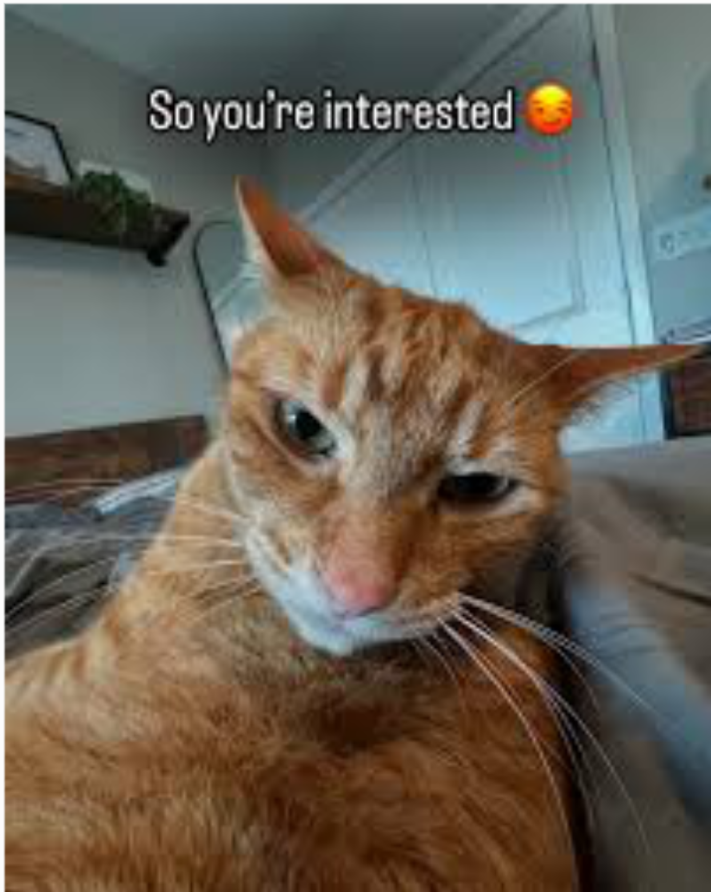


3. Process API



How to Create and Control processes

Difference between : return vs. exit, outside of main()

- return : end that function
- exit : end that process

Fork()

: create a new process

- copy current process
- return value : child's PID
 - $x < 0$: fork failed
 - $x == 0$: no child == this process is child
 - $x > 0$: parent, $x ==$ child's PID
- declared line == child birth point
- Difference between Child vs. Parent : return value of fork(), PID

- Can't guarantee which process(either child or parent) will be executed first

Wait()

: used by a process to wait for a child process.

- return value : Child's PID
- if n children -> wait() for n time

Exec()

: run a program that is different from the calling(parent) program.

system calls that allows a child to break free from its similarity to its parent and execute an entirely new program.

- fork() : run copies of the same(parent) program.
- exec() : run a different program
- memory space (ex. heap, stack) of the program are re-initialized
- Does not create new process => using same PID
- Transform the currently running program into a different program
- successful call to exec() never returns : don't proceed lines after exec() in original code.

int execvp(const char *file, char *const argv[])

- vp : vector path

Isn't it possible to run a completely different program with fork()?

: Since the child and parent have different PIDs, can't we just use the return value of fork() as a branch point with if-else?

=> Technically, yes — but it's inefficient and not recommended.

- fork() can work like a **branch point** like you said, with return value.
- But what about **execvp("wc", args);** ?
- exec() is like a **magic spell** — it replaces the current process with a new, independent one. It's for true execution.

- In theory, you could write the logic of `wc` directly in the child branch. But it would be: redundant, hard to maintain, completely non-portable
- `exec()` is more portable, modular, clean way.
- If `fork()` gives you a two-way branch like a fork in the road, `exec()` is Doraemon's anywhere door — it teleports the process into a different world (program).

Separation of `fork()` and `Exec()`

- run code after the call `fork()` ~ but before the call `exec()`;
 - use for environment setting
1. User : type a command
 2. Shell : figures out where in the file system the executable resides
 3. Shell : calls `fork()` to create a new child process to run the command
 4. Shell : calls `exec()` to run the command
 5. Shell : waits for the command to complete by calling `wait()`
 6. Shell : When the child completes, the shell returns from `wait()`
 7. Shell : prints out a prompt again, ready for your next command

File descriptors assignment in UNIX

- UNIX systems start looking for free file descriptors at 0.
- It's important to close right descriptor -> make empty and use by `open()`
- 0 : STDIN, 1 : STDOUT, 2: STDERR
- Close : empty that slot
- Open : fill that empty slot

Process Control And Users

: for communication, a process should use the **signal()** system call to "catch" various signals : stop, continue, terminate

- Users can only control their own processes
- It is the job of OS to parcel out resources to each user (and their processes) to meet overall system goals