o security research
 - pen testing of CPUs
 - o security training
- Hardware Security Training:
 - "Software Exploitation via Hardware Exploits"
 - "Applied Physical Attacks on x86 Systems"



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Speaker Bio

- Electrical and Computer
 Engineering education, with focus
 on hardware design and test
- 10+ years designing, implementing, and testing SoC silicon debug features
- Hardware and firmware pentesting



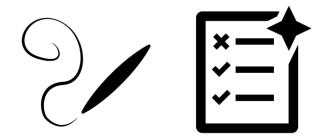
Matt King
@syncsrc
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5 wires, 5 ways to get root

Yeah, we get that part

But what's this?

















IEEE Standard Test Access Port and Boundary-Scan Architecture

IEEE Std 1149.1:

OSI Model

data unit layers application
Network Process to Application data Layers presentation data Data Representation & Encryption Host session data Interhost Communication transport End-to-End Connections segments and Reliability network Media Layers packets Path Determination & Logical Addressing (IP) data link frames Physical Addressing (MAC & LLC) physical Media, Signal and Binary Transmission bits

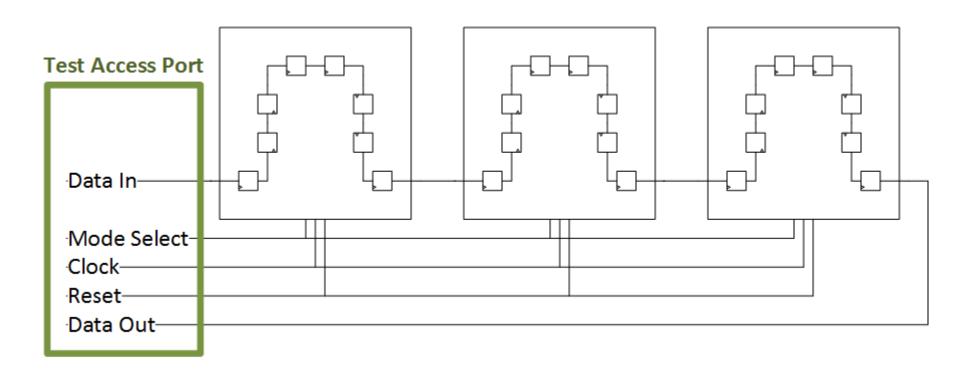
Remember This?

OSI Model data unit layers application Network Process to Application data Layers presentation data Data Representation & Encryption Host session data Interhost Communication transport End-to-End Connections segments and Reliability network Media Layers packets Path Determination & Logical Addressing (IP) data link frames Physical Addressing (MAC & LLC) physical Media, Signal and Binary Transmission bits

JTAG Model

TDI, TDO, TMS, TCK, TRST

Physical Layer: Test Access Port



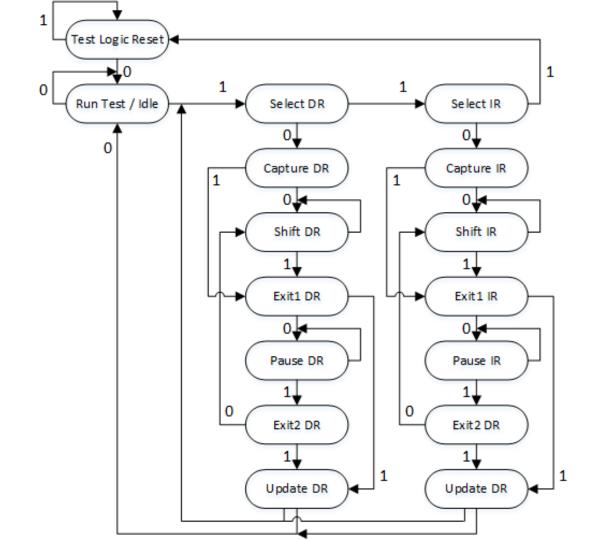
OSI Model data unit layers application Network Process to Application data ayers presentation data Data Representation & Encryption Host session data Interhost Communication transport End-to-End Connections segments and Reliability network Media Layers packets Path Determination & Logical Addressing (IP) data link frames Physical Addressing (MAC & LLC) physical Media, Signal and Binary Transmission bits

JTAG Model

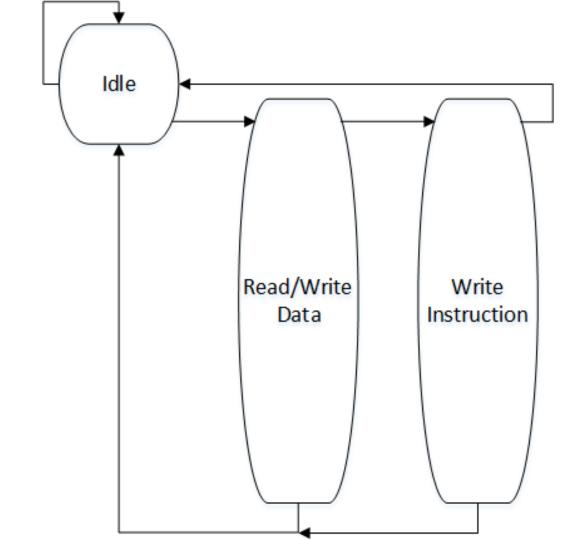
TAP FSM

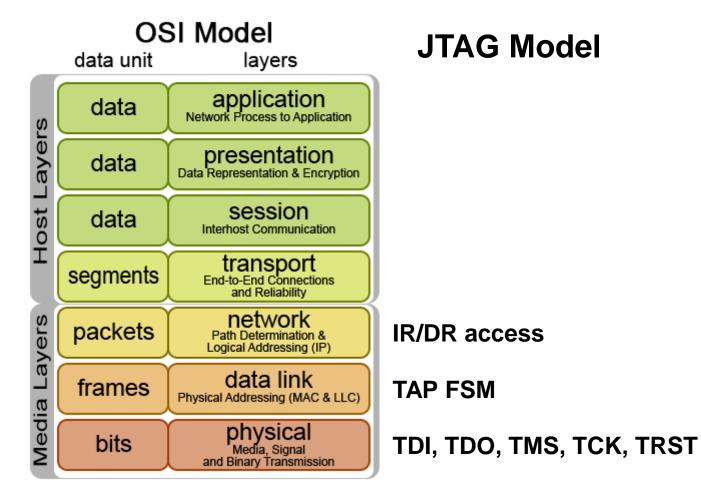
TDI, TDO, TMS, TCK, TRST

Data Link: TAP FSM

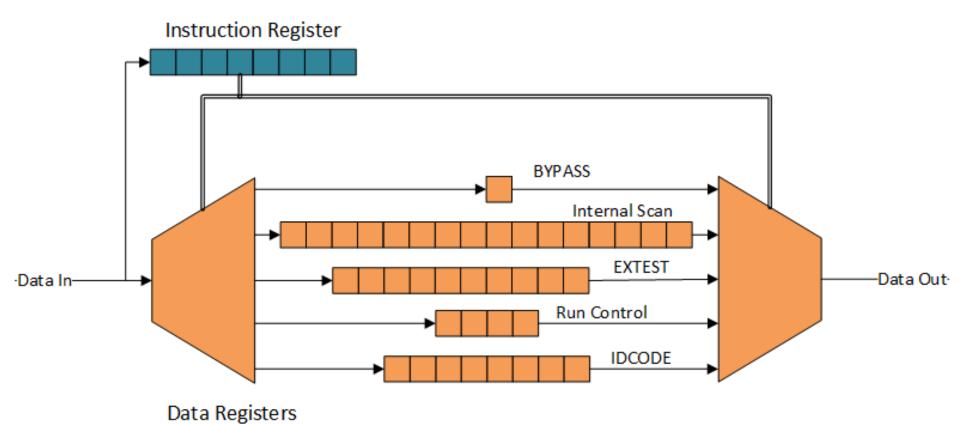


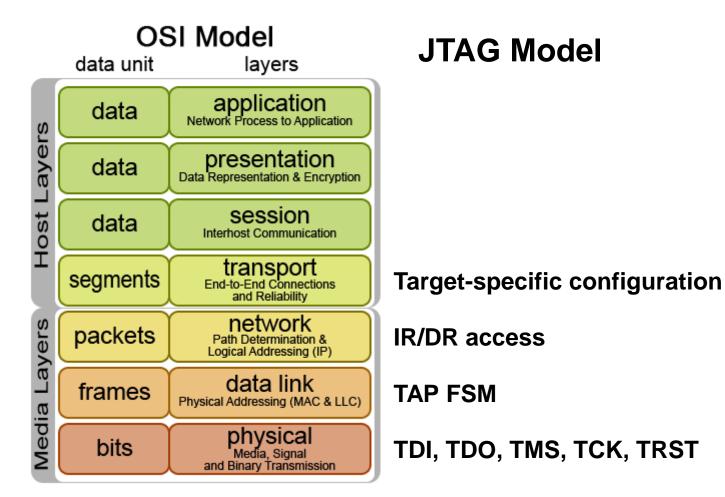
Data Link: TAP FSM





Network Layer: IRs & DRs





Transport Layer: Target-Specific

Table 6-1 TAP Instruction Overview

Code	Instruction	Function		
All 0's	(Free for other use)	Free for other use, such as JTAG boundary scan		
0x01	IDCODE	Selects Device Identification (ID) register		
0x02	(Free for other use)	Free for other use, such as JTAG boundary scan		
0x03	IMPCODE	Selects Implementation register		
0x04 - 0x07	(Free for other use)	Free for other use, such as JTAG boundary scan		
0x08	ADDRESS	Selects Address register		
0x09	DATA	Selects Data register		
0x0A	CONTROL	Selects EJTAG Control register		
0x0B	ALL	Selects the Address, Data and EJTAG Control registers		
0x0C	EJTAGBOOT	Makes the processor take a debug exception after reset		
0x0D	NORMALBOOT	Makes the processor execute the reset handler after reset		

That's just MIPS.

That's just MIPS.

X86 is different ARM is different Each SOC is different

OSI Model data unit layers application Network Process to Application data ayers presentation data Data Representation & Encryption Host session data Interhost Communication transport End-to-End Connections segments and Reliability network -ayers packets Path Determination & Logical Addressing (IP) data link frames Physical Addressing (MAC & LLC) Media physical bits Media, Signal and Binary Transmission

JTAG Model

- --- (no one uses this crap)
- --- N/A sessionless...
- **Target-specific configuration**

IR/DR access

TAP FSM

TDI, TDO, TMS, TCK, TRST

OSI Model data unit layers application Network Process to Application data ayers presentation data Data Representation & Encryption Host session data Interhost Communication transport End-to-End Connections segments and Reliability network -ayers packets Path Determination & Logical Addressing (IP) data link frames Physical Addressing (MAC & LLC) Media physical bits Media, Signal and Binary Transmission

JTAG Model

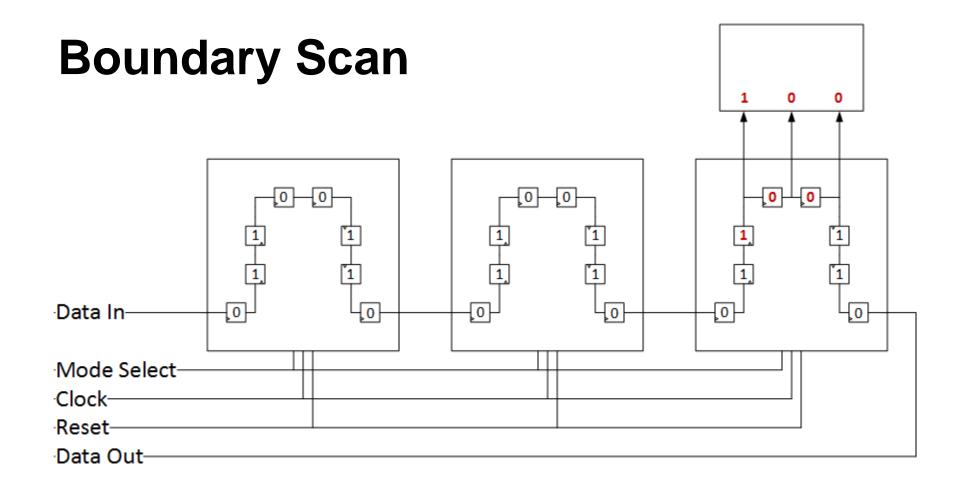
BScan, Memory & Register Access

Target-specific configuration

IR/DR access

TAP FSM

TDI, TDO, TMS, TCK, TRST

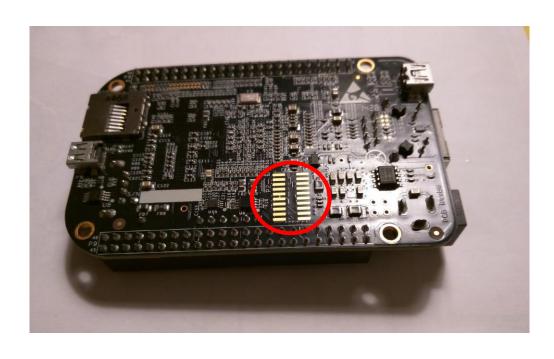


#1: Access Non-Volatile Storage

via Boundary Scan

JTAG on the Beaglebone Black

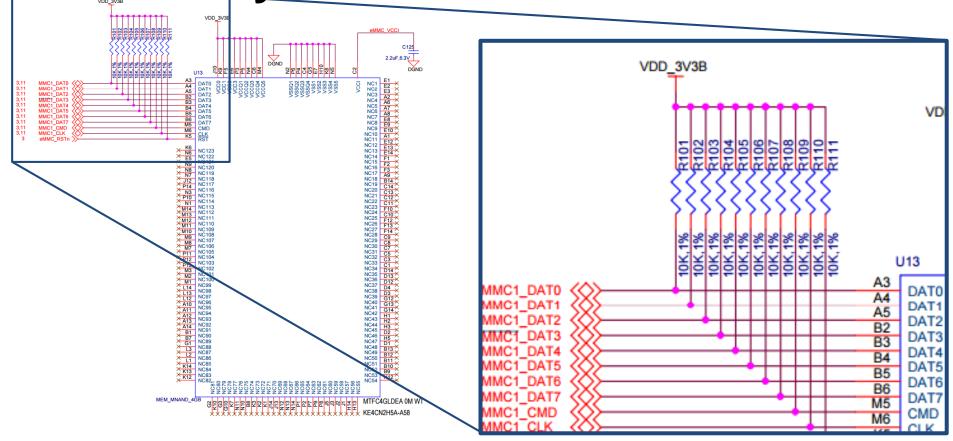




Boundary Scan on the BBB



Boundary Scan on the BBB

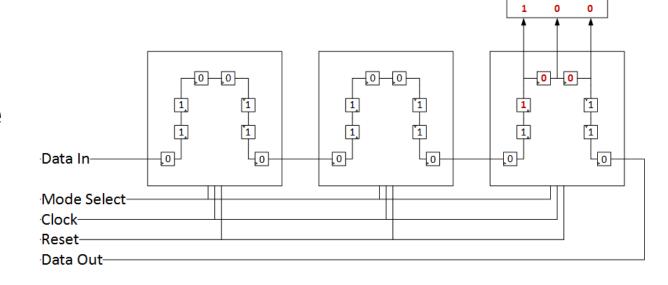


UrJTAG

Initbus: initialize a bus on the selected part

"via BSR": use boundary scan

Pins defined in a .bsdl file



UrJTAG

Detectflash: detect parameters of attached memory device

peek/poke:
read/write a single word

readmem/writemem: dump to file or load from file

Access Non-Volatile Storage

Cons:

- Really slow
- Requires mapping out the full boundary scan or a BDSL file
- Really slow
- Requires emulating the flash interface protocol
- Really, really slow

Pros:

- Access to BGA pins without desoldering anything
- Works regardless of CPU functionality
- You can visit a park, go out for a beer, and sometimes take a week holiday while it works

Run Control

"Optional" vendor-specific registers which:

Allows insertion of commands

- read/write memory, registers, or I/O ports

Control of execution

- step through instructions
- step into other code
- breakpoints

analysis

#2 Scrape memory for offline

Memory Access

```
> mdw 0x00 0x40
0x00000000: 68bba82d 5256e25d 48a6268c c1019709 1337c0de 880999b6 a0a047f
0x00000020: 997d1c46 4d348585 0b94b2e3 ab6b0040 a23dee32 07f1a4b6 b941410
0x00000040: a284c82b c711082f bf2bda21 c2507e77 f035ceca 45d1727d 30c7f4f
0x00000060: b8a7fc33 c3dacc16 265e9ab0 c36d397d d654f8e9 30ab86d8 5fd1cfd
0x00000080: 3aa60e7e 69e72dcc dccc5163 8d115177 68834721 7025e8cb 3e09b7b
0x000000a0: ea92aa2e 468f00f4 22af5680 cee32148 16dac22c 8c8e2372 17d1a38
0x000000c0: bf6e0eb7 254b71dd eca4d4e0 bae09034 83f1413d 998bba3b 8314070
0x000000e0: 7ffe52db 84453273 0fd7fc6e 17209711 4202f0d6 01fb48a0 367ec63
>
> mem2array tcl variable name 32 0xd0000000 0x1000
>
> dump image mydump.bin 0x8000000 0x2000000
>
```

Simple Memory Analysis

strings and grep

- if you know just what you're looking for

binwalk

- it's meant for firmware images, but
- it recognizes commons structures like:
 - files in memory
 - keys and certificates
 - some code blocks

Advanced Memory Analysis

dmesg log of a system:

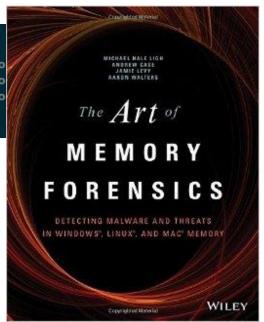
```
AppleThunderboltHAL::earlyWake - complete - took 0 milliseconds

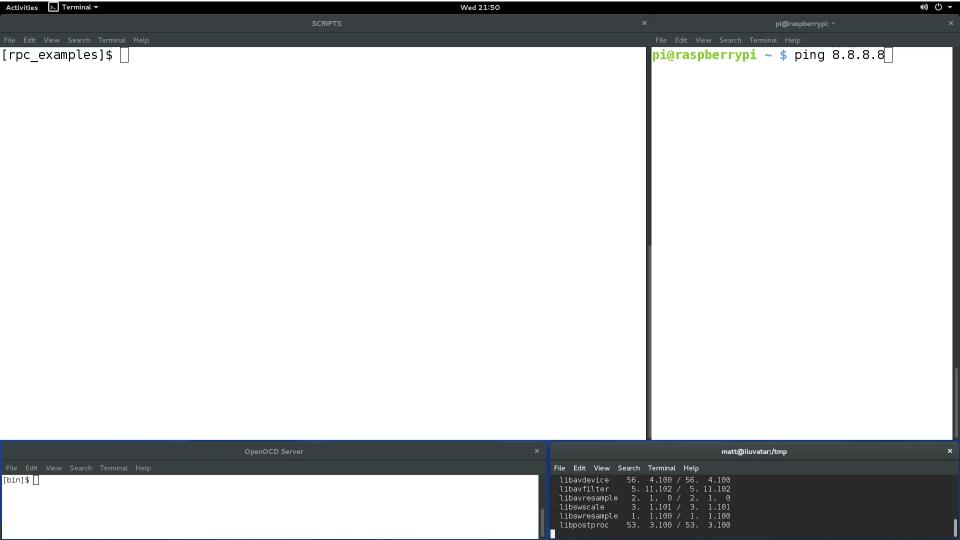
Thunderbolt Self-Reset Count = 0xedefbe00

IOThunderboltSwitch<0xffffff8013f40400>(0x1)::listenerCallback - Thunderbolt HPD packet fo
IOThunderboltSwitch<0xffffff8013f40400>(0x1)::listenerCallback - Thunderbolt HPD packet fo
IOThunderboltSwitch<0xffffff8013f40400>(0x1)::listenerCallback - Thunderbolt HPD packet fo
[ PCI configuration begin ]
[ PCI configuration end, bridges 12, devices 14 ]
```

process list of a system:

Name	Pid	Uid
kernel_task	0	0
.launchd	1	0
com.apple.IconSe	36773	-
com.apple.hiserv	36755	501
UserEventAgent	11	0
kextd	12	0
notifyd	14	0
securityd	15	0
diskarbitrationd	16	0
powerd	17	0
configd	18	0







Memory Scraping & Analysis

Cons:

- Still slow
- May Will almost certainly crash the target
- Not fast
- Need to be careful about paging and MMUs
- Slow

Pros:

- Minimal target knowledge required
- Existing forensic analysis tools
- Inception over Jtag =

SLOWCEPTION

 Plenty of time to take your wife out for dinner

#3: Patching Boot Arguments

Boot Arguments

Commands passed to the kernel (and beyond) at boot:

```
-bootargs=console=ttyS0,115200 root=ubi0:rootfs ubi.mtd=4,2048 rootfstype=ubifs
```

Boot Arguments can tell us to boot in single user mode:

```
-bootargs=console=ttyS0,115200 root=ubi0:rootfs ubi.mtd=4,2048 rootfstype=ubifs 1
```

Boot Arguments

Sometimes, they're hard-coded in the kernel:

```
020bc20: 6c78 0a00 0000 0000 0000 0000 0000 0000
                                                1x..........
020bc30: 636f 6e73 6f6c 653d 7474 7953 302c 3131
                                                console=ttyS0,11
020bc40: 3532 3030 2072 6f6f 743d 3331 3a32 2072
                                                5200 \text{ root} = 31:2 \text{ r}
020bc50: 6f6f 7466 7374 7970 653d 7371 7561 7368
                                                ootfstype=squash
020bc60: 6673 2069 6e69 743d 2f73 6269 6e2f 696e
                                                fs init=/sbin/in
020bc70: 6974 206d 7464 7061 7274 733d 6174 682d
                                                it mtdparts=ath-
020bc80: 6e6f 7230 3a31 3238 6b28 752d 626f 6f74
                                                nor0:128k(u-boot
020bc90: 292c 3130 3234 6b28 6b65 726e 656c 292c
                                                ),1024k(kernel),
020bca0: 3238 3136 6b28 726f 6f74 6673 292c 3634
                                                2816k (rootfs), 64
020bcb0: 6b28 636f 6e66 6967 292c 3634 6b28 6172
                                                k(config),64k(ar
020bcc0: 7429 206d 656d 3d33 324d 0000
                                                t) mem=32M.....
                                     0000
        0000 0000 0000 0000 0000
020bcd0:
                                     0000
```

Boot Arguments

- 1. Set a breakpoint or watchpoint
- 2. Wait for the kernel to be loaded in memory
- 3. Halt
- 4. Patch kernel
- 5. Allow patched kernel to boot

```
File Edit View Search Terminal Help
                                                              File Edit View Search Terminal Help
WASP ----> S27 PHY
                                                             telnet>
GMAC: cfa1 0x5 cfa2 0x7114
eth0: ba:be:fa:ce:08:41
s27 reg init
athrs27_phy_setup ATHR_PHY_CONTROL 4: 0x1000
athrs27 phy setup ATHR PHY SPEC STAUS 4: 0x10
eth0 up
WASP ----> $27 PHY
                                                              oefitz@linUX303: ~/Documents/classes/rpi-jtag
GMAC: cfg1 0xf cfg2 0x7214
                                                              File Edit View Search Terminal Help
                                                             openocd>
eth1: ba:be:fa:ce:08:41
s27 reg init lan
ATHRS27: resetting s27
ATHRS27: s27 reset done
athrs27 phy setup ATHR PHY CONTROL 0: 0x1000
athrs27 phy setup ATHR PHY SPEC STAUS 0: 0x10
athrs27 phy setup ATHR PHY CONTROL 1: 0x1000
athrs27 phy setup ATHR PHY SPEC STAUS 1: 0x10
athrs27 phy setup ATHR PHY CONTROL 2: 0x1000
athrs27 phy setup ATHR PHY SPEC STAUS 2: 0x10
athrs27 phy setup ATHR PHY CONTROL 3: 0x1000
athrs27 phy setup ATHR PHY SPEC STAUS 3: 0x10
eth1 up
eth0, eth1
Autobooting in 1 seconds
File Edit View Search Terminal Help
```



Boot Patch

Cons:

- Has to be done at boot time
- Single user mode means manually mounting everything

Pros:

- Can be done after a kernel signature is checked
- Doesn't depend on persistent storage or root file system contents
- Rebooting doesn't take several days

#4: Kernel Patching

Linux File System ACL Enforcement

```
int generic_permission(struct inode *inode, int mask) {
          int ret;
           * Do the basic permission checks.
           ret = acl_permission_check(inode, mask);
          if (ret != -EACCES)
                     return ret;
           * Searching includes executable on directories, else just read.
           mask &= MAY_READ | MAY_WRITE | MAY_EXEC;
           if (mask == MAY_READ)
                      if (capable_wn_inode_uidgid(inode, CAP_DAC_READ_SEARCH))
                                 return 0:
           return -EACCES;
```

Locating Kernel Functions

```
$ cat /proc/kallsyms
c1000000 T startup 32
c1000000 T text
c1001000 T wakeup_pmode_return
c100104c t bogus_magic
c100104e t save registers
c10adec0 T page_follow_link_light
c10adef0 T page_readlink
c10adf40 T generic_permission
c10ae0f0 t follow dotdot rcu
c10ae270 t follow dotdot
c10ae340 t handle dots
c10ae380 T full name hash
c10ae3c0 T final_putname
c10ae400 t getname flags
c10ae4f0 T getname
c10ae500 T __inode_permission
```

Identifying Patch Point

```
loc 0x000003b:
83 c4 08
                             add
                                     esp,0x8
5b
                                     ebx
                             pop
5e
                                     esi
                             pop
5f
                                     edi
                             pop
5d
                                     ebp
                             pop
c3
                             ret
. . .
ba 02 00 00 00
                                     edx,0x2
                             mov
89 f0
                                     eax,esi
                             mov
e8 dd 4c f8 ff
                             call
                                     0xfff84da0
84 c0
                             test
                                     al,al
Of 85 6e ff ff ff
                             jne
                                     0 \times 00000039
90
                             nop
8d 74 26 00
                             lea
                                     esi, [esi+eiz*1+0x0]
b8 f3 ff ff ff
                                     eax, 0xfffffff3
                                                          ; mov eax, -EACCESS
                             mov
e9 61 ff ff ff
                                     0 \times 0000003b
                                                          ; goto function return
                             jmp
```

Delivery Options

Convert JTAG sequences into a standard format (SVF/XSVF)

Enables replay of debug performed in OpenOCD

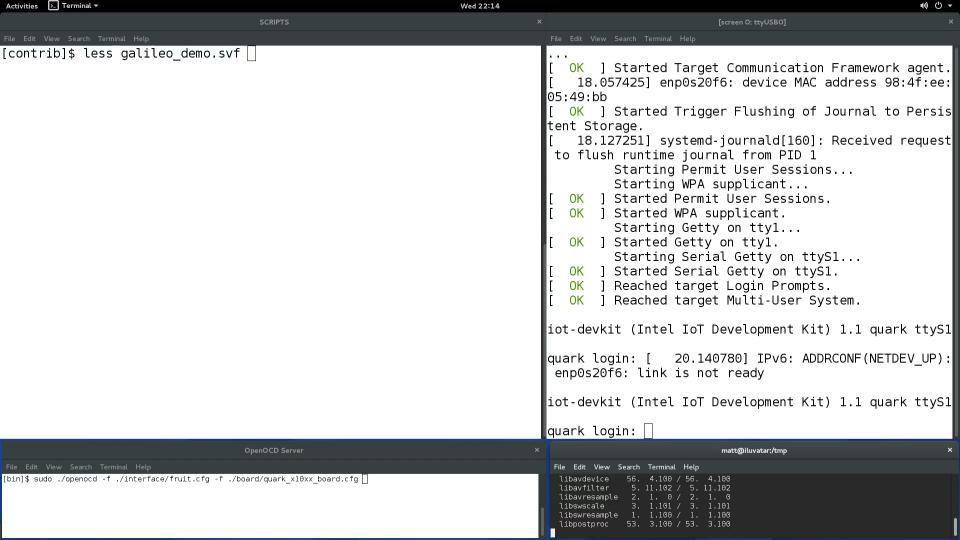
- mww 0xc10ae011 0

!Begin Test Progra TRST OFF: **ENDIR IDLE**; **ENDDR IDLE**; HIR 8 TDI (00); HDR 16 TDI (FFF TIR 16 TDI (0000) TDR 8 TDI (12); SIR 8 TDI (41); SDR 32 TDI (ABC

STATE DRPAUSE

RUNTEST 100 TO

IEnd Tost Program



Kernel Patch

Cons:

- Target & kernel specific
- SVF does not support datadependant control flow
- Need to understand and manage implicit debugger actions
- Effort required to find fixedlocation kernel functions and patches for them

Pros:

- Can be applied to running target
- Fast
- Can be implemented as an SVF
- Many hardware and software options for SVF playback

#5: Patching a Process

getty Parameters

```
$ xxd /sbin/getty
. . .
00006770
          1b 5b 48 1b 5b 4a 00 25
                                     73 25 73 20 28 61 75 74
                                                                |.[H.[J.%s%s (aut|
00006780
          6f 6d 61 74 69 63 20 6c
                                     6f 67 69 6e 29 0a 00 25
                                                                |omatic login)..%|
00006790
          73 3a 20 72 65 61 64 3a
                                     20 25 6d 00 25 73 3a 20
                                                                |s: read: %m.%s: |
000067a0
          69 6e 70 75 74 20 6f 76
                                        72 72 75 6e 00 63 68
                                                                |input overrun.ch|
000067b0
          65 63 6b 6e 61 6d 65 20
                                     66 61 69 6c 65 64 3a 20
                                                                eckname failed: |
000067c0
          25 6d 00 2d 68 00 2d 66
                                     00 2d 2d 00 25 73 3a 20
                                                                |%m.-h.-f<mark>.</mark>--]%s:
000067d0
          63 61 6e 27 74 20 65
                                78
                                     65 63 20 25 73 3a 20 25
                                                                |can't exec %s: %|
          6d 00 38 62 69 74 73
000067e0
                                     61 75 74 6f 6c 6f 67 69
                                                                |m.8bits.autologi|
                                00
```

Changing '--' to '-f' results in user being pre-authenticated

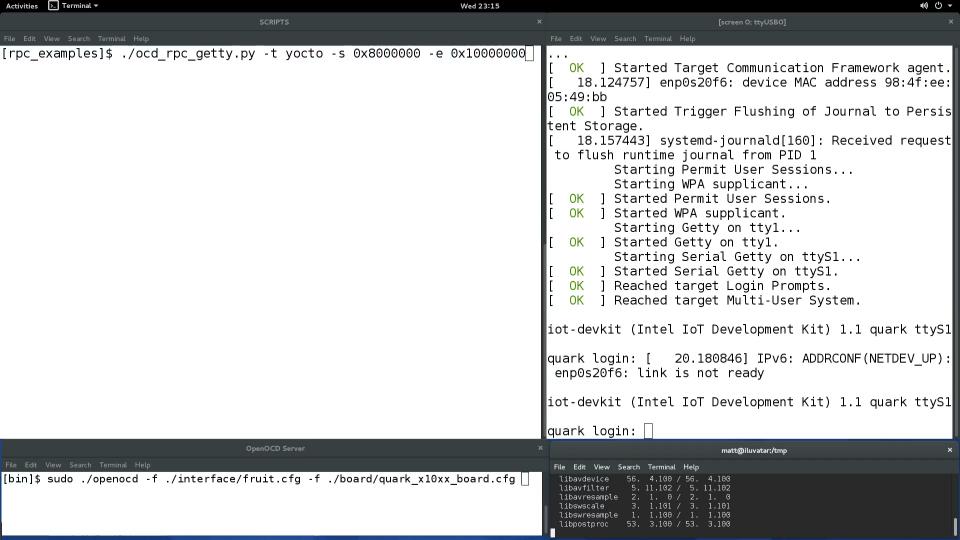
Searching Memory

OpenOCD has a scripting interface

- and examples in python

Check every page in memory for target process

- Read 8 bytes at offset 0x7c9 (512x speedup)
- Compare to signature, patch if match found





Patch a Process

Cons:

- Needs a small signature at a known page offset (to find process in memory)
- Requires knowledge of processes running on the target system

Pros:

- Can be applied to running target
- Arbitrarily change any process
- Still relatively fast
- More resilient to changes in target software

Summary

The CPU runs the show, but JTAG calls the shots via:

- I/O control
- Run control
- Memory access

https://github.com/syncsrc/jtagsploitation

Questions?