ISCAS NIST

循序渐进，学习开发一个 RISC-V 上的操作系统

第 9 章 上下文切换和协作式多任务

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本章内容

- 多任务与上下文
- 协作式多任务的设计与实现

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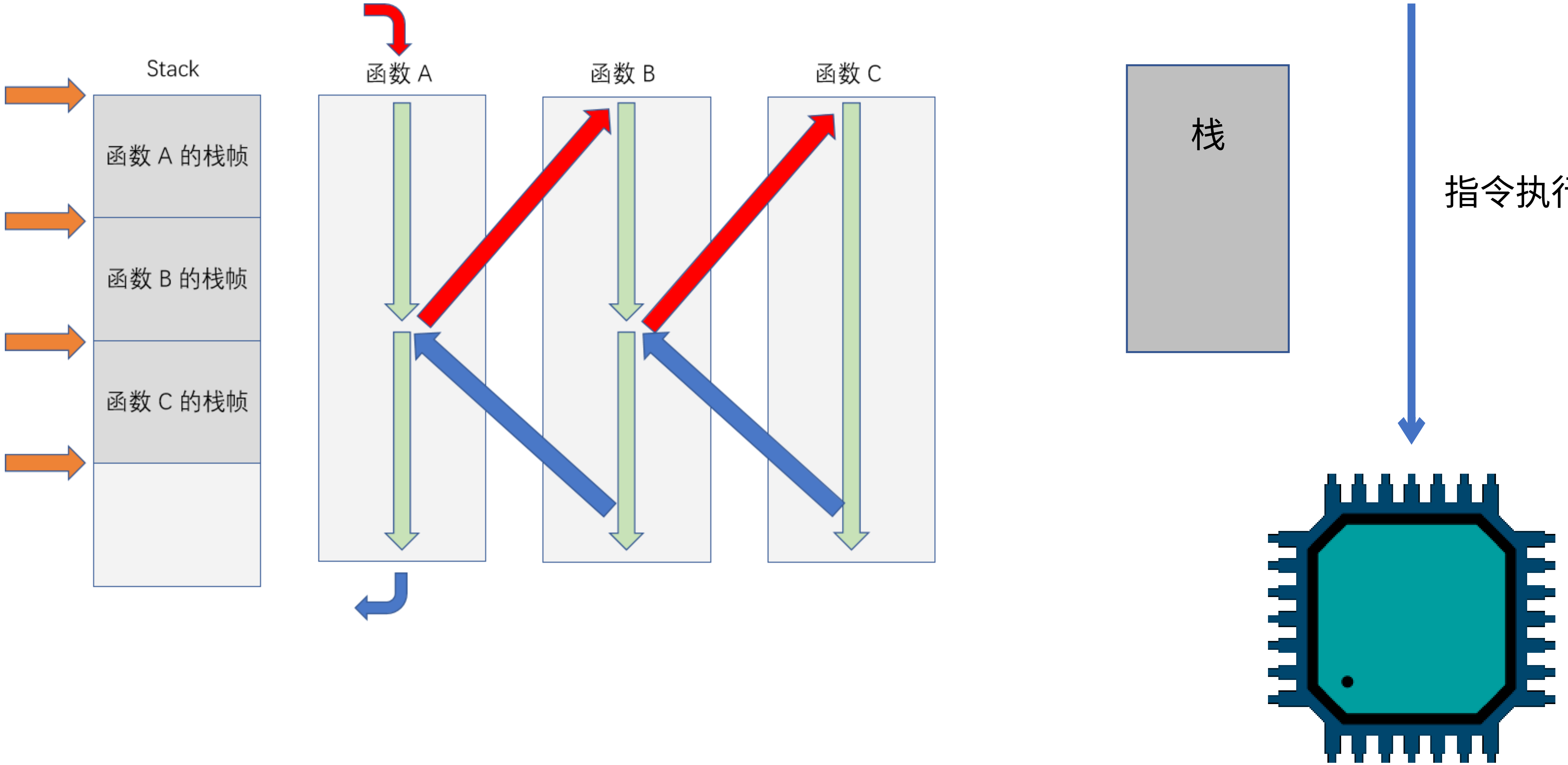
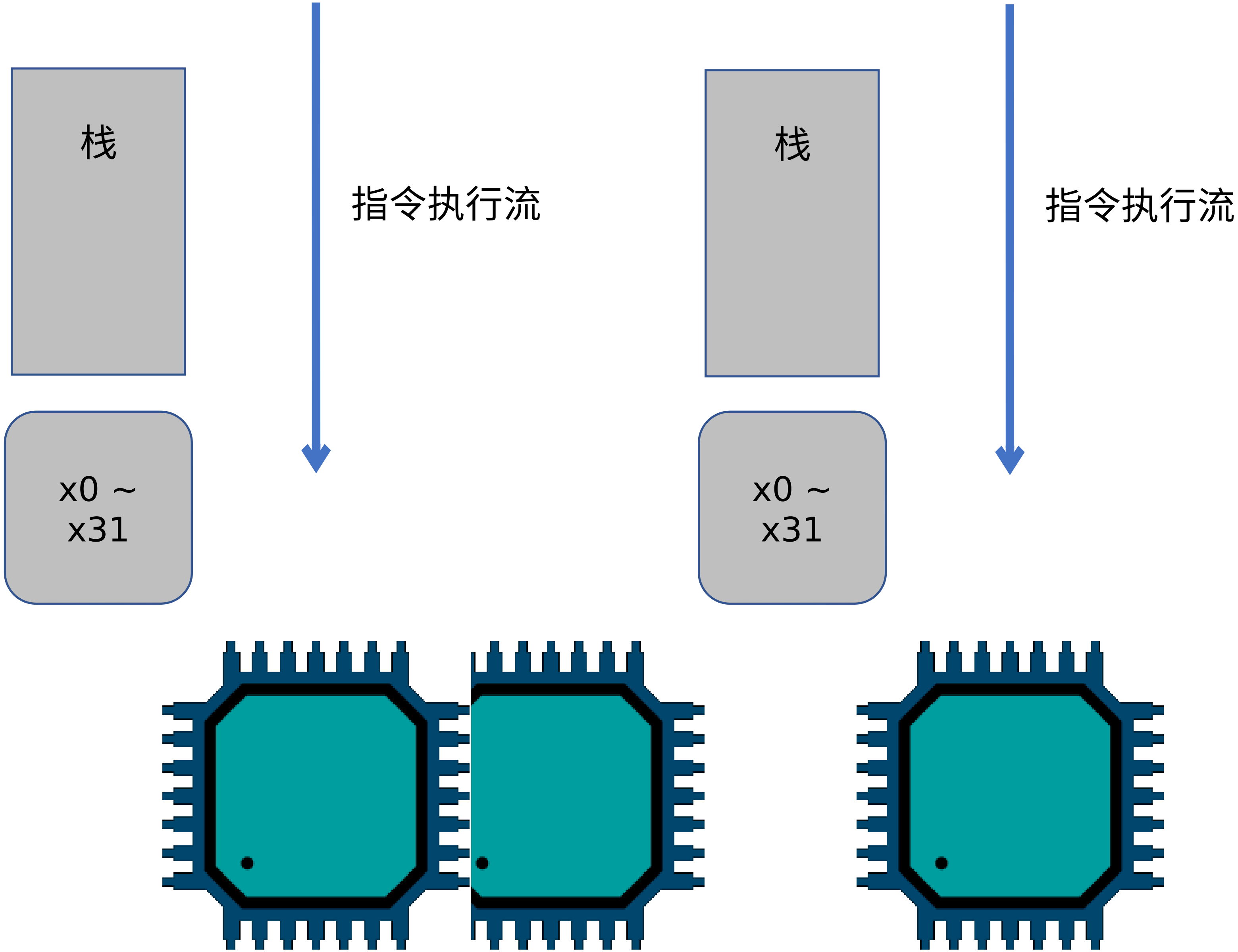
- **多任务与上下文**
  - 任务的概念
  - 多任务的概念
  - 任务上下文的概念

- **协作式多任务的设计与实现**

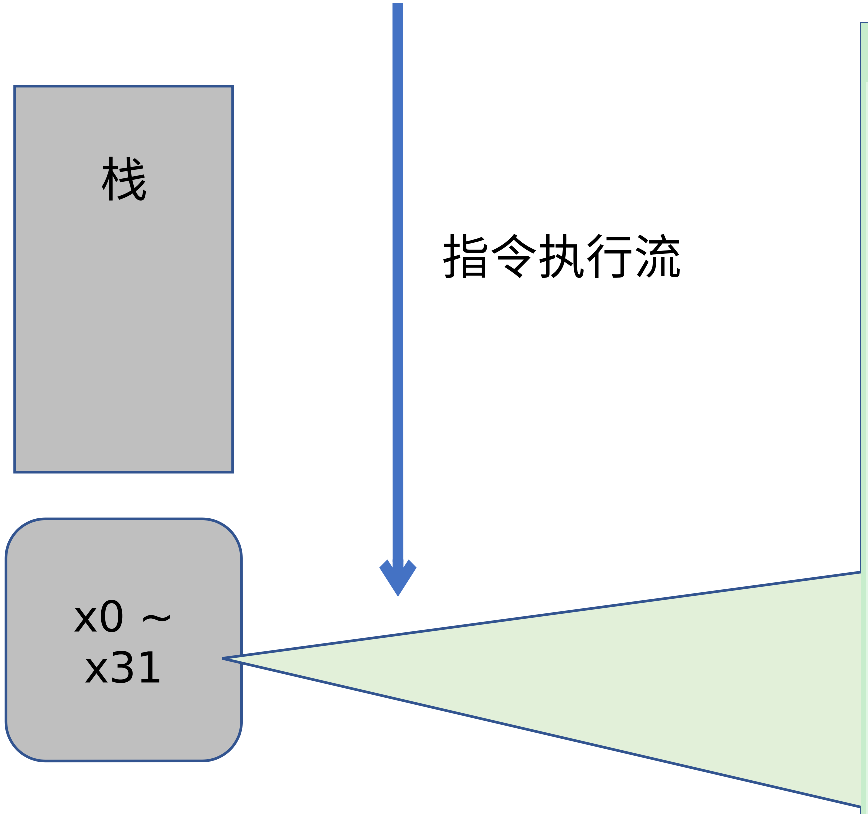
任务（task）



多任务 （Multitask）



任务上下文（Context）




```c
struct context {
        /* ignore x0 */
        reg_t ra;
        reg_t sp;
        reg_t gp;
        reg_t tp;
        .......
        reg_t s10;
        reg_t s11;
        reg_t t3;
        reg_t t4;
        reg_t t5;
        reg_t t6;
};
```

本章内容

- 多任务与上下文

- **协作式多任务的设计与实现**
  - 协作式多任务和抢占式多任务
  - 协作式多任务的设计思路
  - 协作式多任务的关键实现

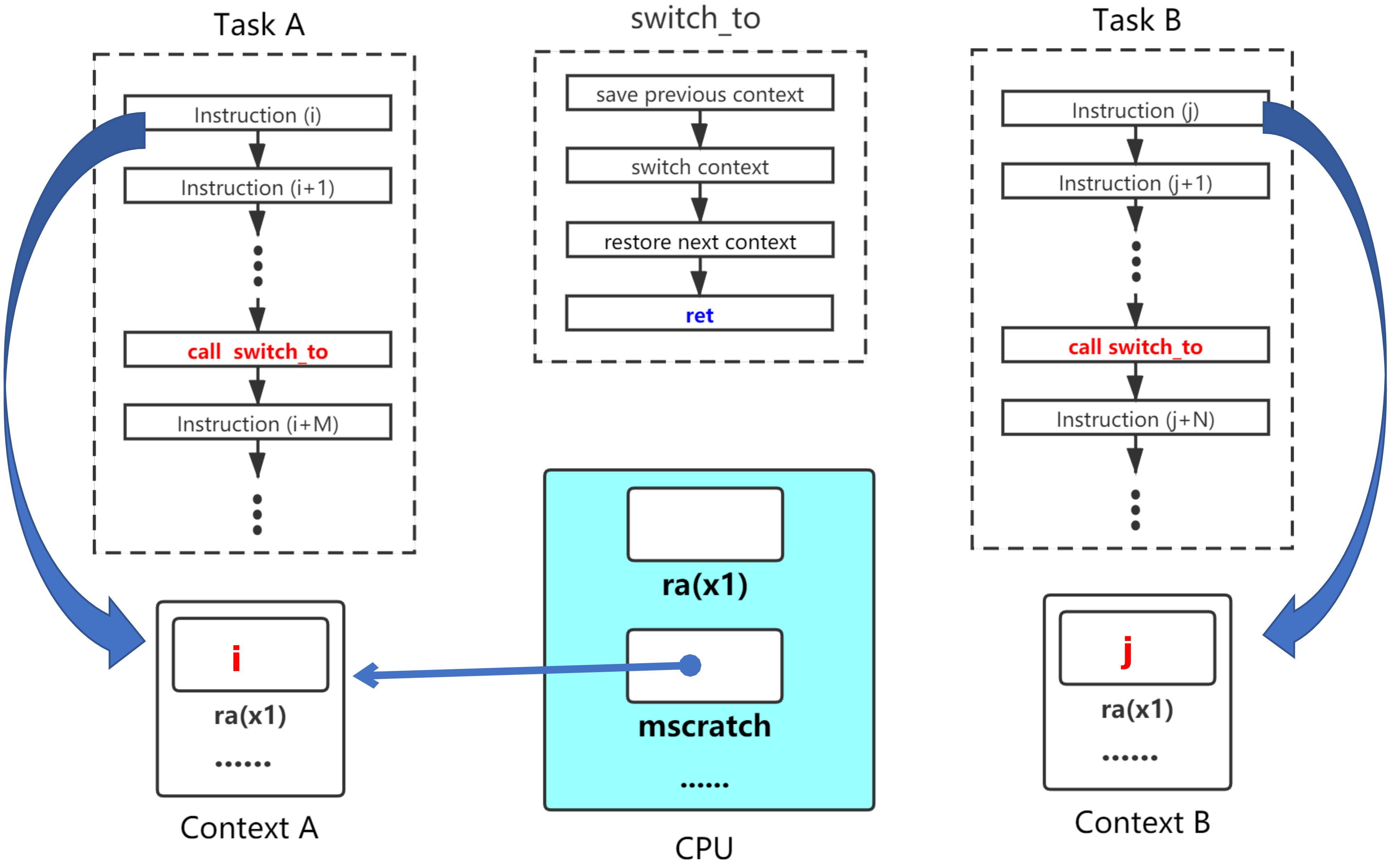
多任务系统的分类

- **协作式多任务（Cooperative Multitasking）**：协作式环境下，下一个进程被调度的前提是当前进程主动放弃时间片。
- **抢占式多任务（Preemptive Multitasking）**：抢占式环境下，操作系统完全决定进程调度方案，操作系统可以剥夺耗时长的进程的时间片，提供给其它进程。

协作式多任务

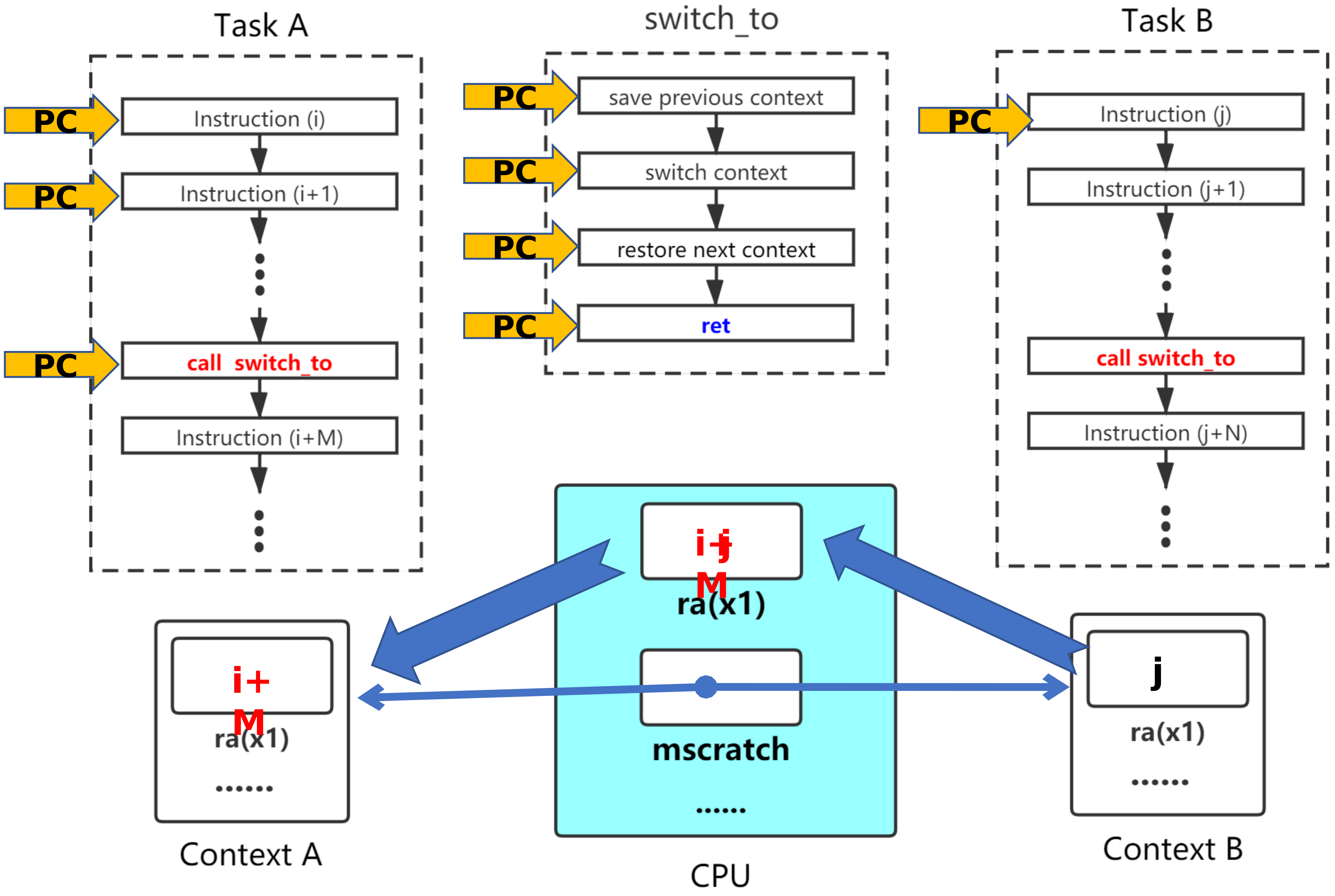
Task A
Instruction (i)
Instruction (i+1)
call switch_to
Instruction (i+M)
Task B
Instruction (j)
Instruction (j+1)
call switch_to
Instruction (j+N)

协作式多任务



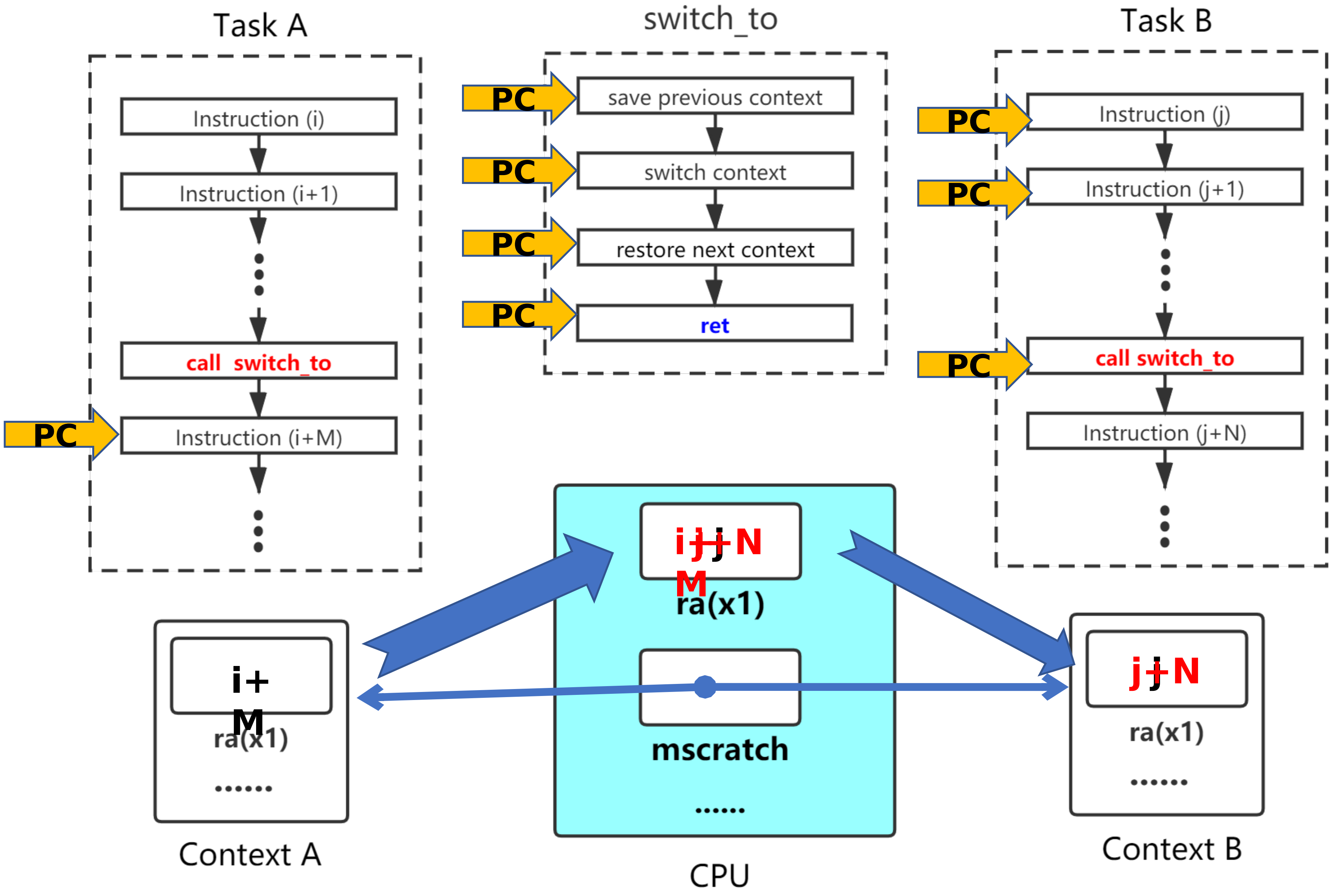
Task A
Instruction (i)
Instruction (i+1)
call switch_to
Instruction (i+M)
switch_to
save previous context
switch context
restore next context
ret
Task B
Instruction (j)
Instruction (j+1)
call switch_to
Instruction (j+N)
i
ra(x1)
......
Context A
ra(x1)
mscratch
......
CPU
j
ra(x1)
......
Context B

协作式多任务



Task A
PC
Instruction (i)
PC
Instruction (i+1)
PC
call switch_to
Instruction (i+M)
switch_to
PC
save previous context
PC
switch context
PC
restore next context
PC
ret
Task B
PC
Instruction (j)
Instruction (j+1)
call switch_to
Instruction (j+N)
i+M
ra(x1)
......
Context A
i+j
M
ra(x1)
mscratch
......
CPU
j
ra(x1)
......
Context B

协作式多任务



Task A
Instruction (i)
Instruction (i+1)
call switch_to
PC
Instruction (i+M)
switch_to
PC
save previous context
PC
switch context
PC
restore next context
PC
ret
Task B
PC
Instruction (j)
PC
Instruction (j+1)
PC
call switch_to
Instruction (j+N)
i+M
ra(x1)
......
Context A
i+M j+N
ra(x1)
mscratch
......
CPU
j+N
ra(x1)
......
Context B

关键函数（switch_to）

```
# void switch_to(struct context *next);
# a0: pointer to the context of the next task
.globl switch_to
.align 4
switch_to:
        # We use mscratch to hold a pointer to context of previous task and swap
        # with a user register (t6).
        # We use t6 as the 'base' for reg_save/reg_load, because it is the
        # very bottom register (x31) and would not be overwritten during loading.
        csrrw   t6, mscratch, t6        # swap t6 and mscratch
        beqz    t6, 1f                  # Notice: previous task may be NULL
        reg_save t6                     # save context of prev task

1:
        # switch mscratch to point to the context of the next task
        csrw    mscratch, a0

        # Restore all GP registers
        # Use t6 to point to the context of the new task
        mv      t6, a0
        reg_load t6

        # do actual context switching
        ret

.end
```

context 数据结构设计

```c
#define STACK_SIZE 1024
uint8_t task_stack[STACK_SIZE];
struct context ctx_task;
```

```c
w_mscratch(0);

ctx_task.sp = (reg_t) &task_stack[STACK_SIZE - 1];
ctx_task.ra = (reg_t) user_task0;
```

```c
struct context *next = &ctx_task;
switch_to(next);
```

```c
struct context {
	/* ignore x0 */
	reg_t ra;
	reg_t sp;
	reg_t gp;
	reg_t tp;
	.......
	reg_t s10;
	reg_t s11;
	reg_t t3;
	reg_t t4;
	reg_t t5;
	reg_t t6;
};
```

谢 谢

欢迎交流合作