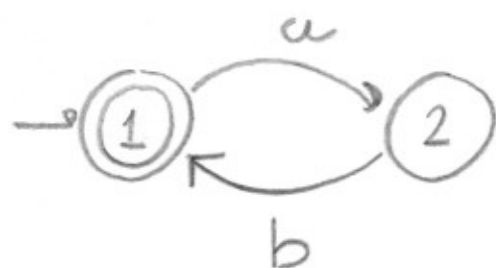


Prodotto di Automi $\Sigma = \{a, b\}$ ①

$$R_1 = (ab)^*$$

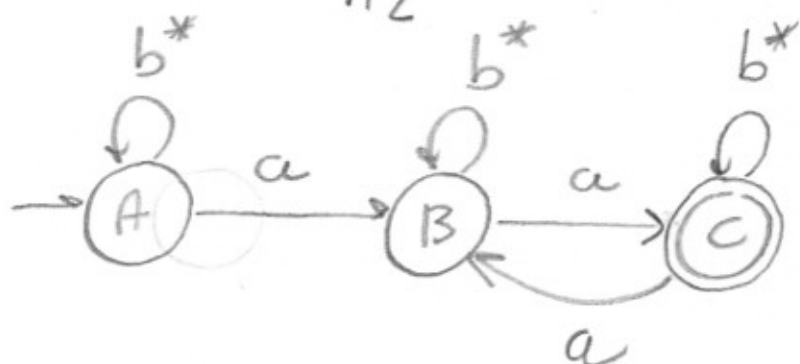
$$R_2 = (\Sigma^* a \Sigma^* a \Sigma^*)^+$$

A_1



$$Q_1 = \{1, 2\}$$

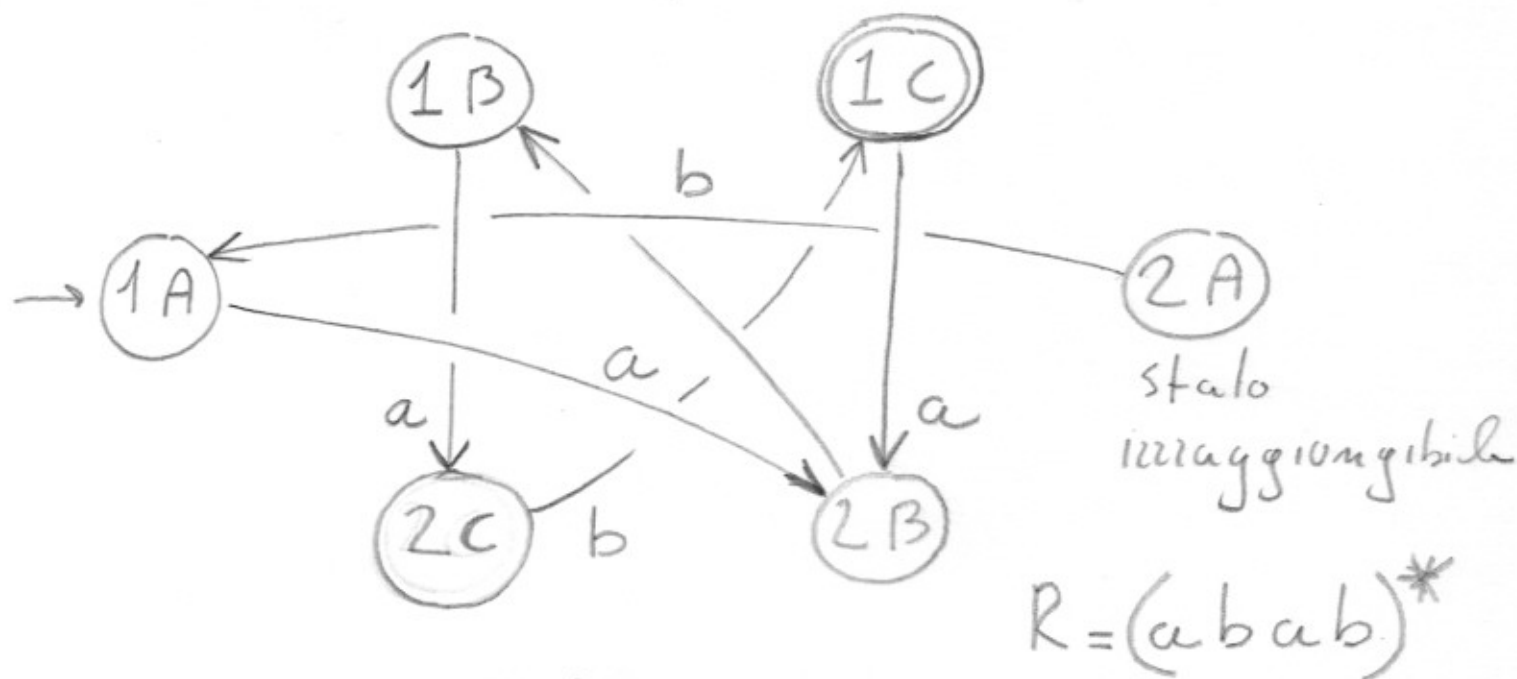
A_2



$$Q_2 = \{A, B, C\}$$

$$Q = Q_1 \times Q_2 = \{(1, A), (1, B), (1, C), \dots\}$$

$$A = A_1 \times A_2 \quad L(A) = L(A_1) \cap L(A_2)$$



poi pulire l'automata ($2A$ irraggiungibile)

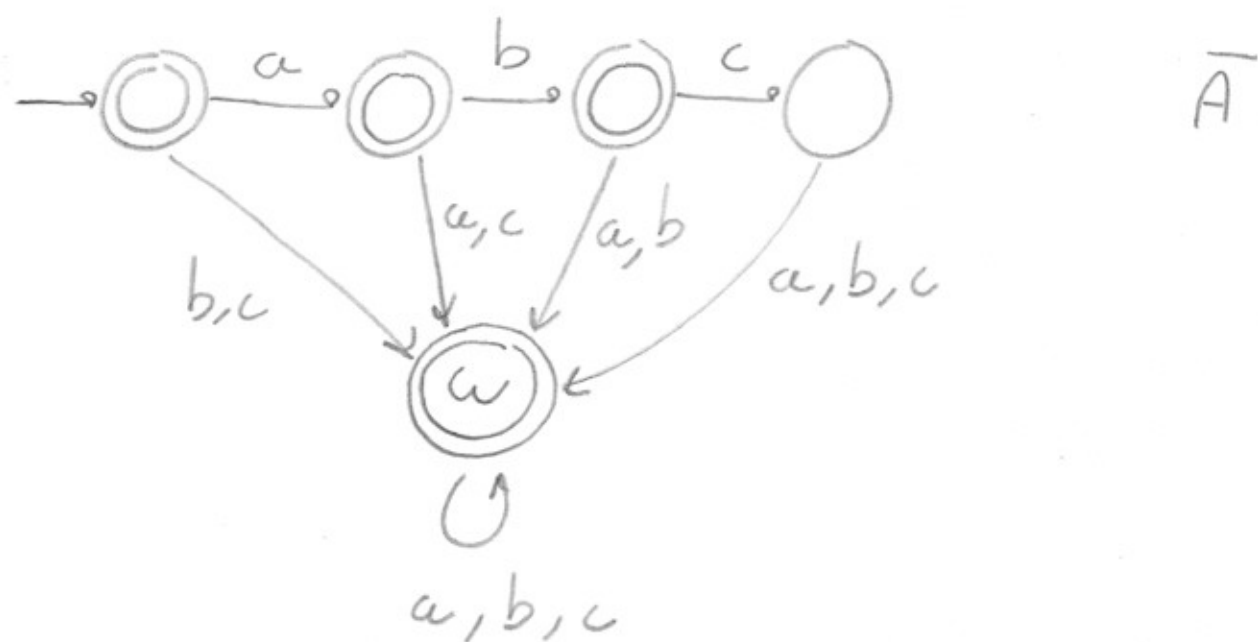
Funziona anche con automi 7 det.

Complemento di automa DET ! (1)

$\Sigma = \{a, b, c\}$



$L(A) = \{abc\}$ (Una sola stringa)



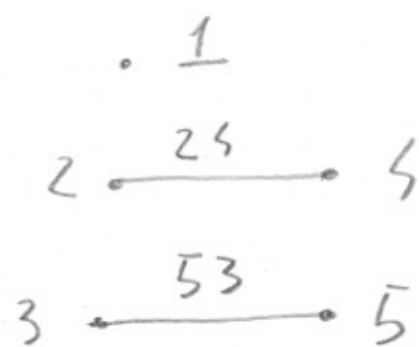
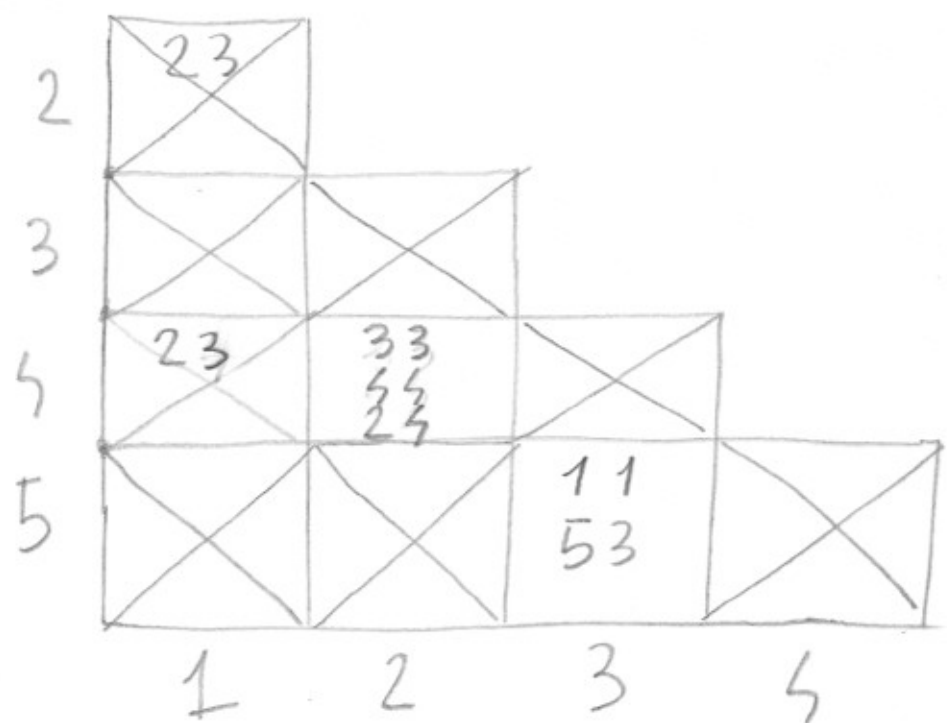
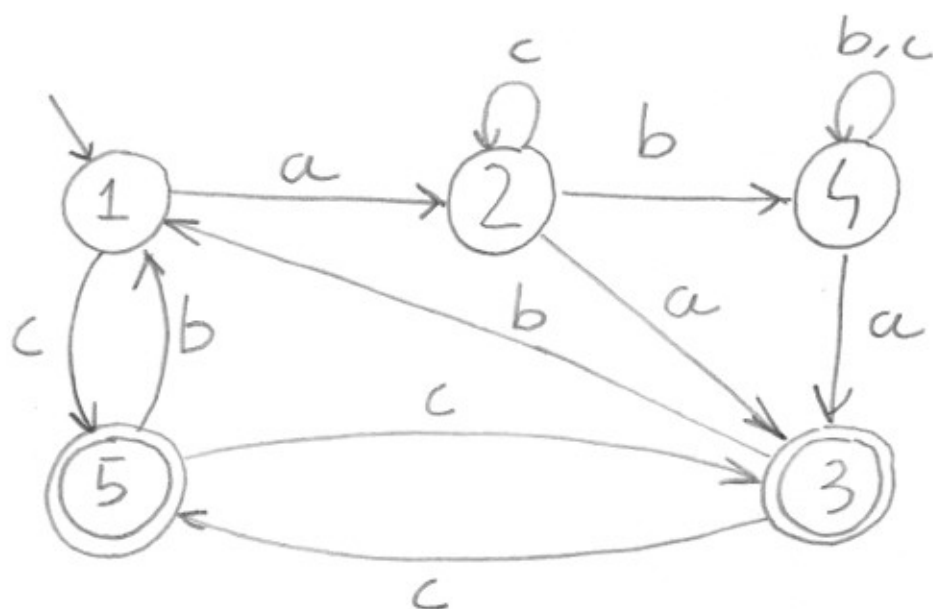
$$\overline{L(A)} = L(\bar{A}) \quad L(\bar{A}) = \Sigma^* - \{abc\}$$

Attenzione: \neg funziona se
l'automato \bar{A} \neg deterministico

Riduzione stati

①

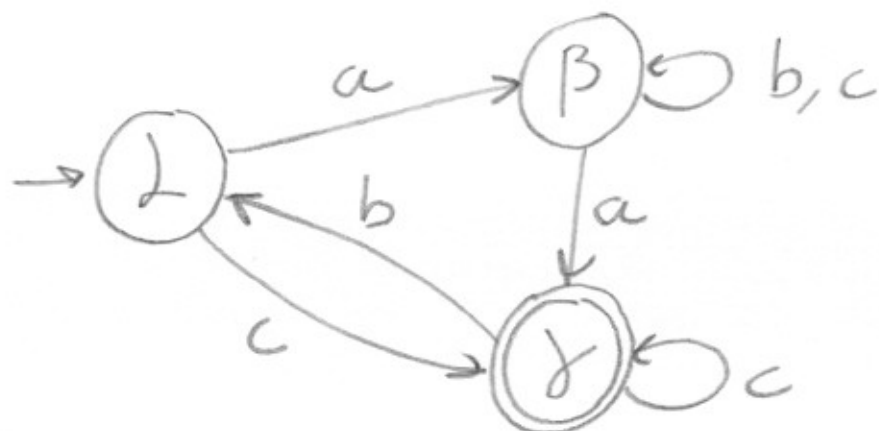
(x automa deterministico) $\Sigma = \{a, b, c\}$



$$\alpha = 1$$

$$\beta = 2, 4$$

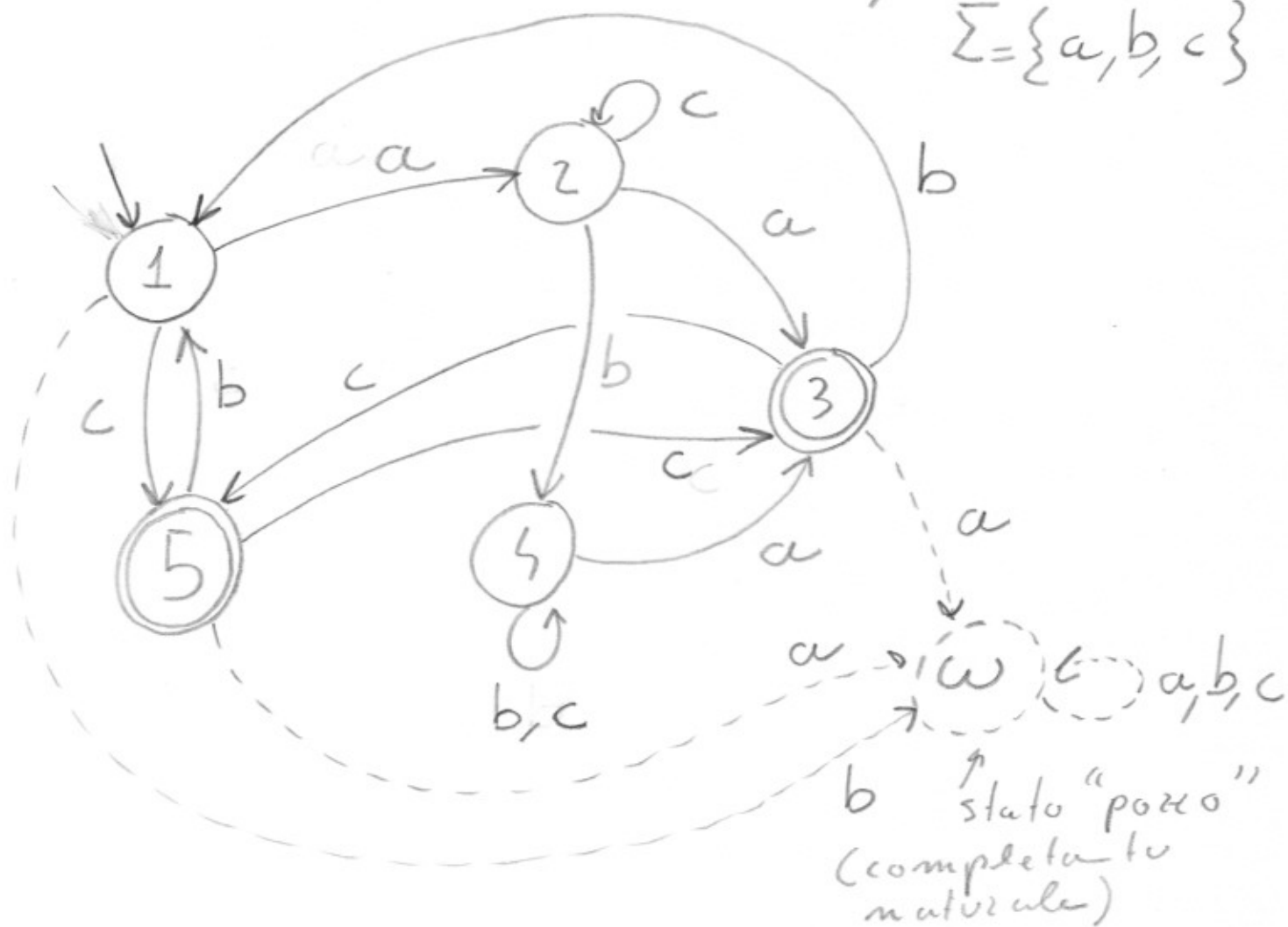
$$\gamma = 3, 5$$



Riduzione stati
(x automa deterministico)

①

$\Sigma = \{a, b, c\}$



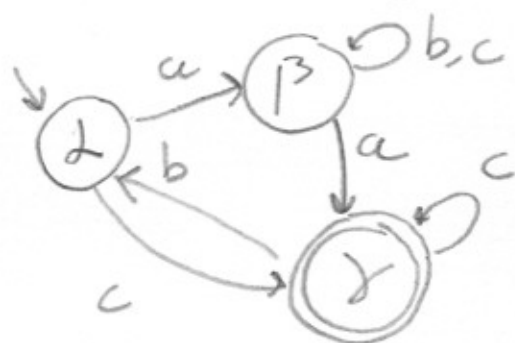
2	$\begin{array}{ c c } \hline 2 & 3 \\ \hline \omega & 4 \\ \hline 5 & 2 \\ \hline \end{array}$			
3	$\begin{array}{ c c } \hline 2 & \omega \\ \hline a & 1 \\ \hline 5 & 5 \\ \hline \end{array}$	$\begin{array}{ c c } \hline 3 & \omega \\ \hline 4 & 1 \\ \hline 2 & 5 \\ \hline \end{array}$		
4	$\begin{array}{ c c } \hline 2 & 3 \\ \hline a & 4 \\ \hline 5 & 4 \\ \hline \end{array}$	$\begin{array}{ c c } \hline 3 & 3 \\ \hline 4 & 4 \\ \hline 2 & 4 \\ \hline \end{array}$	$\begin{array}{ c c } \hline \omega & 3 \\ \hline 1 & 4 \\ \hline 5 & 4 \\ \hline \end{array}$	
5	$\begin{array}{ c c } \hline 2 & \omega \\ \hline a & 1 \\ \hline 5 & 3 \\ \hline \end{array}$	$\begin{array}{ c c } \hline 3 & \omega \\ \hline 4 & 1 \\ \hline 2 & 5 \\ \hline \end{array}$	$\begin{array}{ c c } \hline a & \omega \\ \hline 1 & 1 \\ \hline 5 & 3 \\ \hline \end{array}$	$\begin{array}{ c c } \hline 3 & \omega \\ \hline 4 & 1 \\ \hline 4 & 3 \\ \hline \end{array}$
	1	2	3	4

1

2 $\xrightarrow{24}$ 4

3 $\xrightarrow{53}$ 5

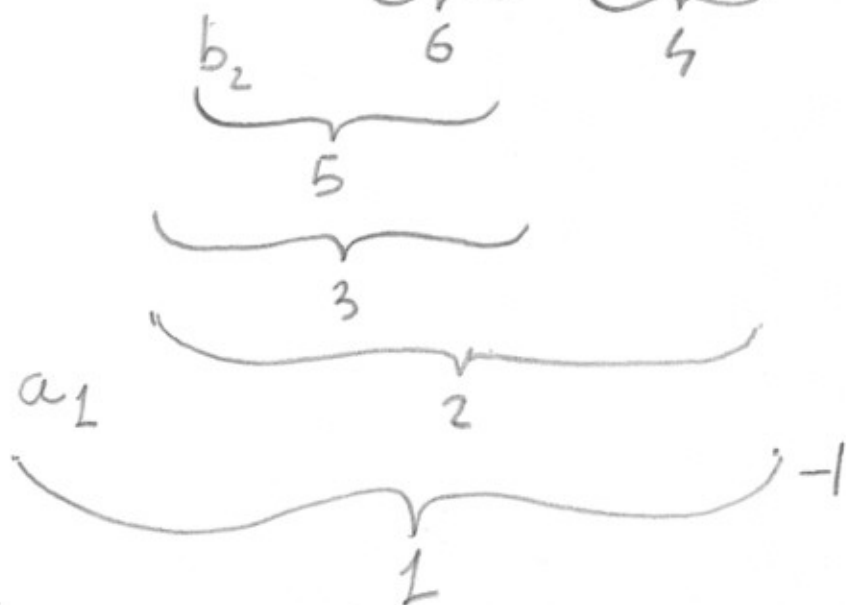
$\alpha = 1 \quad \beta = 2, 4 \quad \gamma = 3, 5$



$$\Sigma = \{a, b, c\}$$

$$R = a(b|cc)^*cb \quad (2)$$

$$R = a_1 (b_2 | c_3 c_4)^* c_5 b_6 \quad -$$



$$R = 1 \cdot -$$

$$1 = a_1 \cdot 2$$

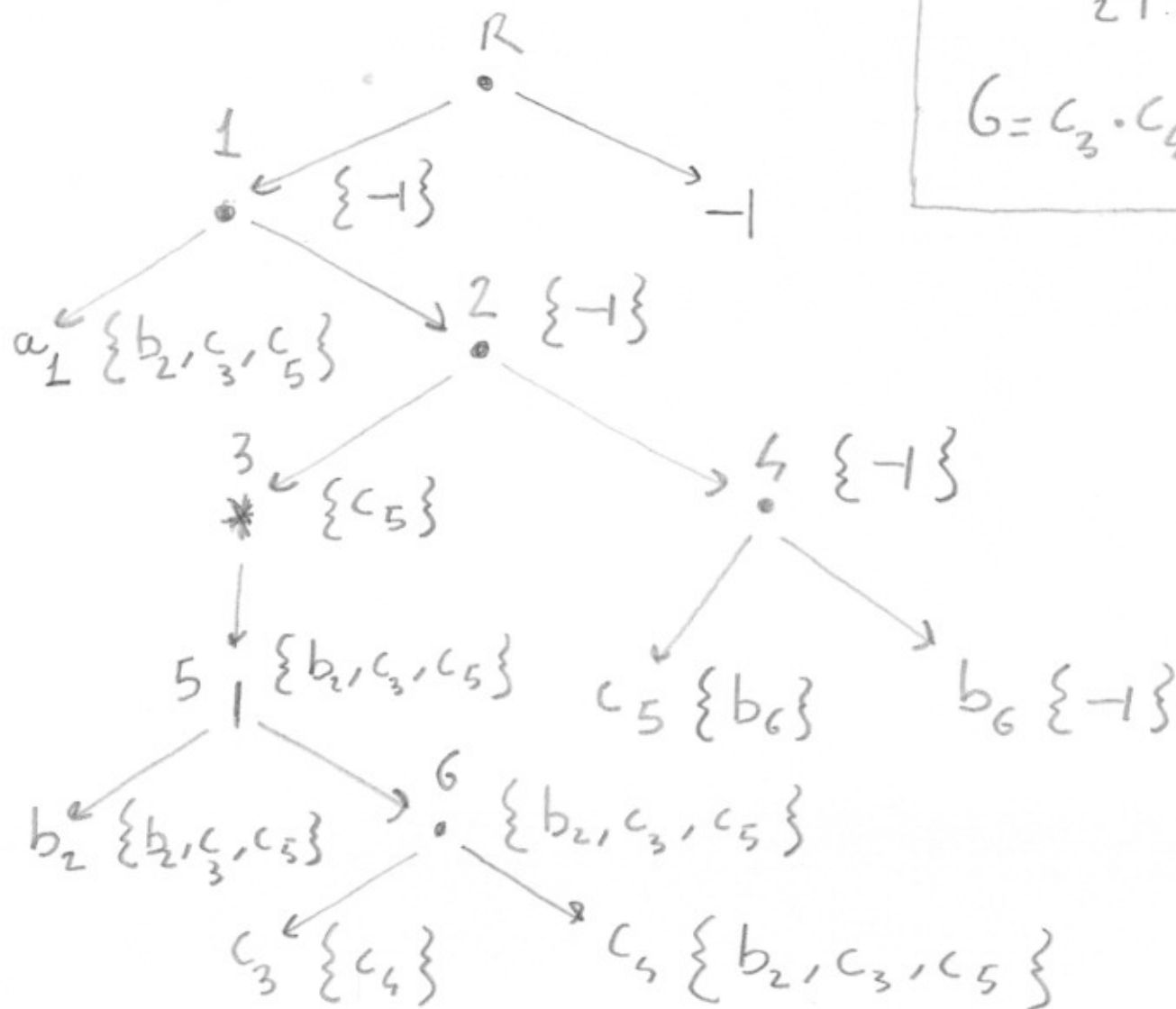
$$2 = 3 \cdot 4$$

$$3 = 5^*$$

$$4 = c_5 \cdot b_6$$

$$5 = b_2 | 6$$

$$6 = c_3 \cdot c_4$$

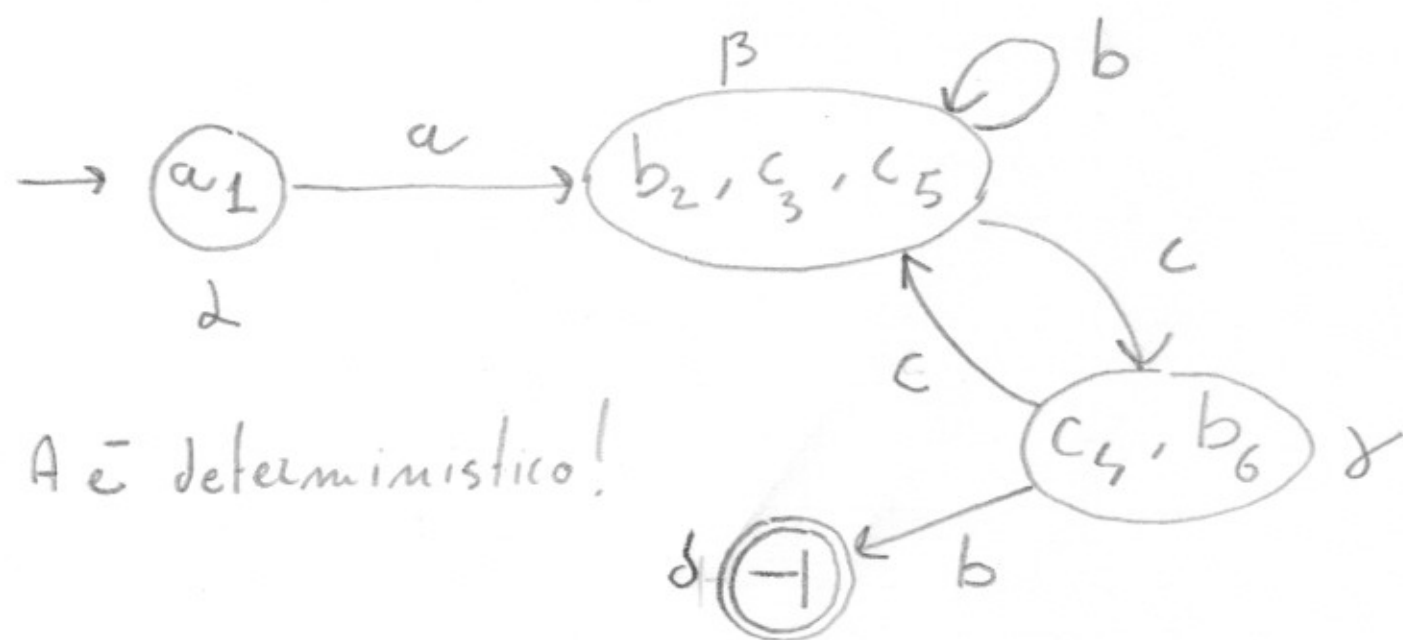


$$A = (\Sigma, Q, q_0, Q_F, \delta)$$

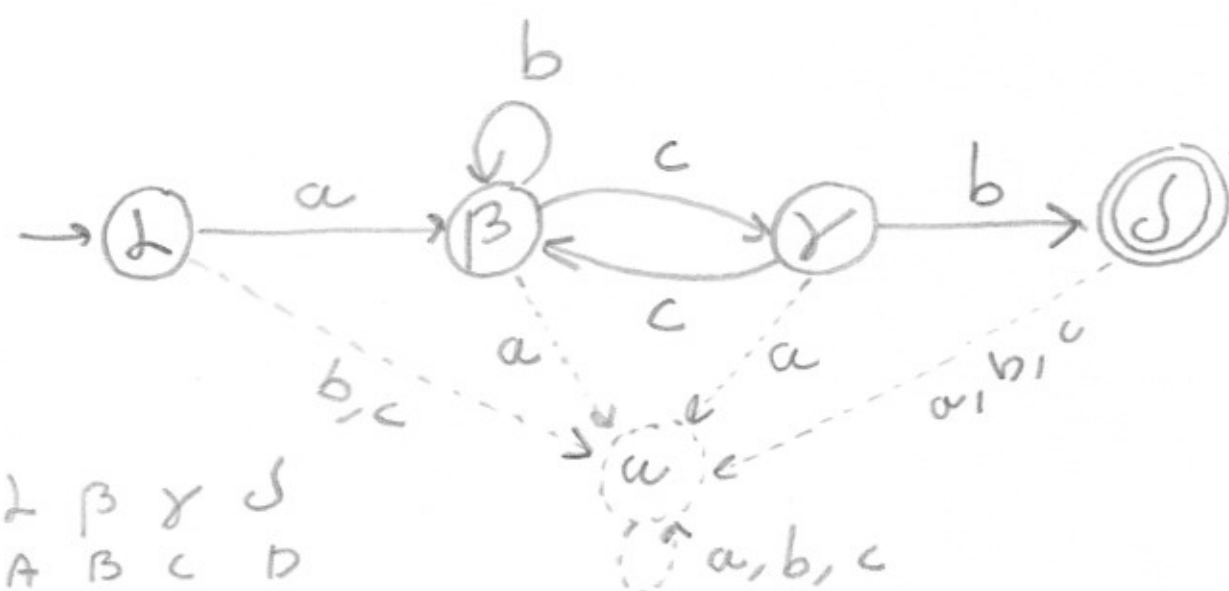
(2)

$$q_0 = \{a \mid a \text{ è carattere iniziale di } R-1\}$$

$$\text{vale a dire } q_0 = \{a_1\}$$



A è deterministico!



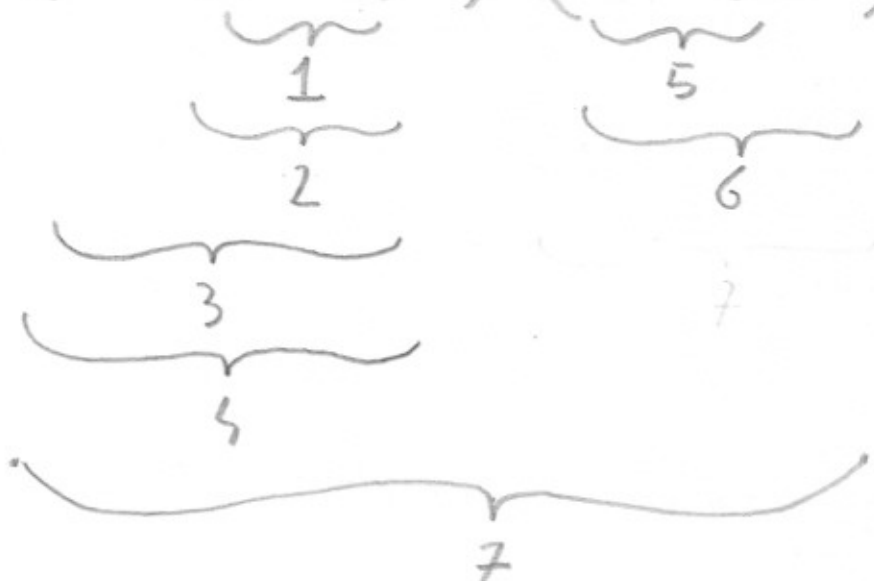
$$G \begin{cases} A \rightarrow a B \\ B \rightarrow b B \mid c C \\ C \rightarrow c B \mid b D \\ D \rightarrow \epsilon \end{cases}$$

G è unilineare
(a destra)

A è l'assioma

$$\Sigma = \{a, b\} \quad R = (a(bb)^*)^*(aa|\epsilon)^+ \quad \textcircled{L}$$

$$R = (a_1(b_2b_3)^*)^*(a_4a_5|\epsilon)^+$$



$$R = 7 \cdot +$$

$$7 = 4 \cdot 6$$

$$6 = 5 | \epsilon$$

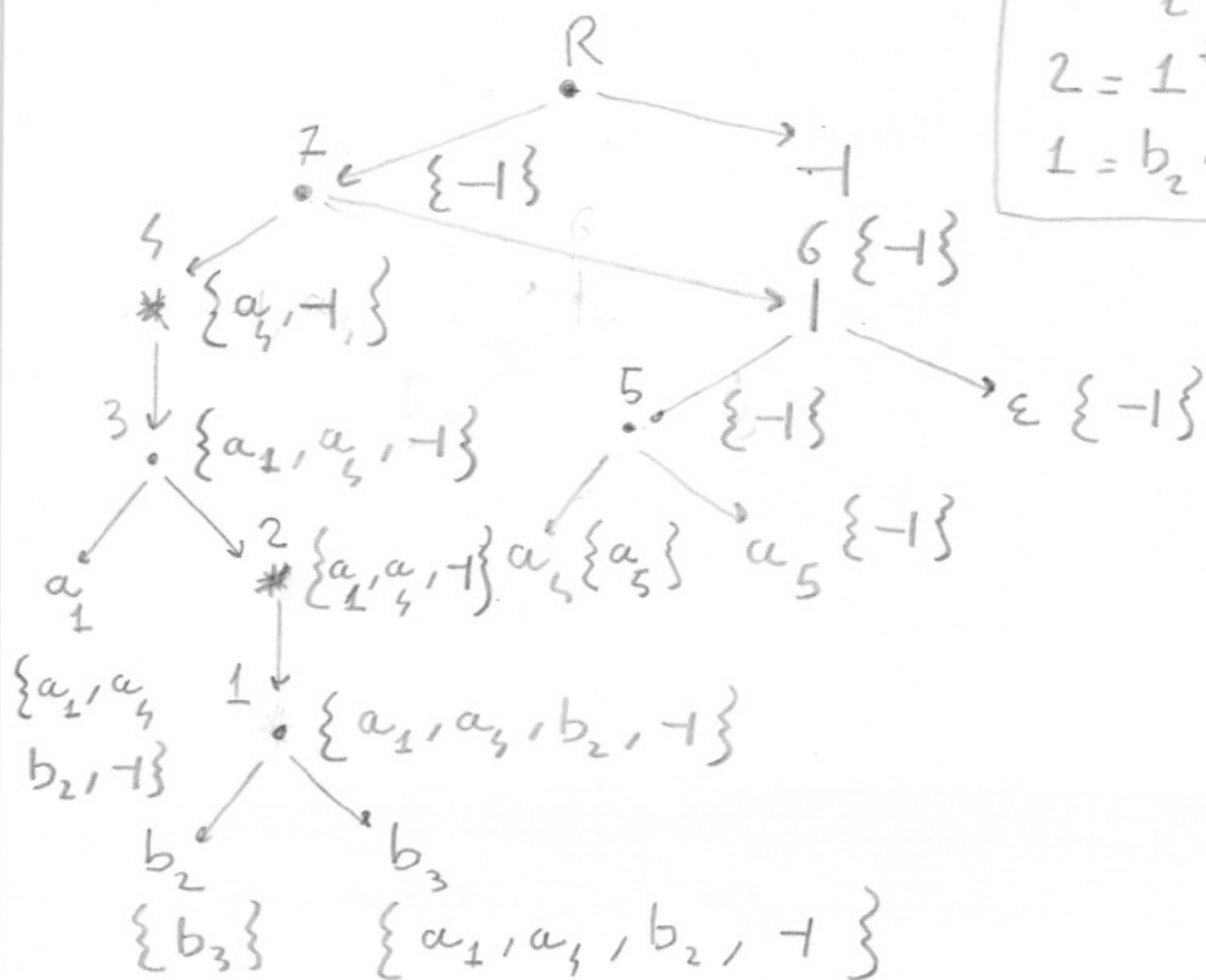
$$5 = a_4 \cdot a_5$$

$$4 = 3^*$$

$$3 = a_1 \cdot 2$$

$$2 = 1^*$$

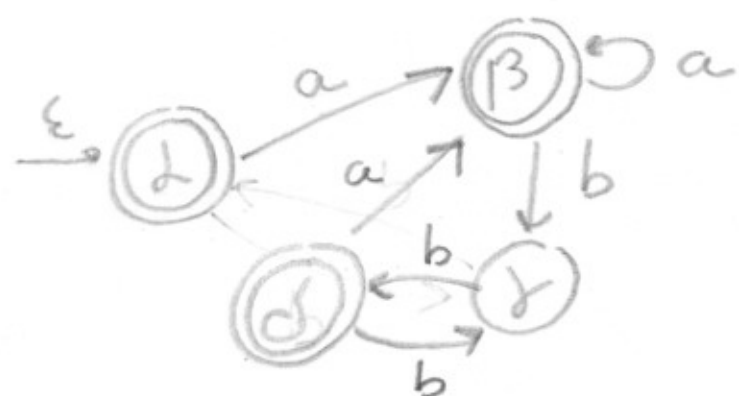
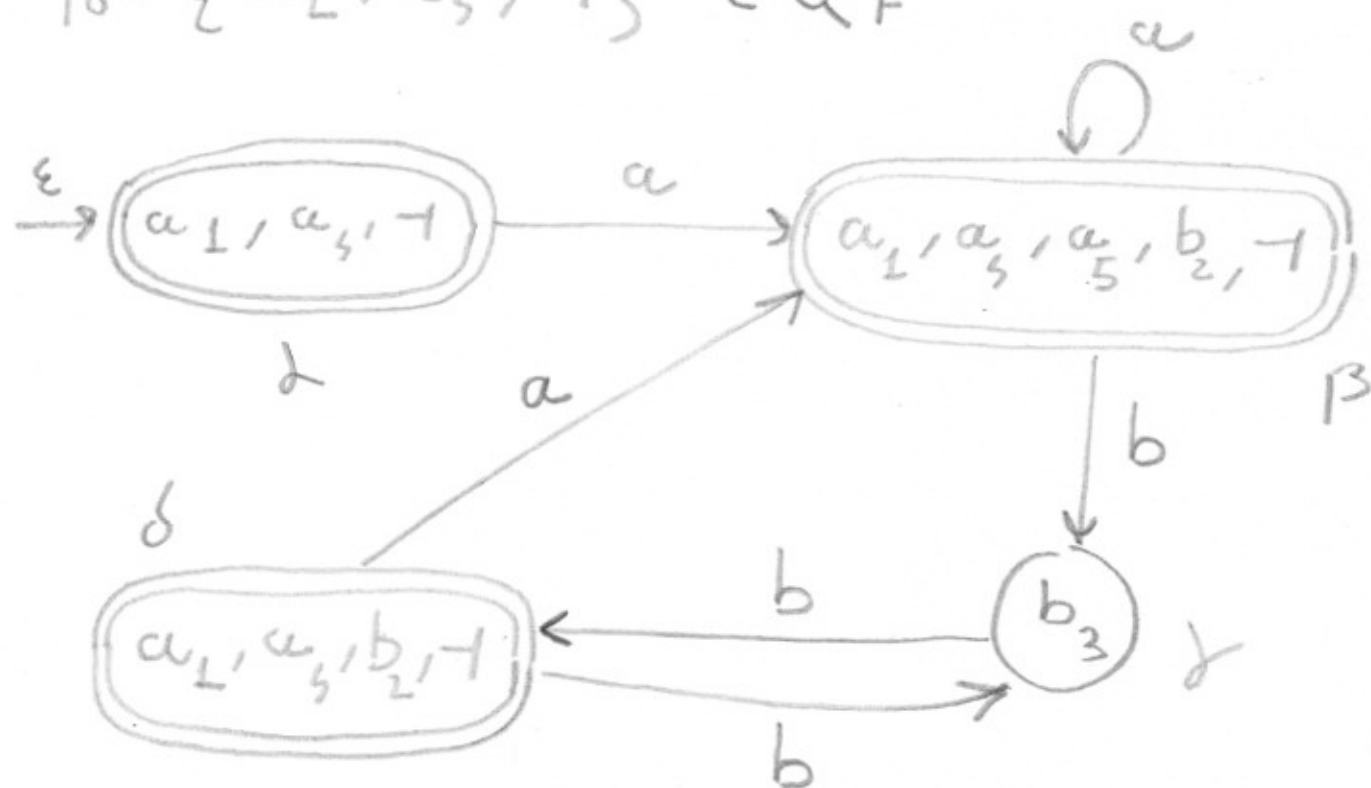
$$1 = b_2 \cdot b_3$$



$$A = (\bar{L}, Q, q_0, Q_F, \delta)$$

(2)

$$q_0 = \{a_1, a_3, \neg\} \in Q_F$$



$\alpha \quad \beta \quad \gamma \quad \delta$
 $A \quad B \quad C \quad D$

$$G \begin{cases} A \rightarrow \epsilon / aB \\ B \rightarrow \epsilon / aB / bC \\ C \rightarrow bD \\ D \rightarrow \epsilon / aB / bC \end{cases}$$

$G \neq \bar{e}$ ambigua

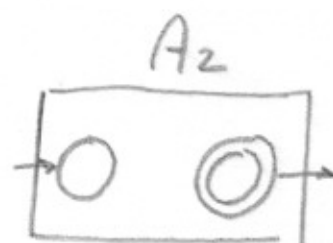
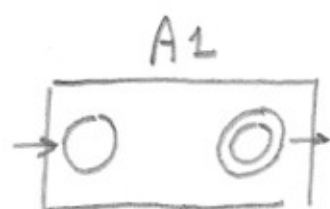
(invece $R \neq \bar{e}$ ambigua:

$$a_1 a_1 \in L(R)$$

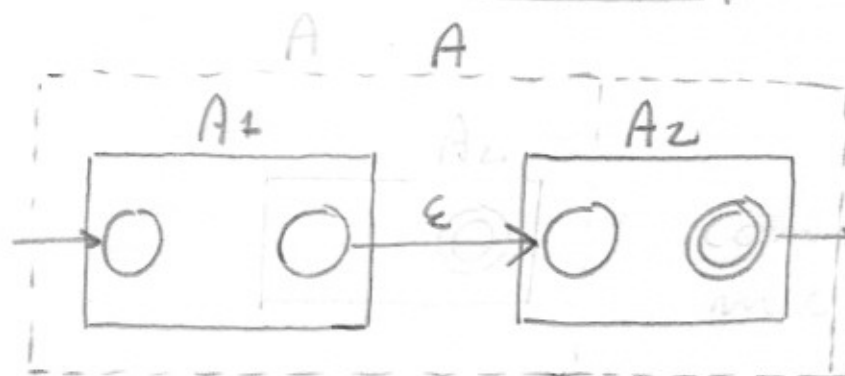
$$a_3 a_3 \in L(R)$$

Composizione modulare (Brezowski)

1



$$A_1 \cdot A_2 = A$$

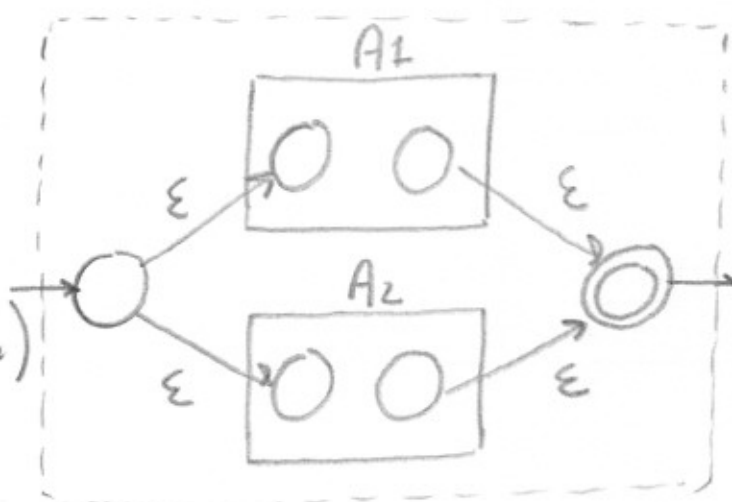


Concatenamento
(1 nuovo ϵ -arco)

$$L(A) = L(A_1) \cdot L(A_2)$$

$$A_1 \cup A_2$$

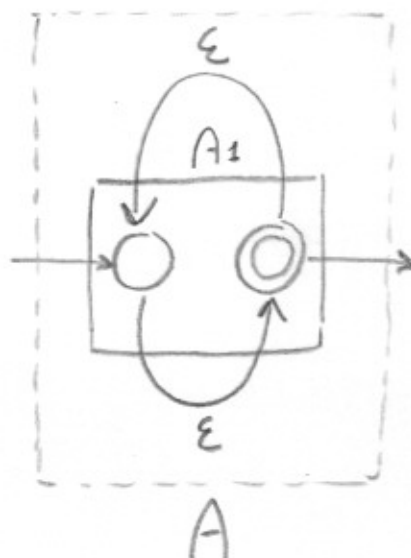
$$L(A) = L(A_1) \cup L(A_2)$$



Unione
(2 nuovi stati)
(3 nuovi ϵ -archi)

$$A_1^*$$

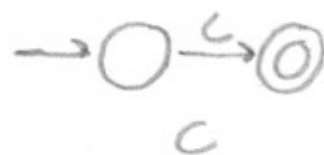
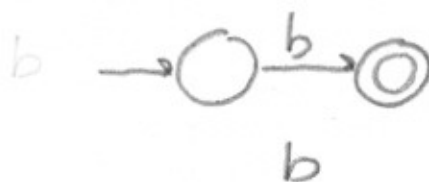
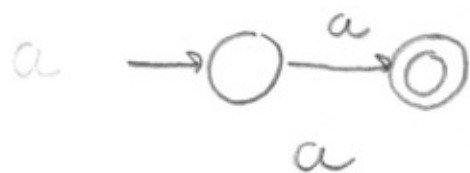
$$L(A) = L(A_1)^*$$



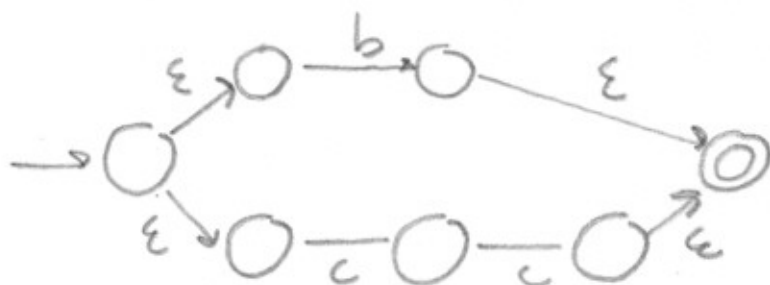
Iterazione
(2 nuovi ϵ -archi)

$$R = a(b|cc)^*cb$$

Calcolo automatico
7 deterministico (1)



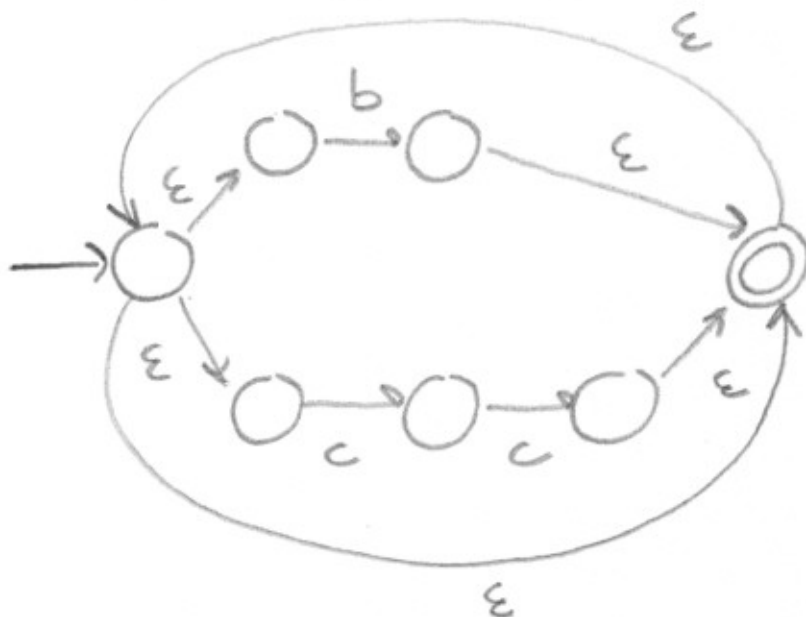
6



5

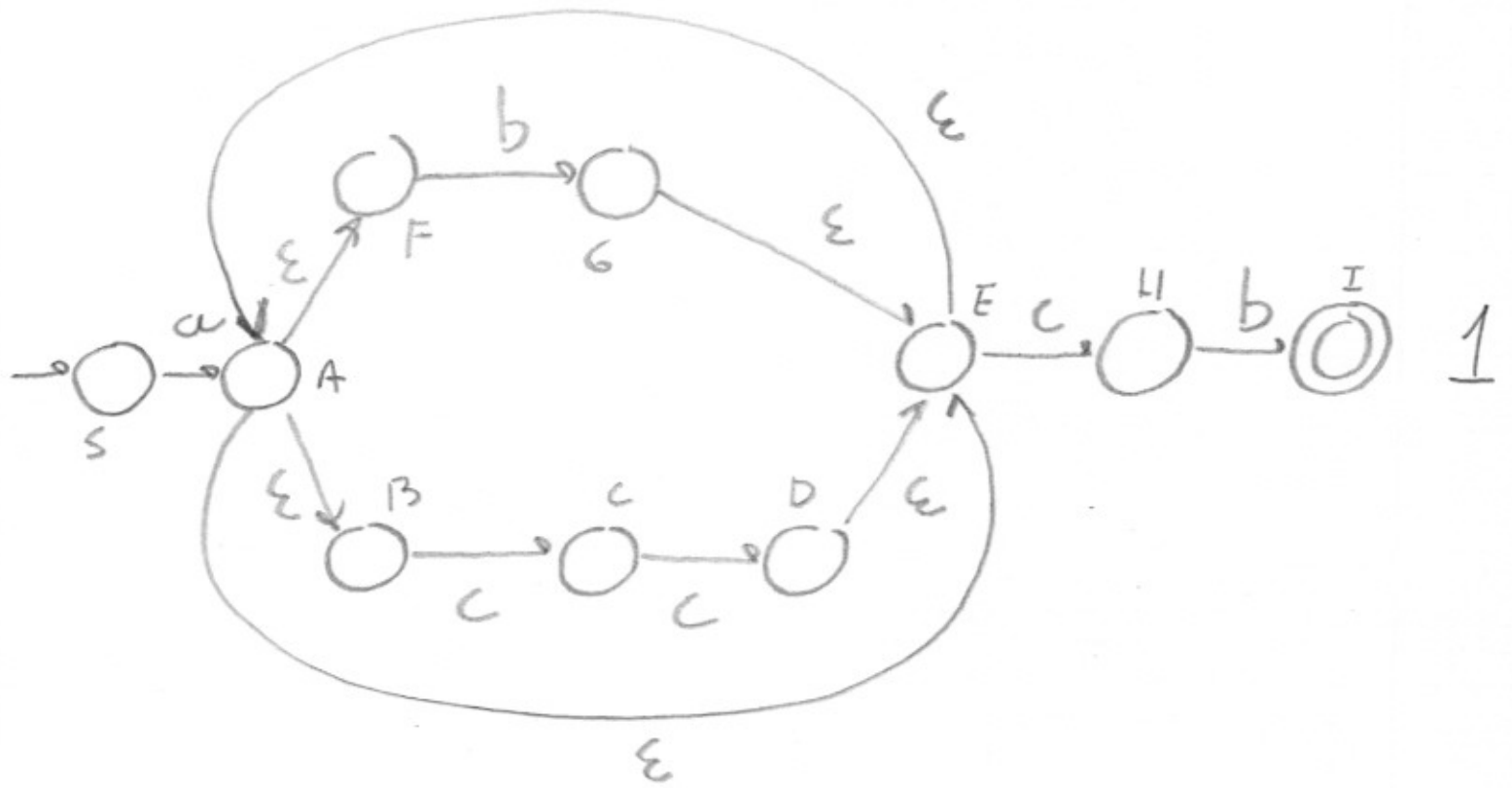
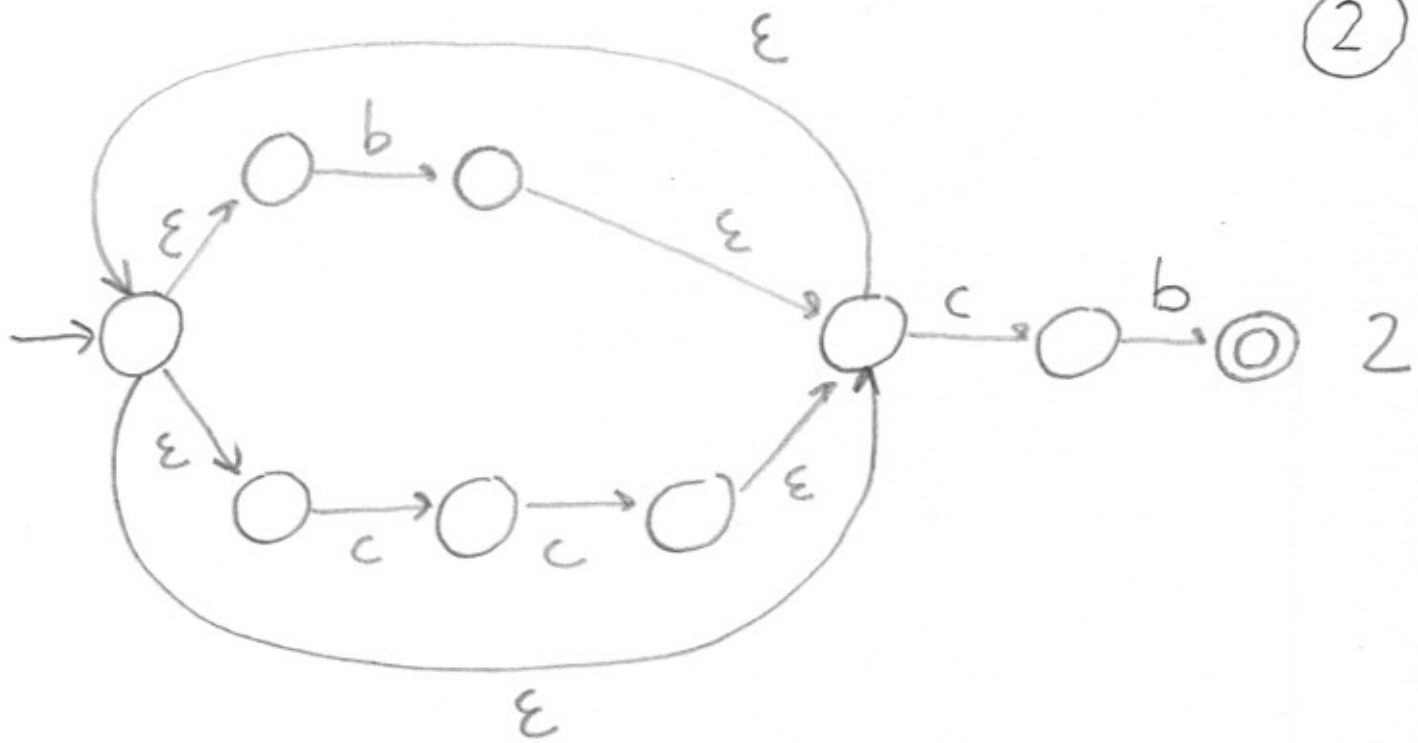


4



3

2



$S \rightarrow a A$
 $A \rightarrow B | F | E$
 $B \rightarrow c C$
 $C \rightarrow c D$

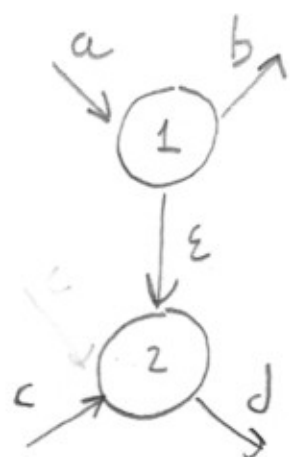
$D \rightarrow E$
 $F \rightarrow b G$
 $G \rightarrow E$
 $E \rightarrow A | c H$

$H \rightarrow b I$
 $I \rightarrow \epsilon$

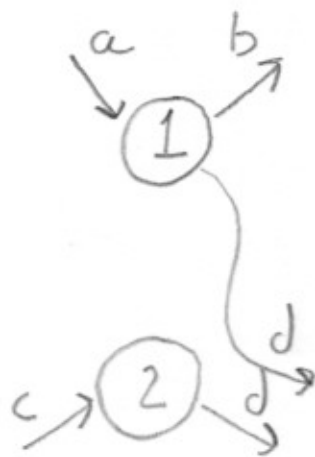
Eliminazione ϵ -transizioni

(1)

- Retrazione archi



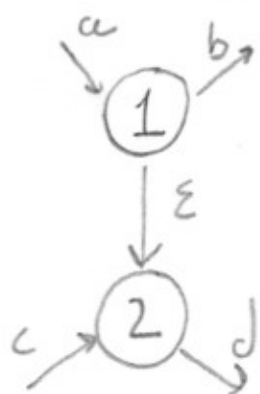
\Rightarrow



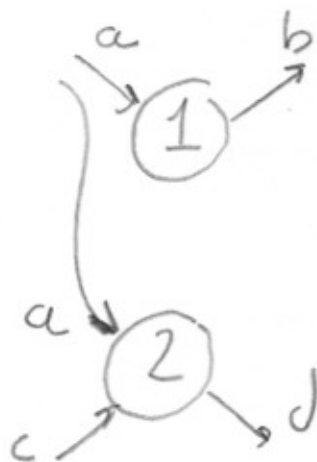
con $d \in \Sigma \cup \epsilon$

Inoltre, se
2 $\bar{\epsilon}$ finale,
1 diventa finale

- Propagazione archi



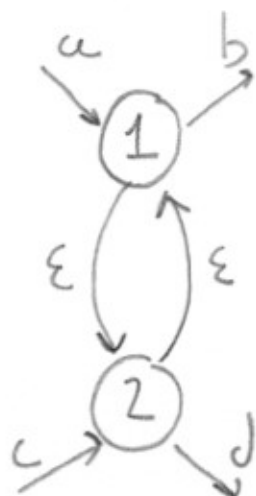
\Rightarrow



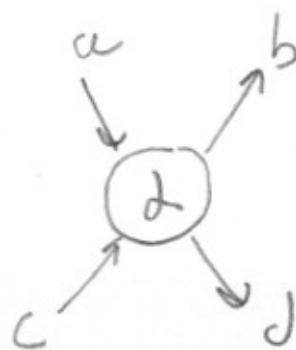
con $a \in \Sigma \cup \epsilon$

Inoltre, se
1 $\bar{\epsilon}$ iniziale,
2 diventa iniziale

- Collasso degli ϵ -cicli



\Rightarrow



Se 1 iniziale

\Downarrow

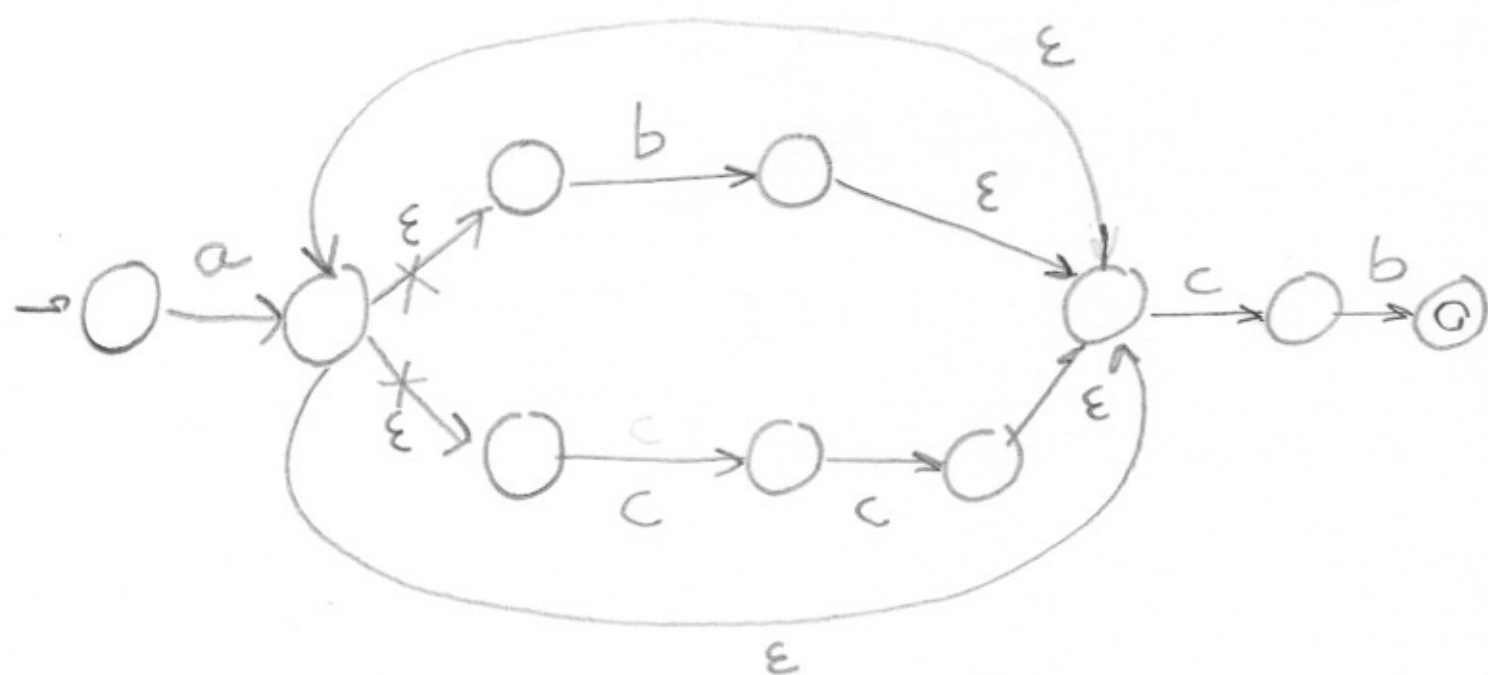
2 iniziale

Se 1 finale

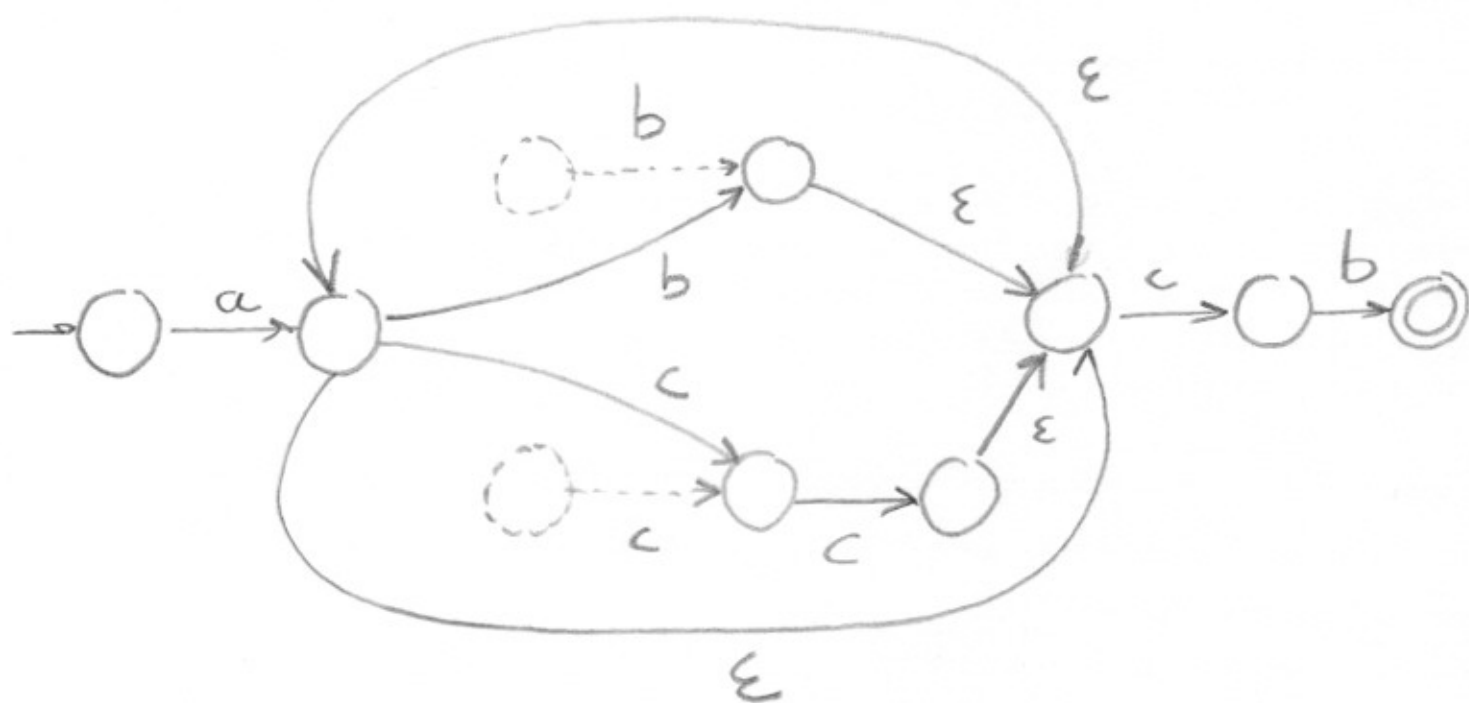
\Downarrow

2 finale

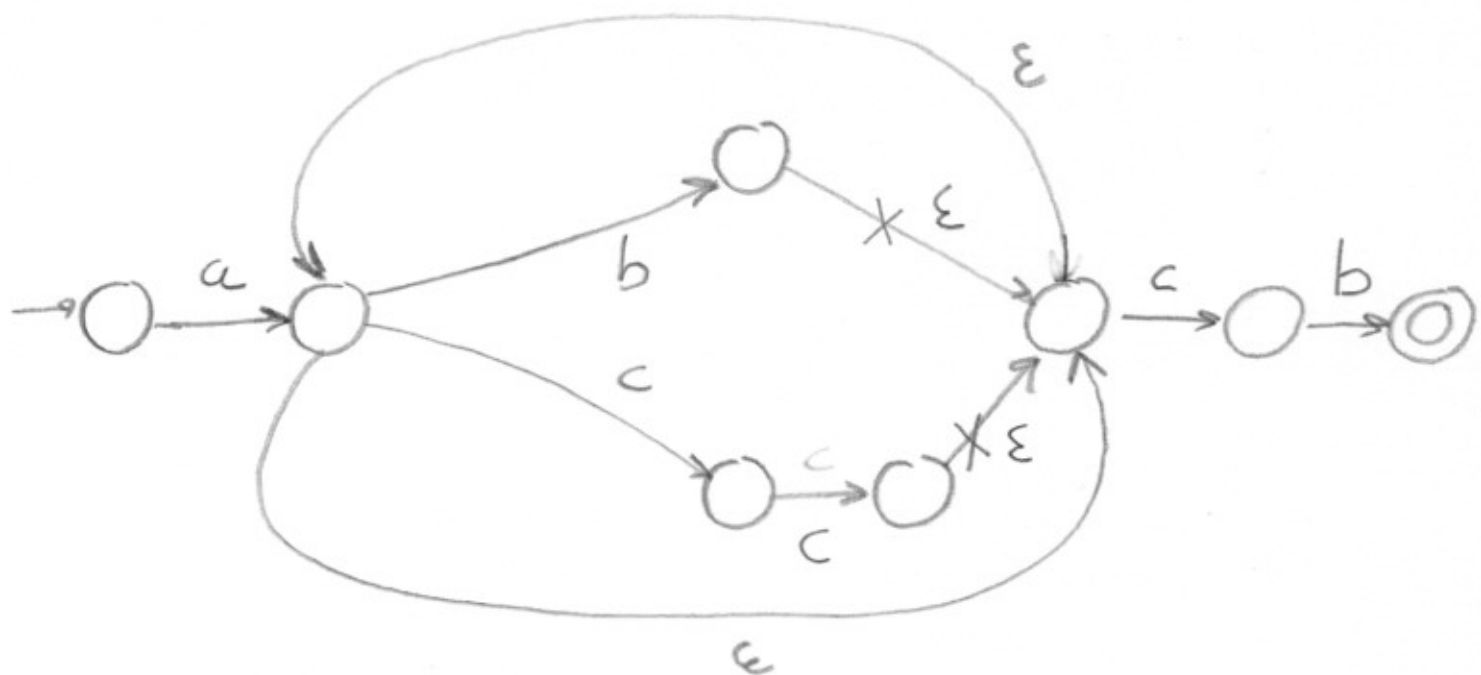
Determinizzazione Automata (1)



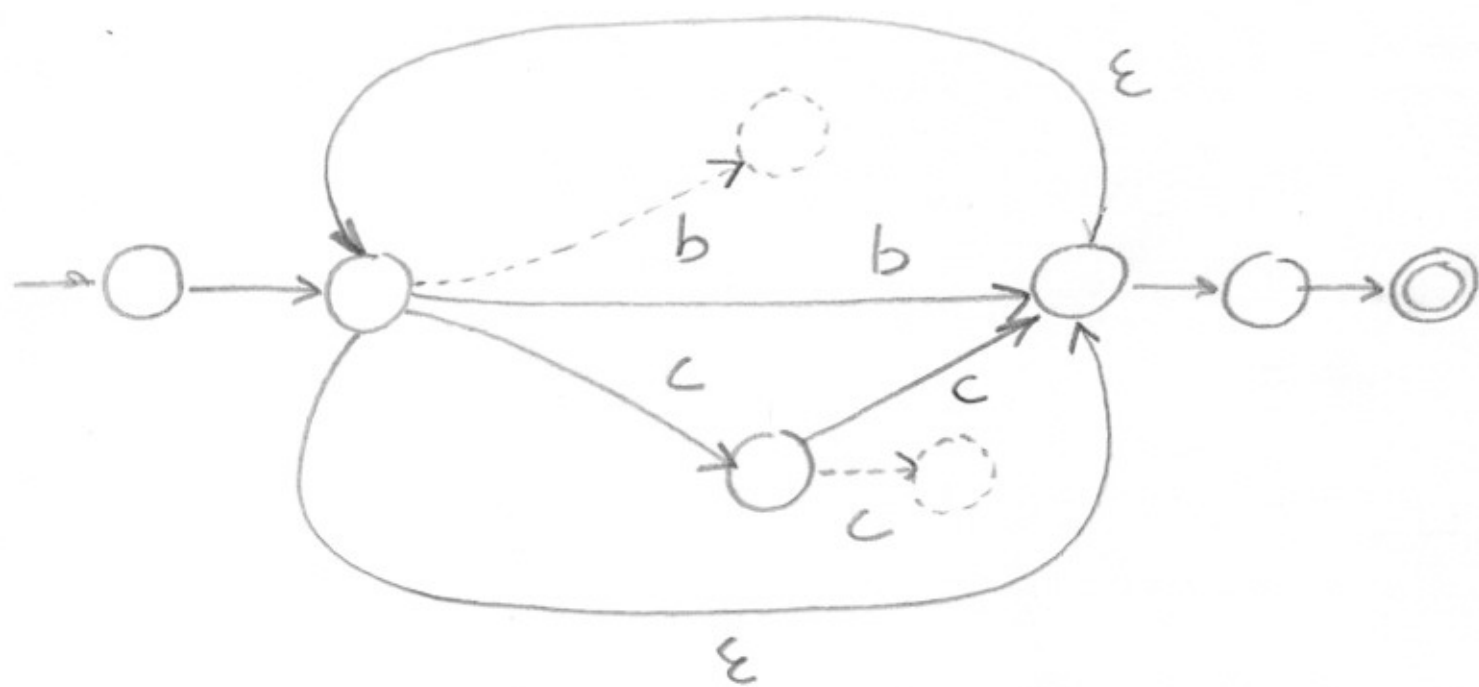
\Downarrow propagazione



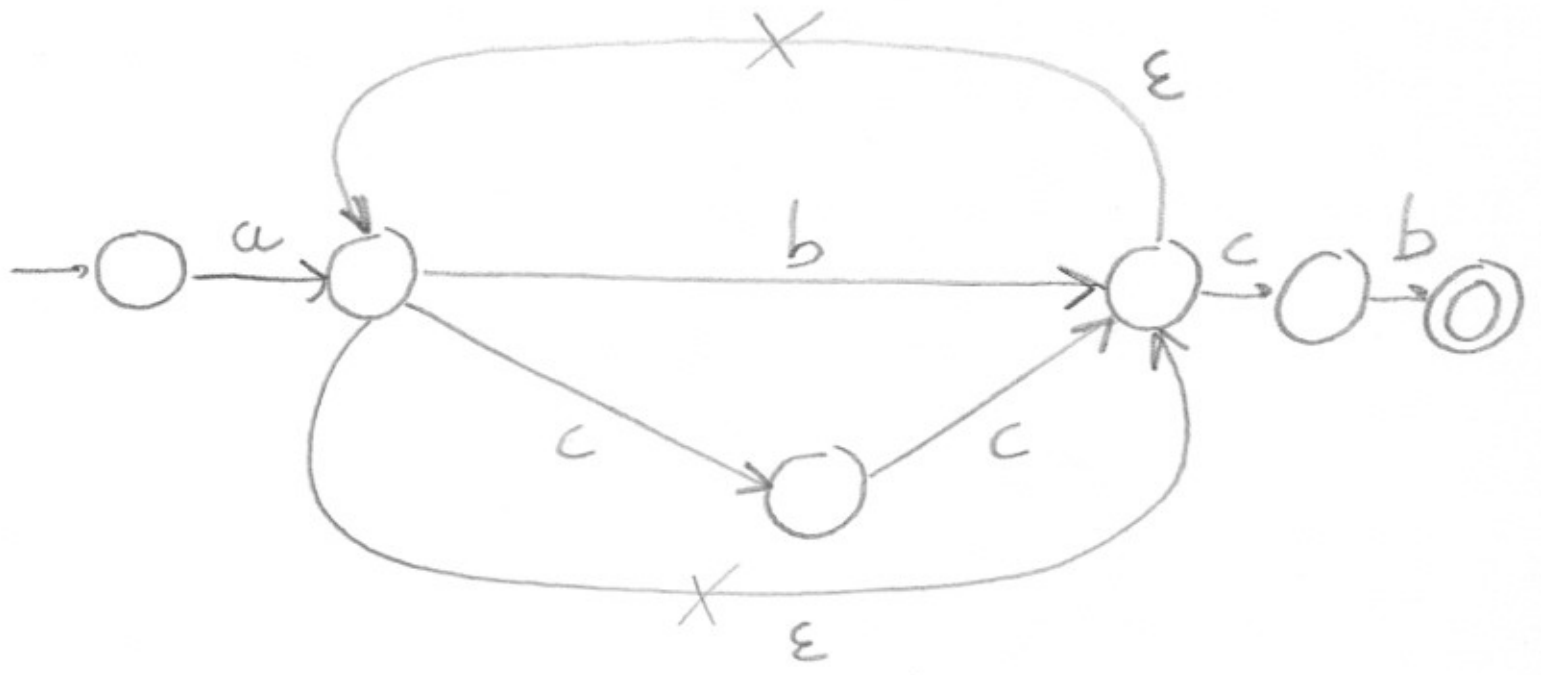
2



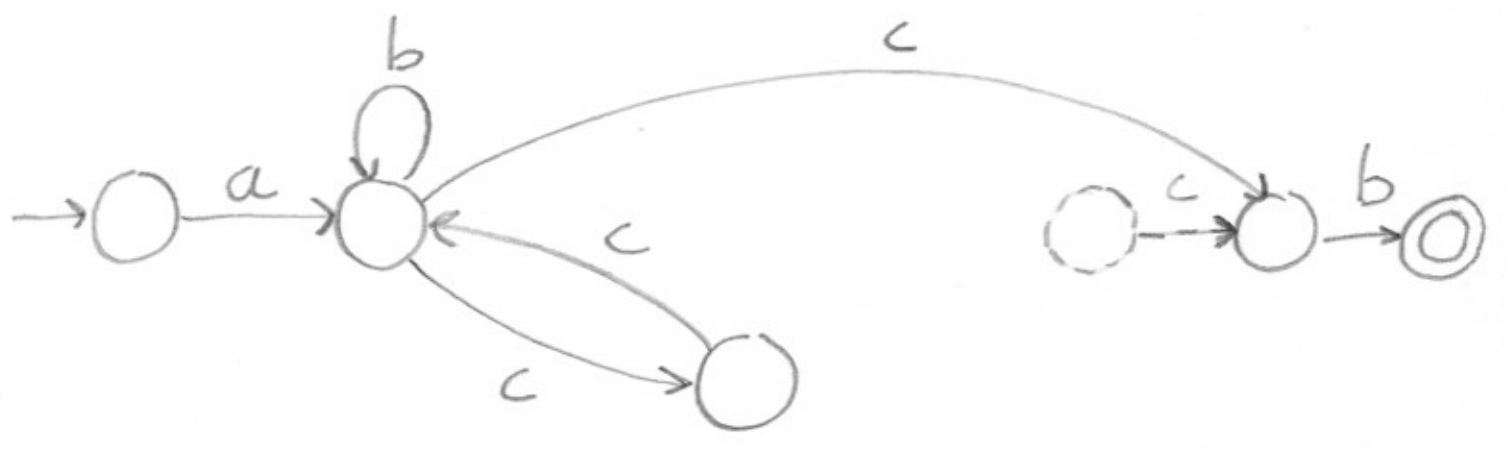
⇓ propagazione



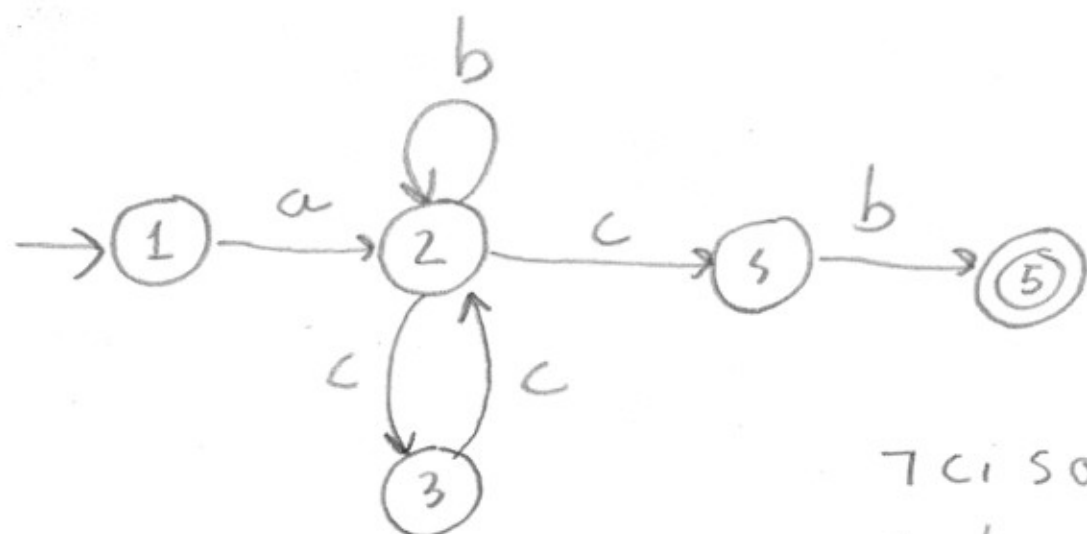
3



⇓
collasso
ε-ciclo

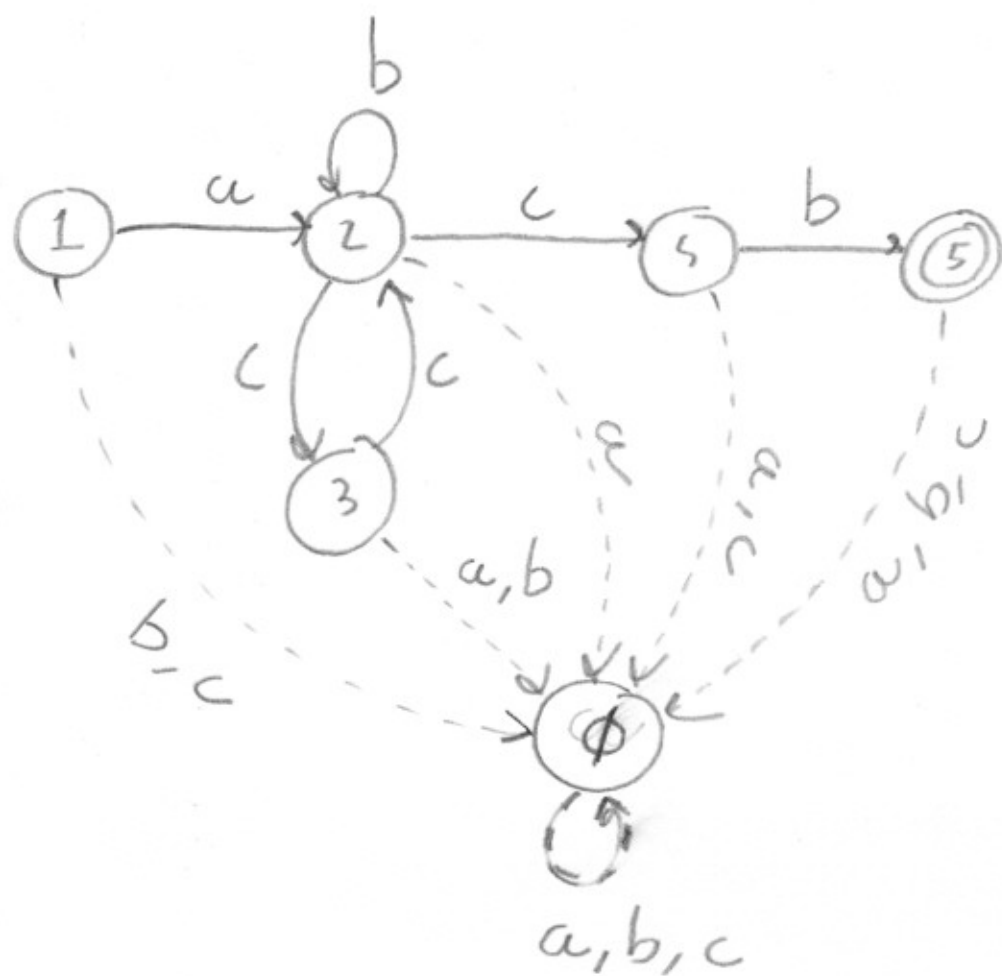


(4)

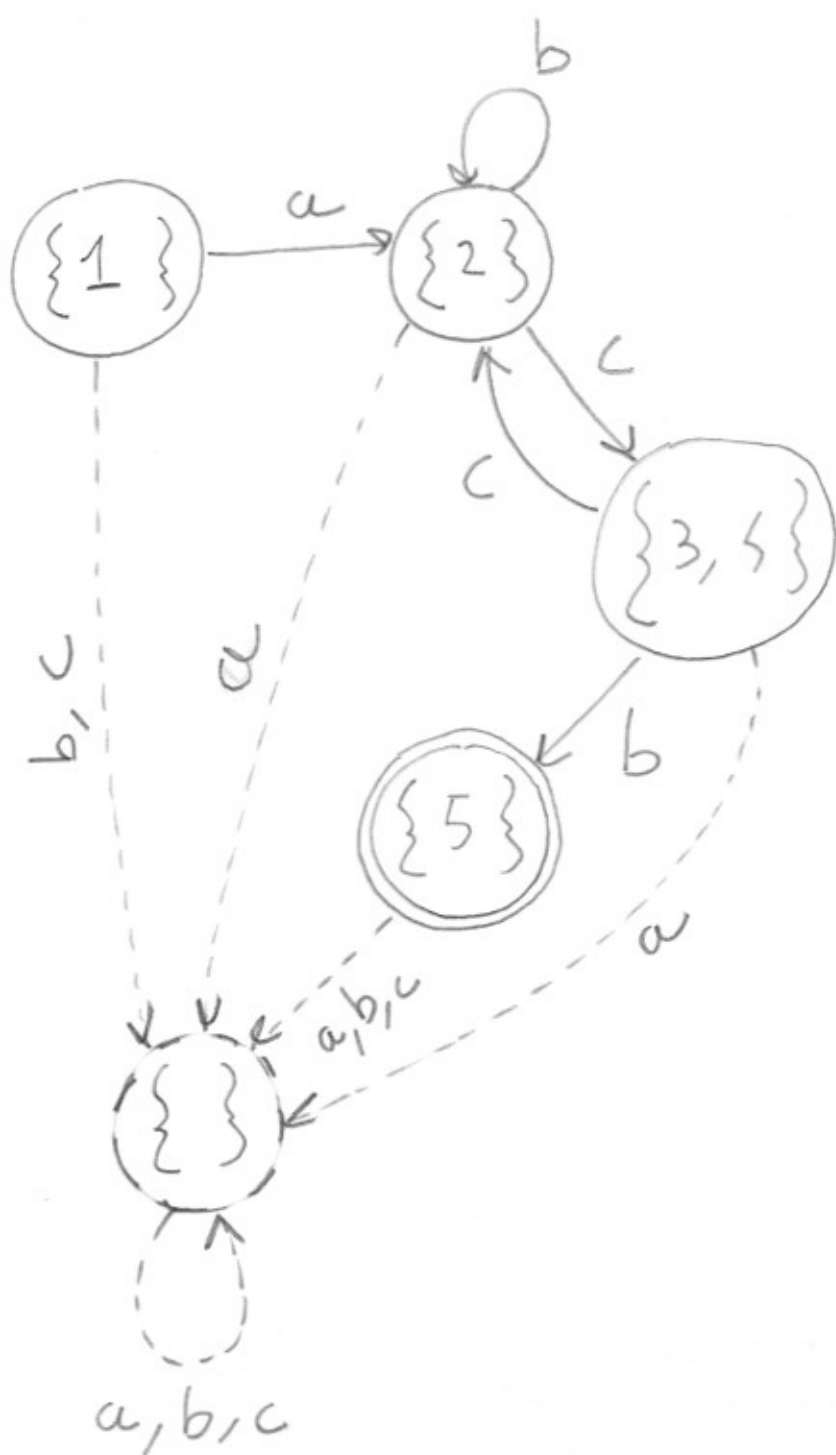
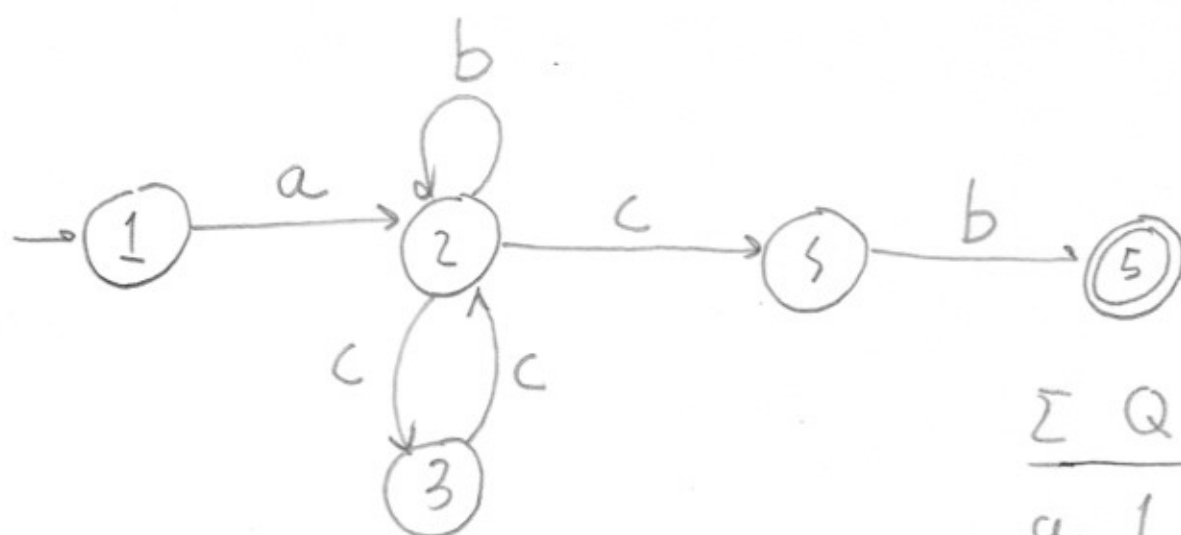


7 ci sono più
ε-transizioni

Se si preferisce . . .



①



Σ	Q	δ	$\delta(q)$
a	1		{2}
b	1		$\emptyset = \{\}$
c	1		\emptyset
a	2		\emptyset
b	2		{2}
c	2		{3,4}
a	3		\emptyset
b	3		\emptyset
c	3		{2}
a	4		\emptyset
b	4		{5}
c	4		\emptyset
a	5		\emptyset
b	5		\emptyset
c	5		\emptyset
a	\emptyset		\emptyset
b	\emptyset		\emptyset
c	\emptyset		\emptyset