

Ex.No.8  
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## IMPLEMENTATION OF ELLIPTIC CURVE CRYPTOSYSTEM

### AIM:

To implement elliptic curve key exchange and elliptic curve encryption algorithms.

### THEORY

#### Affine points:

In elliptic curve cryptography (ECC), an affine point is represented as a pair of coordinates  $(x, y)$  that lie on the elliptic curve. These points are defined over a finite field and follow the curve's equation, typically  $y^2 = x^3 + ax + b$ . Affine points are used in ECC operations.

#### Elliptic Curve Discrete logarithms:

The elliptic curve discrete logarithm problem (ECDLP) involves finding the integer  $k$  given two points  $P$  and  $Q$  on an elliptic curve such that  $Q = k \cdot P$ . ECDLP is computationally hard, making ECC secure for encryption, digital signatures, and key exchange.

#### Point Addition:

Point addition in ECC refers to adding two distinct points  $P$  and  $Q$  on an elliptic curve to get another point  $R$ . This operation follows specific algebraic rules that depend on the curve's equation. If  $P=Q$ , it becomes "point doubling."

#### Scalar Multiplication:

Scalar multiplication in ECC is the process of repeatedly adding a point  $P$  to itself  $k$  times, denoted as  $k \cdot P$ . This is the fundamental operation in elliptic curve cryptography, used in key generation, encryption, and signatures.

#### Inverse of an affine point:

The inverse of an affine point  $P = (x, y)$  on an elliptic curve is  $-P = (x, -y)$ . This is used in elliptic curve operations, such as point subtraction, where  $P - Q$  is defined as  $P + (-Q)$ .

**ALGORITHM:****Elliptic Curve Key Exchange****1. Elliptic Curve Parameters Selection:**

Both parties agree on the public parameters:

- a) A prime number  $p$  defining the finite field  $F_p$ .
- b) Coefficients  $a$  and  $b$  defining the elliptic curve equation  $y^2 = x^3 + ax + b \pmod{p}$ .
- c) A base point  $G = (x_g, y_g)$  on the curve, and its order  $n$ .

**2. Key Generation by Alice:**

- i) Alice chooses a random private key  $d_a$ , where  $d_a \in [1, n-1]$ .
- ii) Alice computes her public key:  $P_a = d_a * G$  (scalar multiplication).
- iii) Alice sends  $P_a$  to Bob.

**3. Key Generation by Bob:**

- i) Bob chooses a random private key  $d_b$ , where  $d_b \in [1, n-1]$ .
- ii) Bob computes his public key:  $P_b = d_b * G$ .
- iii) Bob sends  $P_b$  to Alice.

**4. Shared Secret Calculation:**

- i) Alice computes the shared secret:  $S_b = d_b * P_a$ .
- ii) Bob computes the shared secret:  $S_b = d_b * P_a$ .
- iii) Both  $S_a$  and  $S_b$  are the same point on the curve due to the commutative property of scalar multiplication:  $S = d_a * d_b * G$ .

**5. Shared Secret Extraction:**

Both parties extract the  $x$  coordinate of the shared point  $S$  as the shared secret key.

**Elliptic Curve Encryption:****1. Elliptic Curve Parameters:**

Both parties agree on public parameters:

A prime  $p$  defining the finite field  $F_p$ .

Coefficients  $a$  and  $b$  for the elliptic curve equation  $y^2 = x^3 + ax + b \pmod{p}$ .

A base point  $G = (x_G, y_G)$  on the curve, and its order  $n$ .

**2. Key Generation (By the Receiver):**

The receiver (Bob) chooses a private key  $d_B$  where  $d_B$  in  $[1, n-1]$ .

The receiver computes their public key  $P_B = d_B * G$  and shares  $P_B$  with the sender (Alice).

### 3. Message Encoding (By the Sender):

Alice represents the plaintext message  $M$  as a point on the elliptic curve  $Min E(F_p)$ . If the message is too large, a hashing or encoding function is used to map it to the curve.

### 4. Encryption (By the Sender):

Alice chooses a random integer  $k$  in  $[1, n-1]$ .

Alice computes two values:

$C_1 = k * G$  (a point on the curve).

$C_2 = M + k * P_B$  (the encrypted message point).

The ciphertext is  $(C_1, C_2)$ , and Alice sends this to Bob.

### 5. Decryption (By the Receiver):

Upon receiving  $(C_1, C_2)$ , Bob computes the shared secret  $S = d_B * C_1$ .

Bob then recovers the original message point by computing:  $M = C_2 - S$ .

Since  $S = d_B * C_1 = d_B * (k * G) = k * P_B$ , the original message  $M$  is recovered correctly.

## CODING:

### ECC.py

```
import random
import hashlib
from Crypto.Cipher import AES
from Crypto.Util.Padding import pad,unpad

class ECC:
    def __init__(self,a,b,p,g):
        self.a=a
        self.b=b
        self.p=p
        self.g=g
    def is_point_on_curve(self,point):
        x,y=point
        return (y**2)%self.p==(x**3+self.a*x+self.b)%self.p
```

```

def point_addition(self,P,Q):
    if P==Q:
        return self.point_doubling(P)
    x1,y1=P
    x2,y2=Q
    if x1==x2 and y1 != y2:
        return None
    try:
        slope=(y2 - y1) * pow(x2 - x1,-1,self.p) % self.p
    except ValueError:
        raise ValueError("Cannot compute the modular inverse (points are not suitable for
addition).")
    xr=(slope ** 2 - x1 - x2) % self.p
    yr=(slope * (x1 - xr) - y1) % self.p
    return xr,yr
def point_doubling(self,P):
    x,y=P
    slope=(3*x**2+self.a)*pow(2*y,-1,self.p)%self.p
    xr=(slope**2-2*x)%self.p
    yr=(slope*(x-xr)-y)%self.p
    return xr,yr
def scalar_multiplication(self,k,P):
    result=None
    addend=P
    while k:
        if k & 1:
            if result is None:
                result=addend
            else:
                result=self.point_addition(result,addend)

        addend=self.point_doubling(addend)
        k >>= 1
    return result
def generate_private_key(self):
    return random.randint(1,self.p-1)
def generate_public_key(self,private_key):
    return self.scalar_multiplication(private_key,self.g)

```

```

def derive_shared_secret(self,private_key,public_key):
    return self.scalar_multiplication(private_key,public_key)

def hash_shared_secret(self,secret_point):
    secret_x=secret_point[0]
    shared_key=hashlib.sha256(str(secret_x).encode()).digest()
    return shared_key

def encrypt_message(self,public_key,message):
    ephemeral_private_key=self.generate_private_key()
    ephemeral_public_key=self.generate_public_key(ephemeral_private_key)
    shared_secret=self.derive_shared_secret(ephemeral_private_key,public_key)
    shared_key=self.hash_shared_secret(shared_secret)
    cipher=AES.new(shared_key,AES.MODE_CBC)
    ciphertext=cipher.encrypt(pad(message.encode(),AES.block_size))
    return (ciphertext,cipher.iv,ephemeral_public_key)

def decrypt_message(self,private_key,ciphertext,iv,ephemeral_public_key):
    shared_secret=self.derive_shared_secret(private_key,ephemeral_public_key)
    shared_key=self.hash_shared_secret(shared_secret)
    cipher=AES.new(shared_key,AES.MODE_CBC,iv)
    plaintext=unpad(cipher.decrypt(ciphertext),AES.block_size)
    return plaintext.decode()

```

#### main.py

```

from ECC import ECC

def run_ecdh(ecc,alice_private_key,alice_public_key,bob_private_key,bob_public_key):
    alice_shared_secret=ecc.derive_shared_secret(alice_private_key,bob_public_key)
    print("Alice's Shared Secret:",alice_shared_secret)

    bob_shared_secret=ecc.derive_shared_secret(bob_private_key,alice_public_key)
    print("Bob's Shared Secret:",bob_shared_secret)

    if alice_shared_secret == bob_shared_secret:
        print("Key exchange successful! Shared secret is identical for both parties.")
    else:
        print("Key exchange failed! Shared secret mismatch.")

```

```

def run_ecies(ecc,bob_public_key,bob_private_key):
    message=input("Enter the Message to be encrypted.")
    ciphertext,iv,ephemeral_public_key=ecc.encrypt_message(bob_public_key,message)
    print("Ciphertext:",ciphertext)

decrypted_message=ecc.decrypt_message(bob_private_key,ciphertext,iv,ephemeral_public_
key)
    print("Decrypted Message:",decrypted_message)

def main():
    a=2
    b=3
    p=13
    g=(3,6)
    ecc=ECC(a,b,p,g)
    alice_private_key=ecc.generate_private_key()
    alice_public_key=ecc.generate_public_key(alice_private_key)
    bob_private_key=ecc.generate_private_key()
    bob_public_key=ecc.generate_public_key(bob_private_key)
    print("Alice's Private Key:",alice_private_key)
    print("Alice's Public Key:",alice_public_key)
    print("Bob's Private Key:",bob_private_key)
    print("Bob's Public Key:",bob_public_key)
    choices={'1': run_ecdh,'2': run_ecies}
    print("\nChoose an option:")
    print("1: Elliptic Curve Diffie-Hellman Key Exchange (ECDH)")
    print("2: Elliptic Curve Integrated Encryption Scheme (ECIES)")
    choice=input("Enter your choice (1 or 2): ")
    if choice == '1':
        choices[choice](ecc,alice_private_key,alice_public_key,bob_private_key,bob_public_key)

    elif choice == '2':
        choices[choice](ecc,bob_public_key,bob_private_key)
    else:
        print("Invalid choice. Please enter 1 or 2.")

```

```
if __name__ == "__main__":
    main()
```

### SCREEN SHOTS:

```
C:\Users\gokul\AppData\Local\Programs\Python\Python310\python.exe "C:\
Alice's Private Key: 5
Alice's Public Key: (3, 7)
Bob's Private Key: 7
Bob's Public Key: (3, 6)

Choose an option:
1: Elliptic Curve Diffie-Hellman Key Exchange (ECDH)
2: Elliptic Curve Integrated Encryption Scheme (ECIES)
Enter your choice (1 or 2): 1
Alice's Shared Secret: (3, 7)
Bob's Shared Secret: (3, 7)
Key exchange successful! Shared secret is identical for both parties.
```

```
Choose an option:
1: Elliptic Curve Diffie-Hellman Key Exchange (ECDH)
2: Elliptic Curve Integrated Encryption Scheme (ECIES)
Enter your choice (1 or 2): 2
Enter the Message to be encrypted.gokul
Ciphertext: b'z\n\xa9\x1ey\x08\x7f\xd0T\xb0Z\x84\x92PB\xd3'
Decrypted Message: gokul
```

### RESULT:

Thus, we have implemented elliptic curve key exchange and elliptic curve encryption algorithms successfully.

### Evaluation

Parameter	Max Marks	Marks Obtained
Uniqueness of the Code	50	
Completion of experiment on time	10	
Documentation	15	
Total	75	
Signature of the faculty with Date		