

6.035

Introduction to Code Optimization

Instruction Scheduling

Outline

- Modern architectures
- Introduction to instruction scheduling
- List scheduling
- Resource constraints
- Scheduling across basic blocks
- Trace scheduling

Simple Machine Model

- Instructions are executed in sequence
 - Fetch, decode, execute, store results
 - One instruction at a time
- For branch instructions, start fetching from a different location if needed
 - Check branch condition
 - Next instruction may come from a new location given by the branch instruction

Simple Execution Model

- 5 Stage pipe-line



- Fetch: get the next instruction
- Decode: figure-out what that instruction is
- Execute: Perform ALU operation
 - address calculation in a memory op
- Memory: Do the memory access in a mem. Op.
- Write Back: write the results back

Simple Execution Model

time →

Inst 1

IF	DE	EXE	MEM	WB
----	----	-----	-----	----

Inst 2

IF	DE	EXE	MEM	WB
----	----	-----	-----	----

Simple Execution Model

time →

Inst 1

IF	DE	EXE	MEM	WB
----	----	-----	-----	----

Inst 2

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----	----	-----	-----	----

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Inst 2

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Inst 3

IF	DE	EXE	MEM	WB
----	----	-----	-----	----

Inst 4

IF	DE	EXE	MEM	WB
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Inst 5

IF	DE	EXE	MEM	WB
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From a Simple Machine Model to a Real Machine Model

- Many pipeline stages
 - Pentium 5
 - Pentium Pro 10
 - Pentium IV (130nm) 20
 - Pentium IV (90nm) 31
 - Core 2 Duo 14
- Different instructions taking different amount of time to execute
- Hardware to stall the pipeline if an instruction uses a result that is not ready

Real Machine Model cont.

- Most modern processors have multiple cores
 - Will deal with multicores next week
- Each core has multiple execution units (superscalar)
 - If the instruction sequence is efficient, multiple operations will happen in the same cycles
 - Even more important to have the right instruction sequence

Instruction Scheduling

- Reorder instructions so that pipeline stalls are minimized

Constraints On Scheduling

- Data dependencies
- Control dependencies
- Resource Constraints

Data Dependency between Instructions

- If two instructions access the same variable, they can be dependent
- Kind of dependencies
 - True: write \rightarrow read
 - Anti: read \rightarrow write
 - Output: write \rightarrow write
- What to do if two instructions are dependent.
 - The order of execution cannot be reversed
 - Reduce the possibilities for scheduling

Computing Dependencies

- For basic blocks, compute dependencies by walking through the instructions
- Identifying register dependencies is simple
 - is it the same register?
- For memory accesses
 - simple: $\text{base} + \text{offset}_1 \neq \text{base} + \text{offset}_2$
 - data dependence analysis: $a[2i] \neq a[2i+1]$
 - interprocedural analysis: $\text{global} \neq \text{parameter}$
 - pointer alias analysis: $p1 \rightarrow \text{foo} \neq p2 \rightarrow \text{foo}$

Representing Dependencies

- Using a dependence DAG, one per basic block
- Nodes are instructions, edges represent dependencies

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1: $r2 = *(r1 + 4)$

2: $r3 = *(r1 + 8)$

3: $r4 = r2 + r3$

4: $r5 = r2 - 1$

Representing Dependencies

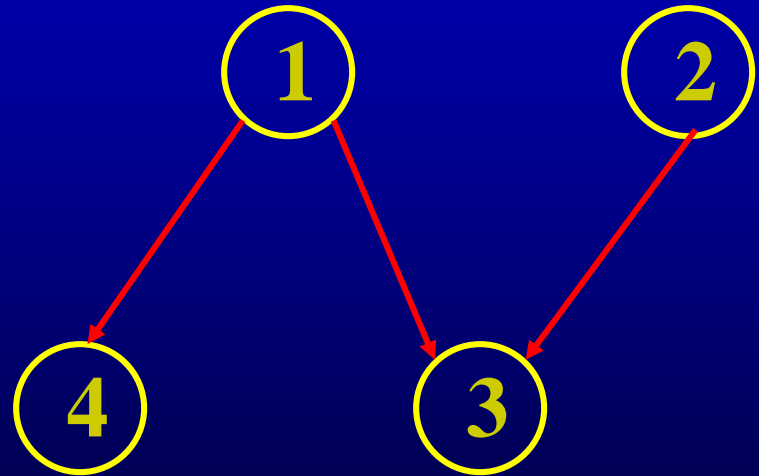
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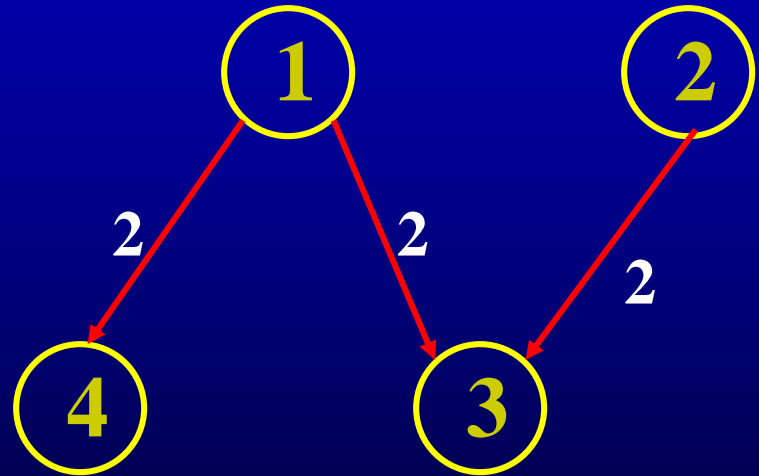
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- Edge is labeled with Latency:
 - $v(i \rightarrow j)$ = delay required between initiation times of i and j minus the execution time required by i

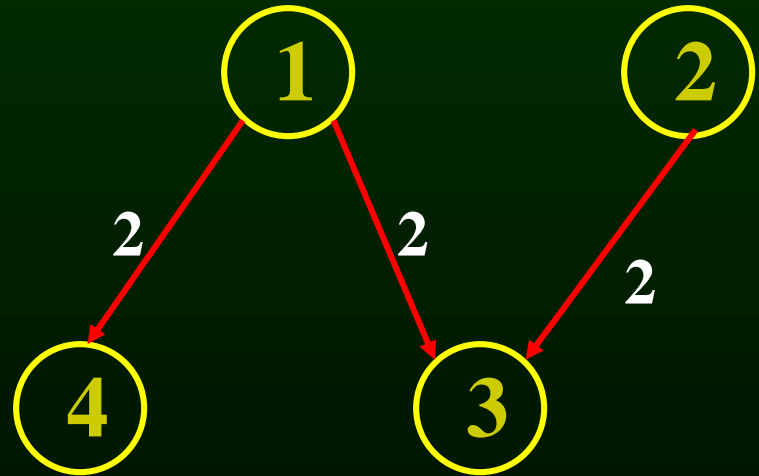
Example

1: $r2 = *(r1 + 4)$

2: $r3 = *(r2 + 4)$

3: $r4 = r2 + r3$

4: $r5 = r2 - 1$



Another Example

1: $r2 = *(r1 + 4)$

2: $*(r1 + 4) = r3$

3: $r3 = r2 + r3$

4: $r5 = r2 - 1$

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②

④

③

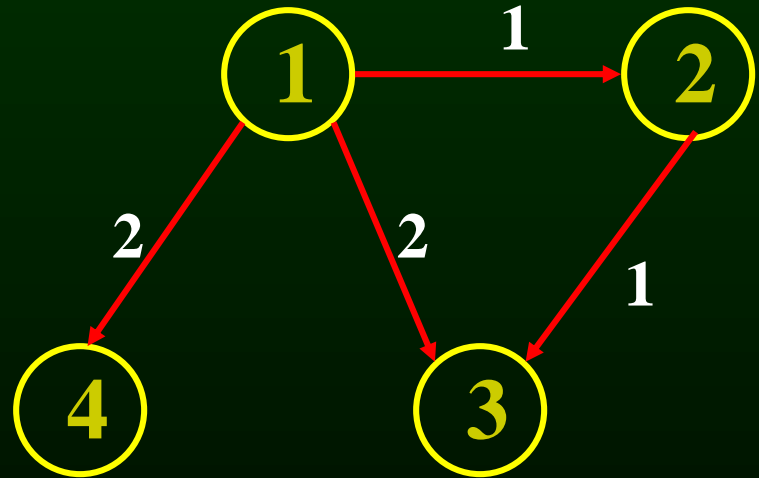
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2: $*(r1 + 4) = r3$

3: $r3 = r2 + r3$

4: $r5 = r2 - 1$



Control Dependencies and Resource Constraints

- For now, let's only worry about basic blocks
- For now, let's look at simple pipelines

Example

```
1: lea    var_a, %rax
2: add    $4, %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: add    %r10, 8(%rsp)
6: and    16(%rsp), %rbx
7: imul   %rax, %rbx
```

Example

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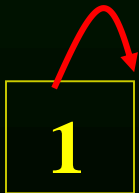
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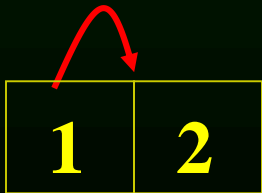
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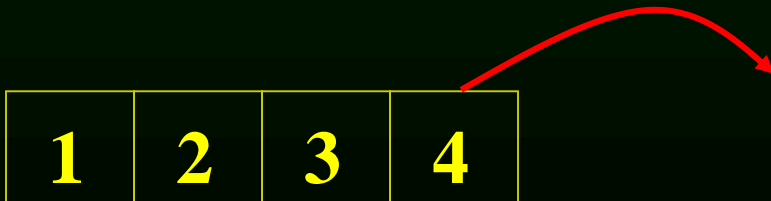
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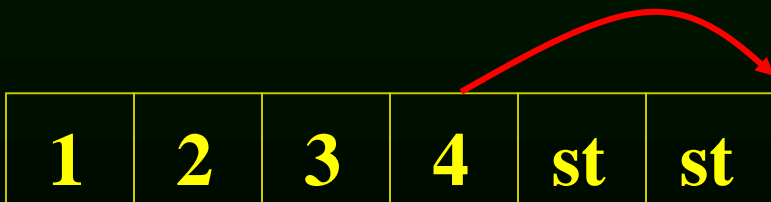
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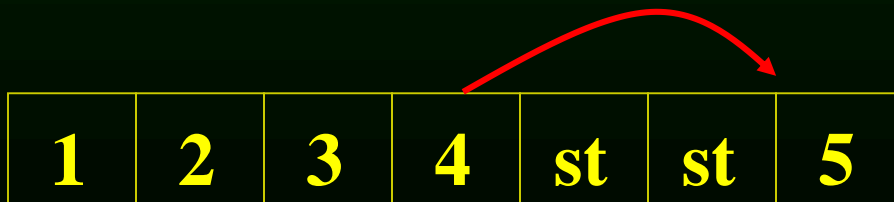
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1	2	3	4	st	st	5
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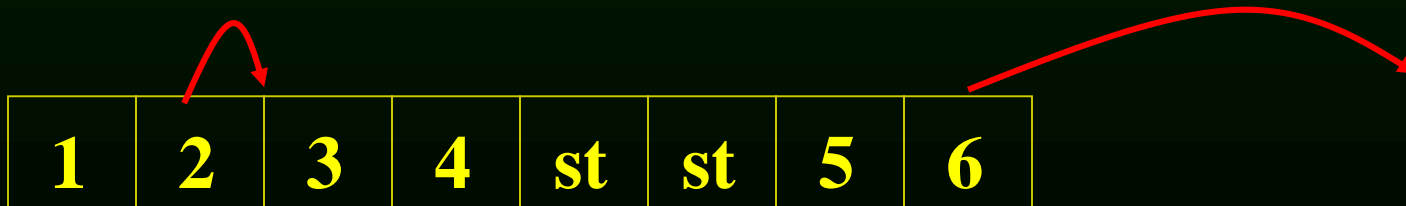
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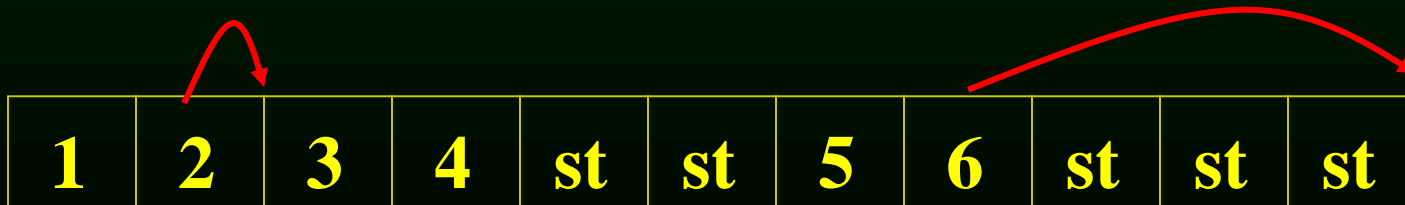
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List Scheduling Algorithm

- Idea
 - Do a topological sort of the dependence DAG
 - Consider when an instruction can be scheduled without causing a stall
 - Schedule the instruction if it causes no stall and all its predecessors are already scheduled
- Optimal list scheduling is NP-complete
 - Use heuristics when necessary

List Scheduling Algorithm

- Create a dependence DAG of a basic block
- Topological Sort

READY = nodes with no predecessors

Loop until READY is empty

 Schedule each node in READY when no stalling

 Update READY

Heuristics for selection

- Heuristics for selecting from the READY list
 - pick the node with the longest path to a leaf in the dependence graph
 - pick a node with most immediate successors
 - pick a node that can go to a less busy pipeline (in a superscalar)

Heuristics for selection

- pick the node with the longest path to a leaf in the dependence graph
- Algorithm (for node x)
 - If no successors $d_x = 0$
 - $d_x = \text{MAX}(d_y + c_{xy})$ for all successors y of x
 - reverse breadth-first visitation order

Heuristics for selection

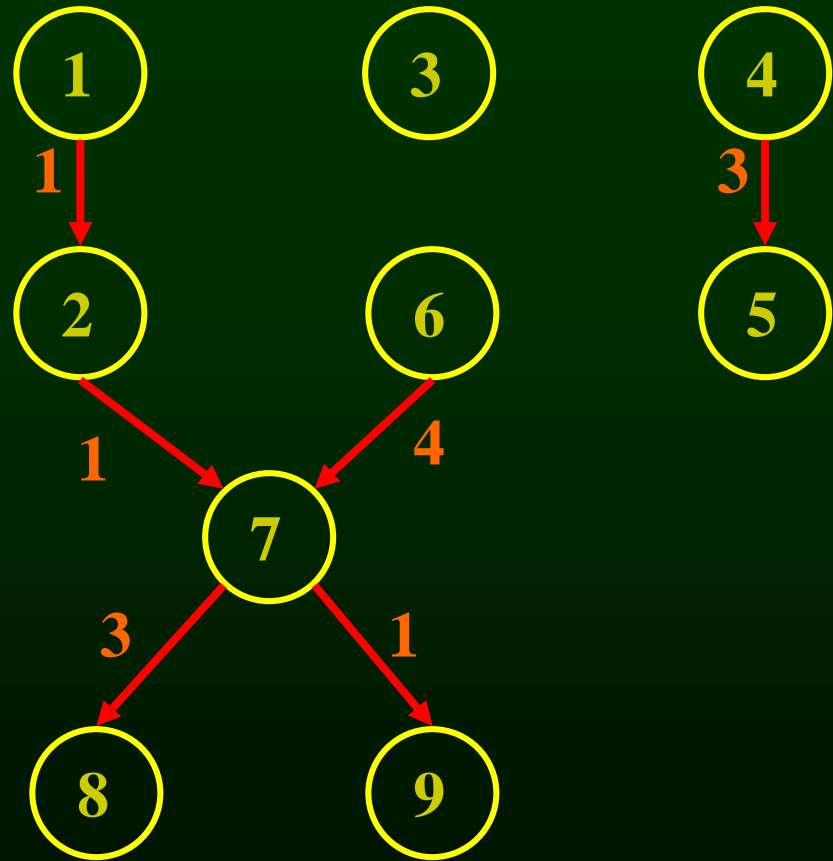
- pick a node with most immediate successors
- Algorithm (for node x):
 - f_x = number of successors of x

Example

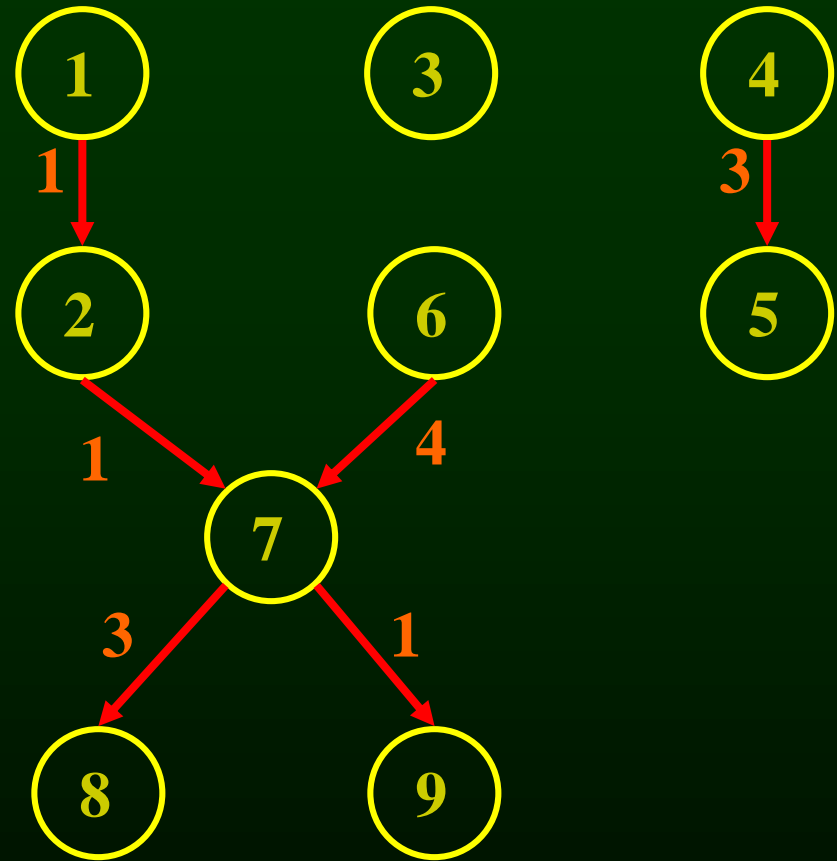
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7:	imul %rax, %rbx	3 cycles
8:	mov %rbx, 16(%rsp)	
9:	lea var_b, %rax	

Example

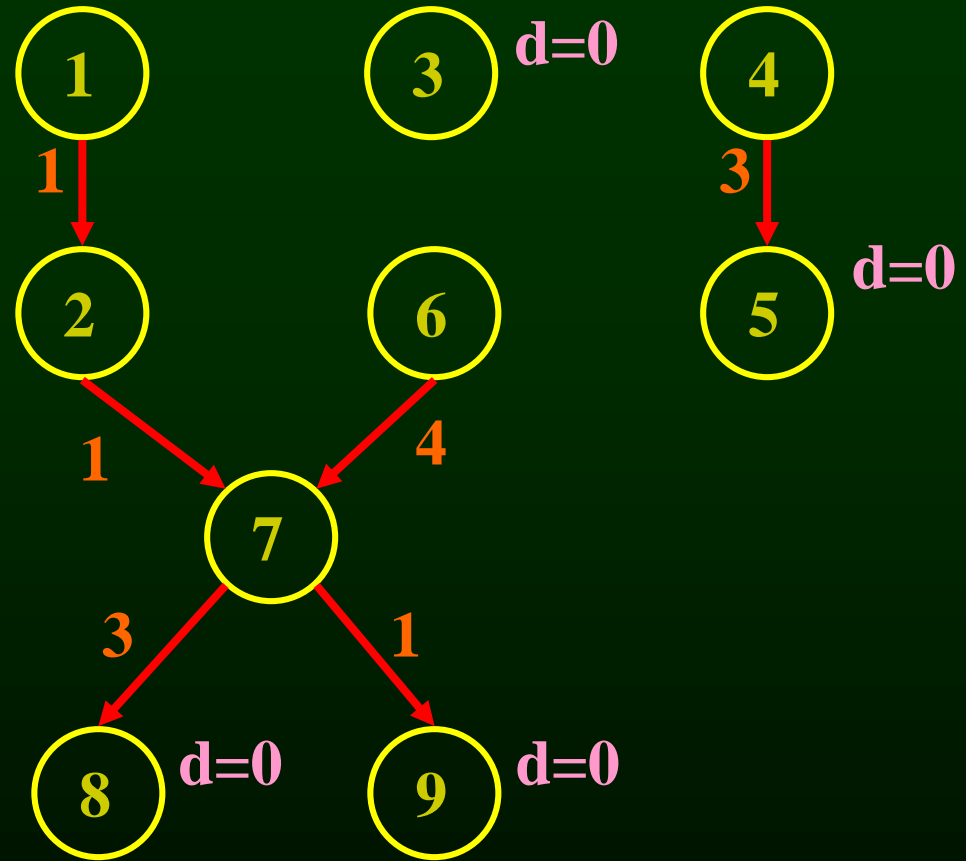
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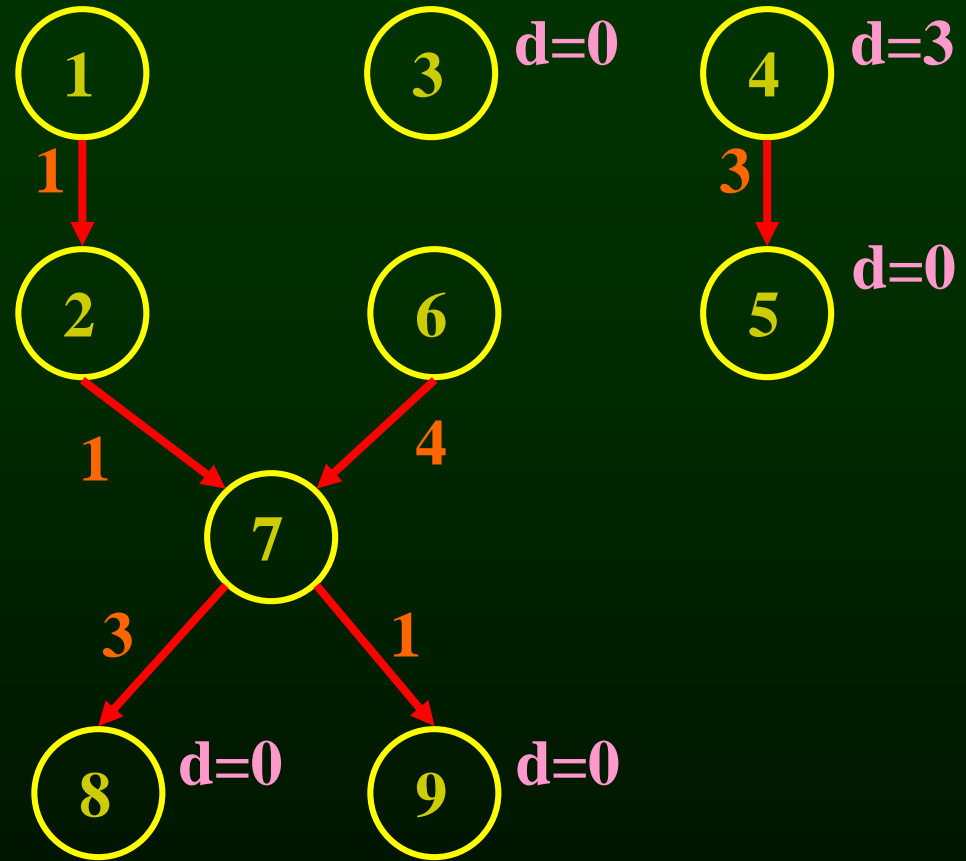
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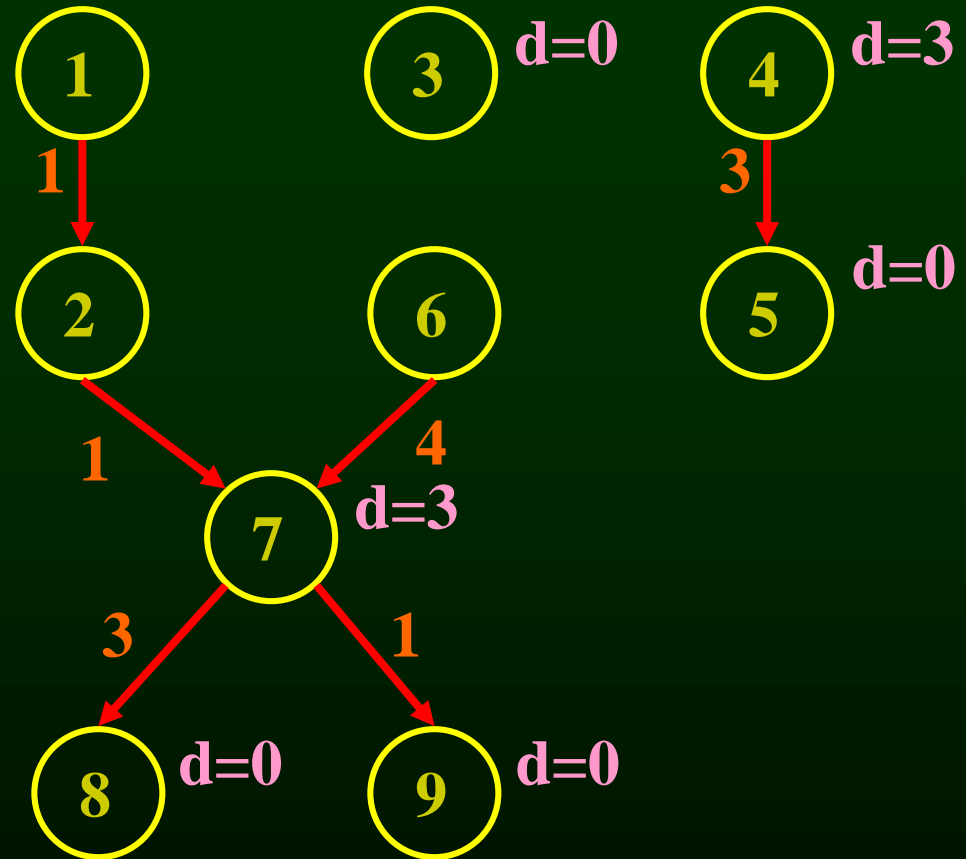
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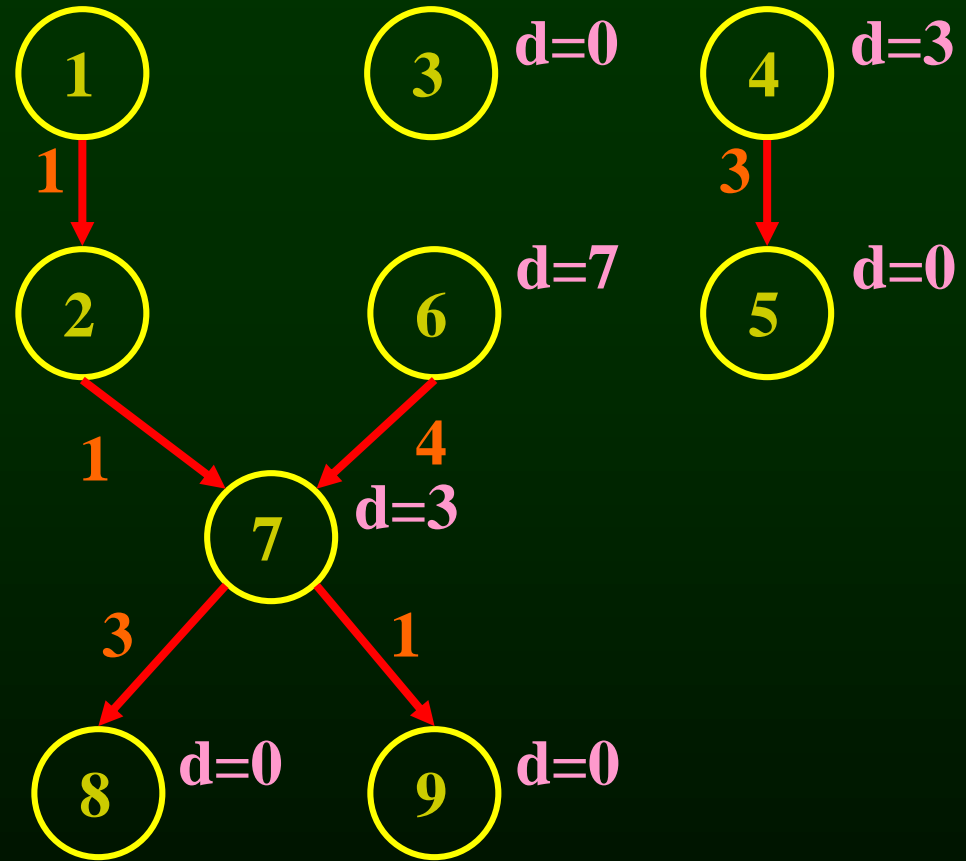
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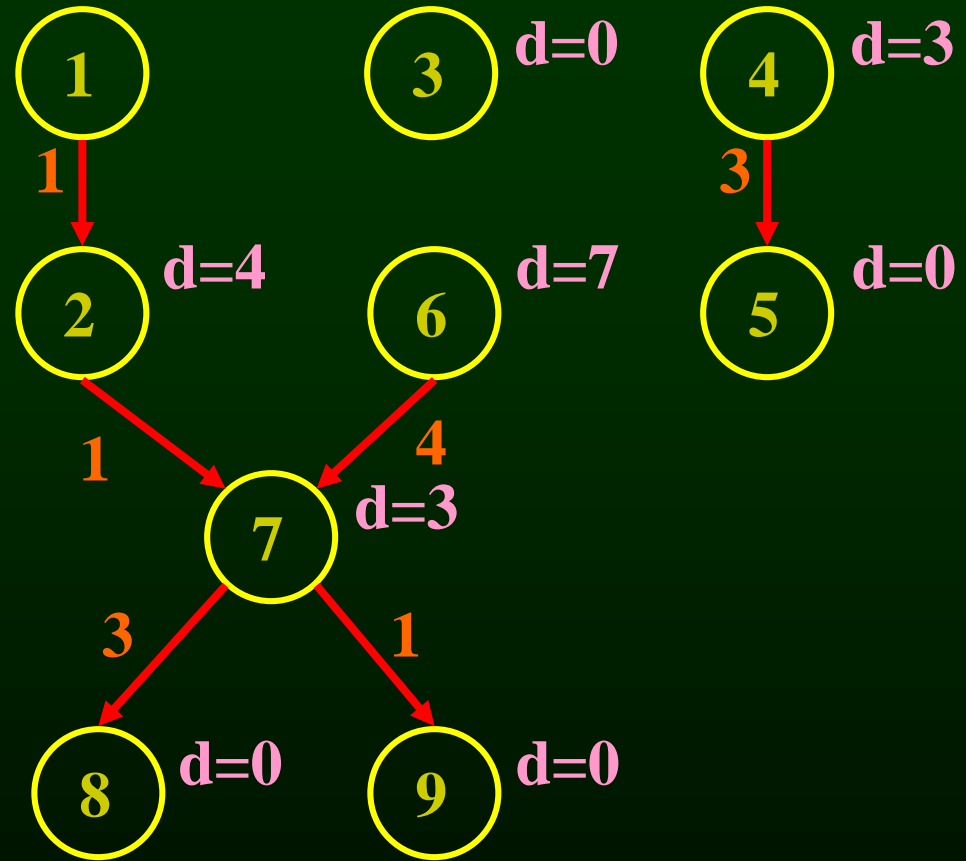
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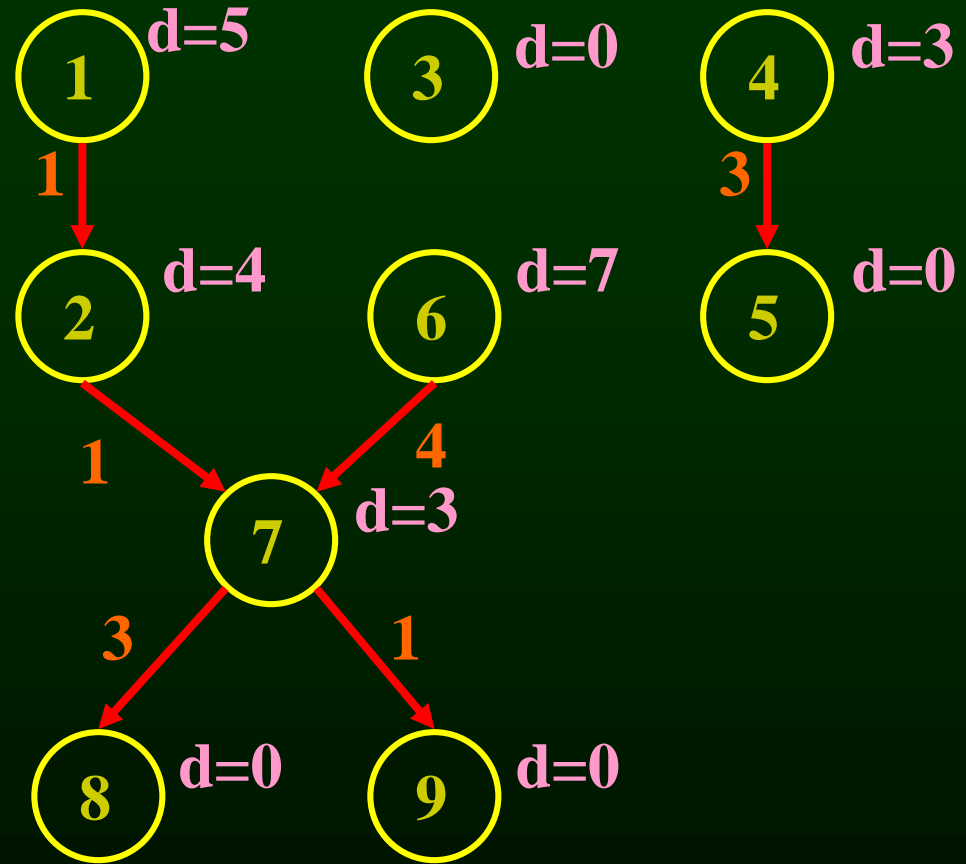
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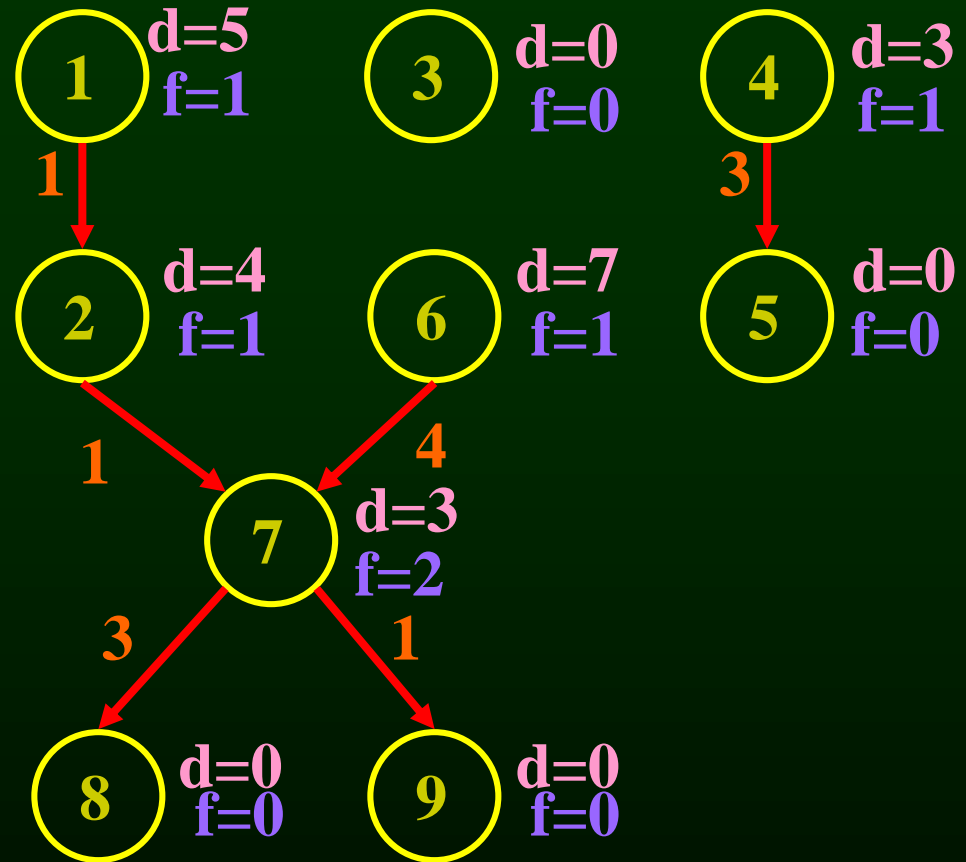
Example



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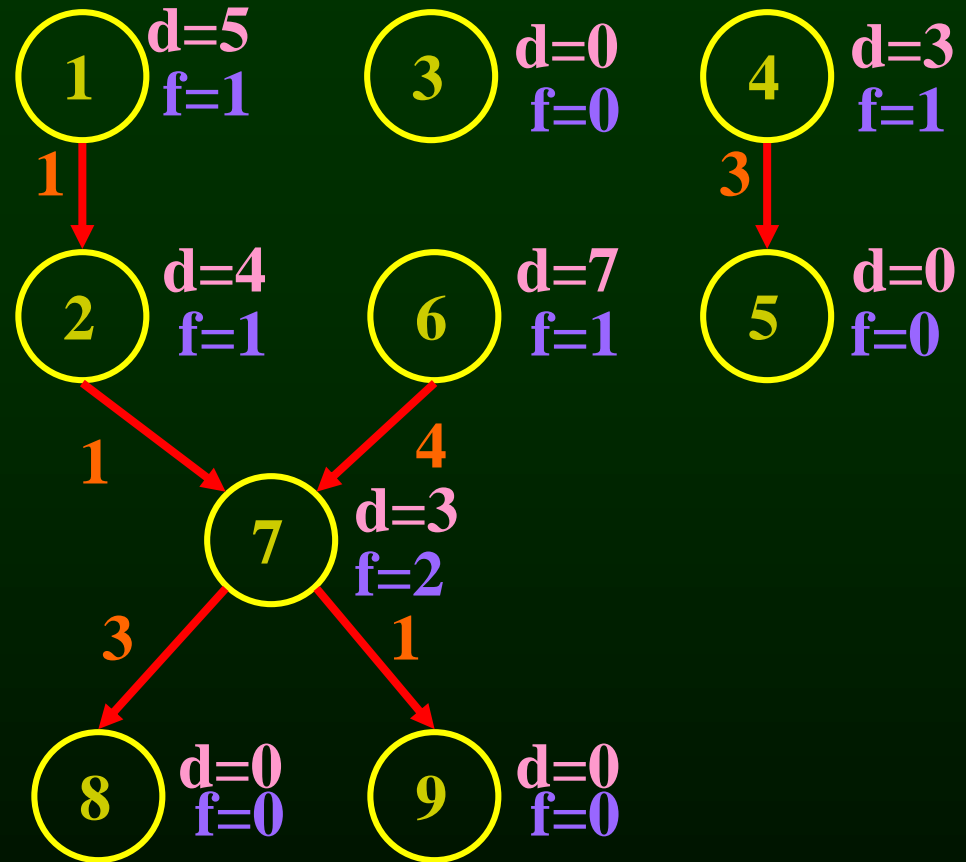


Example



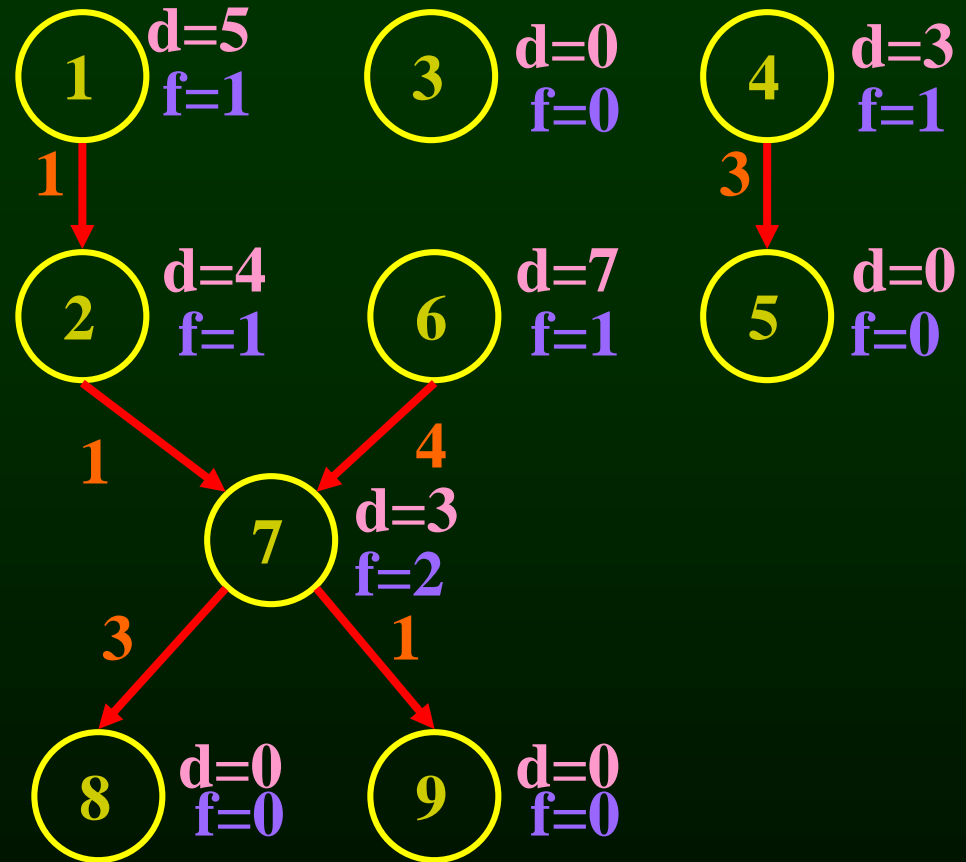
Example

READY = { }



Example

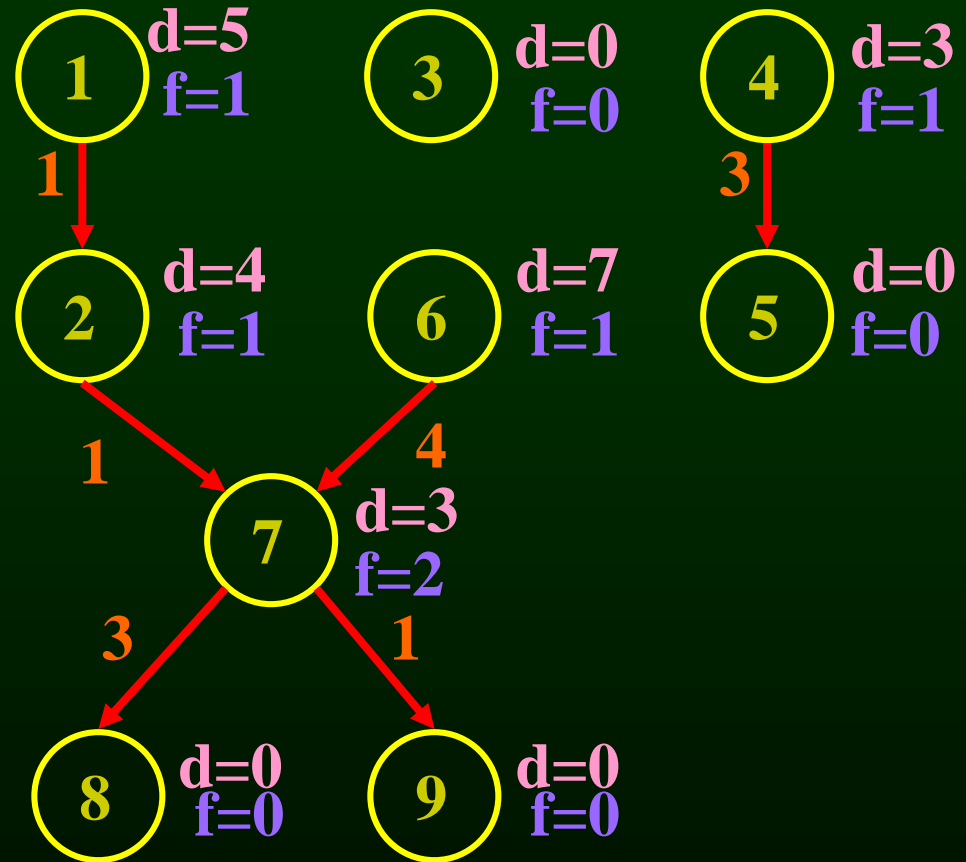
READY = { 1, 3, 4, 6 }



Example

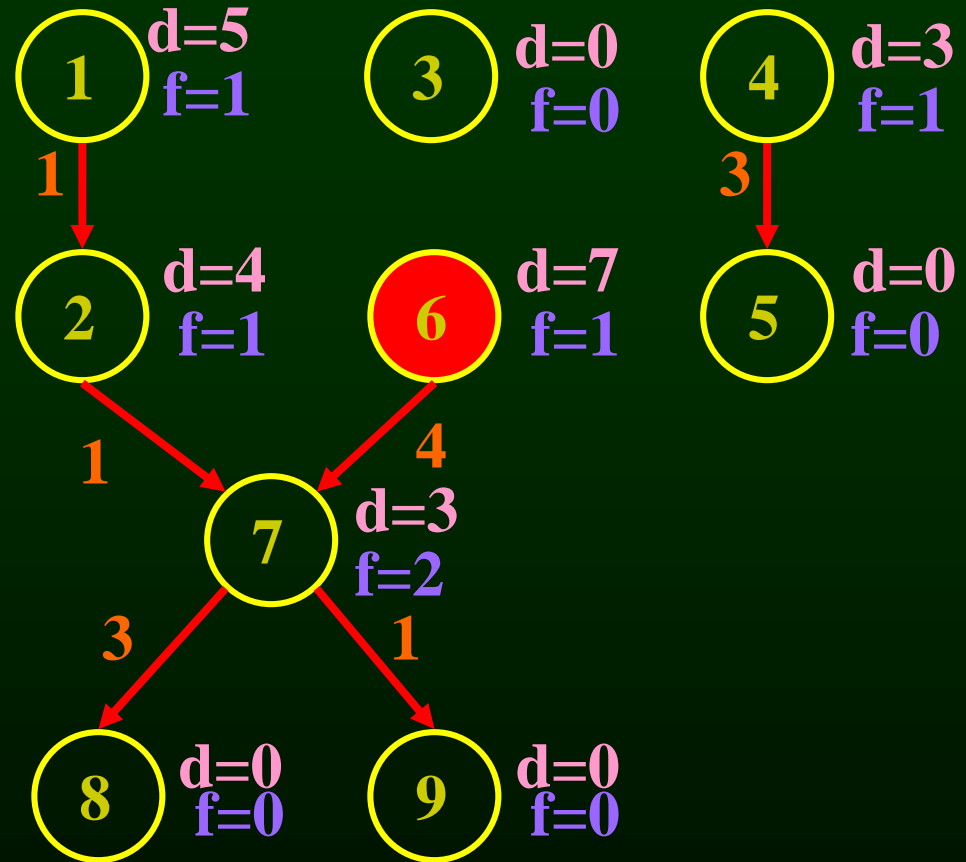
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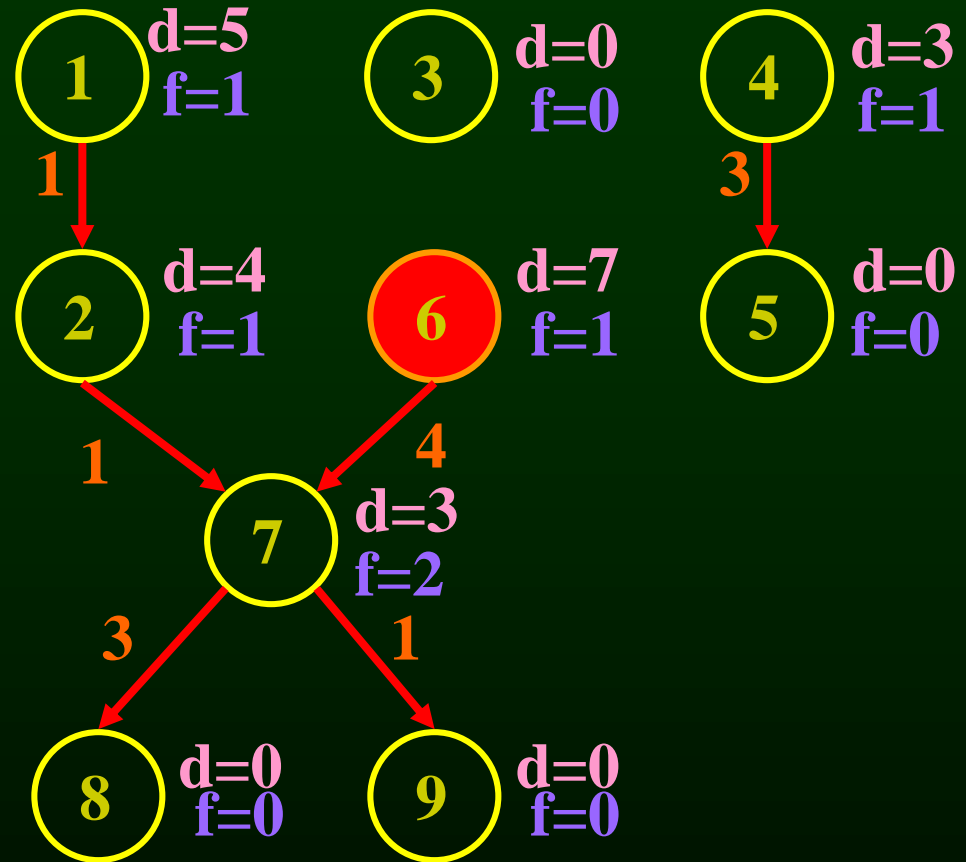
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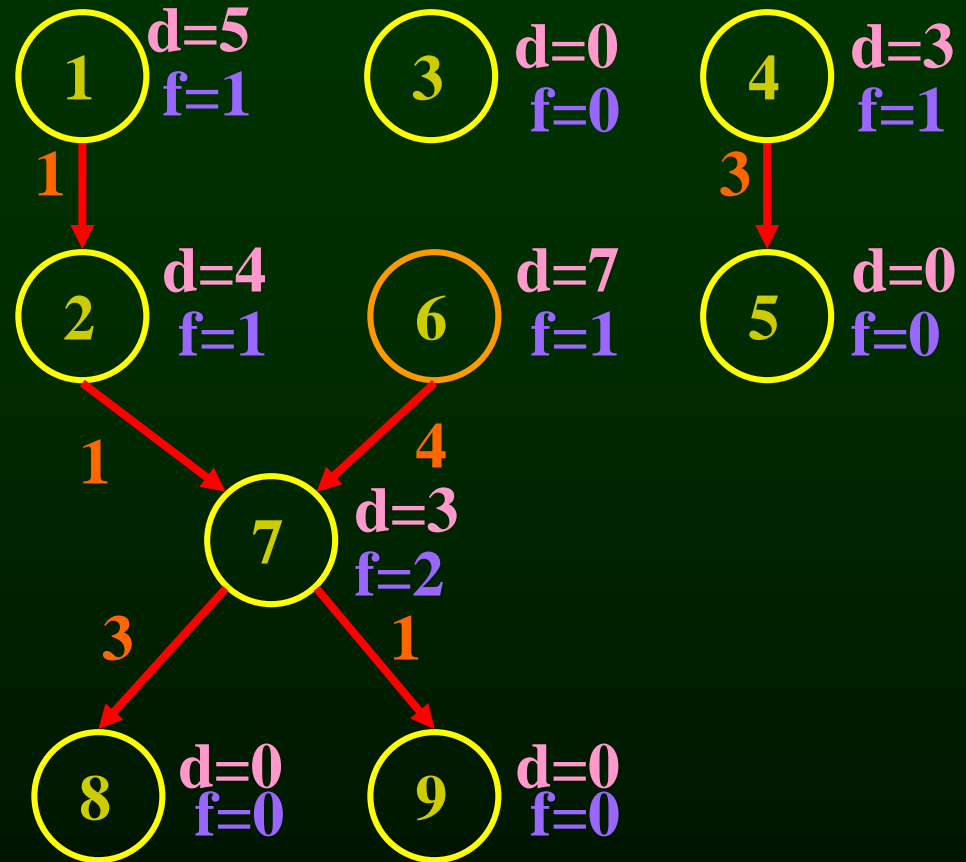
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6

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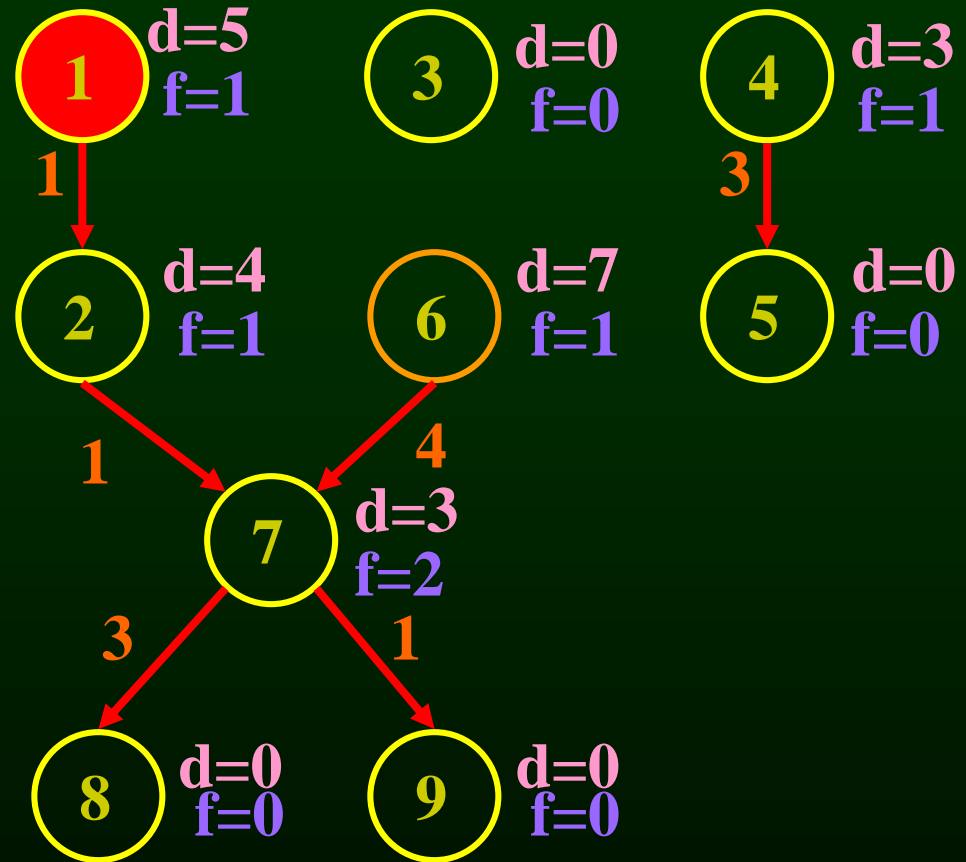
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6

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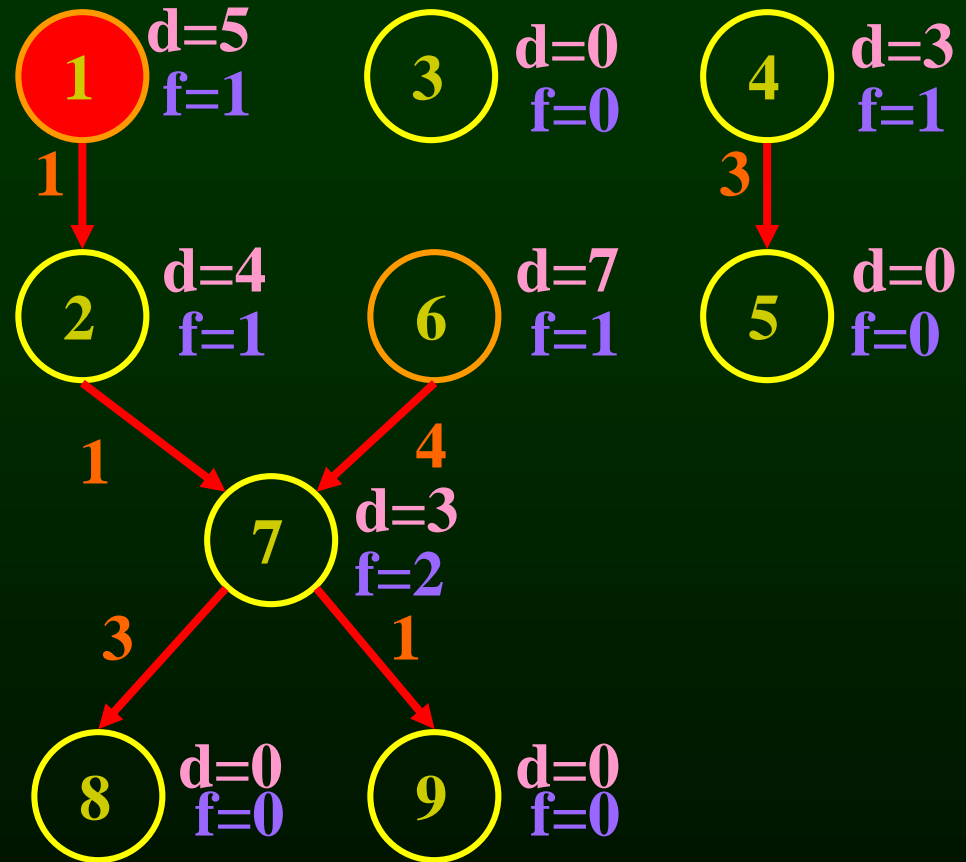
READY = { 1, 4, 3 }



6

Example

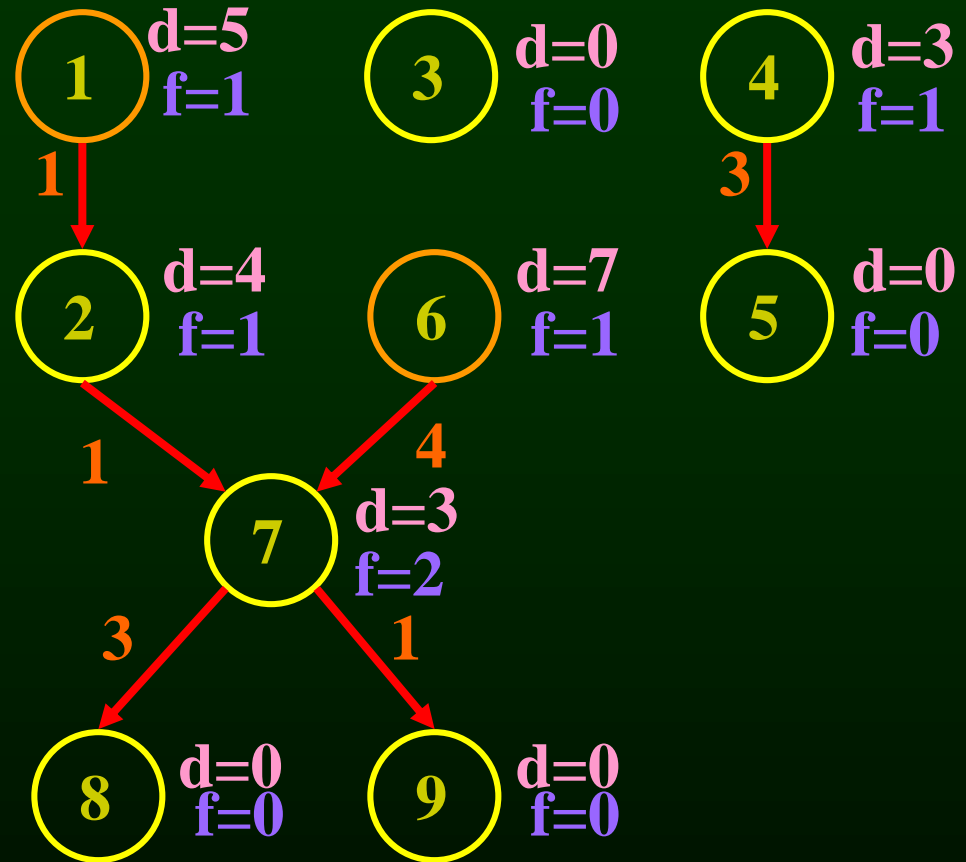
READY = { 1, 4, 3 }



6	1
---	---

Example

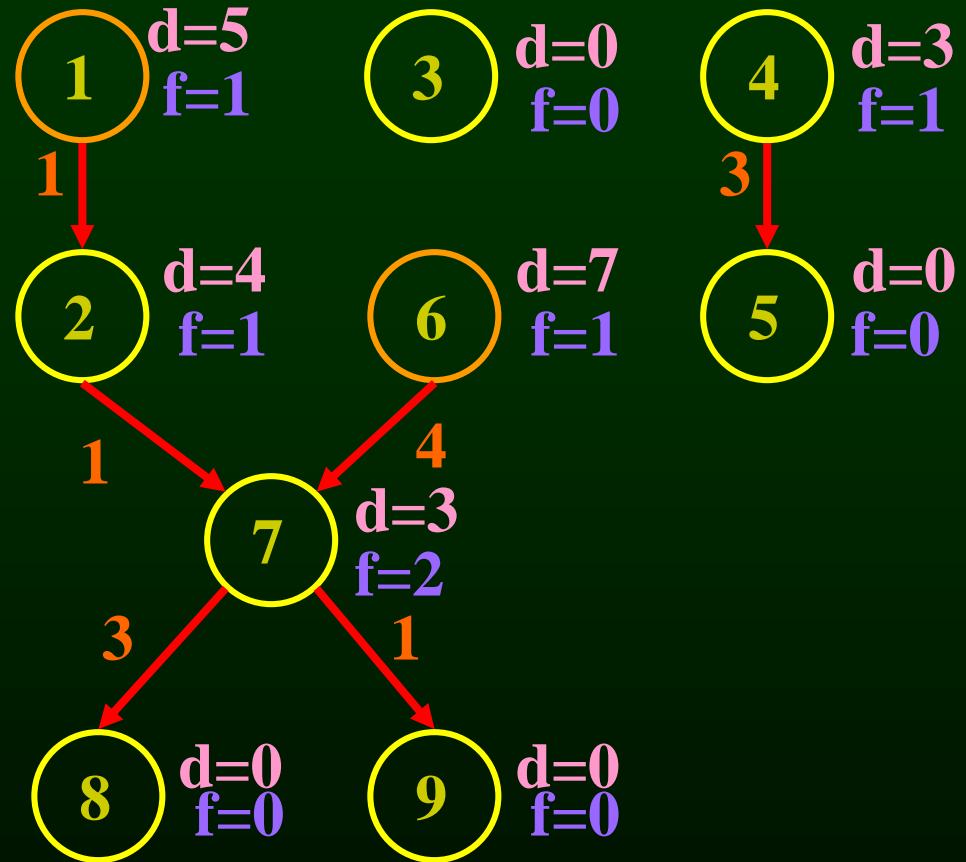
2
READY = { 4, 3 }



6	1
---	---

Example

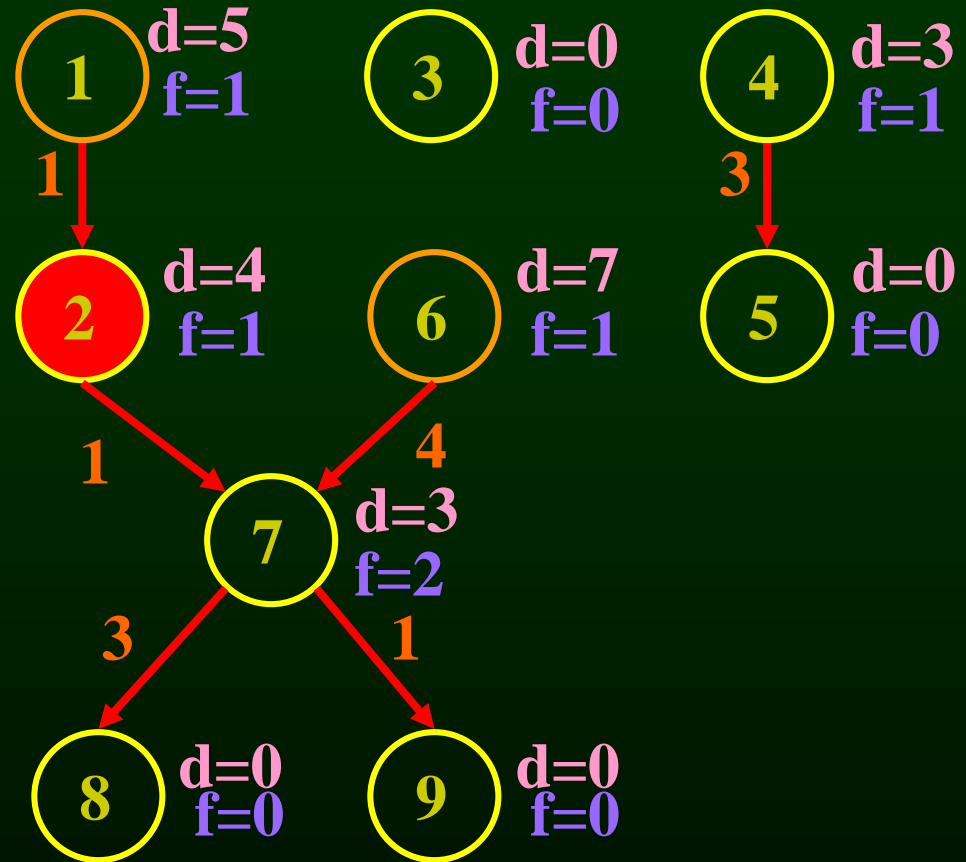
READY = { 2, 4, 3 }



6	1
---	---

Example

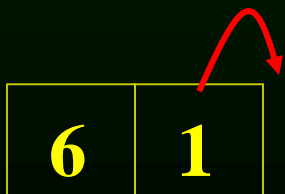
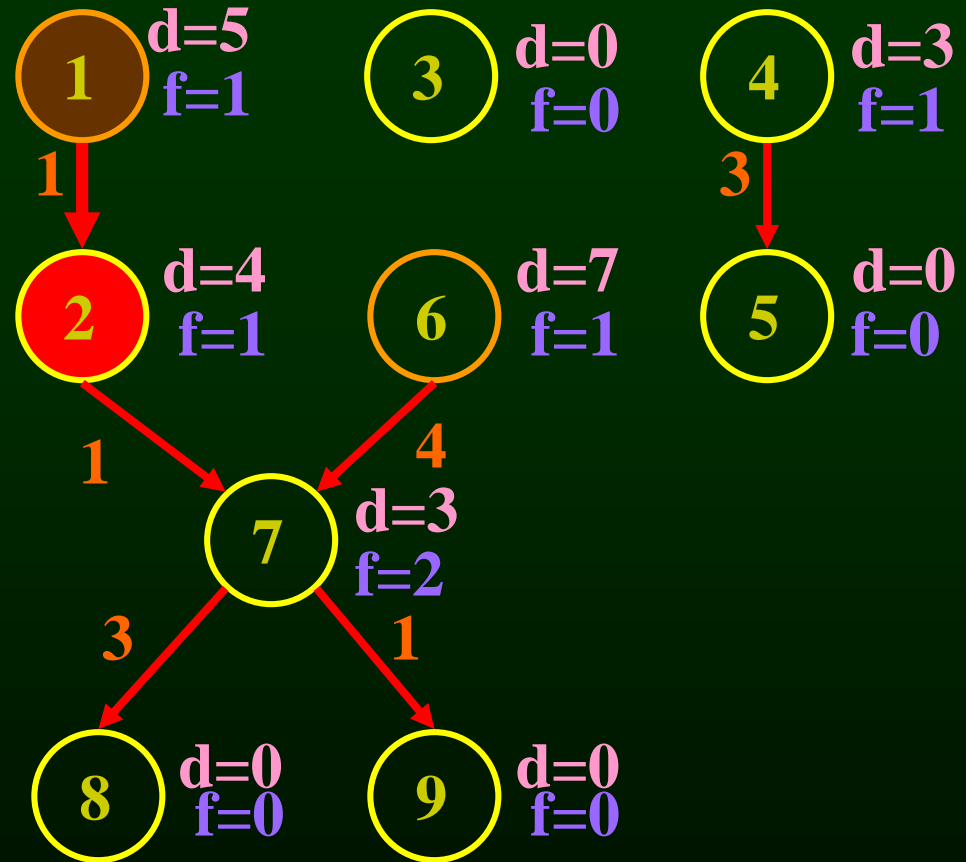
READY = { 2, 4, 3 }



6	1
---	---

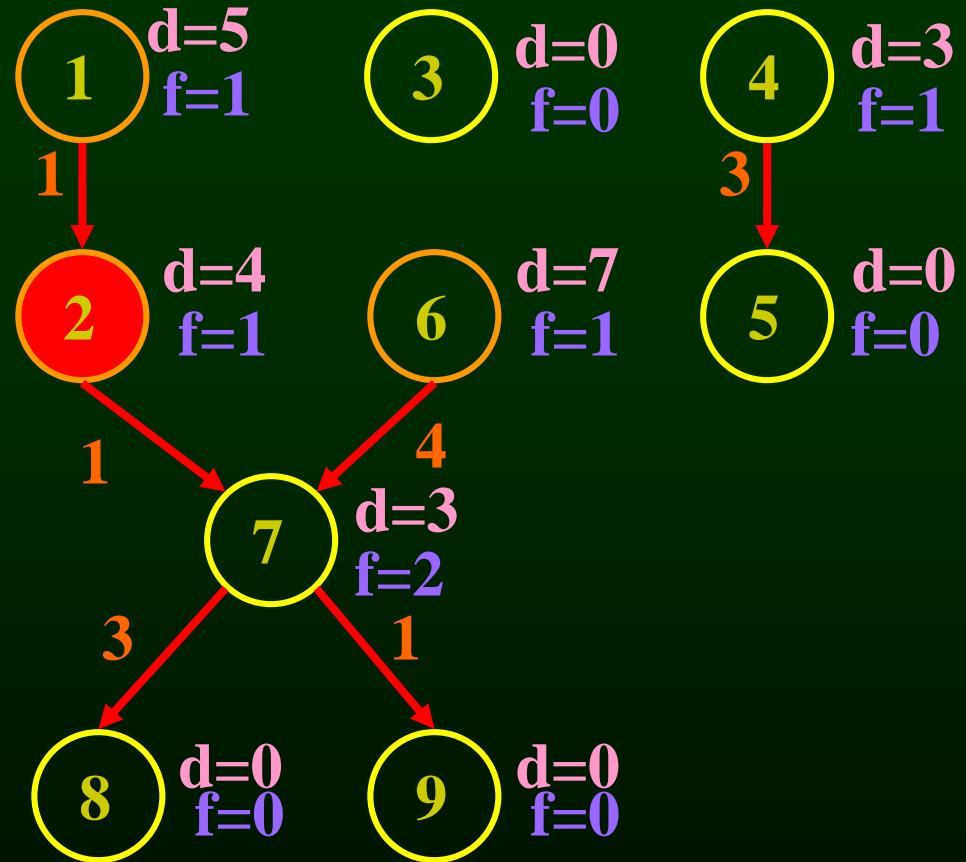
Example

READY = { 2, 4, 3 }



Example

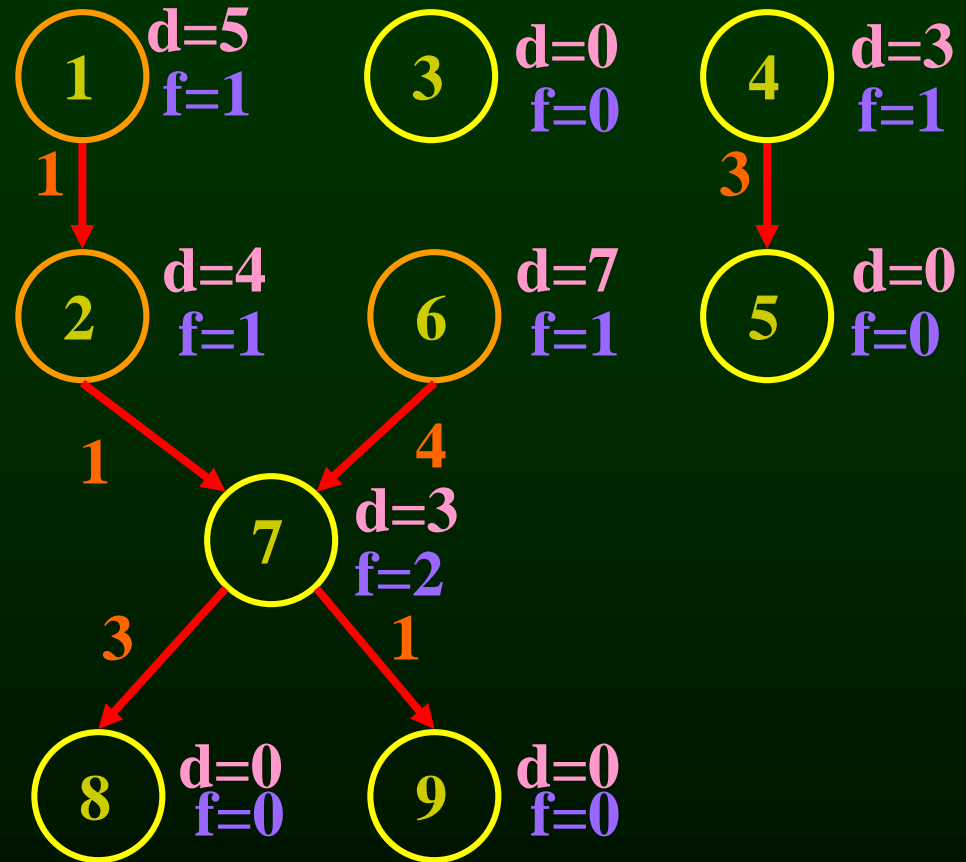
READY = { 2, 4, 3 }



6	1	2
---	---	---

Example

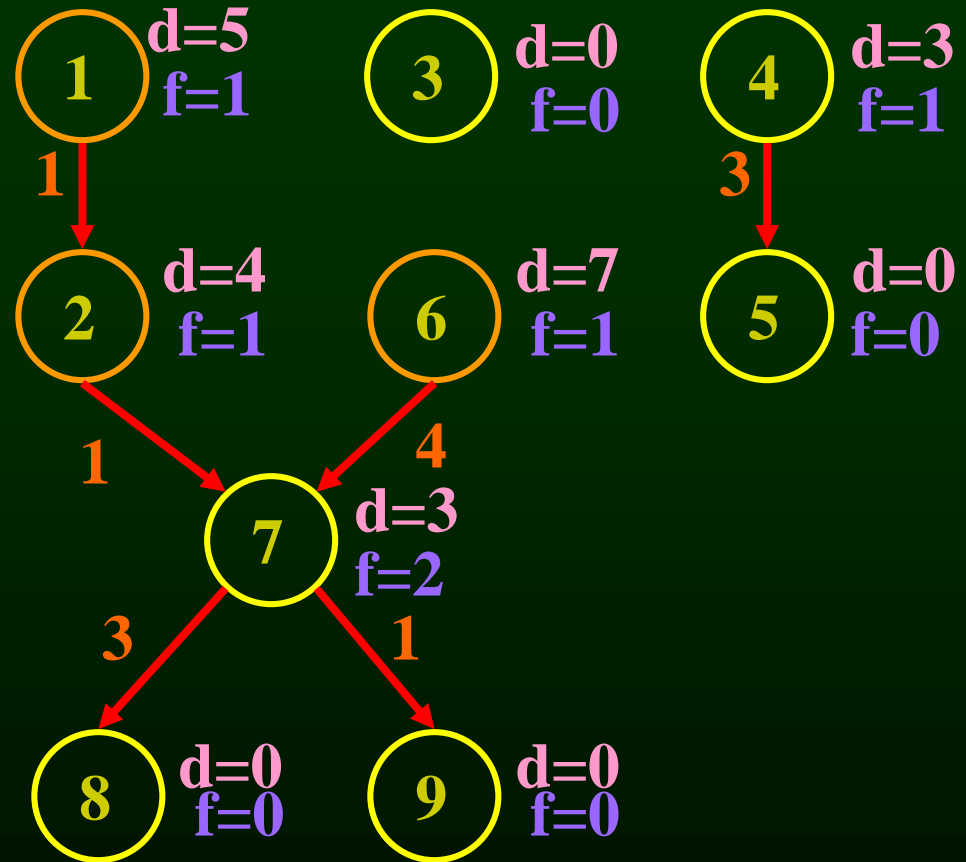
7
READY = { 4, 3 }



6	1	2
---	---	---

Example

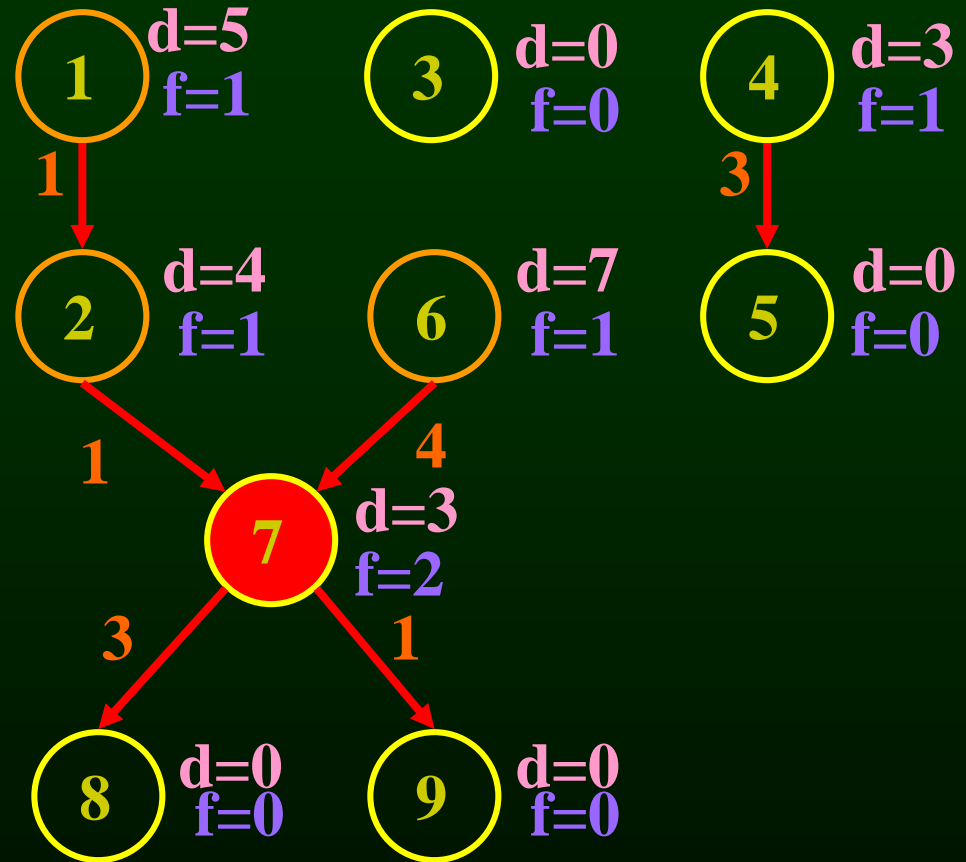
READY = { 7, 4, 3 }



6	1	2
---	---	---

Example

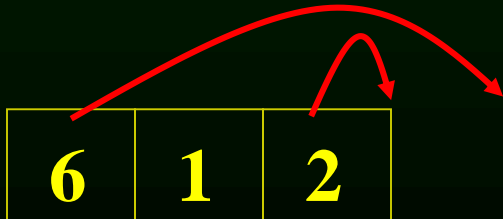
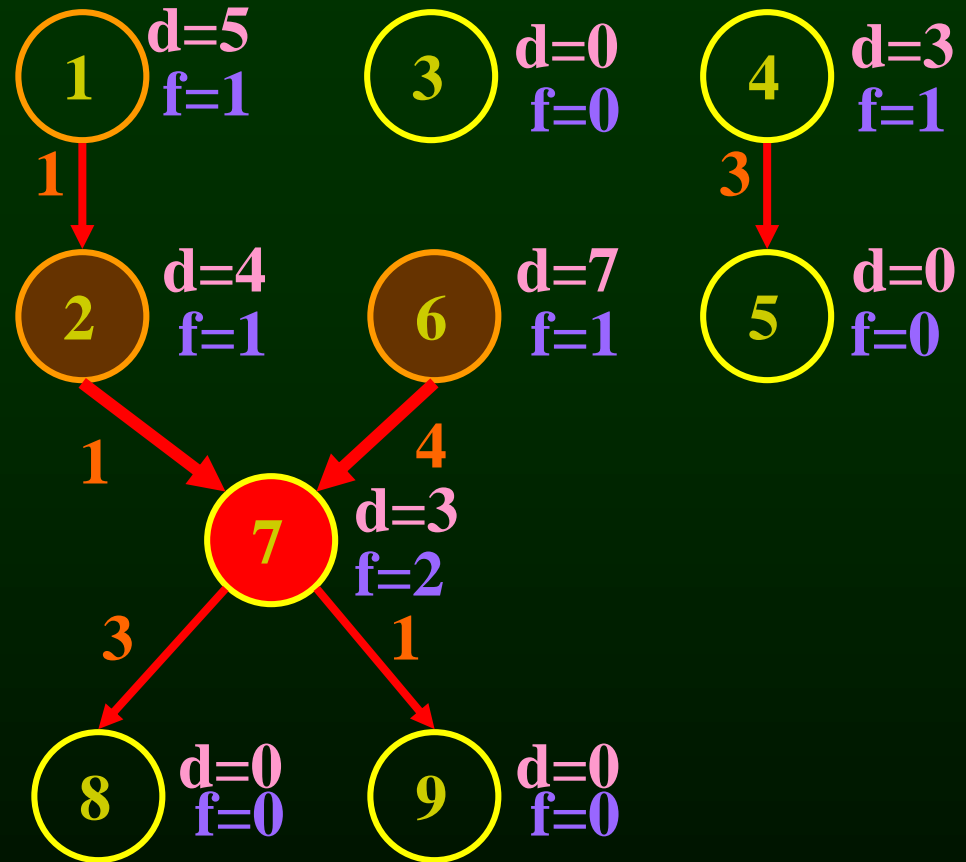
READY = { 7, 4, 3 }



6	1	2
---	---	---

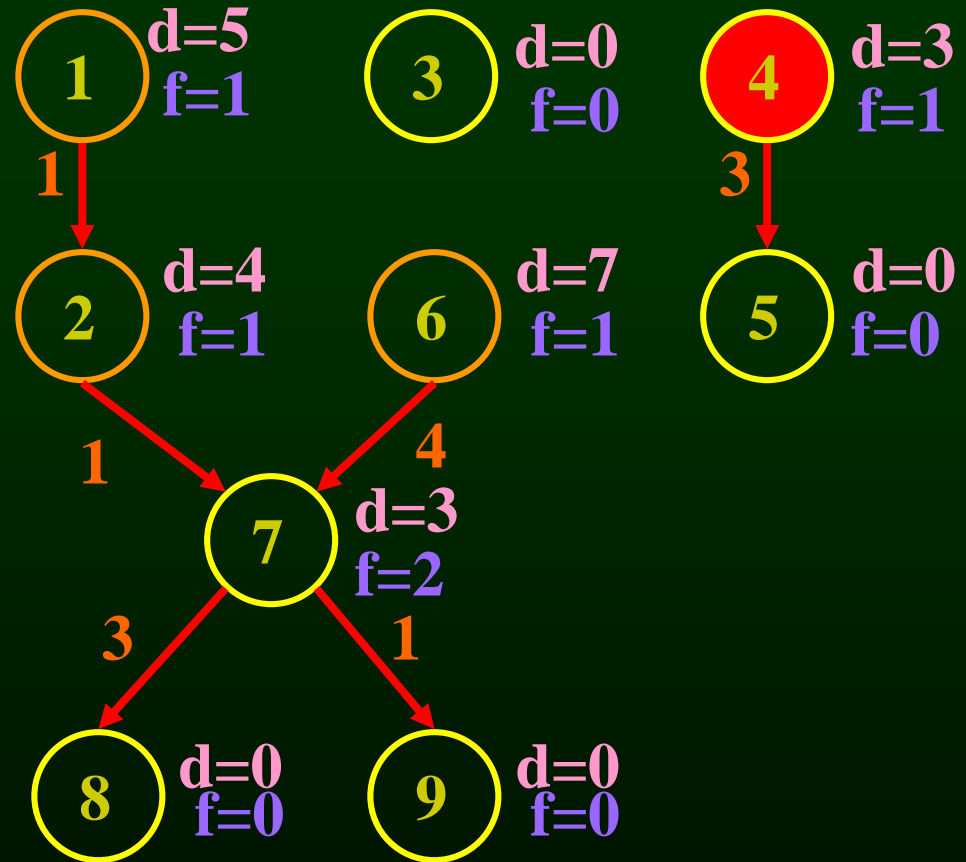
Example

READY = { 7, 4, 3 }



Example

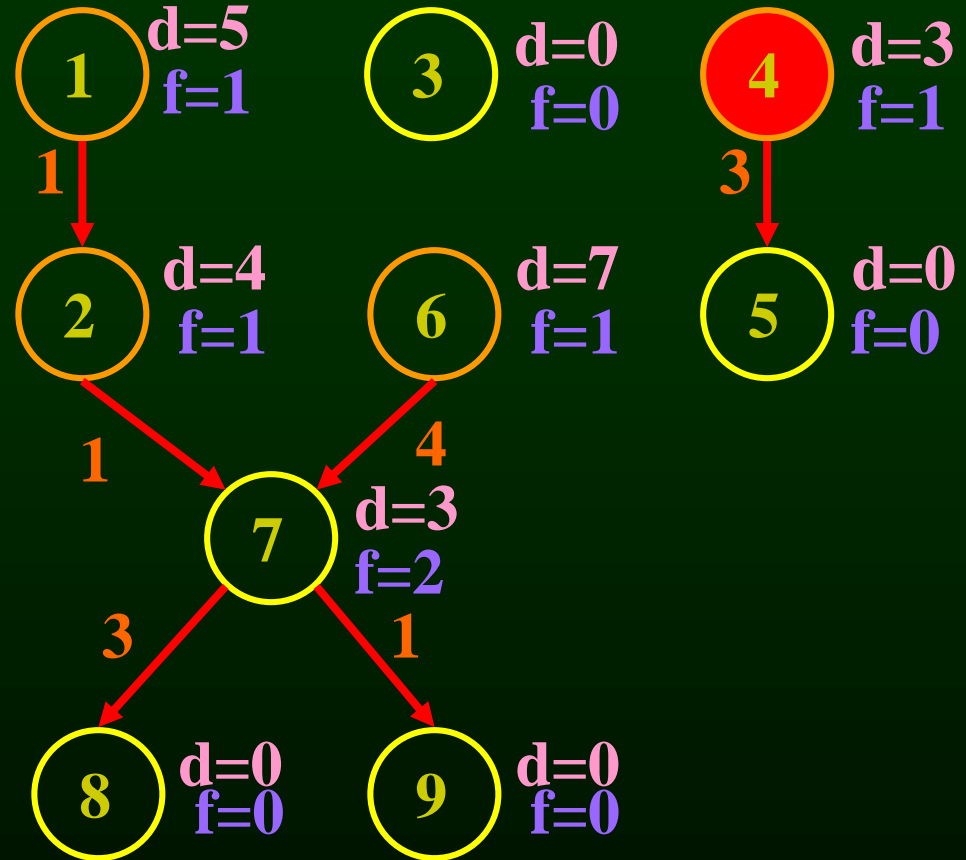
READY = { 7, 4, 3 }



6	1	2
---	---	---

Example

READY = { 7, 4, 3 }

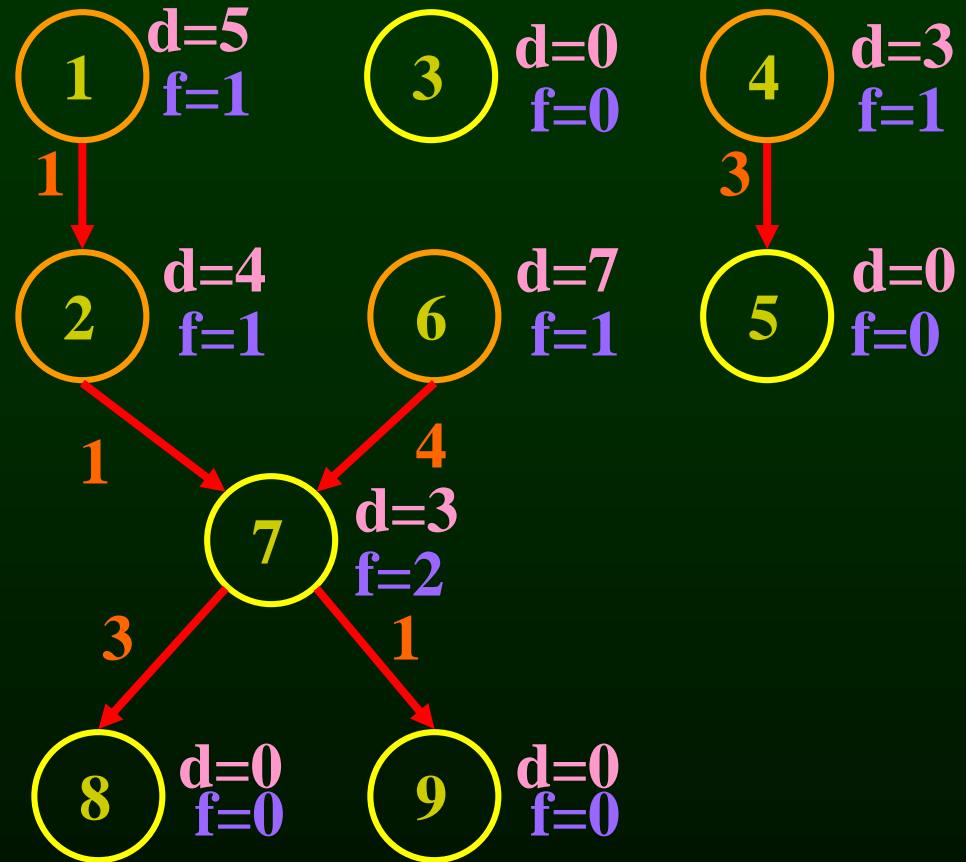


6	1	2	4
---	---	---	---

Example

5

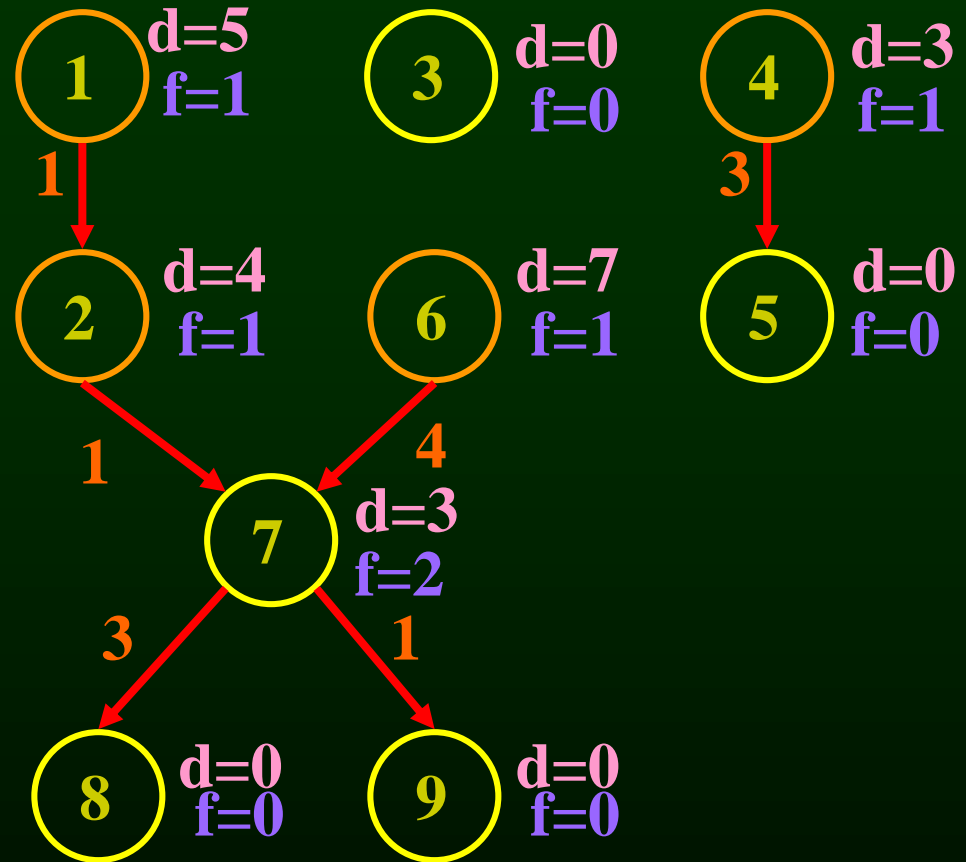
READY = { 7, 3 }



6	1	2	4
---	---	---	---

Example

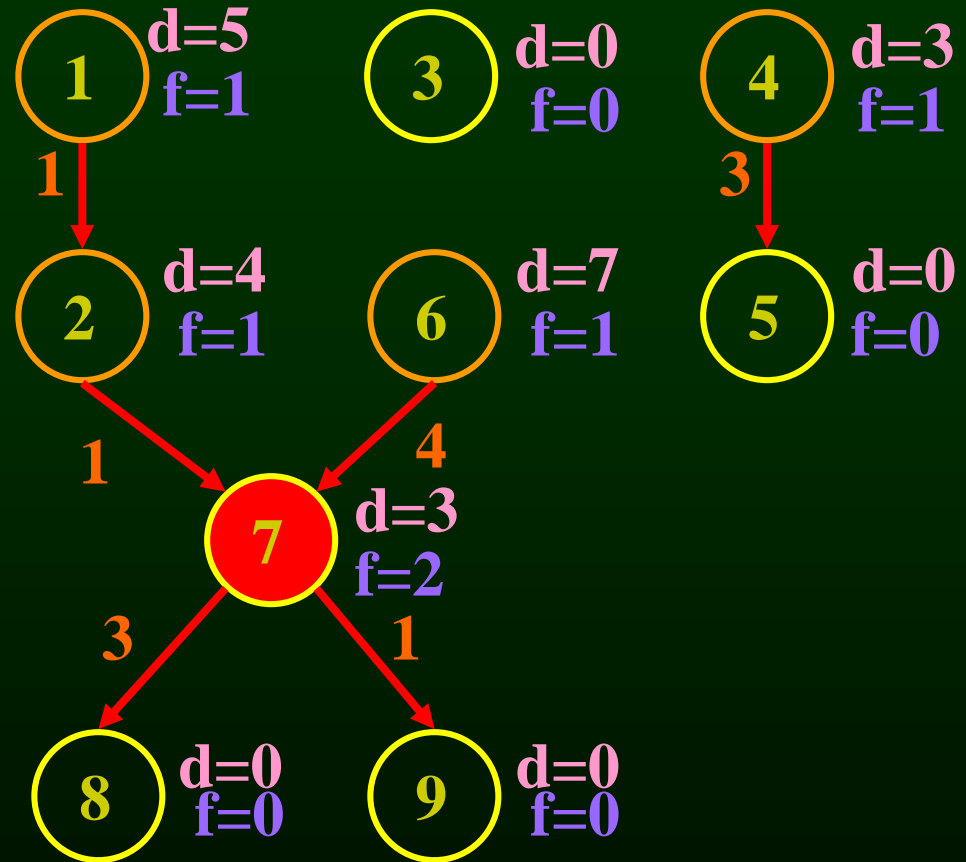
READY = { 7, 3, 5 }



6	1	2	4
---	---	---	---

Example

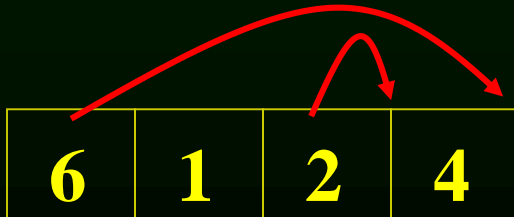
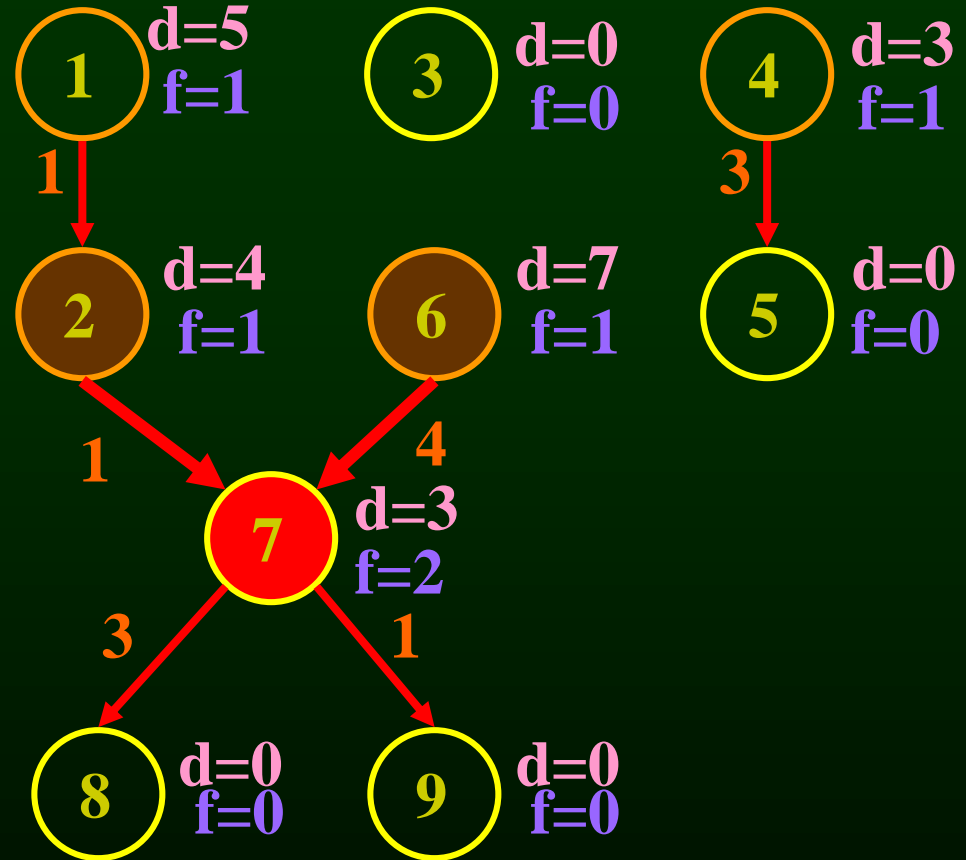
READY = { 7, 3, 5 }



6	1	2	4
---	---	---	---

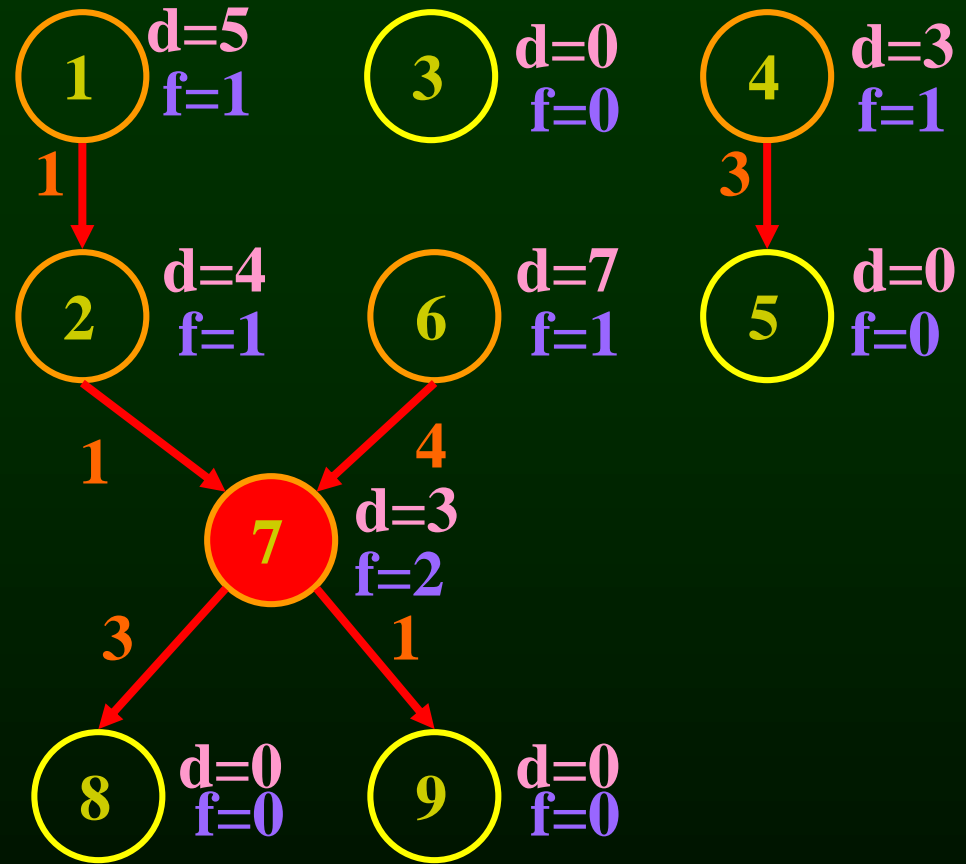
Example

READY = { 7, 3, 5 }



Example

READY = { 7, 3, 5 }

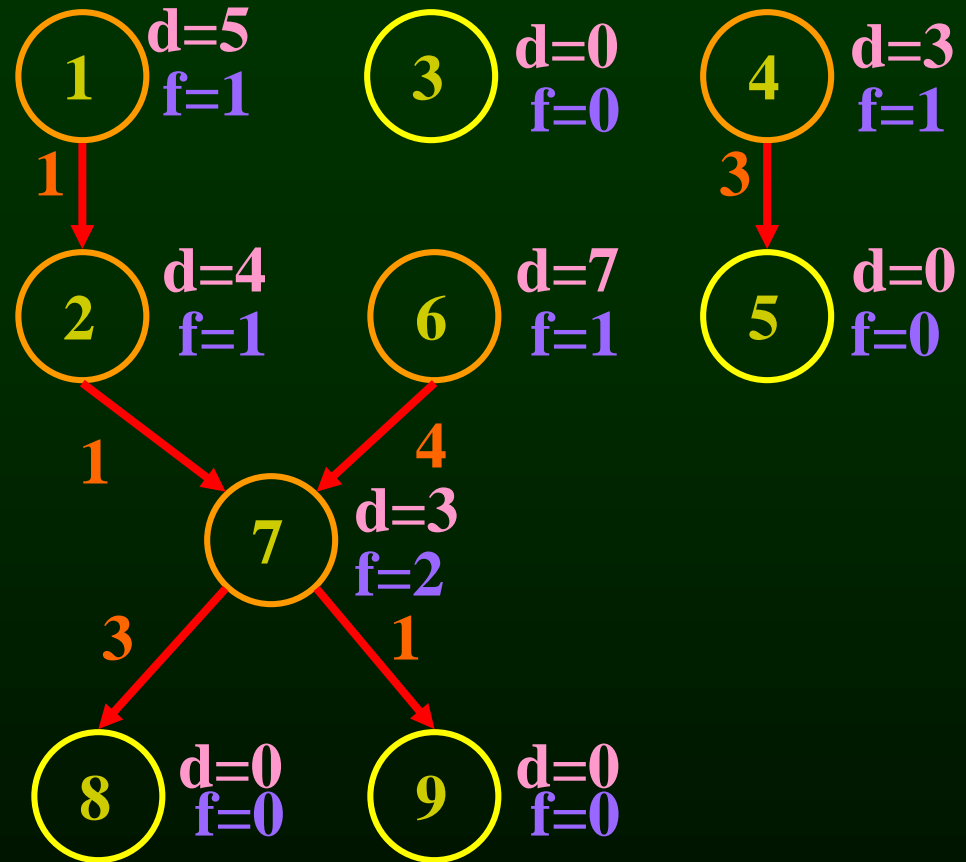


6	1	2	4	7
---	---	---	---	---

Example

8, 9

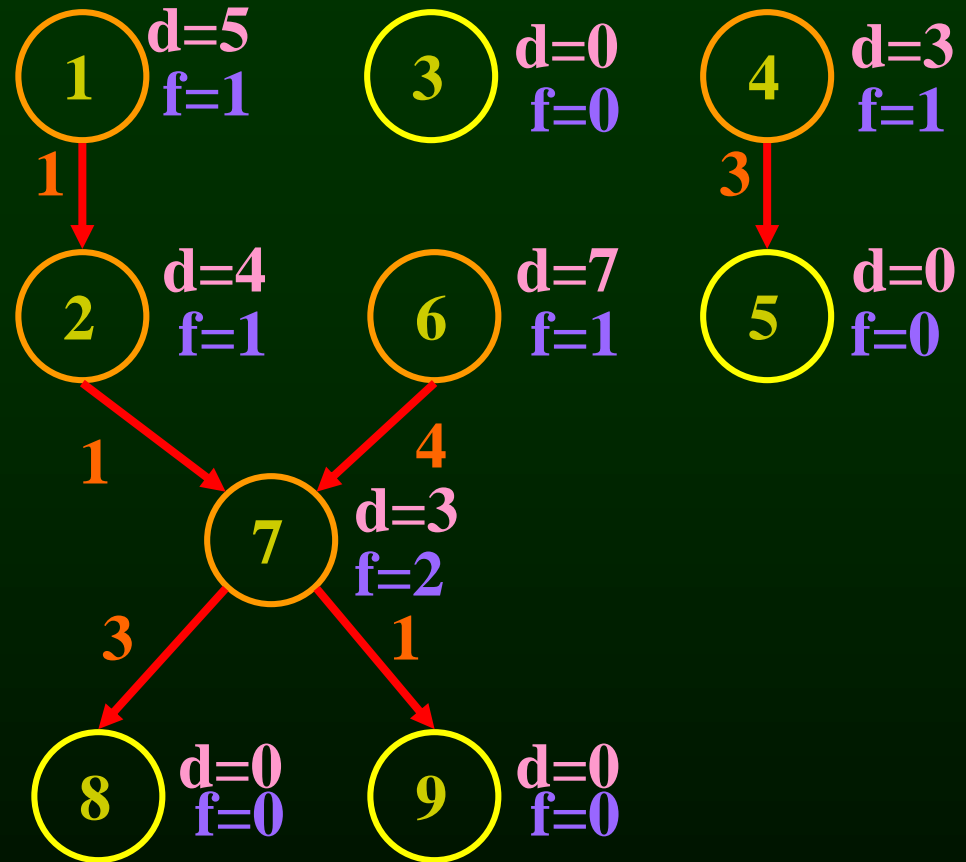
READY = { 3, 5 }



6	1	2	4	7
---	---	---	---	---

Example

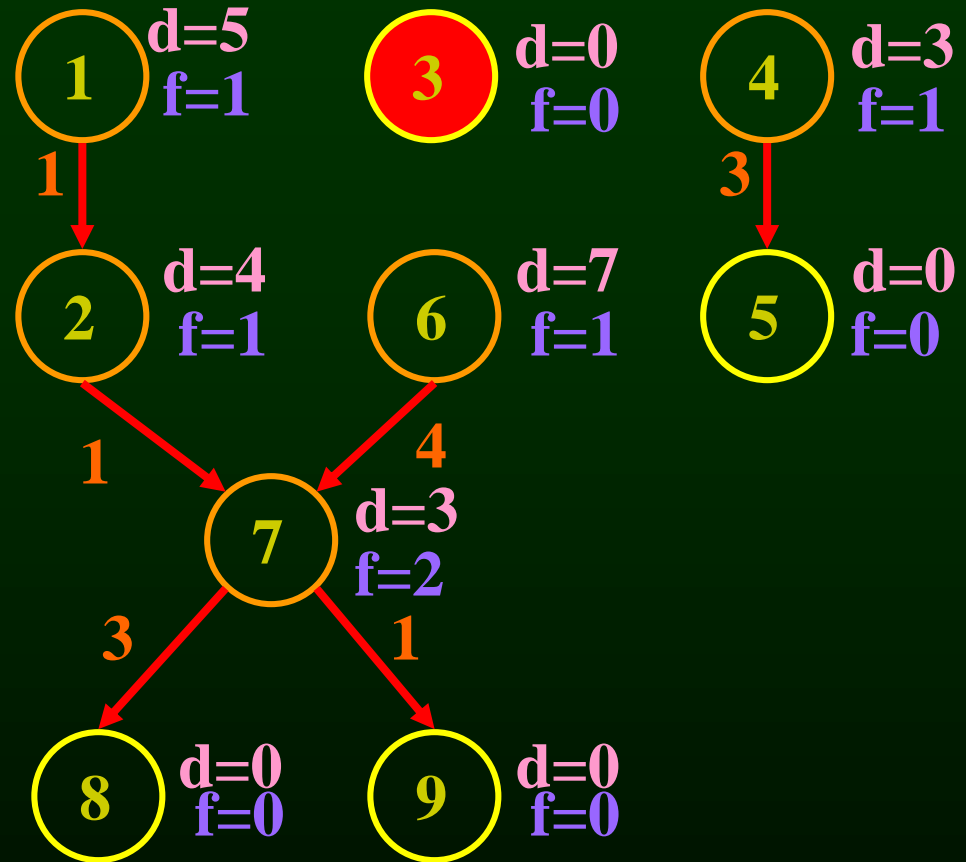
READY = { 3, 5, 8, 9 }



6	1	2	4	7
---	---	---	---	---

Example

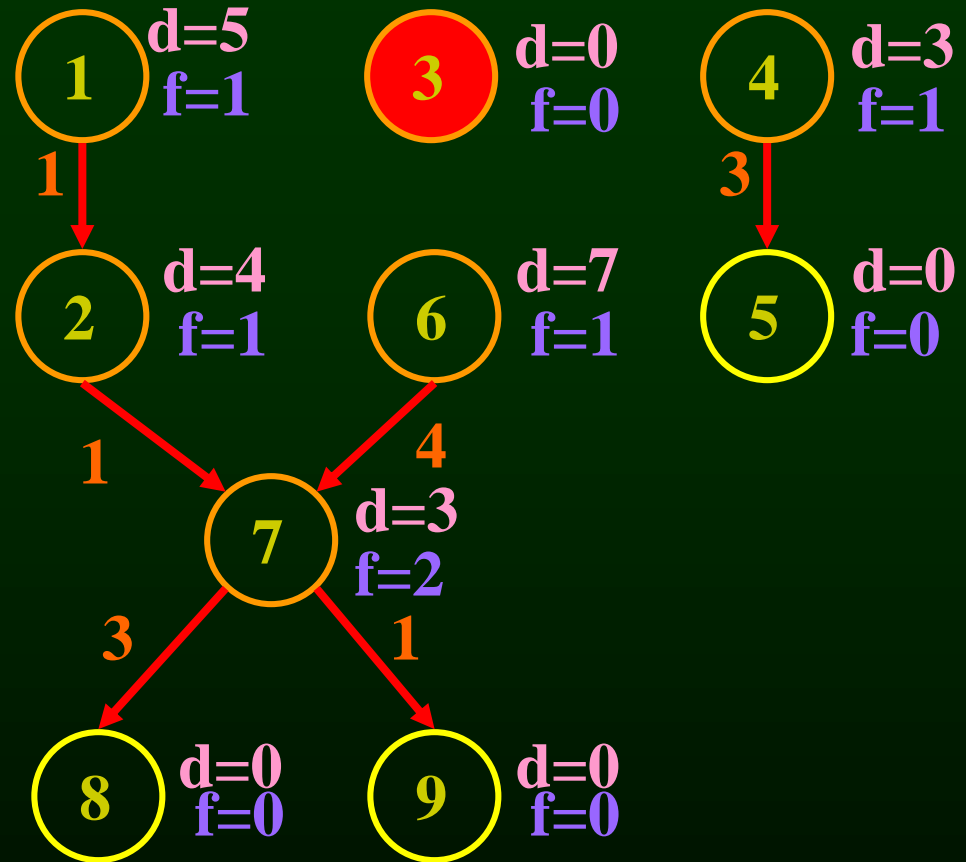
READY = { 3, 5, 8, 9 }



6	1	2	4	7
---	---	---	---	---

Example

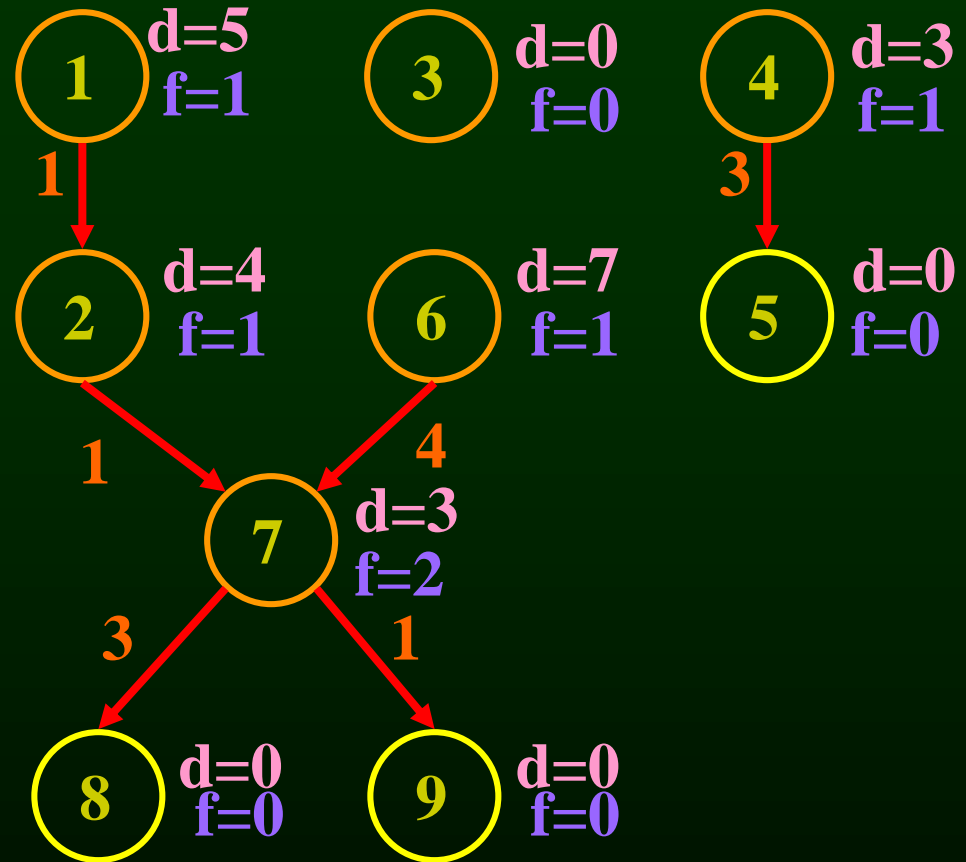
READY = { 3, 5, 8, 9 }



6	1	2	4	7	3
---	---	---	---	---	---

Example

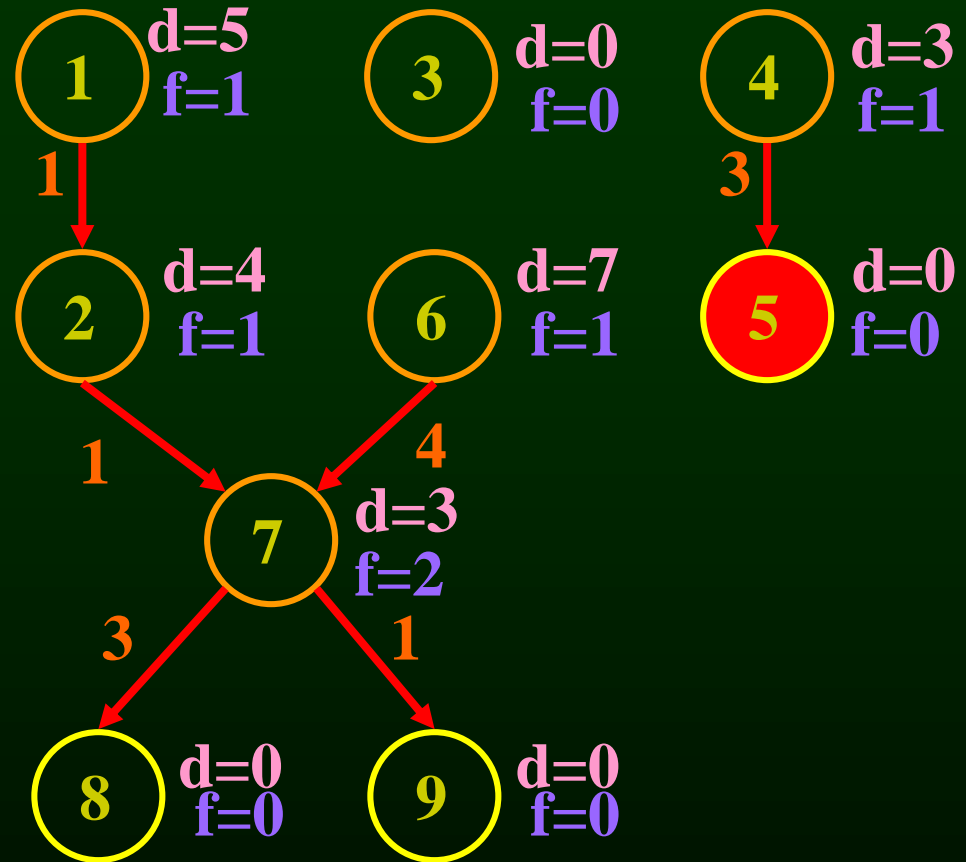
READY = { 5, 8, 9 }



6	1	2	4	7	3
---	---	---	---	---	---

Example

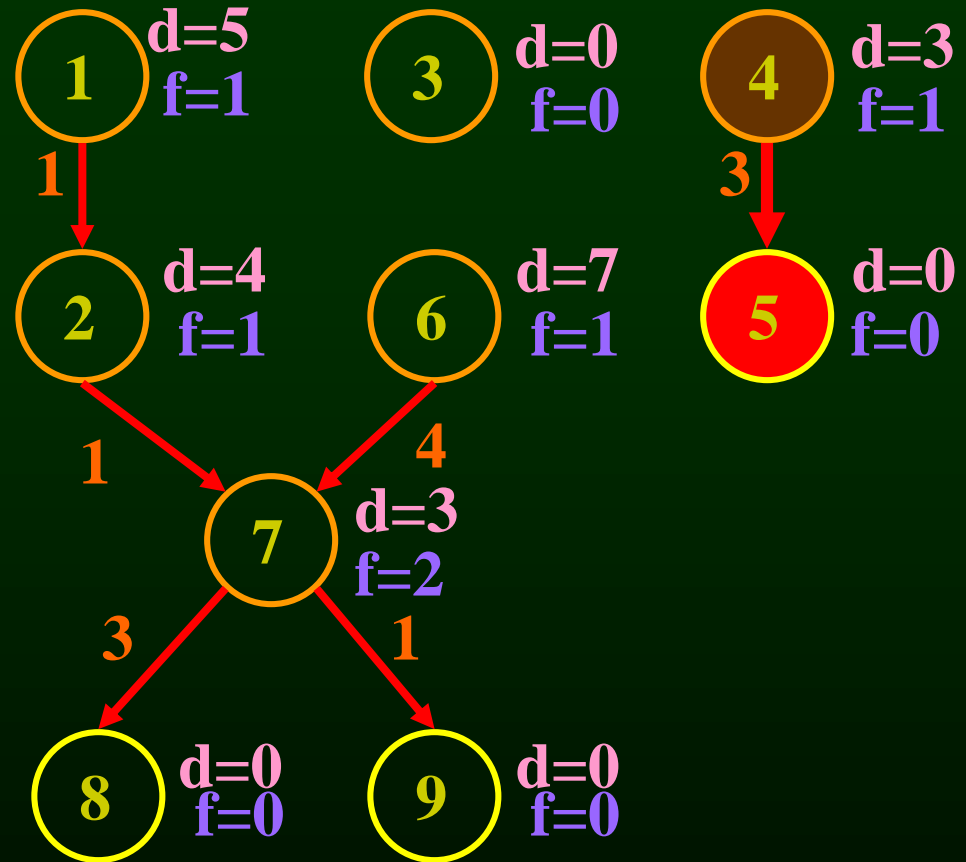
READY = { 5, 8, 9 }



6	1	2	4	7	3
---	---	---	---	---	---

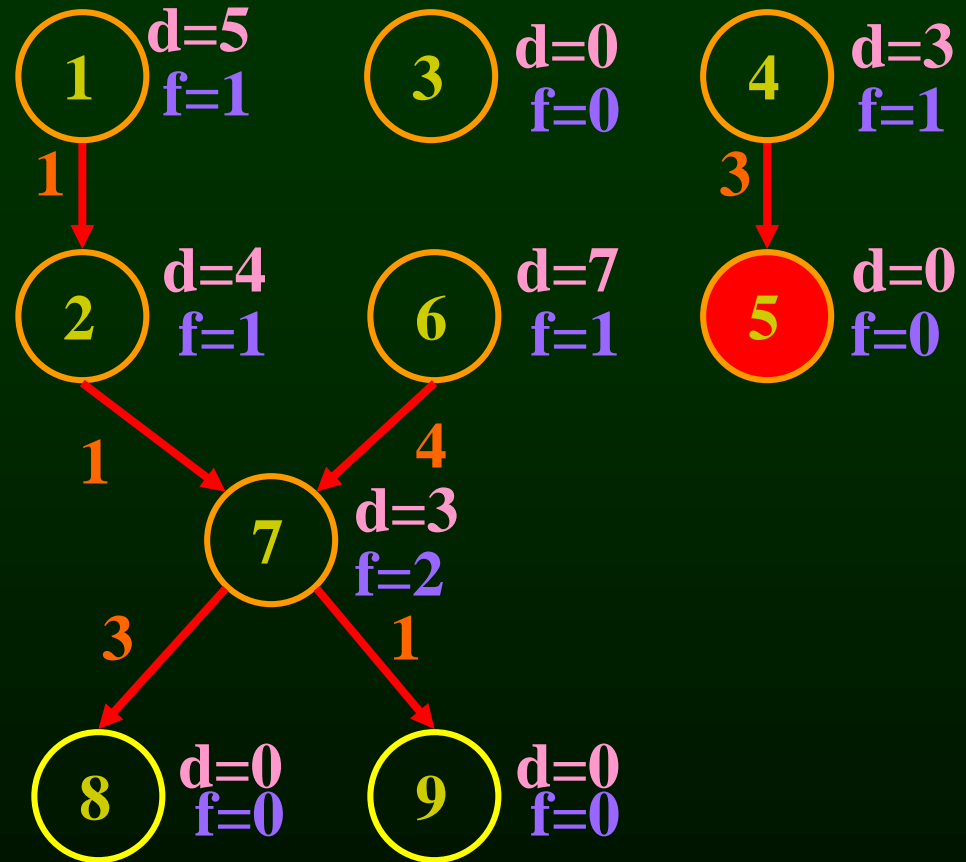
Example

READY = { 5, 8, 9 }



Example

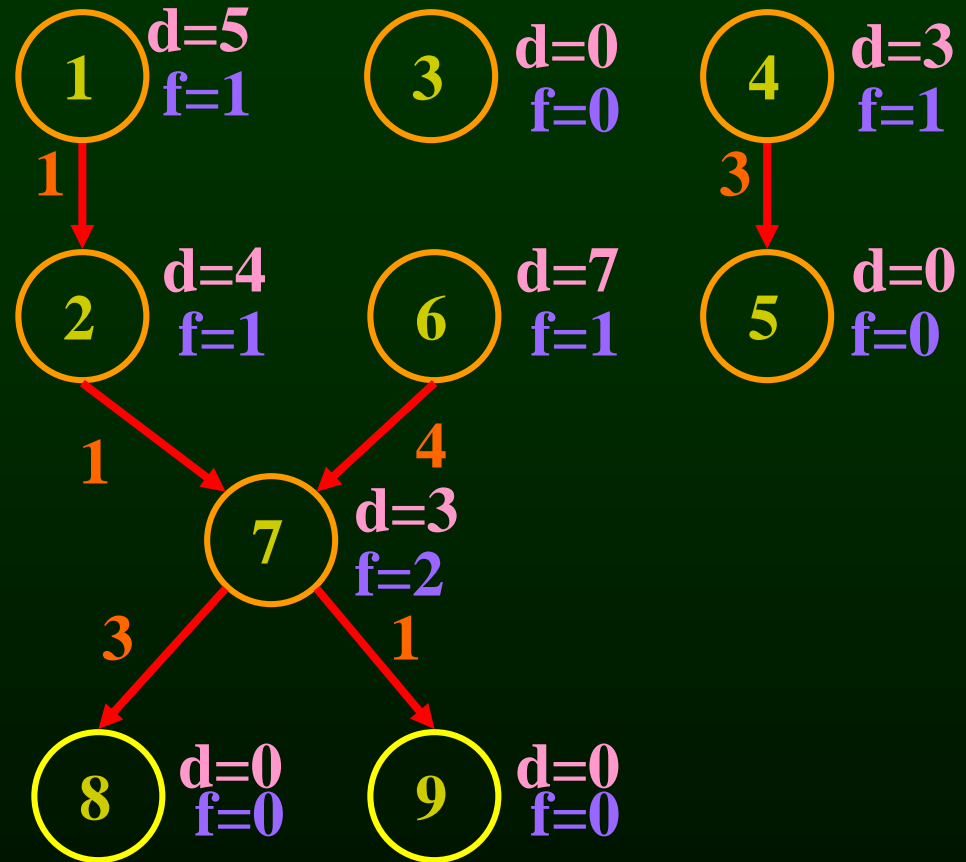
READY = { 5, 8, 9 }



6	1	2	4	7	3	5
---	---	---	---	---	---	---

Example

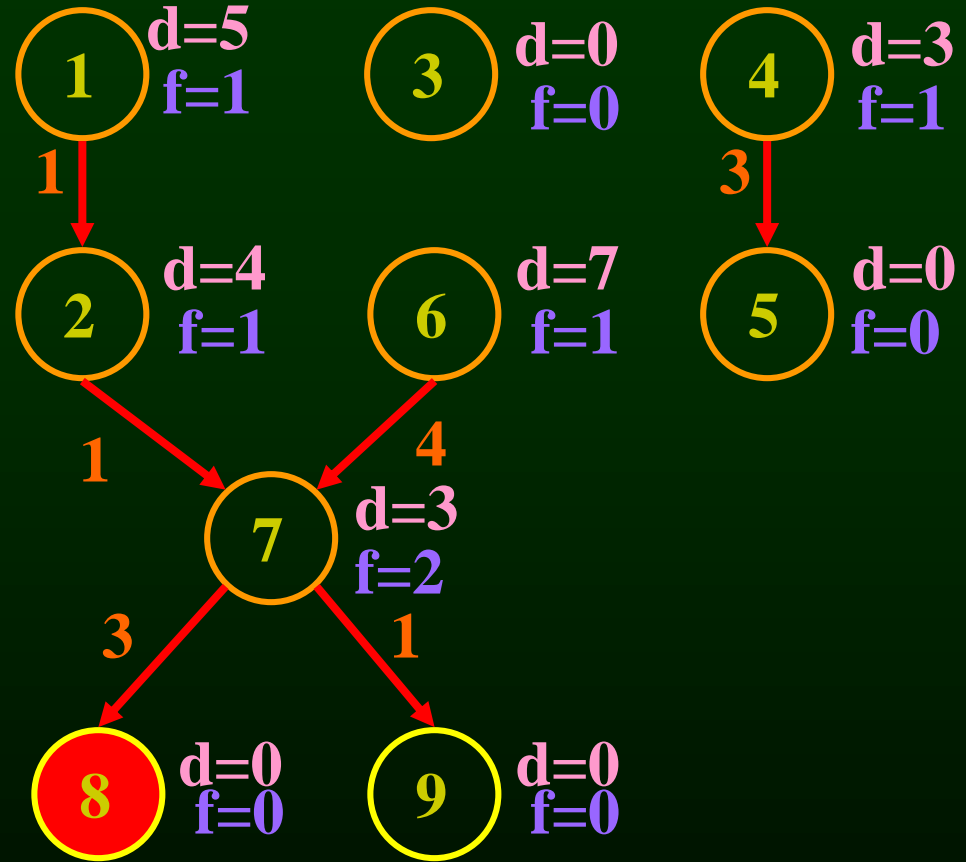
READY = { 8, 9 }



6	1	2	4	7	3	5
---	---	---	---	---	---	---

Example

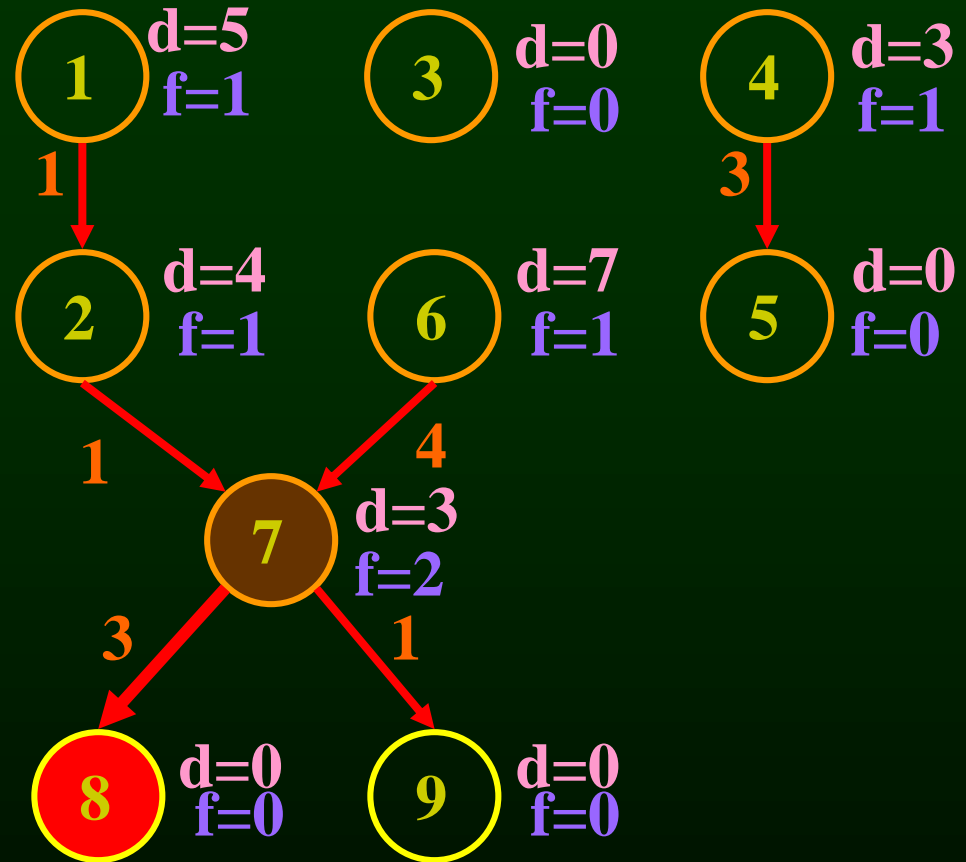
READY = { 8, 9 }



6	1	2	4	7	3	5
---	---	---	---	---	---	---

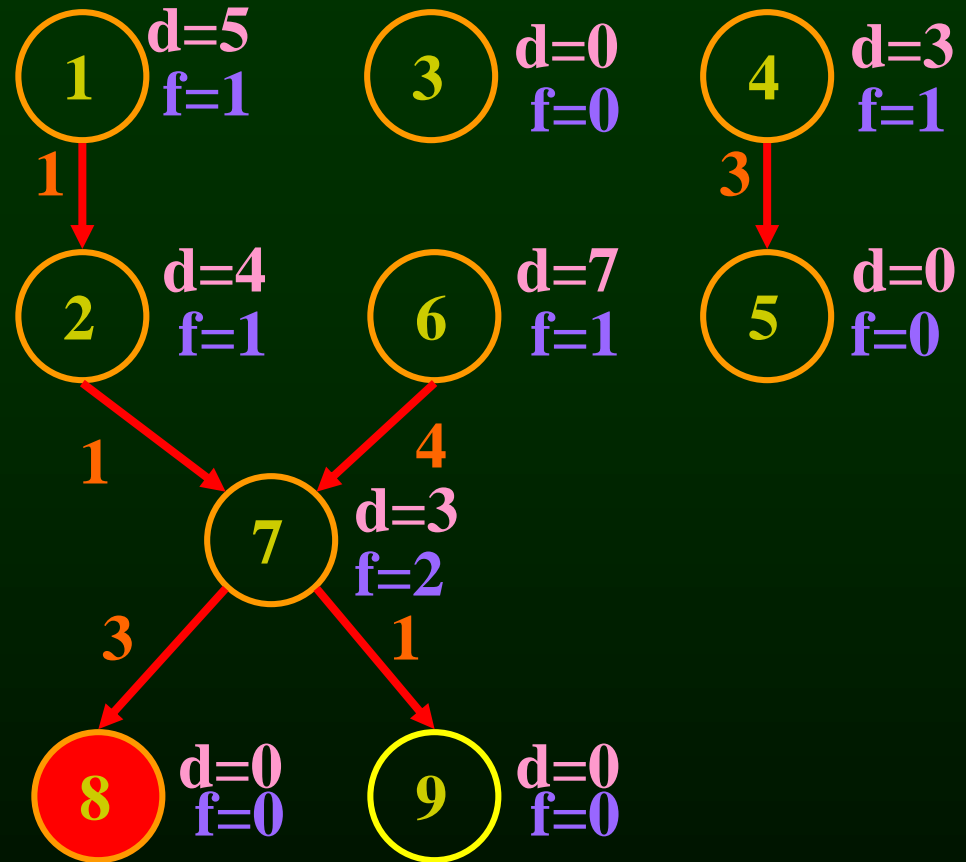
Example

READY = { 8, 9 }



Example

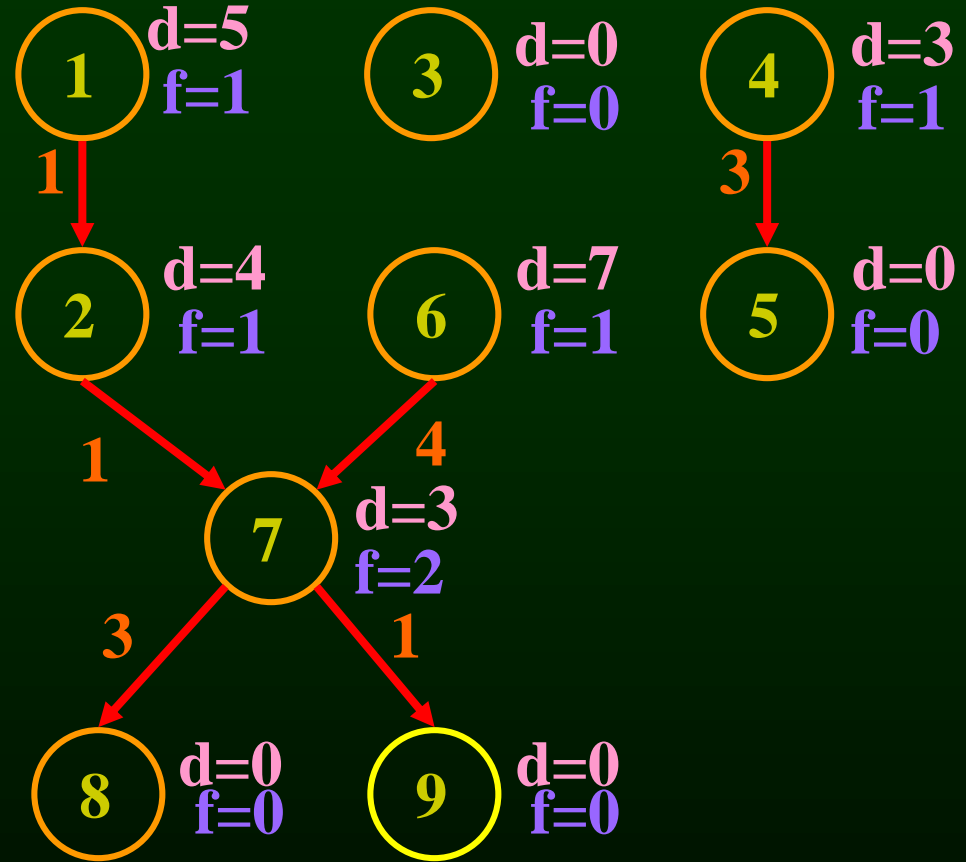
READY = { 8, 9 }



6	1	2	4	7	3	5	8
---	---	---	---	---	---	---	---

Example

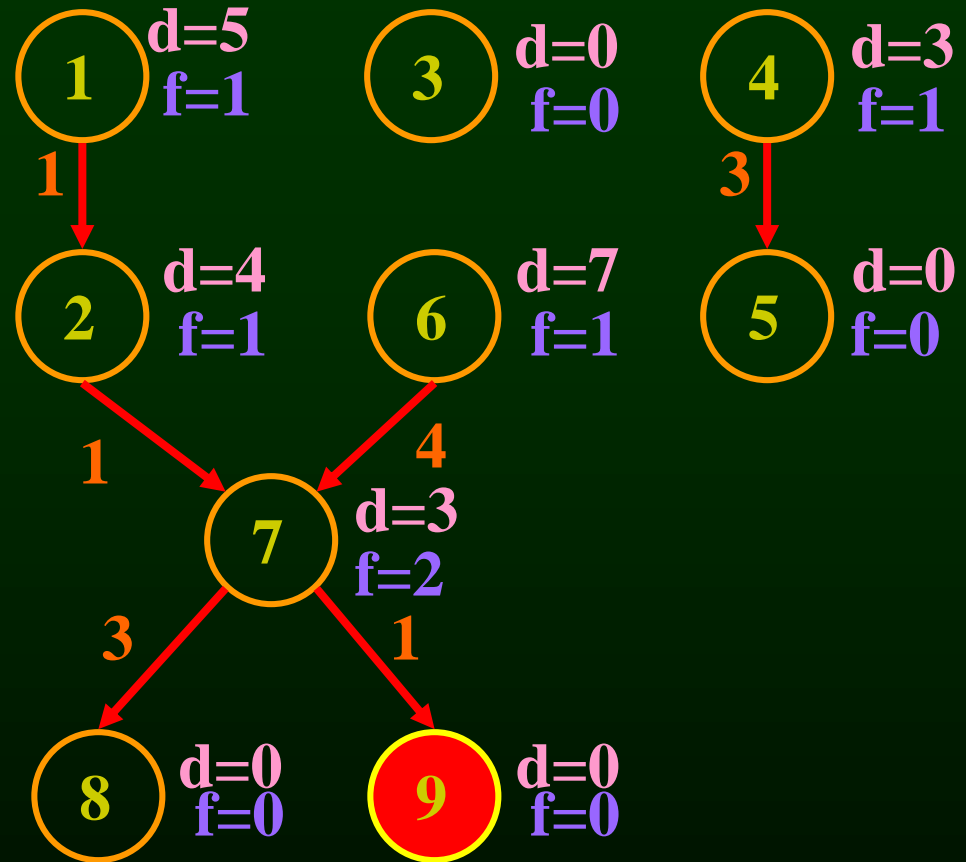
READY = { 9 }



6	1	2	4	7	3	5	8
---	---	---	---	---	---	---	---

Example

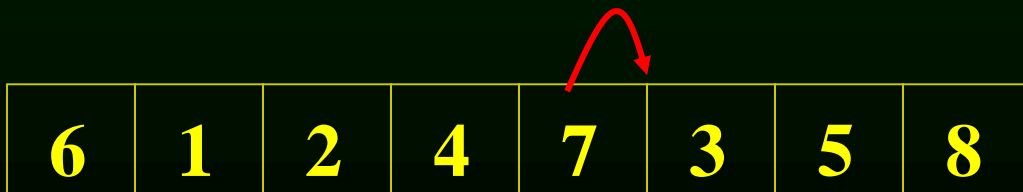
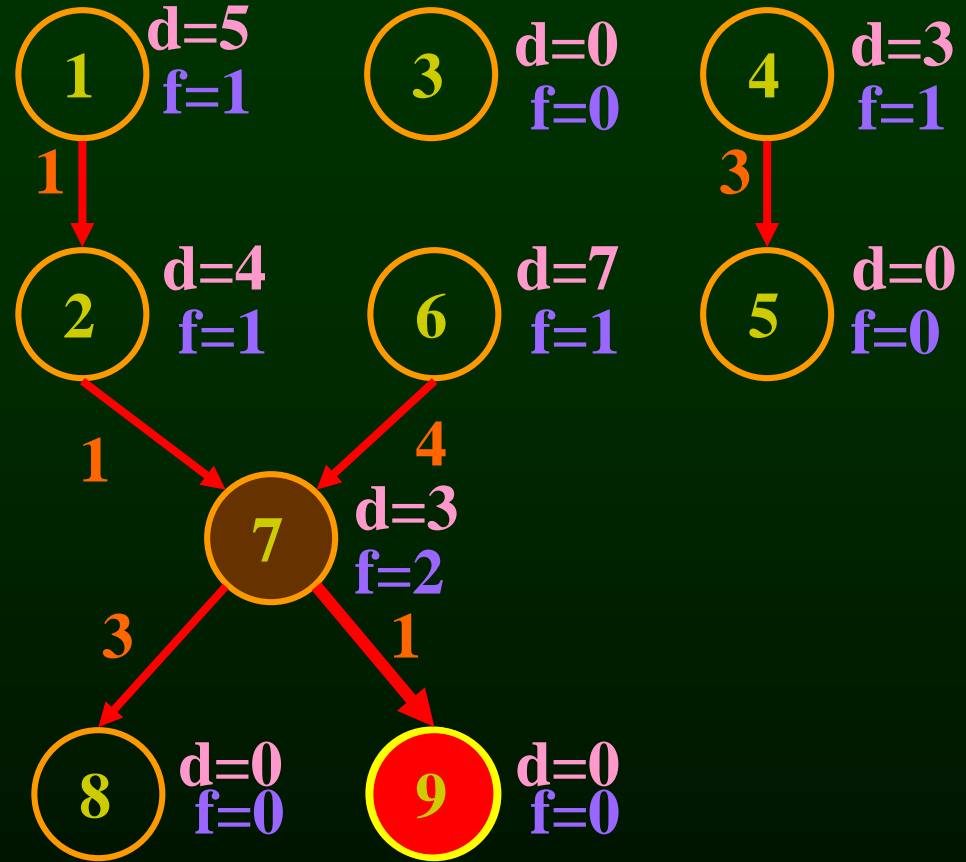
READY = { 9 }



6	1	2	4	7	3	5	8
---	---	---	---	---	---	---	---

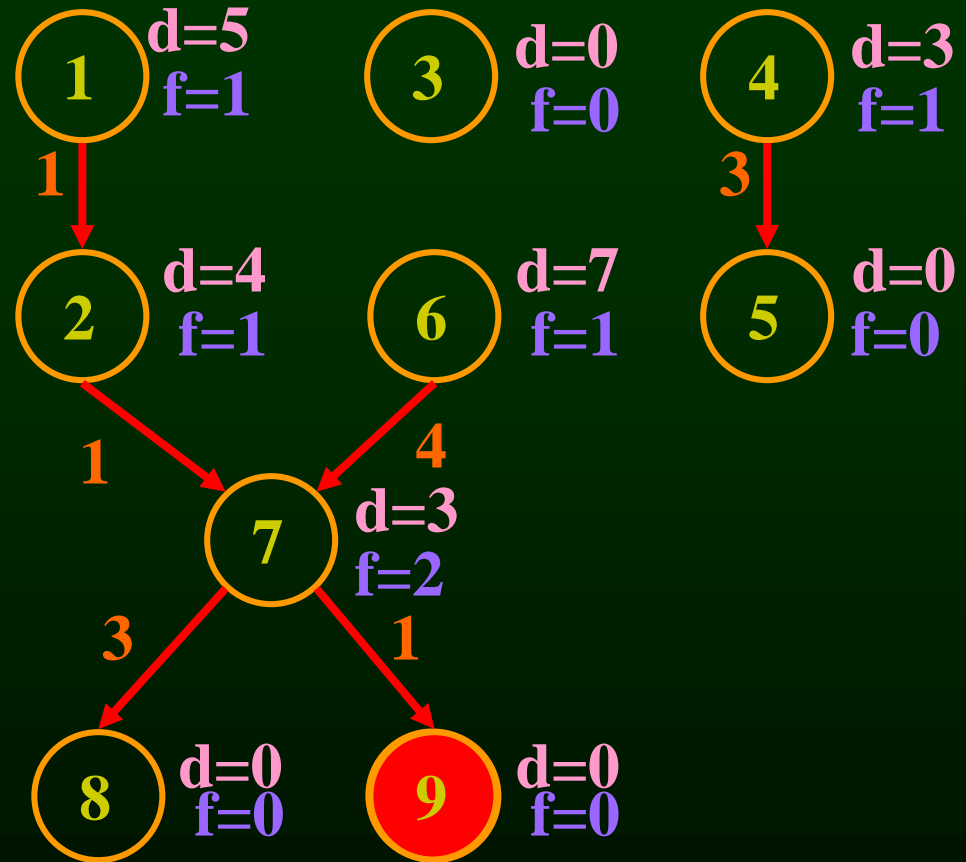
Example

READY = { 9 }



Example

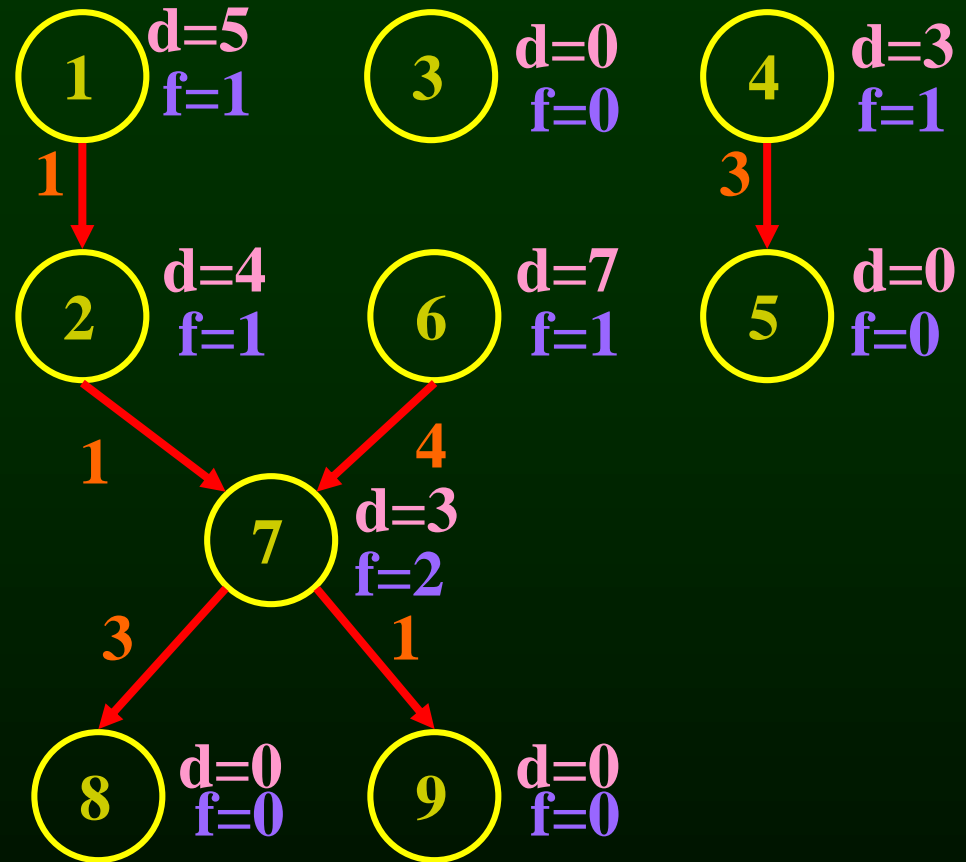
READY = { 9 }



6	1	2	4	7	3	5	8	9
---	---	---	---	---	---	---	---	---

Example

READY = { }



6	1	2	4	7	3	5	8	9
---	---	---	---	---	---	---	---	---

Example

		Results In
1:	lea var_a, %rax	1 cycle
2:	add \$4, %rax	1 cycle
3:	inc %r11	1 cycle
4:	mov 4(%rsp), %r10	3 cycles
5:	add %r10, 8(%rsp)	
6:	and 16(%rsp), %rbx	4 cycles
7:	imul %rax, %rbx	3 cycles
8:	mov %rbx, 16(%rsp)	
9:	lea var_b, %rax	

6	1	2	4	7	3	5	8	9
---	---	---	---	---	---	---	---	---

9 cycles

Example

		Results In
1:	lea var_a, %rax	1 cycle
2:	add \$4, %rax	1 cycle
3:	inc %r11	1 cycle
4:	mov 4(%rsp), %r10	3 cycles
5:	add %r10, 8(%rsp)	
6:	and 16(%rsp), %rbx	4 cycles
7:	imul %rax, %rbx	3 cycles
8:	mov %rbx, 16(%rsp)	
9:	lea var_b, %rax	

1	2	3	4	st	st	5	6	st	st	st	7	8	9
---	---	---	---	----	----	---	---	----	----	----	---	---	---

*14 cycles vs
9 cycles*

6	1	2	4	7	3	5	8	9
---	---	---	---	---	---	---	---	---

Outline

- Modern architectures
- Introduction to instruction scheduling
- List scheduling
- Resource constraints
- Scheduling across basic blocks
- Trace scheduling

Resource Constraints

- Modern machines have many resource constraints
- Superscalar architectures:
 - can run few parallel operations
 - But have constraints

Resource Constraints of a Superscalar Processor

- Example:

- One fully pipelined reg-to-reg unit
 - All integer operations taking one cycle

In parallel with

- One fully pipelined memory-to/from-reg unit
 - Data loads take two cycles
 - Data stores take one cycle

List Scheduling Algorithm with resource constraints

- Represent the superscalar architecture as multiple pipelines
 - Each pipeline represent some resource

List Scheduling Algorithm with resource constraints

- Represent the superscalar architecture as multiple pipelines
 - Each pipeline represent some resource
- Example
 - One single cycle reg-to-reg ALU unit
 - One two-cycle pipelined reg-to/from-memory unit

ALU						
MEM 1						
MEM 2						

List Scheduling Algorithm with resource constraints

- Create a dependence DAG of a basic block
- Topological Sort

READY = nodes with no predecessors

Loop until READY is empty

Let $n \in \text{READY}$ be the node with the highest priority

Schedule n in the earliest slot

that satisfies precedence + resource constraints

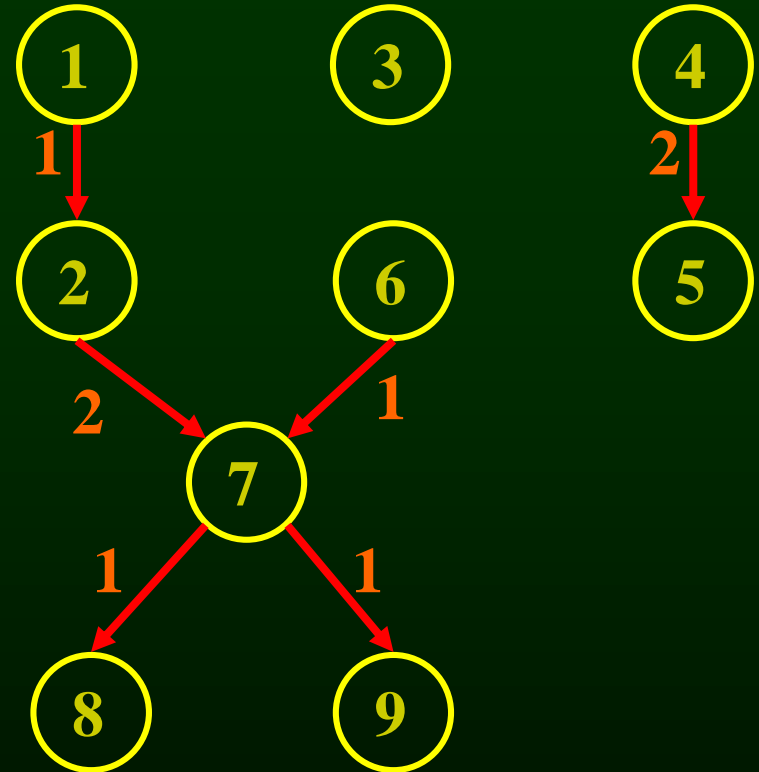
Update READY

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

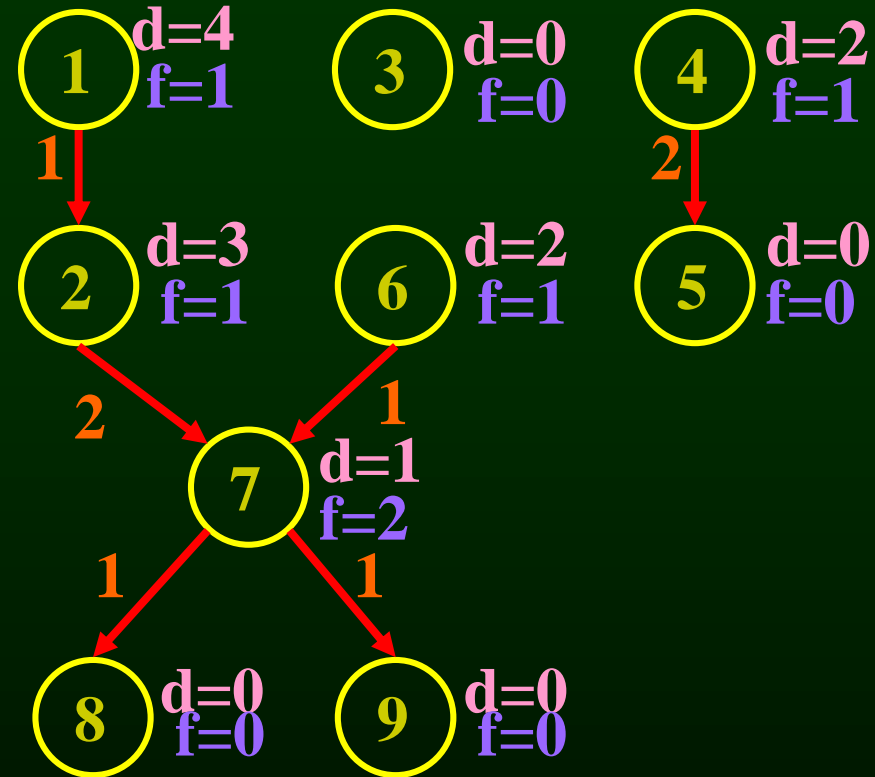
Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```



Example

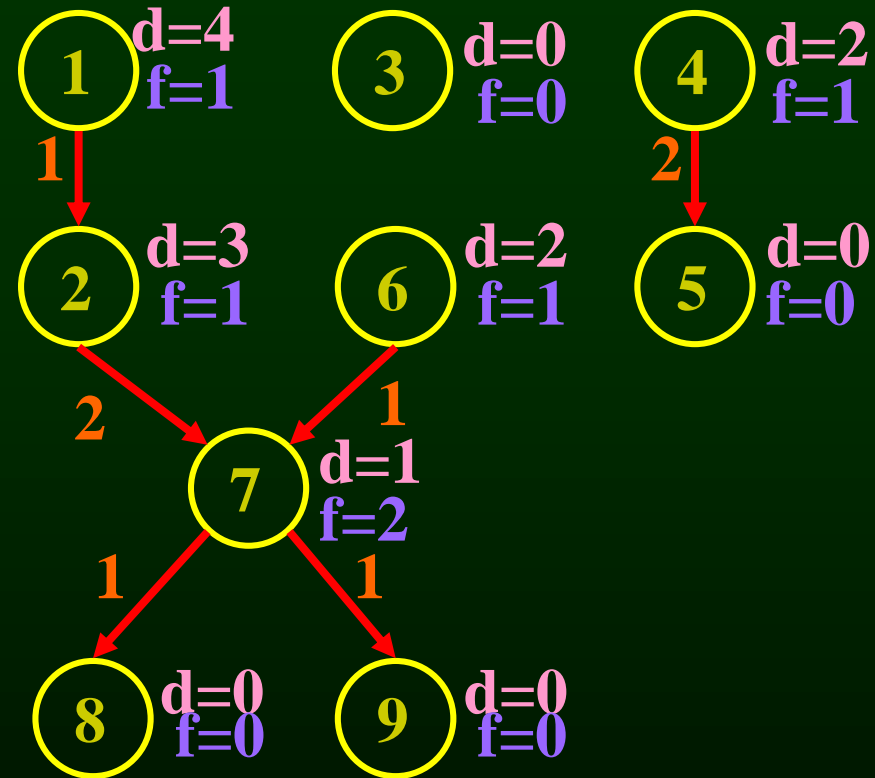
```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```



Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

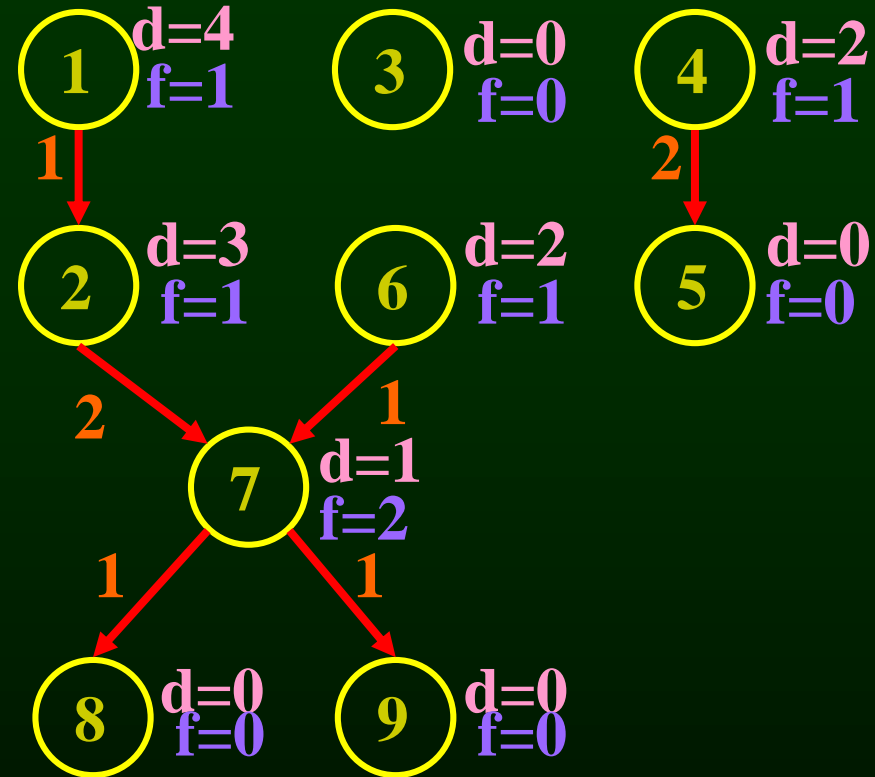
READY = { 1, 6, 4, 3 }



Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 1, 6, 4, 3 }



ALUOp

MEM 1

MEM 2

Example

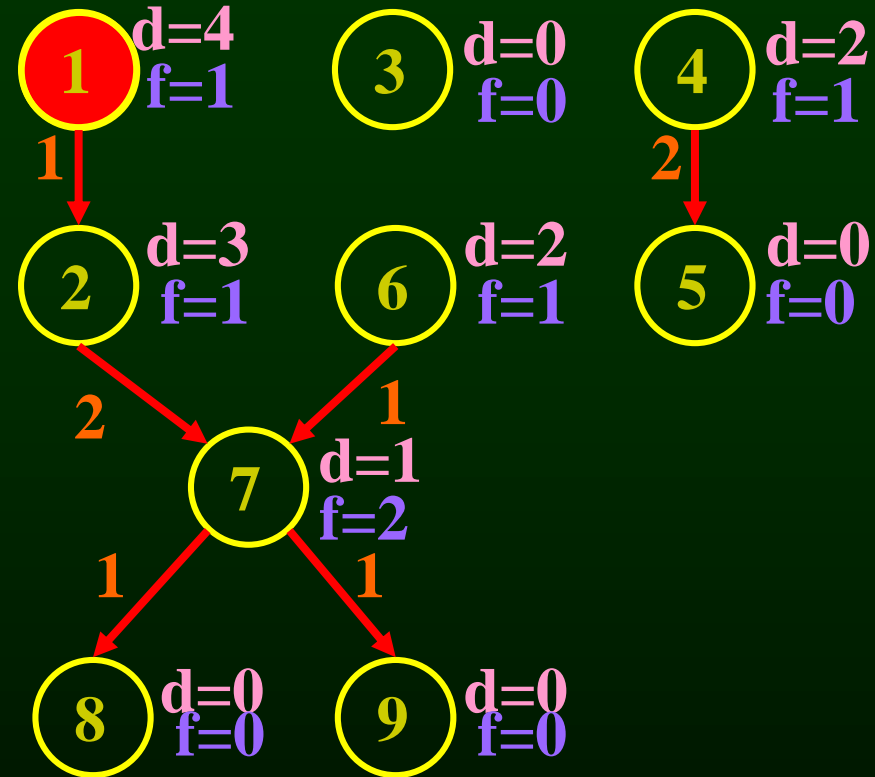
```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov     4(%rsp), %r10
5: mov     %r10, 8(%rsp)
6: and     $0x00ff, %rbx
7: imul    %rax, %rbx
8: lea     var_b, %rax
9: mov     %rbx, 16(%rsp)
```

READY = { 1, 6, 4, 3 }

ALUOp

MEM 1

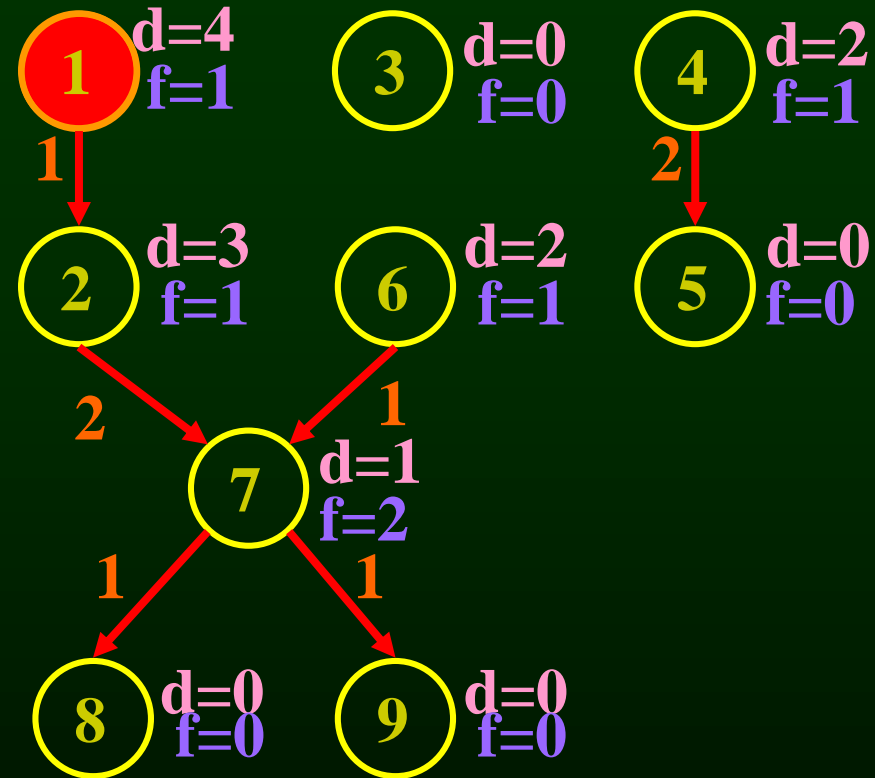
MEM 2



Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 1, 6, 4, 3 }



ALUOp

1

MEM 1

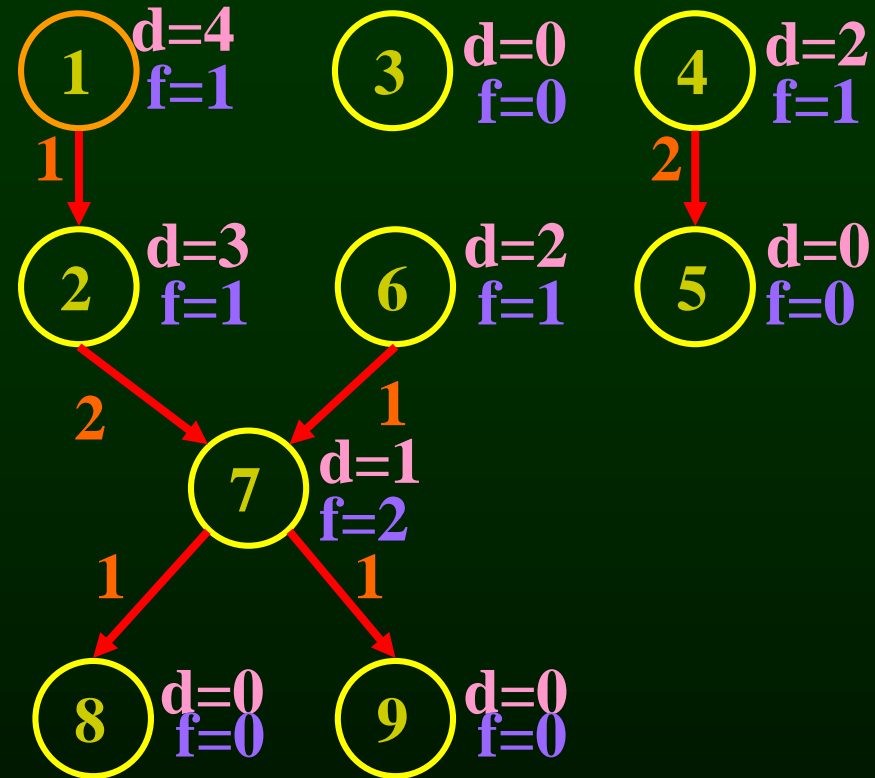
MEM 2

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 6, 4, 3 } ← 2



ALUOp

1

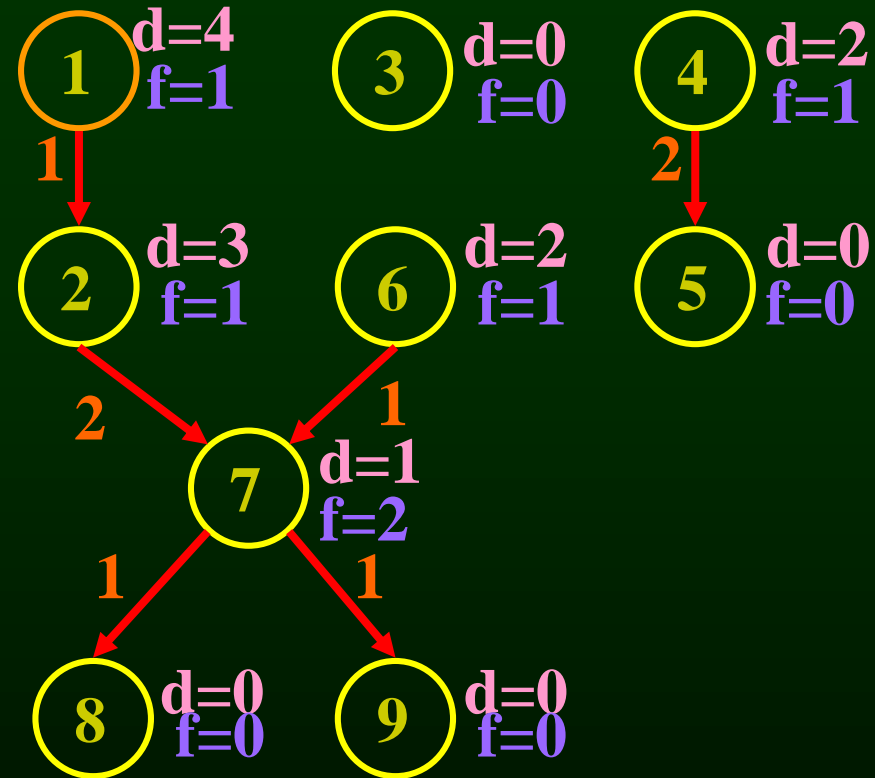
MEM 1

MEM 2

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 2, 6, 4, 3 }



ALUOp

1

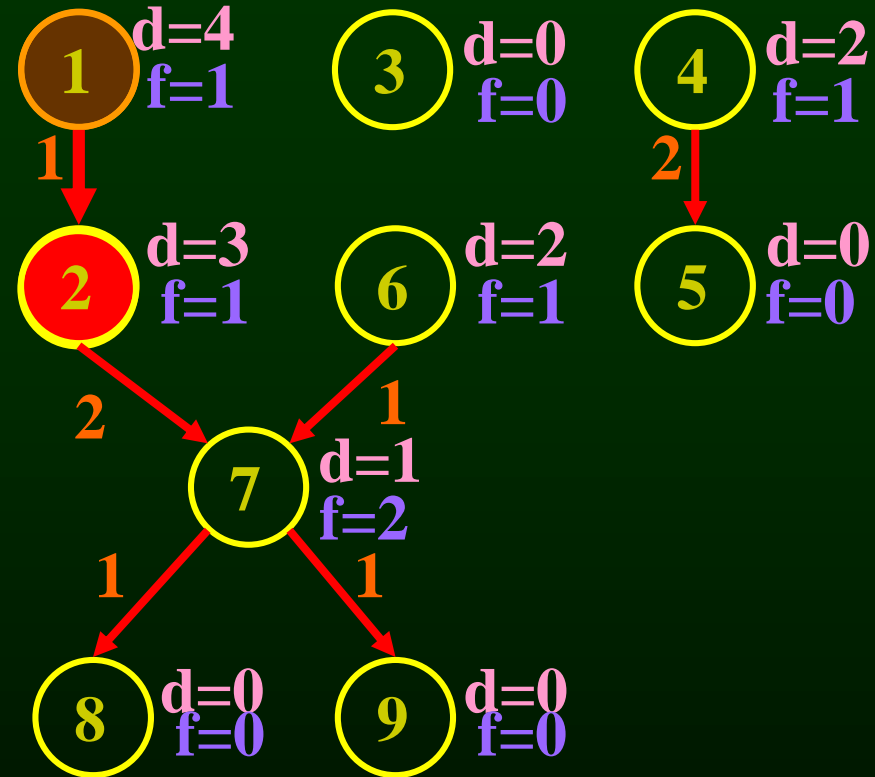
MEM 1

MEM 2

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 2, 6, 4, 3 }



ALUOp

1

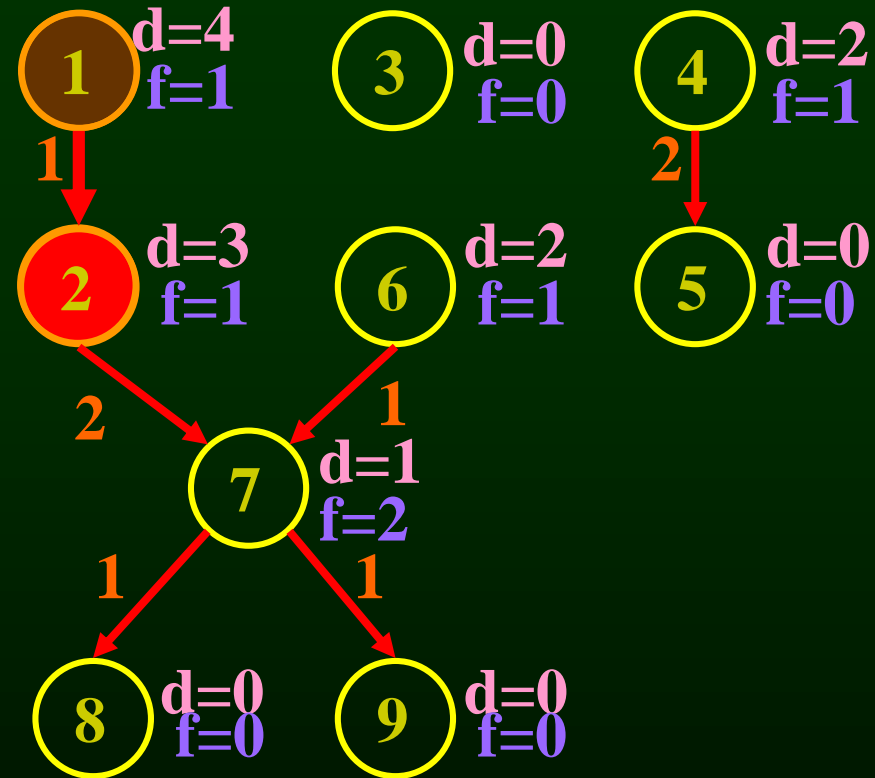
MEM 1

MEM 2

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov     4(%rsp), %r10
5: mov     %r10, 8(%rsp)
6: and     $0x00ff, %rbx
7: imul    %rax, %rbx
8: lea     var_b, %rax
9: mov     %rbx, 16(%rsp)
```

READY = { 2, 6, 4, 3 }



ALUOp

1

MEM 1

2

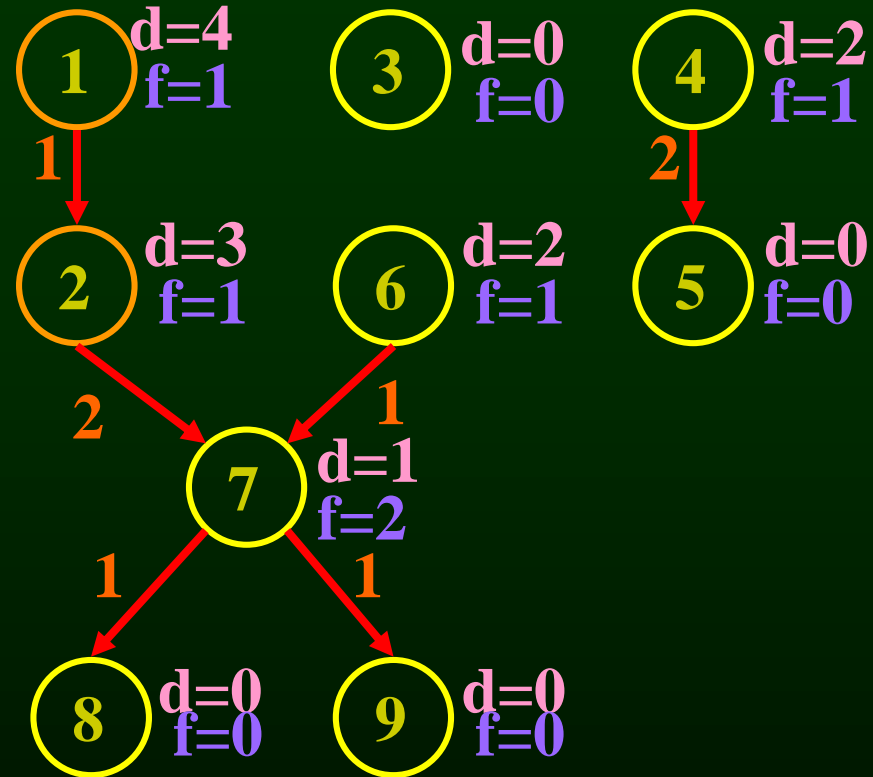
MEM 2

2

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 6, 4, 3 }



ALUOp

1

MEM 1

2

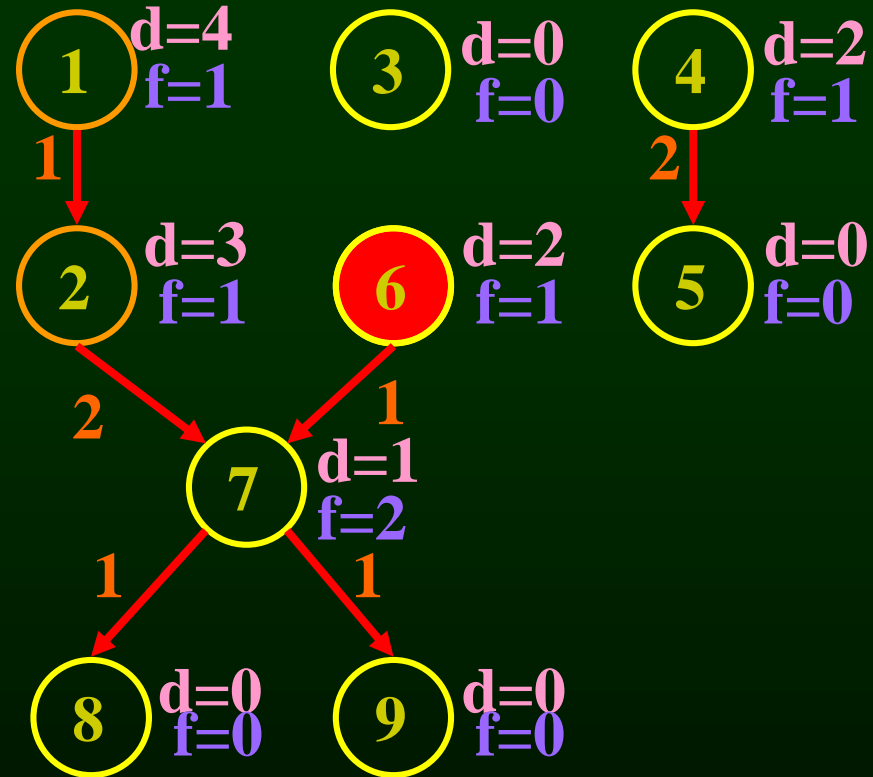
MEM 2

2

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 6, 4, 3 }



ALUOp

1

MEM 1

2

MEM 2

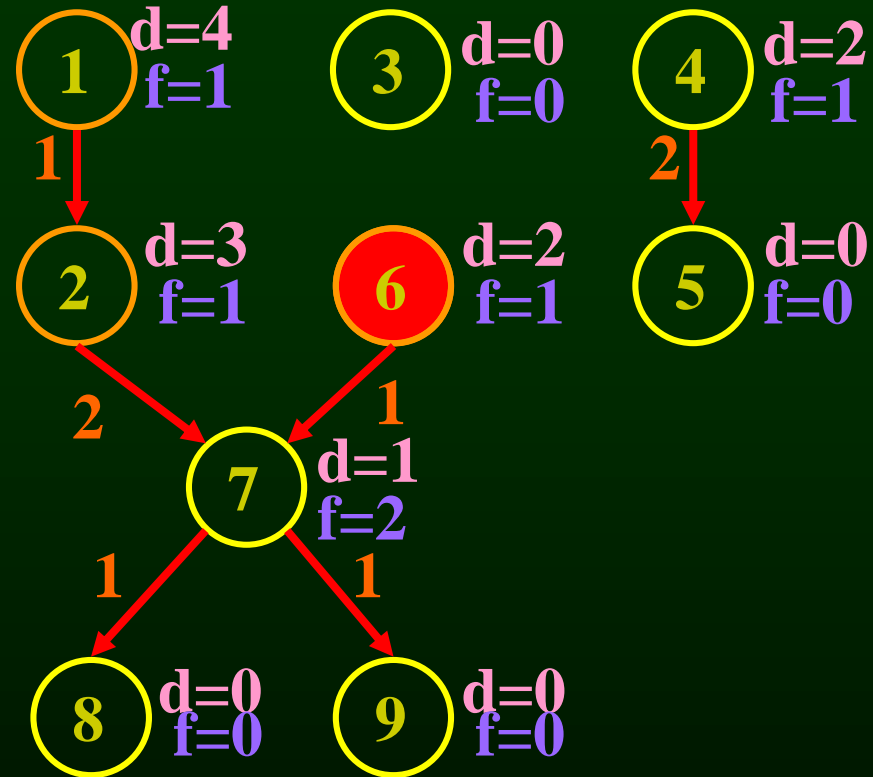
2

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 6, 4, 3 }



ALUOp

1

6

MEM 1

2

MEM 2

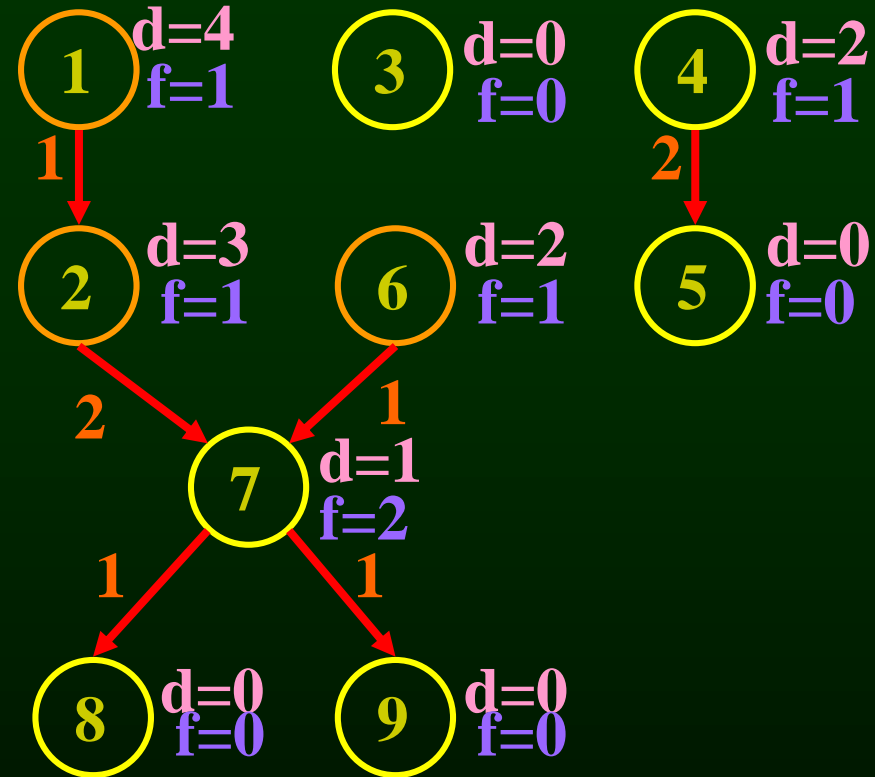
2

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 4, 3 } ← 7

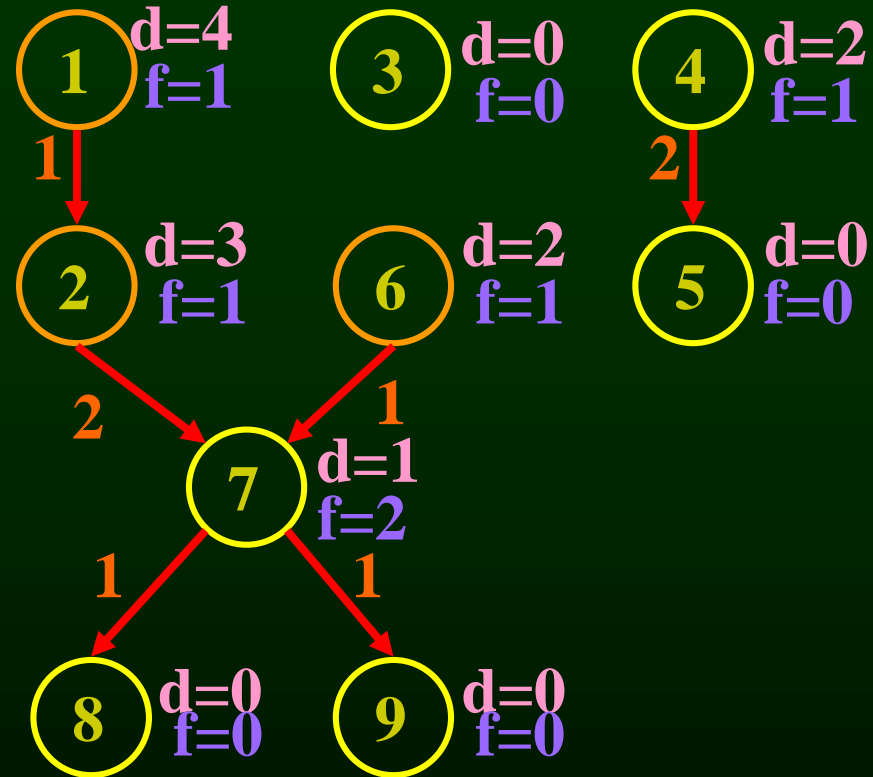


ALUOp	1	6				
MEM 1		2				
MEM 2			2			

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 4, 7, 3 }



ALUOp

1 6

MEM 1

2

MEM 2

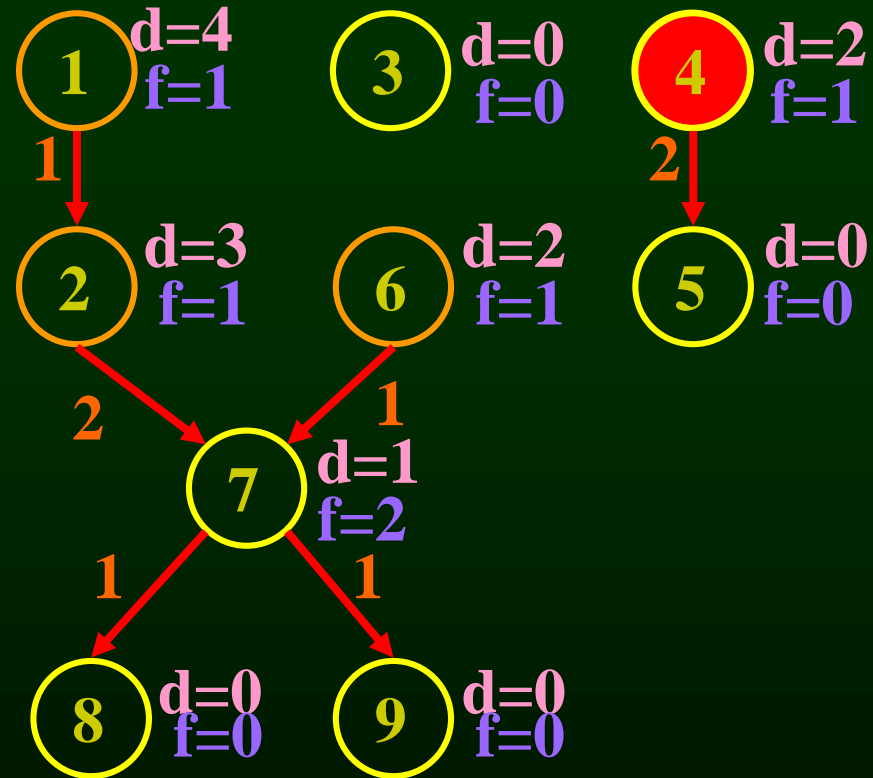
2

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 4, 7, 3 }



ALUOp

1 6

MEM 1

2

MEM 2

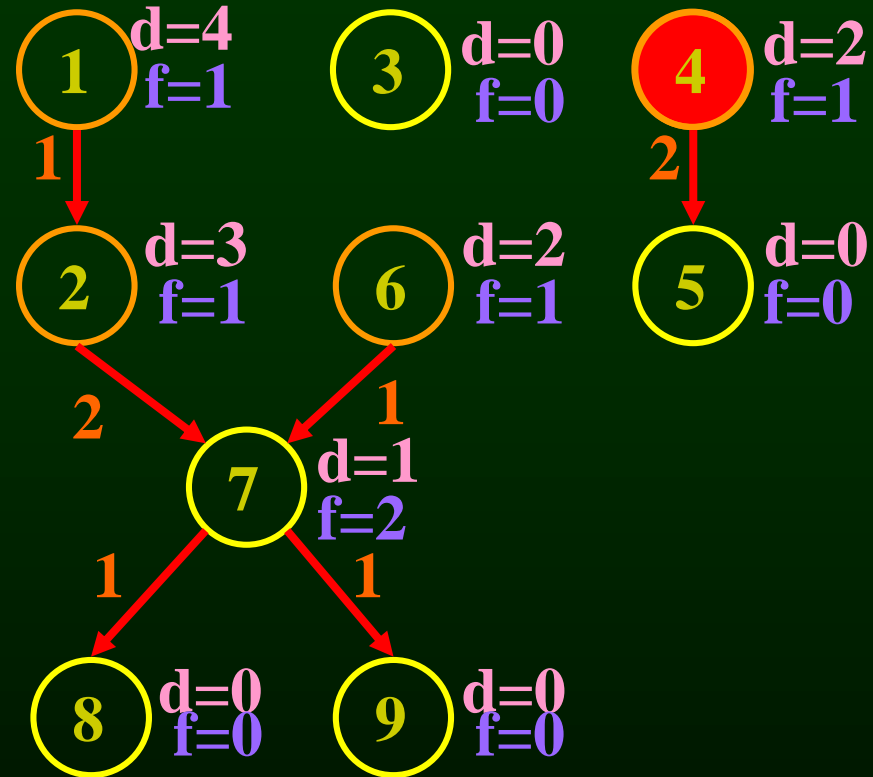
2

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 4, 7, 3 }



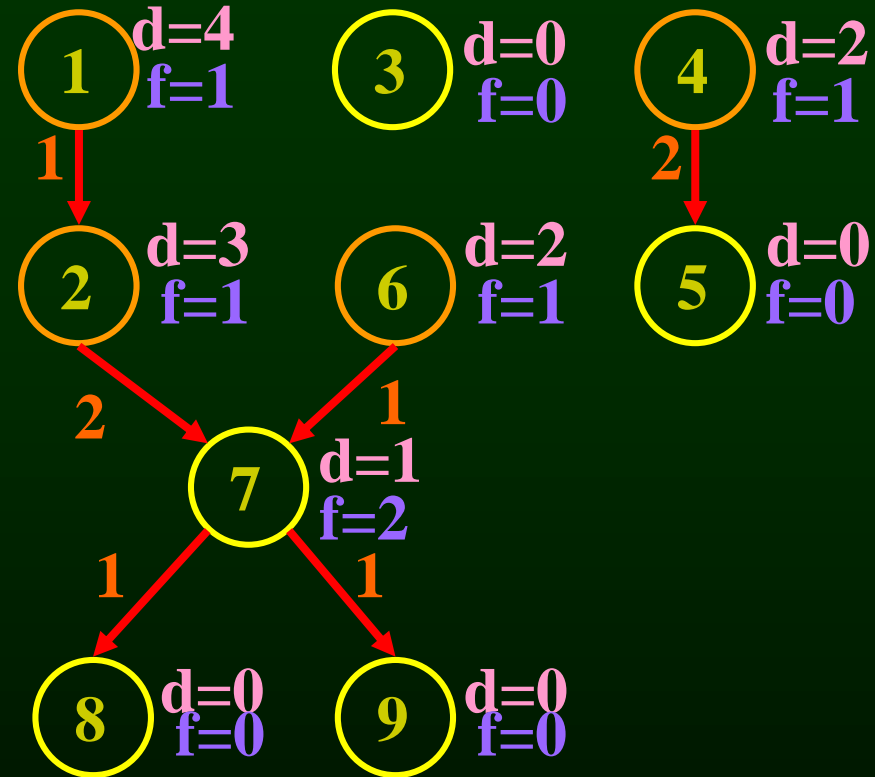
ALUOp	1	6					
MEM 1	4	2					
MEM 2		4	2				

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 7, 3 } ← 5

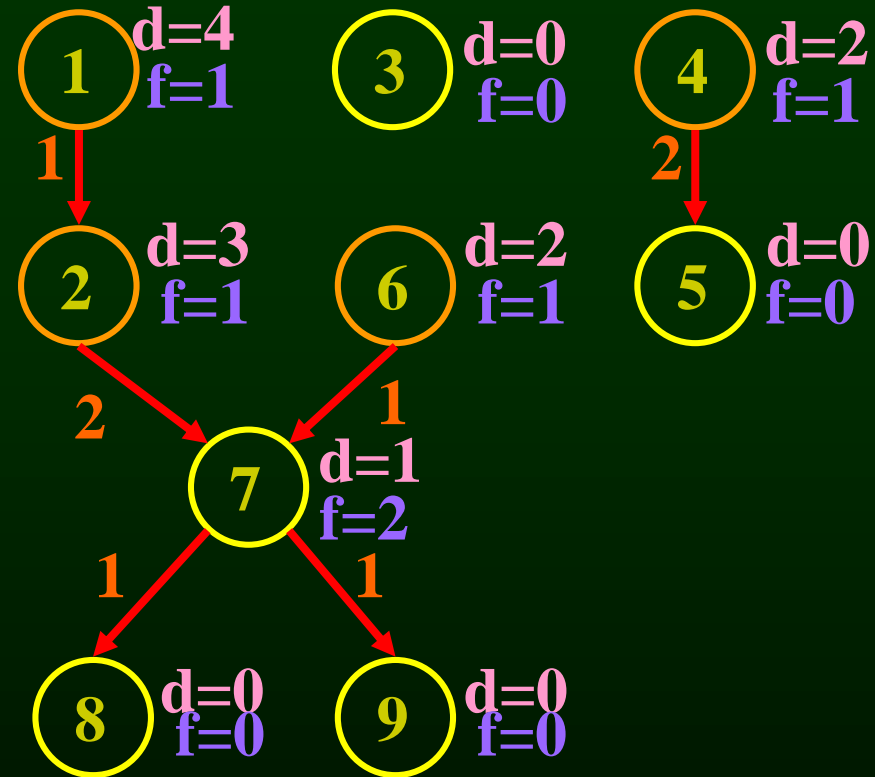


ALUOp	1	6					
MEM 1	4	2					
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 7, 3, 5 }



ALUOp

1 6

MEM 1

4 2

MEM 2

4 2

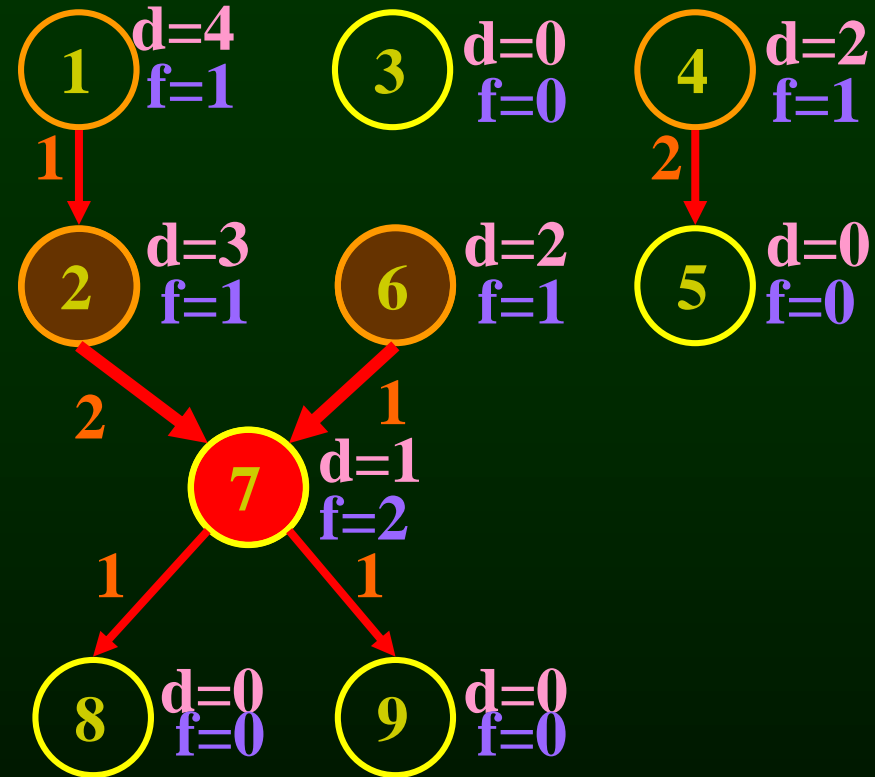
1	6					
4	2					
	4	2				

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 7, 3, 5 }



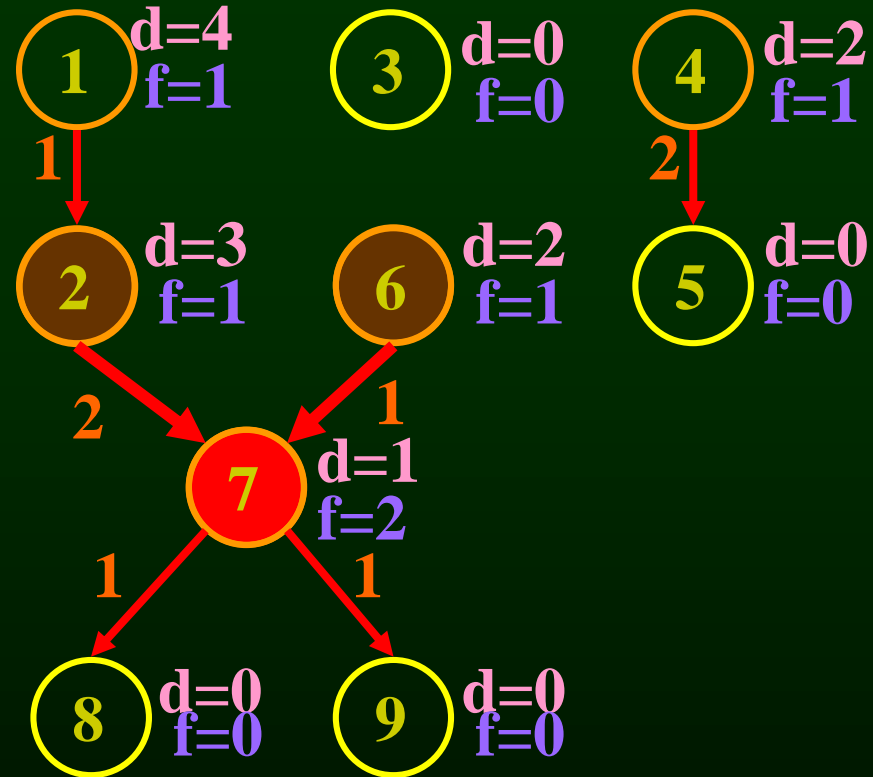
ALUOp	1	6				
MEM 1	4	2				
MEM 2		4	2			

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 7, 3, 5 }



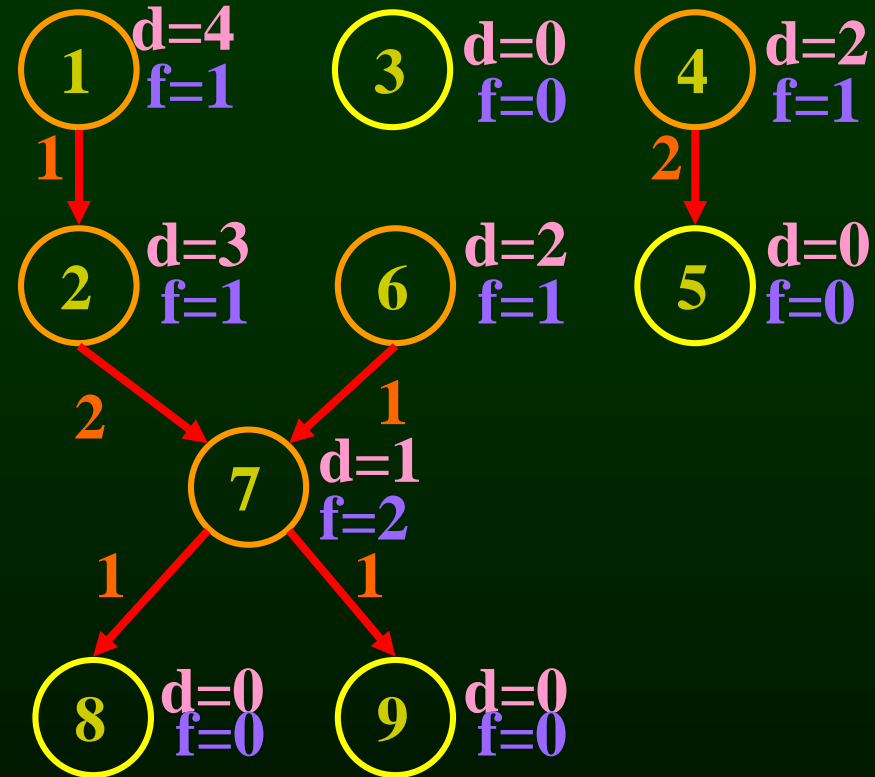
ALUOp	1	6		7		
MEM 1	4	2				
MEM 2		4	2			

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 3, 5 } ← 8, 9

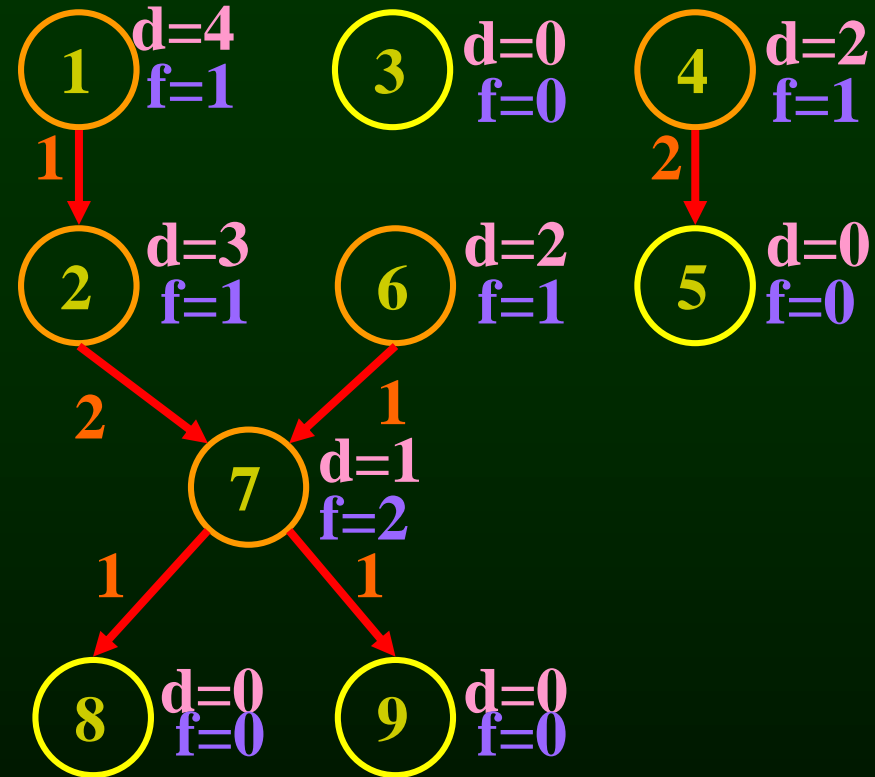


ALUOp	1	6		7			
MEM 1	4	2					
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 3, 5, 8, 9 }



ALUOp

1

6

7

MEM 1

4

2

MEM 2

4

2

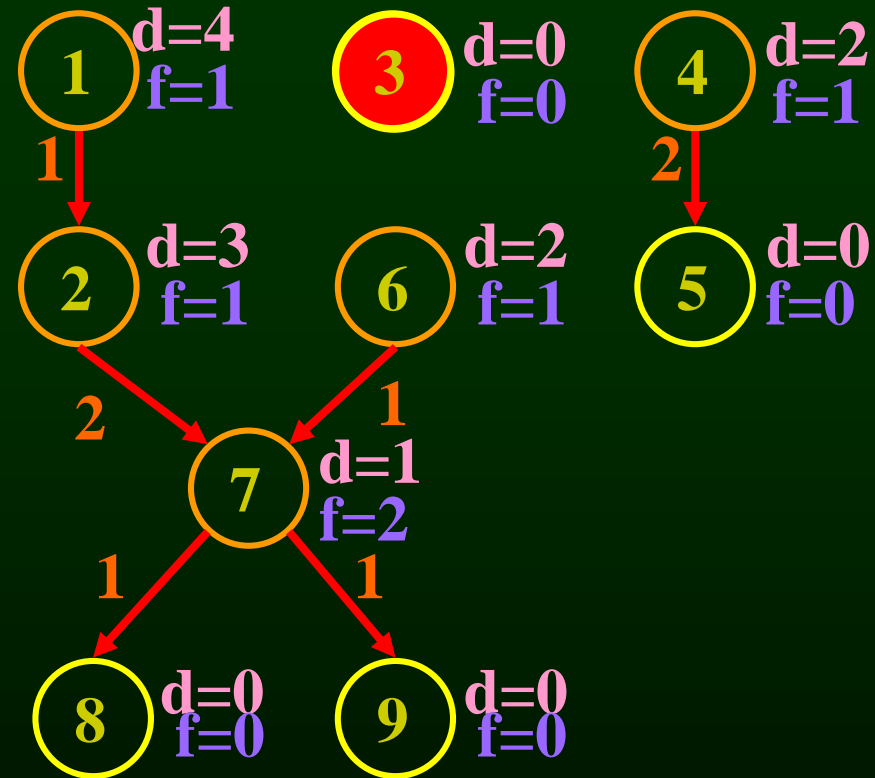
1	6		7			
4	2					
	4	2				

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 3, 5, 8, 9 }

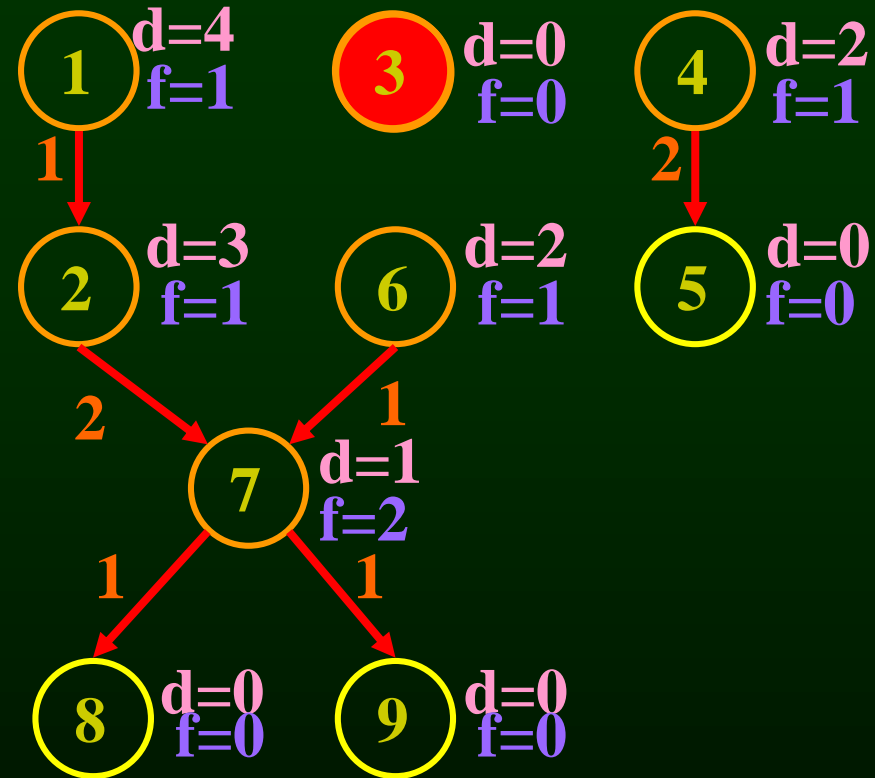


ALUOp	1	6		7			
MEM 1	4	2					
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 3, 5, 8, 9 }

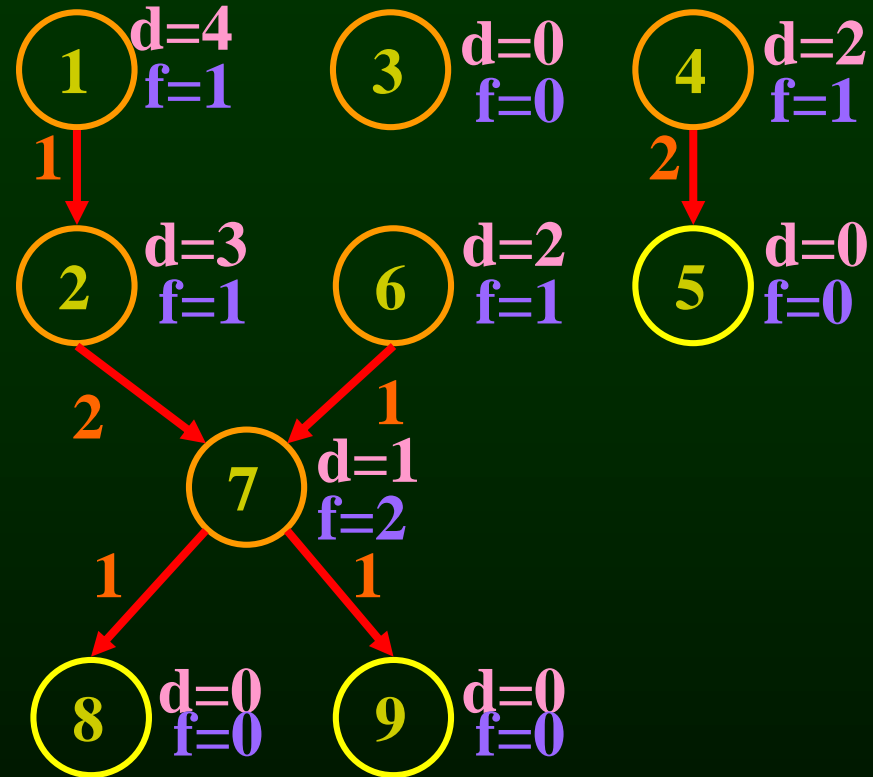


ALUOp	1	6	3	7			
MEM 1	4	2					
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 5, 8, 9 }

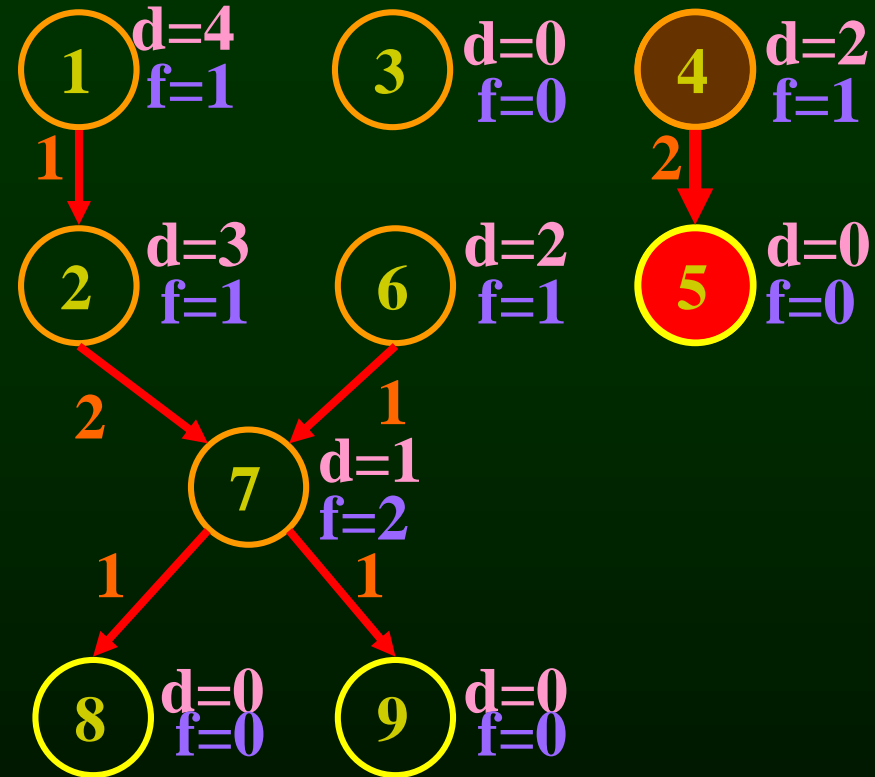


ALUOp	1	6	3	7			
MEM 1	4	2					
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 5, 8, 9 }

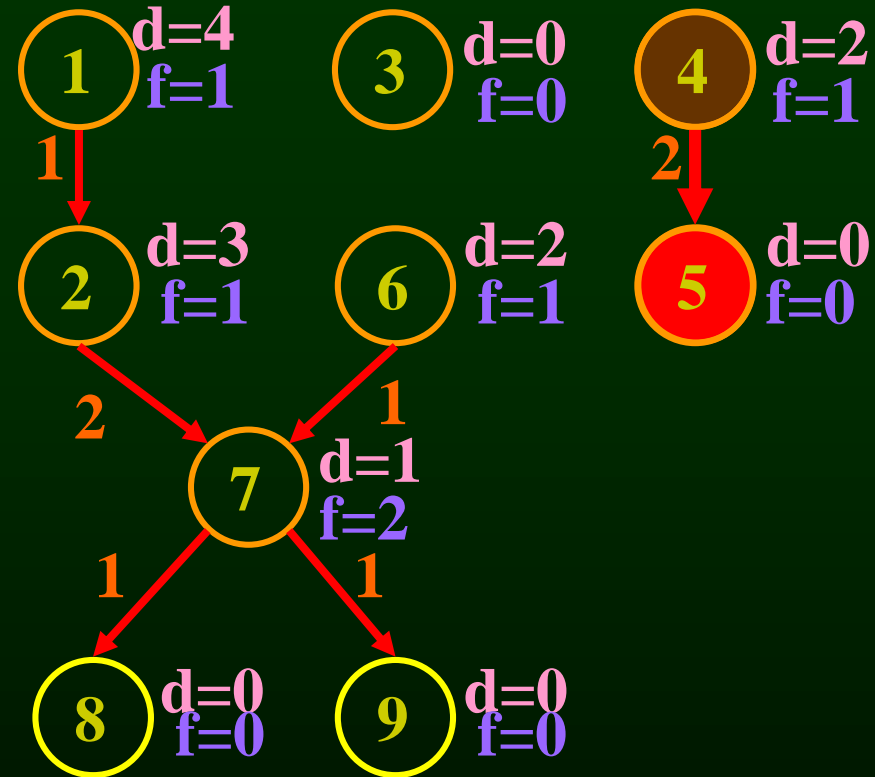


ALUOp	1	6	3	7			
MEM 1	4	2					
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 5, 8, 9 }

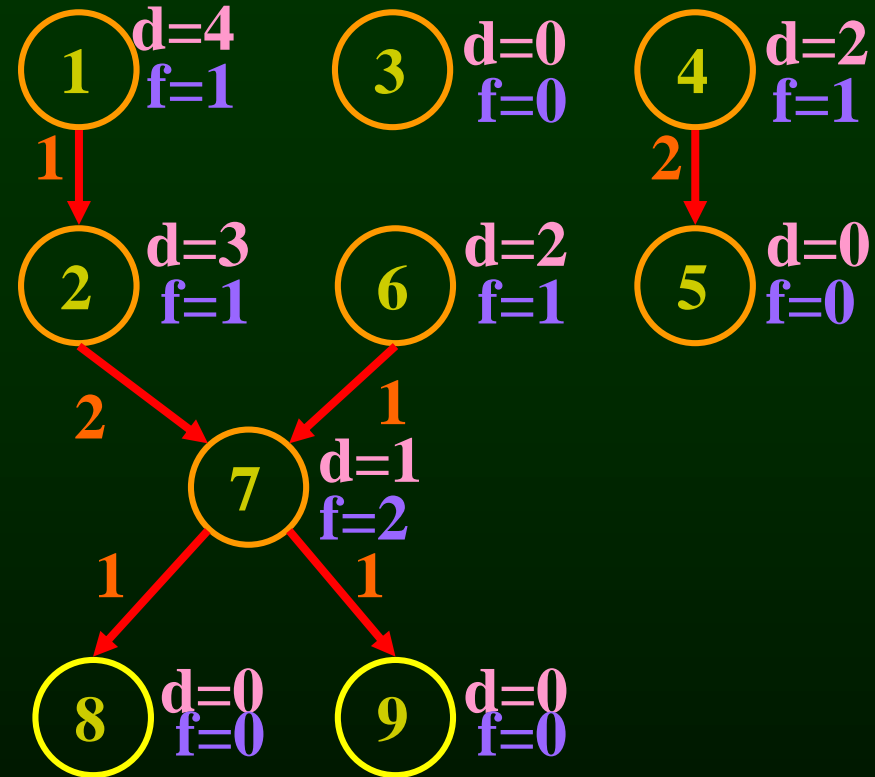


ALUOp	1	6	3	7			
MEM 1	4	2	5				
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 8, 9 }

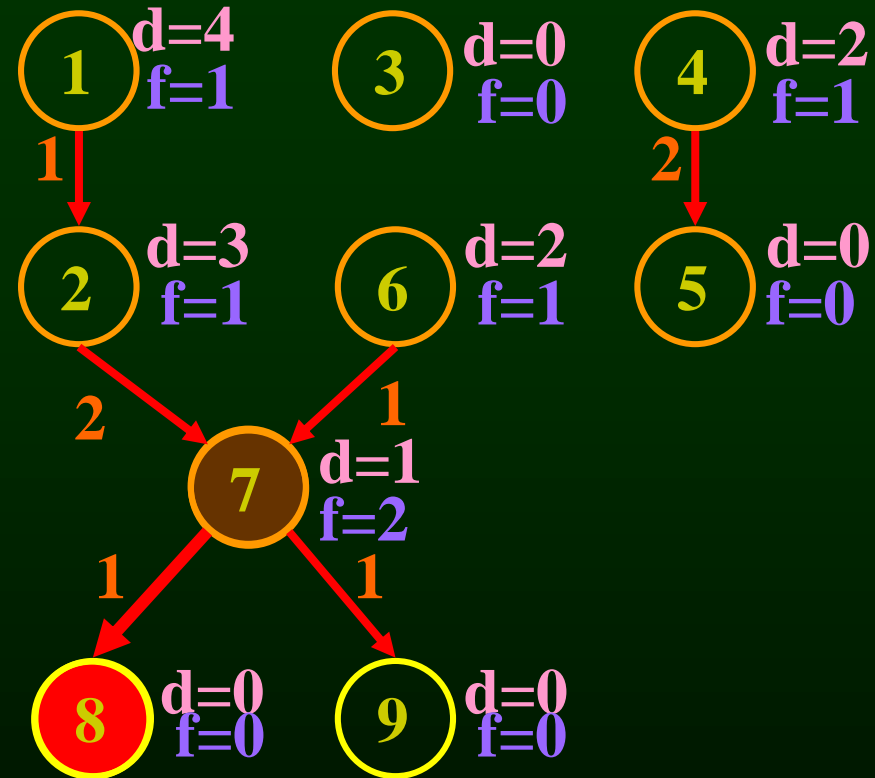


ALUOp	1	6	3	7			
MEM 1	4	2	5				
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 8, 9 }



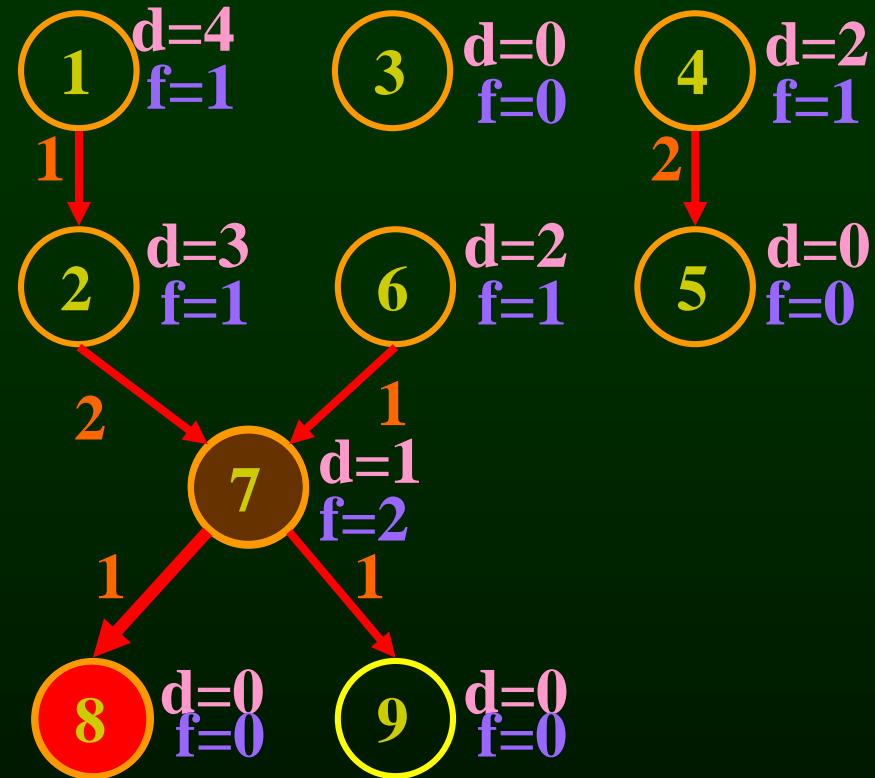
ALUOp	1	6	3	7			
MEM 1	4	2	5				
MEM 2		4	2				

Example

```

1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
    
```

READY = { 8, 9 }

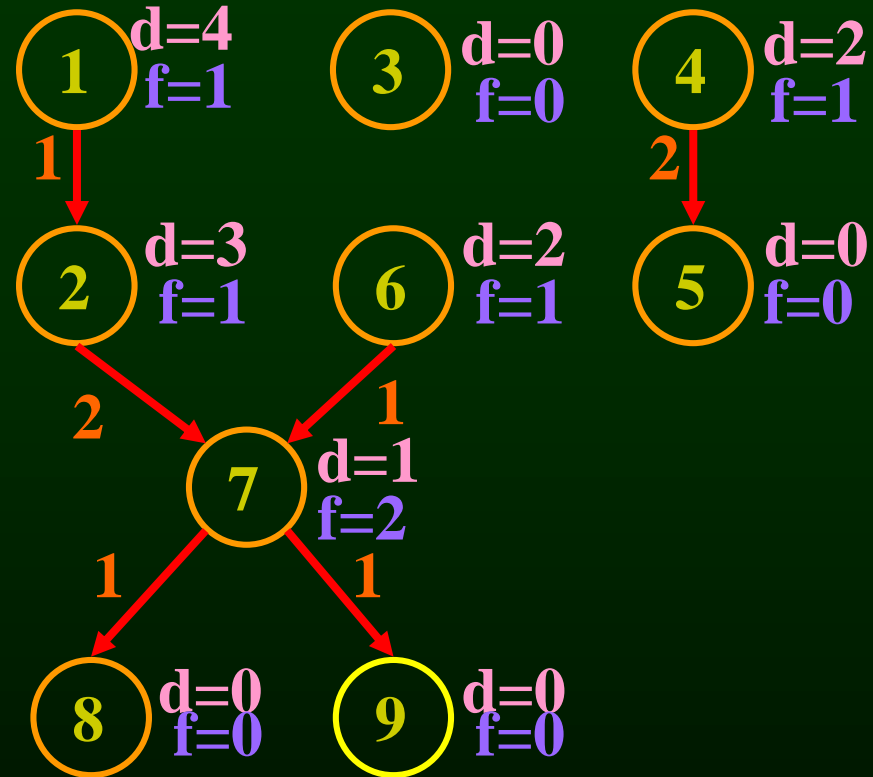


ALUOp	1	6	3	7	8		
MEM 1	4	2	5				
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 9 }

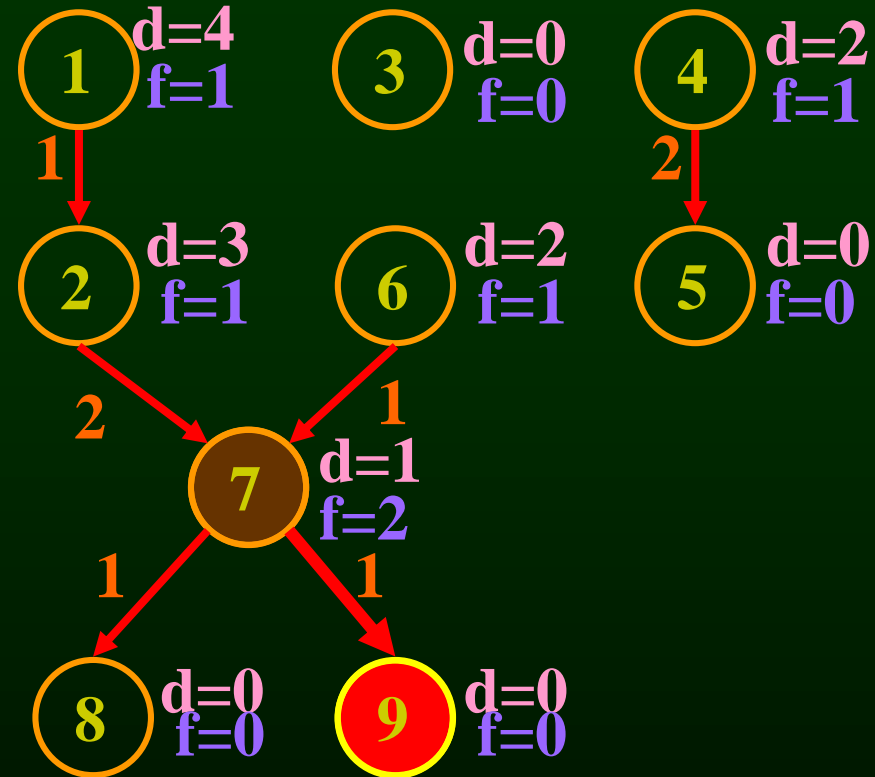


ALUOp	1	6	3	7	8		
MEM 1	4	2	5				
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 9 }

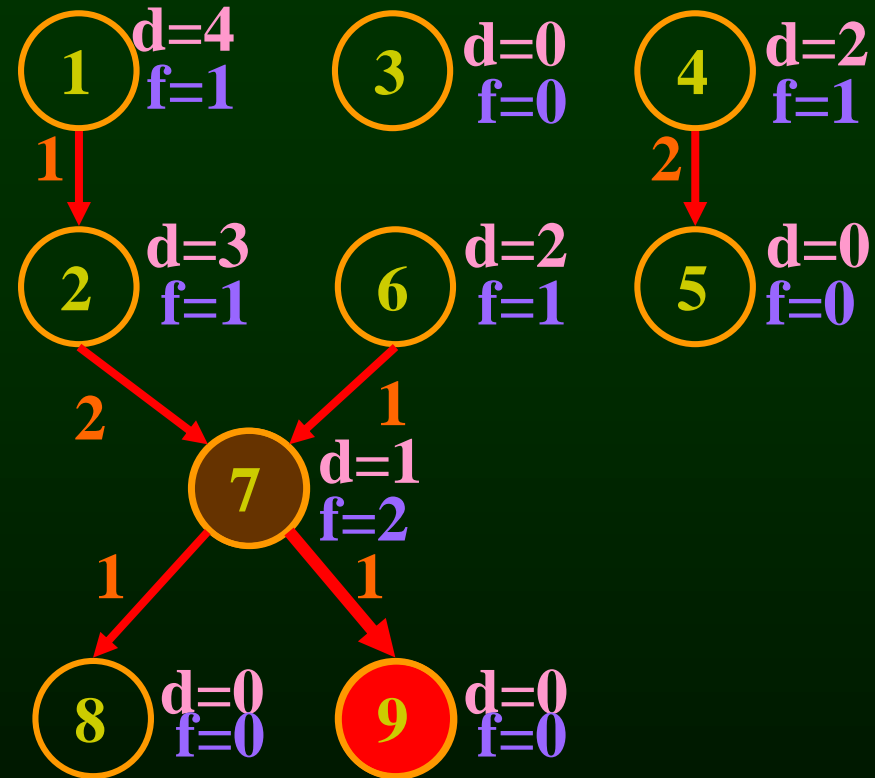


ALUOp	1	6	3	7	8		
MEM 1	4	2	5				
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov     4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and     $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { 9 }

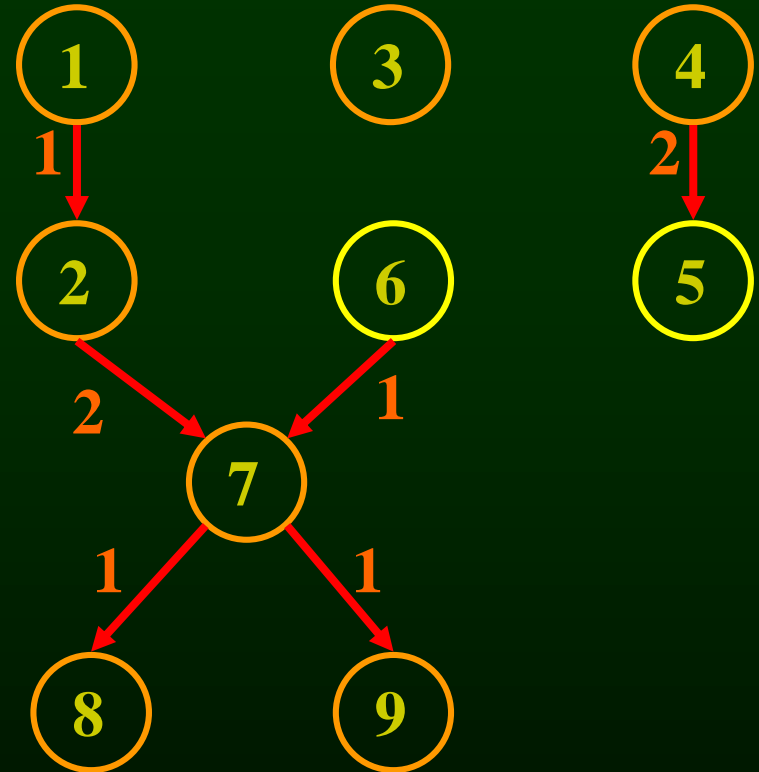


ALUOp	1	6	3	7	8		
MEM 1	4	2	5		9		
MEM 2		4	2				

Example

```
1: lea    var_a, %rax
2: add    4(%rsp), %rax
3: inc    %r11
4: mov    4(%rsp), %r10
5: mov    %r10, 8(%rsp)
6: and    $0x00ff, %rbx
7: imul   %rax, %rbx
8: lea    var_b, %rax
9: mov    %rbx, 16(%rsp)
```

READY = { }



ALUOp	1	6	3	7	8		
MEM 1	4	2	5		9		
MEM 2		4	2				

Outline

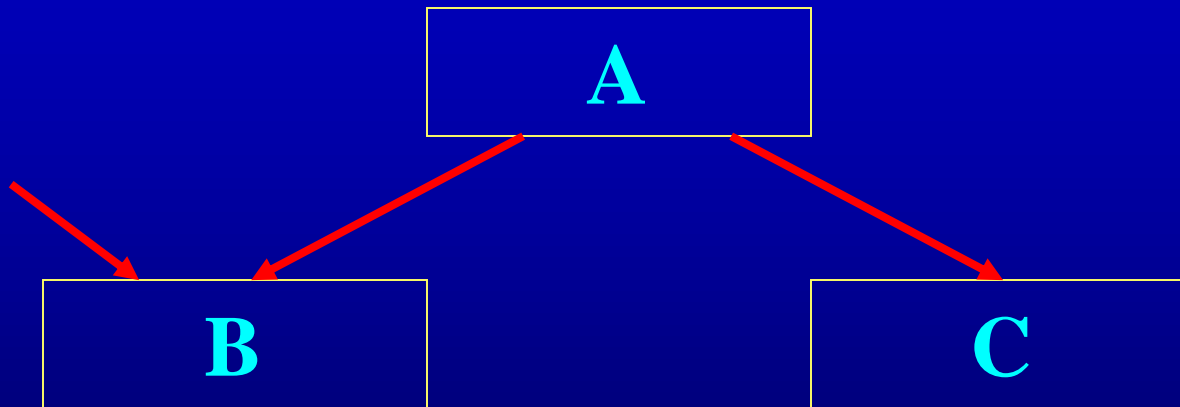
- Modern architectures
- Introduction to instruction scheduling
- List scheduling
- Resource constraints
- Scheduling across basic blocks
- Trace scheduling

Scheduling across basic blocks

- Number of instructions in a basic block is small
 - Cannot keep a multiple units with long pipelines busy by just scheduling within a basic block
- Need to handle control dependence
 - Scheduling constraints across basic blocks
 - Scheduling policy

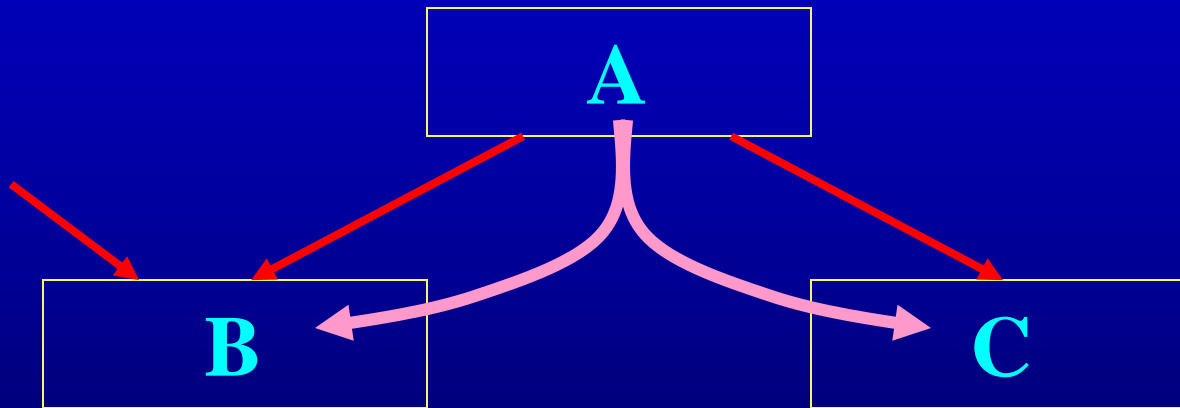
Moving across basic blocks

- Downward to adjacent basic block



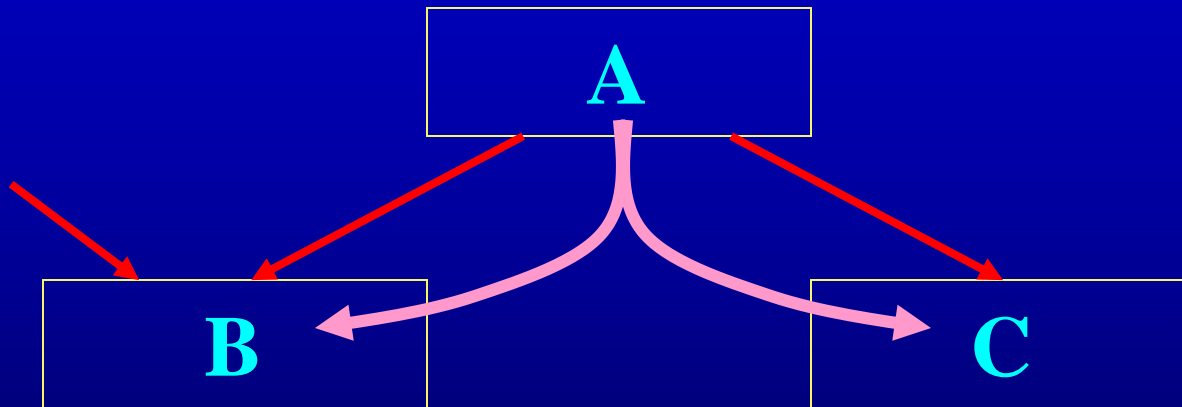
Moving across basic blocks

- Downward to adjacent basic block



Moving across basic blocks

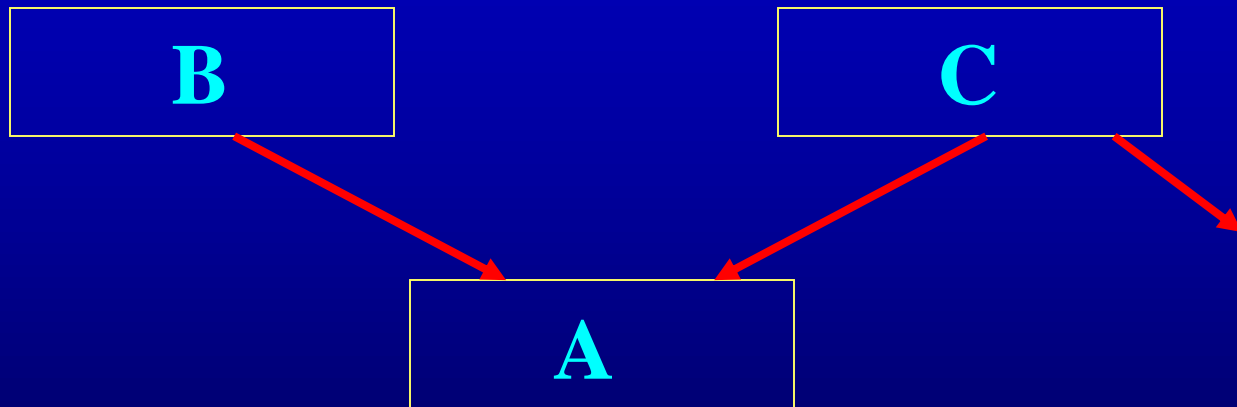
- Downward to adjacent basic block



- A path to B that does not execute A?

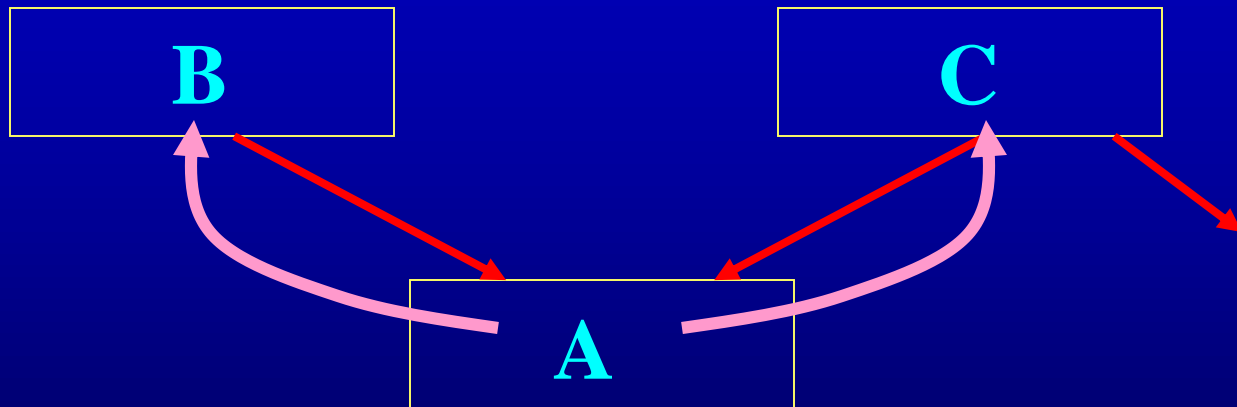
Moving across basic blocks

- Upward to adjacent basic block



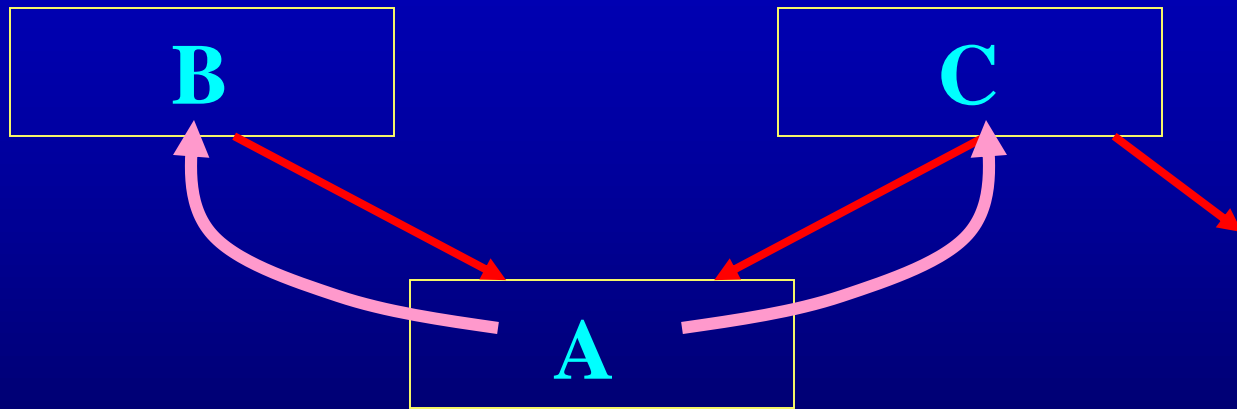
Moving across basic blocks

- Upward to adjacent basic block



Moving across basic blocks

- Upward to adjacent basic block



- A path from C that does not reach A?

Control Dependencies

- Constraints in moving instructions across basic blocks

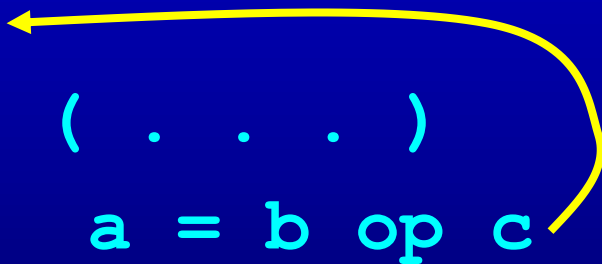
Control Dependencies

- Constraints in moving instructions across basic blocks

```
if ( . . . )  
    a = b op c
```


Control Dependencies

- Constraints in moving instructions across basic blocks

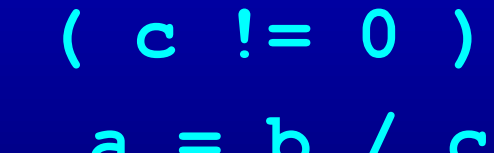


```
if ( . . . )  
    a = b op c
```

A yellow curved arrow originates from the end of the assignment statement `a = b op c` and points back to the opening parenthesis of the `if` statement, illustrating a control dependency where the execution of the assignment is conditional on the `if` statement's condition.

Control Dependencies

- Constraints in moving instructions across basic blocks



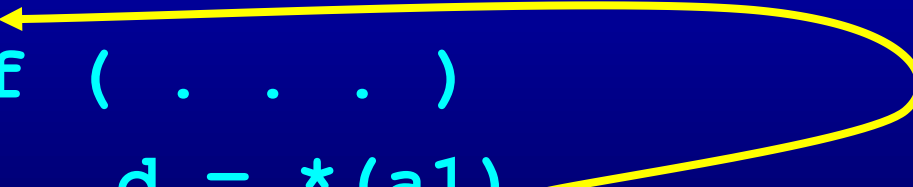
```
if ( c != 0 )  
    a = b / c
```

NO!!!

Control Dependencies

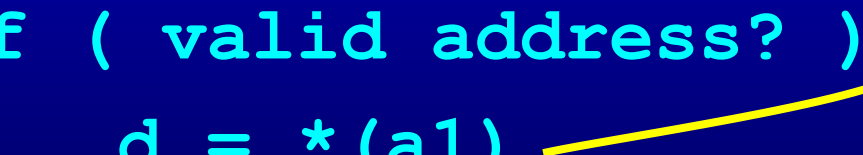
- Constraints in moving instructions across basic blocks

```
If ( . . . )  
    d = *(a1)
```



Control Dependencies

- Constraints in moving instructions across basic blocks



```
If ( valid address? )  
    d = *(a1)
```

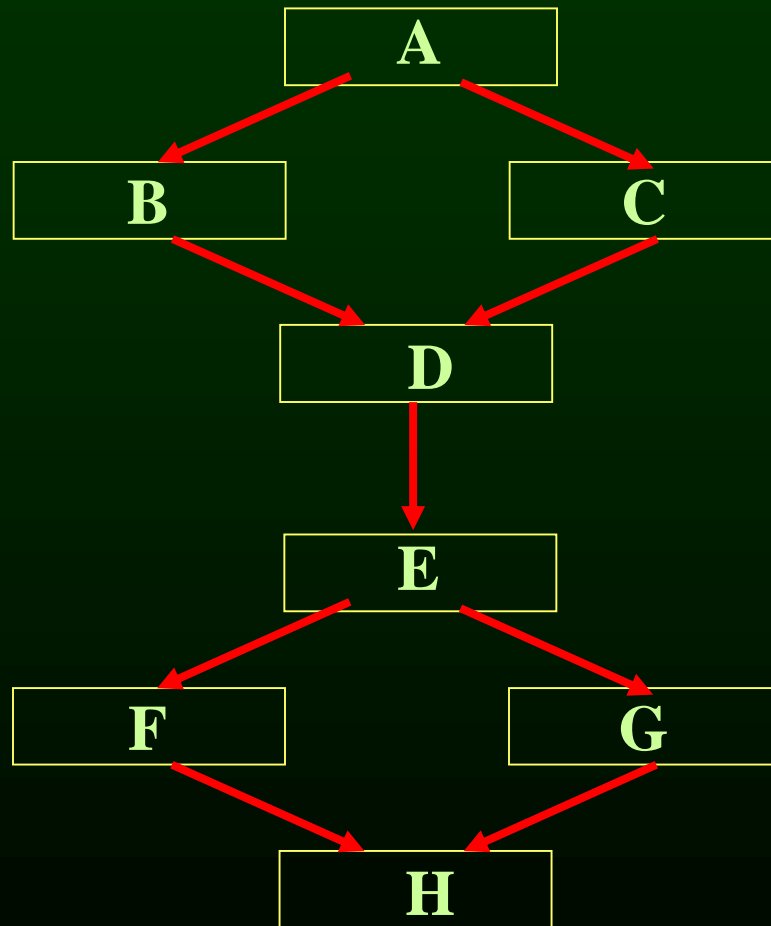
Outline

- Modern architectures
- Introduction to instruction scheduling
- List scheduling
- Resource constraints
- Scheduling across basic blocks
- Trace scheduling

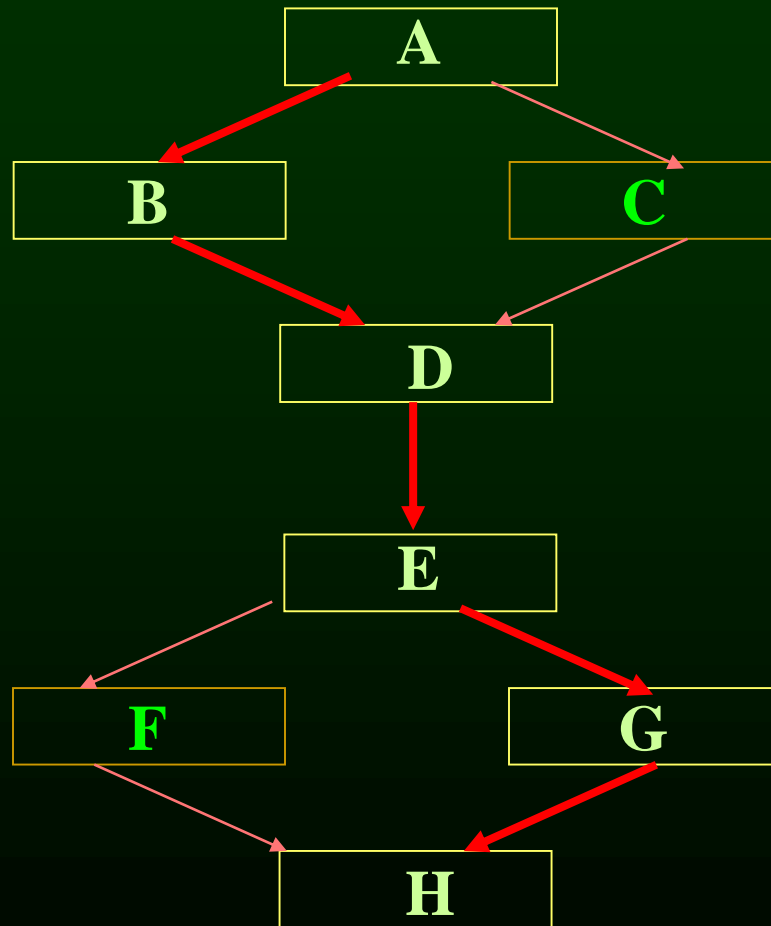
Trace Scheduling

- Find the most common trace of basic blocks
 - Use profile information
- Combine the basic blocks in the trace and schedule them as one block
- Create clean-up code if the execution goes off-trace

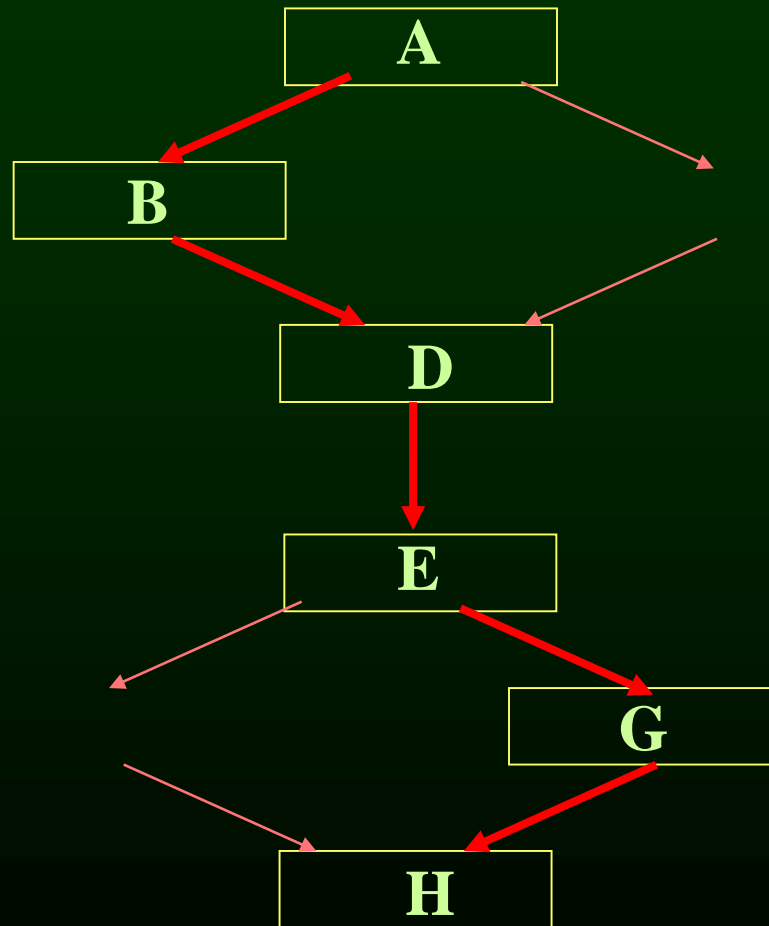
Trace Scheduling



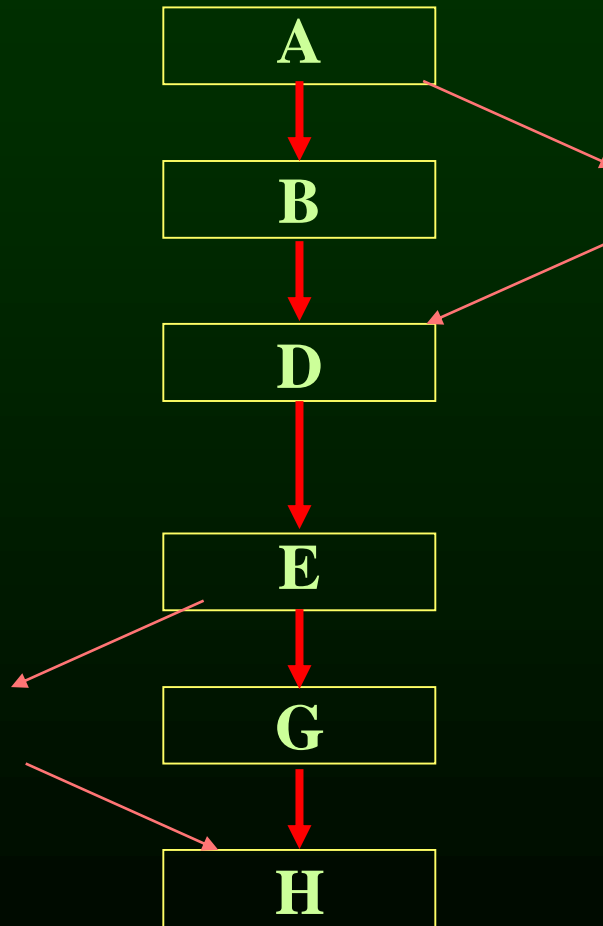
Trace Scheduling



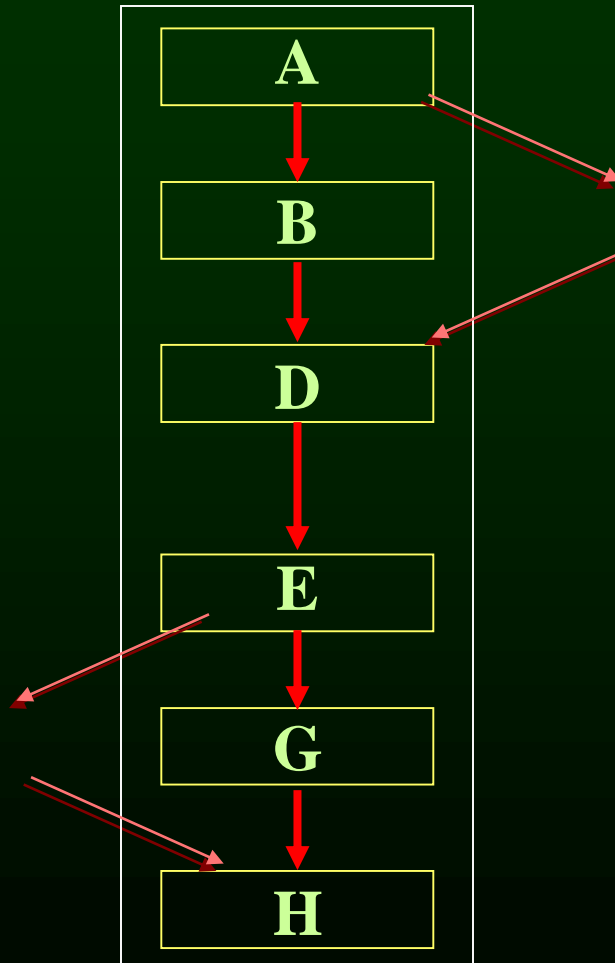
Trace Scheduling



Trace Scheduling

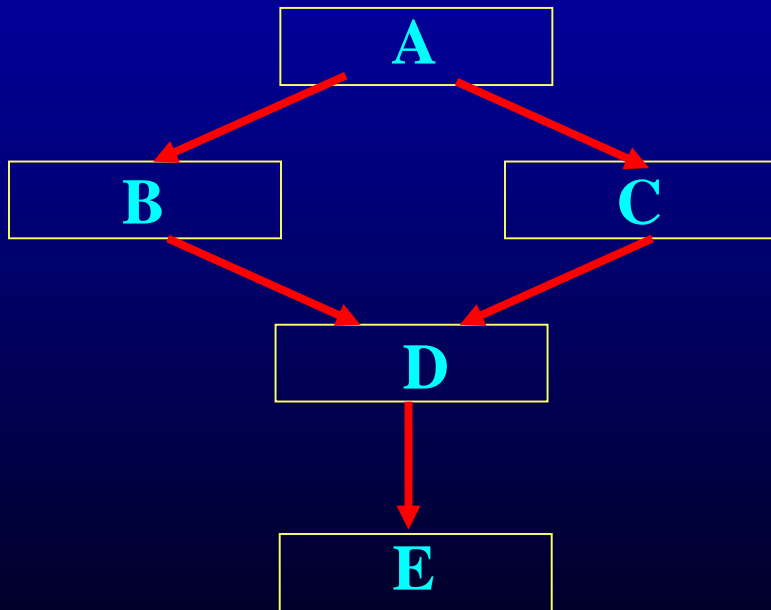


Trace Scheduling



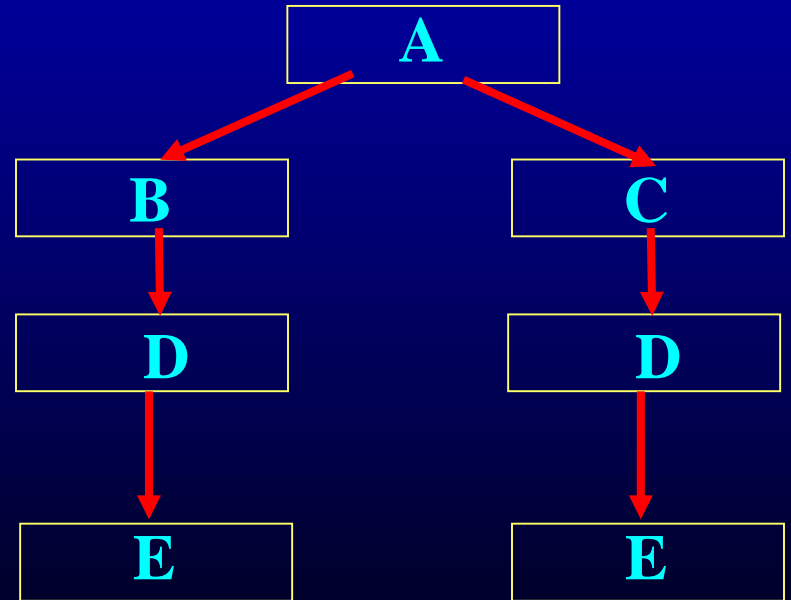
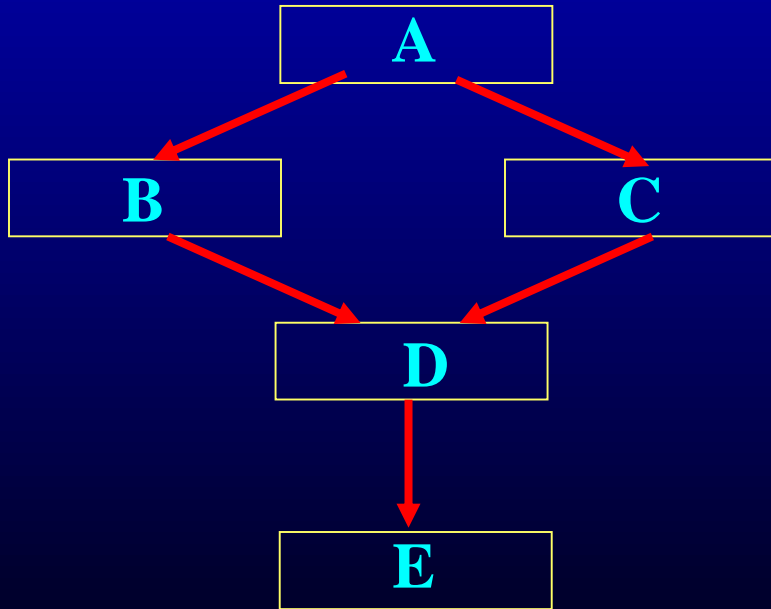
Large Basic Blocks via Code Duplication

- Creating large extended basic blocks by duplication
- Schedule the larger blocks

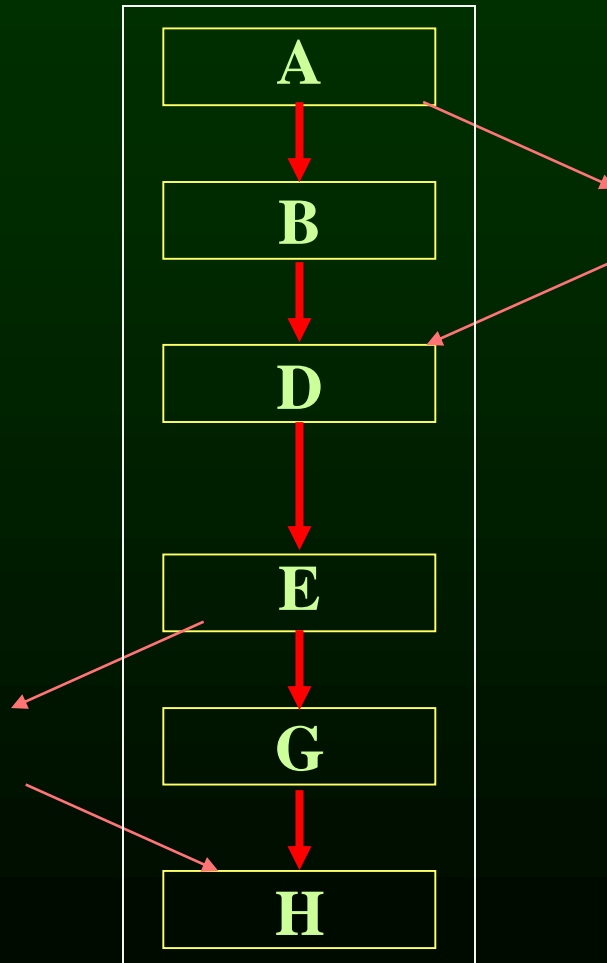


Large Basic Blocks via Code Duplication

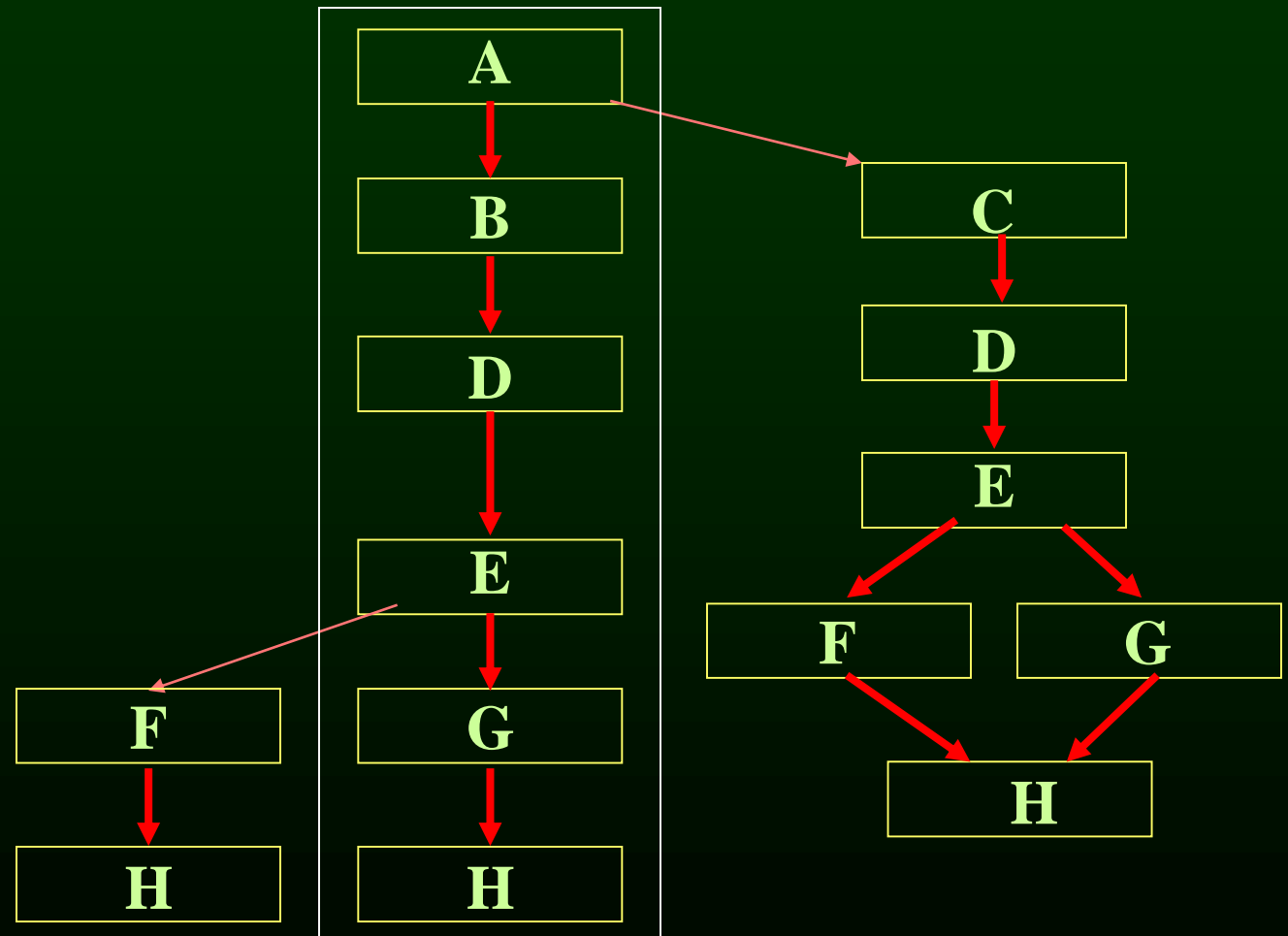
- Creating large extended basic blocks by duplication
- Schedule the larger blocks



Trace Scheduling



Trace Scheduling



Next

- Scheduling for loops
- Loop unrolling
- Software pipelining
- Interaction with register allocation
- Hardware vs. Compiler