### Advanced R Programming - Lecture 6

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### Today

Performant Code

Computational complexity

Parallelism

Improving R code

Parallelism in R

Rcpp

Memoization

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# Questions since last time?

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# Writing fast code

Speed is important!

## Writing fast code

Speed is important!

Time to write code



# Writing fast code

Performant Code

Speed is important!

Time to write code
Time to maintain (understand) code



# Writing fast code

Performant Code

Speed is important!

Time to write code
Time to maintain (understand) code
Time to execute code



### Old Adage About Software

Performant Code

"You can have it Good, Fast, Cheap. Pick any two."



#### Performance

Performant Code

- 1. Performance
- 2. Complexity

Complexity affects performance...



#### Performance

- 1. Performance
- 2. Complexity

Complexity affects performance...

...but performance does'nt affect complexity



### Computational complexity

Theoretical worst case

Big-Oh notation

Basic operations

Relationship: operations to problem size



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# Big Oh

"How fast does a function grow?"

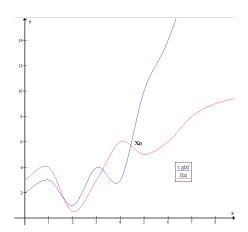
$$f(n)=O(g(n))$$

$$|f(n)| \le C * |g(n)| \forall n > X_0$$

number of operations



### Big Oh



# Big Oh

#### Example

$$f(n) = n^2 + 100n + 100$$



## Big Oh

#### Example

$$f(n) = n^2 + 100n + 100$$
$$f(n) = O(n^2)$$



# Complexities

Big Oh	Name	Example
O(1)	constant	assignments
O(log(N))	logarithmic	binary search (of sorted input)
O(N)	linear	max
$O(N^2)$	quadratic	naive vector-matrix mult.
$O(N^c)$	polynomial	naive matrix-matrix mult.
$O(c^n)$	exponential	brute force cracking of password



```
statement 2
statement c
```

statement 1

O(1)

### Determine complexity

```
if(a)
  statement a
else
  statement b
```

$$\mathsf{max}(\mathsf{O}(\mathsf{a}), \mathsf{O}(\mathsf{b}))$$



## Determine complexity

```
for(i in 1:N)
  statement i
```



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```
for(i in 1:N)
  for (j in 1:M)
                     0?
    statement i,j
```

# Determine complexity

```
for(i in 1:N)
                      O(N * M)
  for (j in 1:M)
    statement i,j
```

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$$g(n) = O(n^2)$$
$$O(n^3)$$

### What is parallelism?

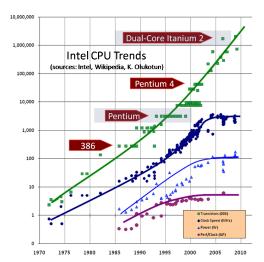
Multiple cores

Each core work with its own part

Cores can exchange information



#### Why parallelism?





# Why parallelism?

Single core limits

Handling larger data

Solving problems faster

More and more important



# Types of parallelism

Multicore systems

Distributed systems

Graphical processing units (GPU)



#### Speedup

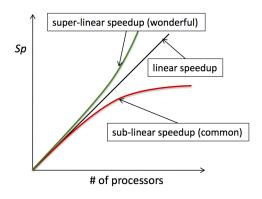


Figure: source



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#### Theoretical limits

#### Strong scaling: Amdahl's law

Deals with fixed problem size, increasing resources

Weak scaling: Gustafsons law

Deals with increasing size problem along with increasing resources



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$$S_p = \frac{1}{f_s + \frac{f_p}{P}}$$

Where:

 $f_s$  = serial fraction of code  $f_p$  = parallel fraction of code P = number of cores

For a fixed size problem!

#### Amdahl's law

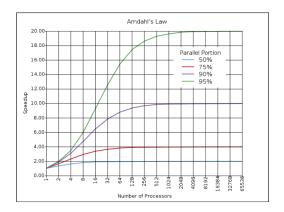


Figure: source



#### Gustafsons law

$$S_p = P - \alpha * (P - 1)$$

Where:

 $\alpha=$  the largest non-parallelizable fraction of any parallel process P= number of cores



### Practical problems

Costs of parallelism communication load balancing scheduling

fine-grained vs embarrassingly parallel



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### Practical problems

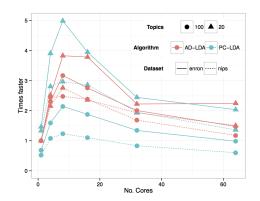


Figure: source



# Donald E. Knuth on Optimization

Programmers waste enormous amounts of time thinking about, or worrying about, the speed of noncritical parts of their programs, and these attempts at efficiency actually have a strong negative impact when debugging and maintenance are considered.

- Donald E. Knuth



#### Performance

#### Depends on many things

- 1. Code
- 2. Complexity
- 3. Compiler
- 4. Hardware
- 5. Language

If you don't measure, you don't optimize!



# How to optimize

- 1. Write code that works with accompanying test suite
- 2. Profile your code for bottlenecks
- 3. Try to eliminate the bottle necks
- 4. Redo 2-3 until fast enough



## **Profiling**

```
Rprof(tmp <- tempfile(),</pre>
  line.profiling = TRUE,
  memory.profiling = TRUE)
test_data <- pxweb::get_pxweb_data(</pre>
   nr1 =
     "http://api.scb.se/OV0104/v1/doris/sv/ssd/BE/BE0101
                                     /BE0101A/BefolkningNy",
   dims = list(Region = c('*'),
     Civilstand = c('*),
     Alder = c('*'),
     Kon = c('*').
     ContentsCode = c('*'),
     Tid = as.character(1970),
   clean = TRUE)
Rprof()
summaryRprof(tmp, lines = "show", memory = "both")
```

# Profiling

\$by.self

	sell.time	sell.pct	total.time	total.pct	mem.total
get_pxweb_data.R#102	1.96	39.2	1.96	39.2	579.2
get_pxweb_data_internal.R#42	1.16	23.2	1.16	23.2	405.0
get_pxweb_data.R#56	0.52	10.4	0.52	10.4	31.3
get_pxweb_data.R#80	0.38	7.6	0.38	7.6	29.1
get_pxweb_data.R#82	0.32	6.4	0.32	6.4	40.7
get_pxweb_data_internal.R#48	0.26	5.2	0.26	5.2	73.2
<pre>get_pxweb_data_internal.R#74</pre>	0.26	5.2	0.26	5.2	29.8
get_pxweb_data.R#83	0.08	1.6	0.08	1.6	17.2
api_catalogue.R#75	0.02	0.4	0.02	0.4	0.0
get_pxweb_data_internal.R#44	0.02	0.4	0.02	0.4	12.6
get_pxweb_data_internal.R#71	0.02	0.4	0.02	0.4	16.0

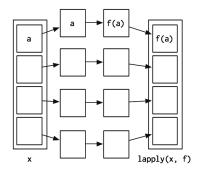
## **Improvements**

- 1. Look for existing solutions
- 2. Do less work
- 3. Vectorise
- 4. Parallelize
- 5. Avoid copies
- 6. Find smarter algorithms



## Parallelism in R

### Based on lapply()



## parallel package

### Two approaches:

- 1. mclapply()
- 2. parLapply()



# mclapply()

#### Pros

Simple to use Low overhead (startup)

#### Cons

Does not work on Windows
Only multi core



# parLapply(type="psock")

#### **Pros**

Works everywhere Good for testing/developing

#### Cons

Slow on multiple nodes



# parLapply(type="mpi")

#### Pros

Good for multiple computers Good for production

#### Cons

Can be used interactively Needs Rmpi package



## Example

example

# Rcpp

Often called interfacing

Similar can be done with Java and Fortran

Extremely fast!

But just handle bottlenecks!



Rcpp

Rcpp

$$f(n) = \begin{cases} n, & \text{if } n < 2 \\ F(n-1) + F(n-2), & \text{otherwise} \end{cases}$$

## Fibonacci R

```
fr <- function(n) {
   if (n < 2) return(n)
   f(n-1) + f(n-2)
}

system.time(fr(30))
user system elapsed
2.246  0.171  2.451</pre>
```

### Fibonacci C++

```
library(Rcpp)
cppFunction(code = '
  int fcpp(int n) {
    if (n < 2) return(n);
    return(fcpp(n-1) + fcpp(n-2));
,)
system.time(fcpp(30))
          system elapsed
user
0.007000000 0.000000000 0.006999999
```

## Memoization

A simple optimization technique Example of a general technique in optimization of trading memory for computation

Memoization stores (caches) results of function calls

If called again, returns old value

Depends on functional programming



```
> library(memoise)
> a <- function(x) runif(1)</pre>
> replicate(3, a())
[1] 0.6709919 0.3490709 0.4772027
> b <- memoise(a)</pre>
> replicate(3, b())
[1] 0.1867441 0.1867441 0.1867441
```



```
> c <- memoise(function(x) {Sys.sleep(1); runif(1)})</pre>
> system.time(print(c()))
[1] 0.7816399
      system elapsed
user
0.003 0.004 1.001
> system.time(print(c()))
[1] 0.7816399
user system elapsed
0.001 0.000 0.000
> forget(c)
[1] TRUE
> system.time(print(c()))
[1] 0.9234995
      system elapsed
user
0.003 0.004
                1.001
```

Memoization