

# SORTING TECHNIQUES

OMKAR BASNET

INCHARGE, CSIT/BCA

TEXAS INTERNATIONAL  
COLLEGE



# SORTING TECHNIQUES

( 8 HRS. )

Analysis of selection sort

Divide and conquer strategy

Quick and merge sort

Heap sort

# ANALYSIS OF SELECTION SORT



Selection sort is an algorithm that selects the smallest element from an unsorted list in each iteration and places that element at the beginning of the unsorted list.



This sorting algorithm is an in-place comparison-based algorithm in which the list is divided into two parts, the sorted part at the left end and the unsorted part at the right end.

20

12

10

15

2

Select first element as minimum

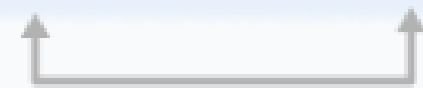
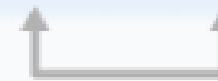
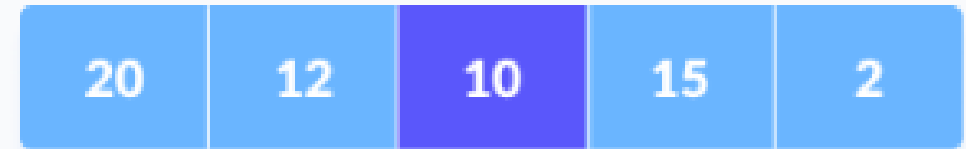
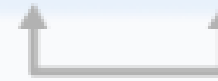
SET THE FIRST ELEMENT AS MINIMUM

# HOW SELECTION SORT WORKS?

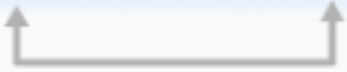
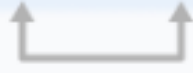
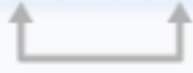
# HOW SELECTION SORT WORKS?

- Compare **minimum** with the second element. If the second element is smaller than **minimum**, assign the second element as **minimum**.

Compare **minimum** with the third element. Again, if the third element is smaller, then assign **minimum** to the third element otherwise do nothing. The process goes on until the last element.



Compare minimum with the remaining elements



Compare minimum with the remaining elements



Swap the first with minimum

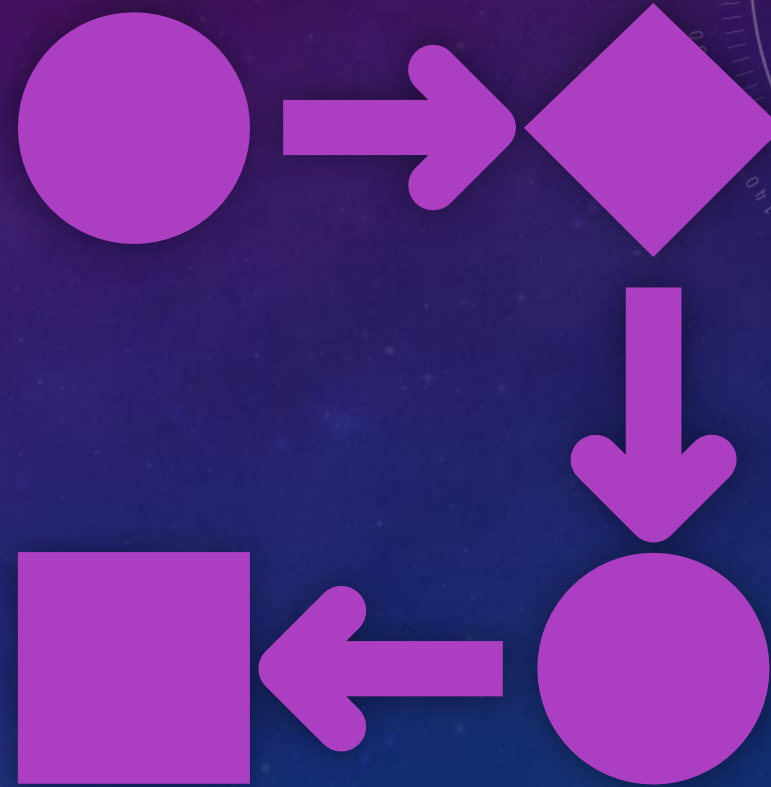
# HOW SELECTION SORT WORKS?

AFTER EACH ITERATION, MINIMUM IS PLACED IN THE FRONT OF THE UNSORTED LIST.



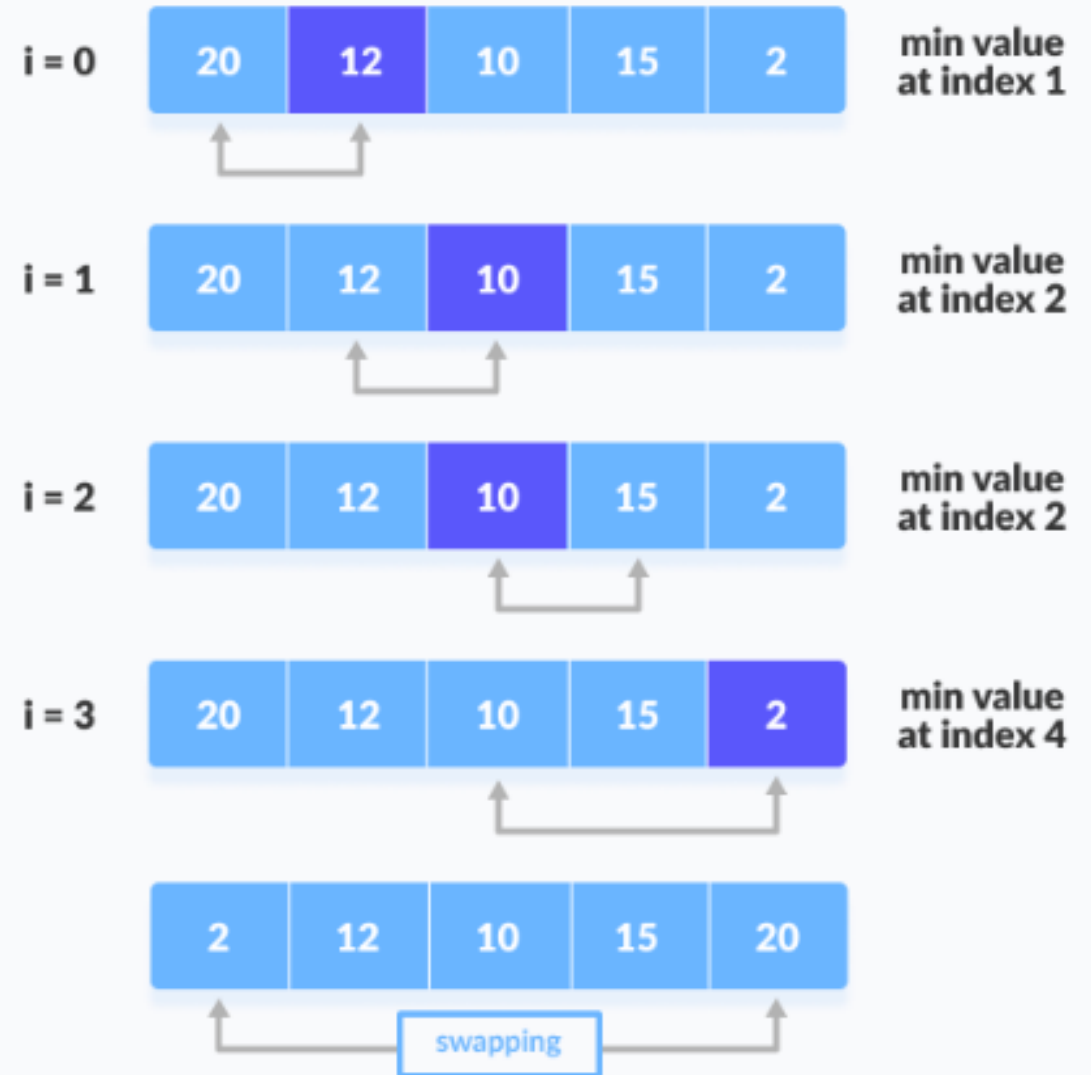
# HOW SELECTION SORT WORKS?

- For each iteration, indexing starts from the first unsorted element. Step 1 to 3 are repeated until all the elements are placed at their correct positions.



# HOW SELECTION SORT WORKS?

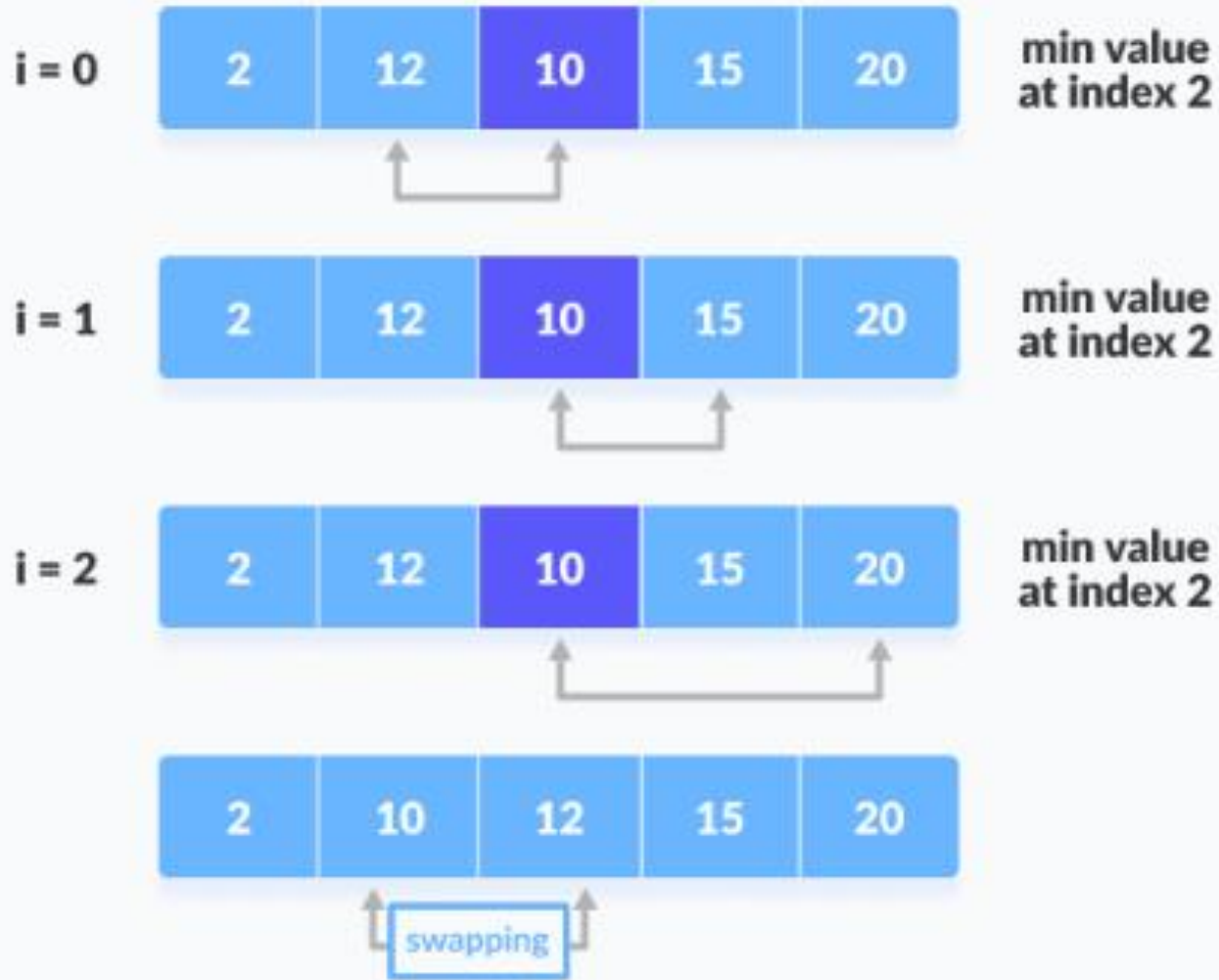
step = 0



The first iteration



step = 1



The second iteration

HOW  
SELECTION  
SORT  
WORKS?

step = 2



The third iteration

HOW  
SELECTION  
SORT  
WORKS?

step = 3

i = 0



min value  
at index 3



already in place

The fourth iteration

HOW  
SELECTION  
SORT  
WORKS?

Cycle	Number of Comparison
1st	$(n-1)$
2nd	$(n-2)$
3rd	$(n-3)$
...	...
last	1

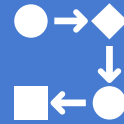
Number of comparisons:  $(n - 1) + (n - 2) + (n - 3) + \dots + 1 = n(n - 1) / 2$

nearly equals to  $n^2$ .

**Complexity** =  $O(n^2)$

# COMPLEXITY

# TIME COMPLEXITIES:



**Worst Case Complexity:**  $O(n^2)$

If we want to sort in ascending order and the array is in descending order then, the worst case occurs.



**Best Case Complexity:**  $O(n^2)$

It occurs when the array is already sorted



**Average Case Complexity:**  $O(n^2)$

It occurs when the elements of the array are in jumbled order (neither ascending nor descending).