

How Facebook's HHVM Uses C++ for Fun and Profit



Drew Paroski

What is HHVM?

- HHVM = HipHop Virtual Machine
- HHVM is a new open-source virtual machine designed to execute programs written in PHP or Hack*

^{*} Hack is a new language that evolved from PHP which adds a bunch of useful features

What is HHVM?

 Uses trace-based JIT compilation to deliver superior performance while maintaining the flexibility and productivity that PHP developers are used to

 Supports virtually all of the PHP 5.6 language including eval()

HHVM at Facebook

- A lot of core logic for a wide range of Facebook's features and products is written in PHP/Hack
- HHVM was developed to deliver better performance for Facebook's PHP/Hack code base
- HHVM evolved from a PHP->C++ transpiler called "HPHPc"

Performance

As of mid-2014, HHVM (and HPHPc before it)
has realized a ~10x increase in throughput and
over a 75% reduction in memory usage for
facebook.com compared with PHP 5.2

* Notes:

- This compares HHVM circa 2014 with a version of vanilla PHP from 2009
- These figures were calculated by taking multiple delta measurements over multiple years and stitching them together – direct comparison is not possible because FB's codebase is no longer compatible with vanilla PHP
- HHVM/HPHPc and FB's codebase have co-evolved over the past 4 years;
 until recently HHVM optimization efforts focused solely on FB's codebase

Facebook Data Center



Performance

- The performance of the stock php.net interpreter has notably improved in recent years
- Still, HHVM produces a boost of ~2x or more for a lot of PHP applications when compared with PHP 5.5 with opcode cache

Productivity

- HHVM supports the flexible "edit, save, run" development workflow that PHP developers are accustomed to
- When a PHP source file gets modified, HHVM detects that the file changed and recompiles PHP code as needed

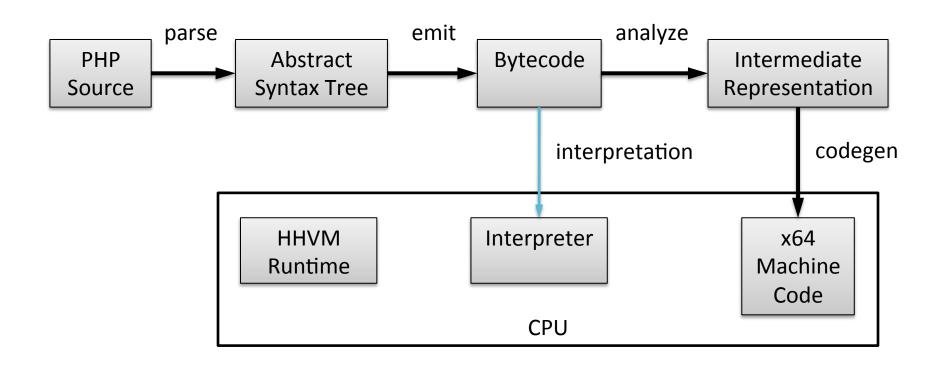
What can HHVM run?

- facebook.com
- Symfony
- MediaWiki
- WordPress
- Laravel
- Magento
- Drupal
- And dozens more of the most popular PHP frameworks and applications

How does HHVM work?

- PHP is parsed and then converted into untyped bytecode and metadata
- Basic blocks of bytecode (tracelets) are analyzed using live type information and converted to a typed IR
- Typed IR is then optimized and converted into x64 machine code

Basic Compilation Pipeline

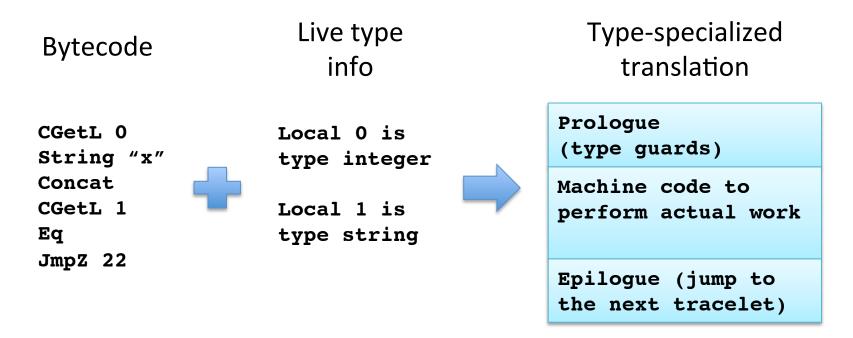


PHP -> HHBC compilation

```
0: EmptyM <L:0 EL:1>
 1: <?php
                                                          13: JmpNZ 23 (96)
 2: $a = array("hello" => "world");
 3: f($a, 42);
                                                          18: String "hit: "
                                                          23: String "\n"
 4: f($a, "hello");
                                                          28: CGetL2 1
 5: function f($a, $k) {
                                                          30: Concat
       if (!empty($a[$k])) {
 6:
                                                          31: Concat
                                                          32: Print
         print "hit: ".$k."\n";
                                                          33: PopC
 8:
        return;
                                                          34: Null
 9:
     } else {
                                                          35: RetC
         print "miss: ".$k."\n";
10:
                                                          36: String "miss: "
11:
         return;
                                                          41: String "\n"
12: }
                                                          46: CGetL2 1
                                                          48: Concat
13: }
                                                          49: Concat
                                                          50: Print
                                                          51: PopC
                                                          52: Null
                                                          53: RetC
```

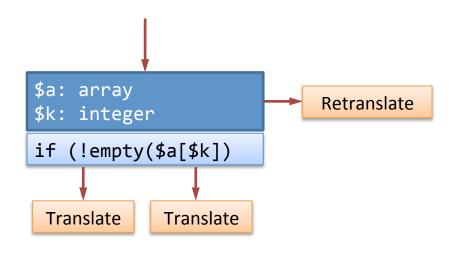
Tracelets

 HHVM's JIT compiler translates small amounts of bytecode at a time, using run time type information to drive predictions and generate type-specialized machine code:



```
function f($a, $k) {
   if (!empty($a[$k])) {
     echo "hit $k\n";
     return;
   } else {
     echo "miss $k\n";
     return;
   }
}
```

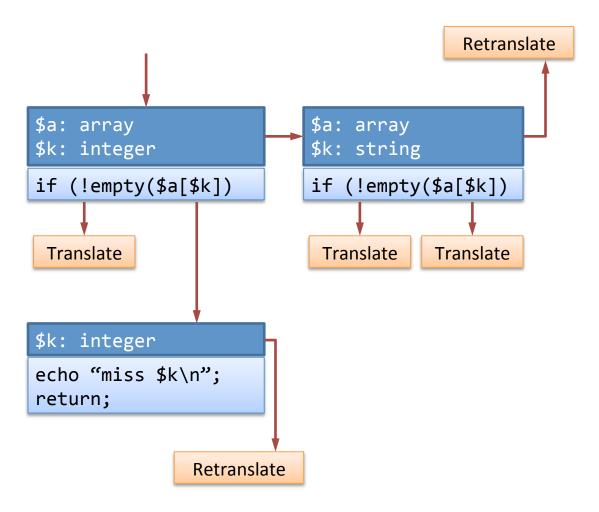
```
function f($a, $k) {
   if (!empty($a[$k])) {
     echo "hit $k\n";
     return;
   } else {
     echo "miss $k\n";
     return;
   }
}
```



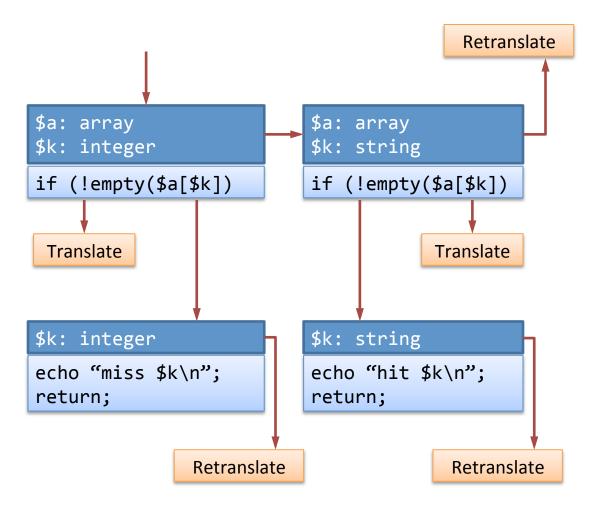
```
function f($a, $k) {
   if (!empty($a[$k])) {
     echo "hit $k\n";
     return;
   } else {
     echo "miss $k\n";
     return;
   }
}
```

```
$a: array
                             Retranslate
$k: integer
if (!empty($a[$k])
 Translate
$k: integer
echo "miss $k\n";
return;
                Retranslate
```

```
function f($a, $k) {
  if (!empty($a[$k])) {
    echo "hit $k\n";
    return;
  } else {
    echo "miss $k\n";
    return;
  }
}
```



```
function f($a, $k) {
  if (!empty($a[$k])) {
    echo "hit $k\n";
    return;
  } else {
    echo "miss $k\n";
    return;
  }
}
```



Another translation example

```
$n = 3 * $n + 1;
```



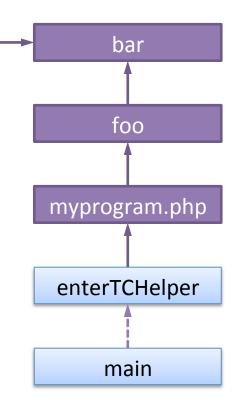
Int 3
CGetL 0
Mul
Int 1
Add
SetL 0
PopC
...



```
cmpl
       $0xa,-0x8(%rbp)
         retranslate
jne
       -0x10(%rbp),%rax
mov
       %rax,%rcx
mov
       %rcx
shl
       %rcx,%rax
add
       $0x1,%r13d
mov
       %rax,%r13
add
```

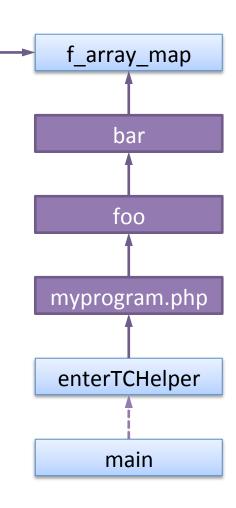
HHVM's Execution Model

- HHVM models the flow of execution using a stack of frames referred to as the *call stack*
- Each frame represents a function invocation
- VM frames correspond to PHP function invocations (shown in purple)
- C++ frames correspond to C++ function invocations (shown in blue)



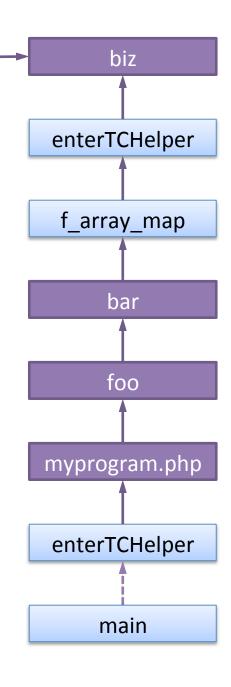
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- PHP code can call into C++ code



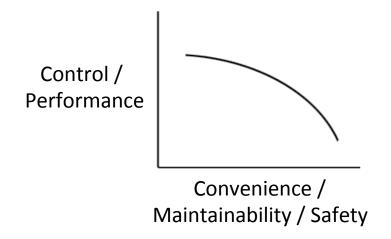
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- C++ frames correspond to C++ function invocations (shown in blue)
- PHP code can call into C++ code
- C++ code can call back into PHP code



C++ hits a sweet spot

- Superb control and performance on one end
- Convenience, maintainability, and safety on the other
- Gives the programmer fine-grained control to choose between these competing concerns as desired in different parts of the codebase



- Performance is equal or better than any other language (aside from assembly)
- Gives the programmer an incredible amount of freedom and control:
 - Manual memory management
 - -Unsafe casting, field size, and layout
 - Flexible, light-weight interop with assembly and machine code

- C++ offers many convenient features that can be used as little or as much as desired
 - Virtual methods
 - Multiple inheritance
 - reinterpret_cast vs. dynamic_cast
 - Plain old data vs. constructors/destructors
 - Raw pointers vs. references vs. smart pointers
 - Stack allocation vs. malloc vs. new
 - Templates and macros

Templates

- Templates are great for maintainability
- More hygienic than preprocessor macros
- For HHVM, templates were particularly useful for critical parts of the engine where we wanted the compiler to do as much inlining as possible to improve perf

Templates

Eliminating branches:

```
void foo(bool b, ...) {
  if (b) {
    bar();
  }
  ...
}

test %rdi, %rdi
jz L1
call 0x4005a0 <bar>
```

L1:

```
template <bool b>
void foo(..) {
  if (b) {
    bar();
call 0x4005a0 <bar>
```

Templates

Reducing indirection:

```
void foo(int(*fn)(), ...)
{
  int x = fn();
  ...
}
```

call *%rdi

```
template <class T>
void foo(...) {
  int x = T::staticMeth();
  ...
}

call 0x400780 <Foo::staticMeth>
```

C++11's Lambdas

- Lambdas are useful because the help keep related pieces of logic together in one place
- For HHVM, we typically use reference capture (i.e. "[&]") and we're careful about making sure lambdas do not outlive the captured variables on the stack

C++11's Lambdas

```
void CodeGenerator::cgCountArray(IRInstruction* inst) {
  ifThenElse(vmain(), vcold(), CC_Z,
    [&](Vout& v) {
      cgCallNative(v, inst);
    },
    [&](Vout& v) {
      v << loadl{baseReg[ArrayData::offsetofSize()],</pre>
                  dstReg};
  );
```

X Macros

- C++'s preprocessor is unhygienic, but it is extremely powerful and it can be very useful
- The X Macro technique can help with maintainability if used judiciously
- HHVM's bytecode definitions use the X Macro technique to make it easy to add, remove, or modify bytecode instructions

X Macros

```
#define OPCODES \
   O(PopC, NA, ONE(CV), NOV, NF) \
   O(PopV, NA, ONE(VV), NOV NF) \
   O(CGetL, ONE(LA), NOV, ONE(CV), NF) \
   O(Add, NA, TWO(CV,CV), ONE(CV), NF) \
   ...
enum class Op : uint8_t {
   #define O(name, ...) name,
    OPCODES
#undef O
};
```

X Macros

```
#define OPCODES \
 O(PopC, NA, ONE(CV), NOV, NF) \
 O(PopV, NA, ONE(VV), NOV NF) \
 O(CGetL, ONE(LA), NOV, ONE(CV),
                                      NF) \
 O(Add, NA, TWO(CV,CV), ONE(CV),
                                      NF) \
#define PUSH_NOV /* nop */
#define PUSH_ONE(t) PUSH_##t
#define PUSH_TWO(t1, t2) PUSH ##t2; PUSH ##t1
#define PUSH CV
#define PUSH_VV
#define O(name, imm, push, pop, flags) \
.. PUSH ##push ..
OPCODES
#undef O
#undef PUSH NOV
```

Unions and field size

- Unions are super useful when dealing with dynamically-typed values
- Unions can also be used to reduce the size of structs that have mutually exclusive fields
- Ability to control field size also comes in handy:
 - 8-bit, 16-bit, 32-bit, or 64-bit integers

TypedValue union

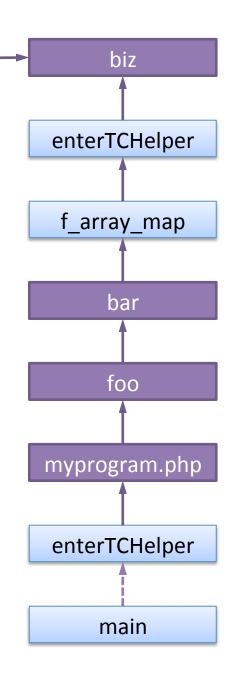
```
struct TypedValue {
                           enum DataType : int8 t {
                             KindOfClass
                                           = -13,
 union {
                             KindOfUninit
                                           = 0x00,
   int64 t
                 num;
   double
                             KindOfNull
                                           = 0x08,
                dbl;
   StringData*
                             KindOfBoolean
                                           = 0x09,
               pstr;
                             KindOfInt64
   ArrayData* parr;
                                           = 0x0a
   ObjectData* pobj;
                             KindOfDouble
                                           = 0x0b
   ResourceData* pres;
                             KindOfString
                                           = 0x14
                             KindOfArray
   Class*
                                           = 0x20,
                pcls;
                pref;
                             KindOfObject
                                           = 0x30,
   RefData*
                             KindOfResource = 0x40,
  } m data;
                                           = 0x50,
                             KindOfRef
 DataType m type;
 AuxUnion m aux;
};
                           };
```

Unsafe casts and bit-stealing

- C++ allows for unsafe casts between integers and different pointer types
- The implementation of malloc used by HHVM always returns chunks of memory aligned to 8-byte boundaries
- Unsafe casts and bit masking can be used to steal the low bits of pointers

Activation Records

- HHVM's ActRec struct is used as the header for each VM frame
- ActRecs store essential information about the PHP function invocation, such as:
 - The name and other metadata pertaining to the current function
 - Where to jump to when the function returns
 - Necessary bookkeeping to support getting a PHP backtrace (i.e. debug_backtrace())



```
struct ActRec {
  ActRec* m savedFp;
  uint64 t m savedRip;
  Func* m func;
  uint32 t m soff;
  uint32 t m numArgsAndFlags;
  union {
    ObjectData* m this;
    Class* m cls;
  };
  union {
    VarEnv* m varEnv;
    ExtraArgs* m extraArgs;
    StringData* m_invName;
 };
};
```

Saved FP	Saved RIP	
Func	Bytecode offset	# args / flags
"this" pointer / late-bound class	VarEnv / ExtraArgs / InvName	

```
struct ActRec {
  ActRec* m savedFp;
  uint64 t m savedRip;
  Func* m func;
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    ObjectData* m this;
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    ExtraArgs* m extraArgs;
    StringData* m invName;
 };
};
```

Saved FP	Saved RIP	
Func	Bytecode offset	# args / flags
"this" pointer / late-bound class	VarEnv / ExtraArgs / InvName	

 For unions of pointers, we steal the low bit(s) and use it as a tag to disambiguate

```
bool hasThis() const {
   return m_this && !((intptr_t)m_this & 1);
}
ObjectData* getThis() {
   return m_this;
}
void setThis(ObjectData* val) {
   m_this = val;
}
bool hasClass() {
   return (intptr_t)m_cls & 1;
}
Class* getClass() {
   return (Class*)((intptr_t)m_cls & 1);
}
void setClass(Class* val) {
   m_cls = (Class*)((intptr_t)val | 1);
}
```

```
struct ActRec {
  ActRec* m savedFp;
  uint64 t m savedRip;
  Func* m func;
  uint32 t m soff;
  uint32_t m_numArgsAndFlags;
  union {
    ObjectData* m this;
    Class* m cls;
  };
  union {
    VarEnv* m varEnv;
    ExtraArgs* m extraArgs;
    StringData* m_invName;
 };
};
```

Saved FP	Saved RIP	
Func	Bytecode offset	# args / flags
"this" pointer / late-bound class	VarEnv / ExtraArgs / InvName	

 For m_numArgsAndFlags, we steal the high bit(s) for flags

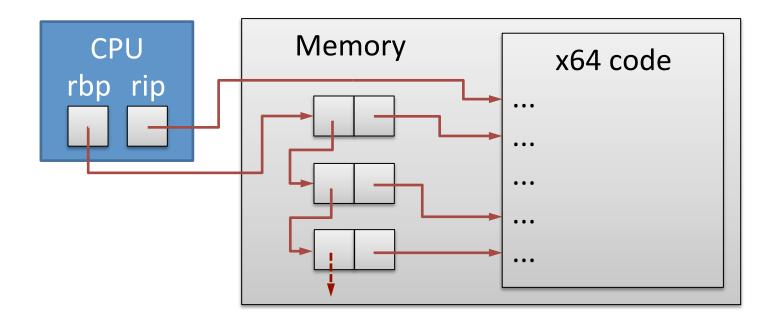
```
int32_t numArgs() {
   return m_numArgsAndFlags & 0x7fffffff;
}
bool isCtorFrame() {
   return m_numArgsAndFlags & (1 << 31);
}
void initNumArgs(uint32_t numArgs) {
   assert(!(numArgs & (1 << 31)));
   m_numArgsAndFlags = numArgs;
}
void initNumArgsCtorFrame(uint32_t numArgs) {
   assert(!(numArgs & (1 << 31)));
   m_numArgsAndFlags = numArgs | (1 << 31);
}</pre>
```

Profiling

- The C++ ecosystem has lots of mature tools for profiling
- Linux's perf tool is awesome
 - Uses a sampling technique to determine which functions a program spends the most CPU time on
 - Keeps track of the full call stack for each sample

Profiling

- On x64, Linux's perf tool works by using the *rbp-chain* to walk the call stack
- C++ code needs be compiled with frame pointers (i.e. -fno-omit-frame-pointer)



Profiling

- ActRec is designed so the first 16 bytes has the same layout as C++ frames
- HHVM sets up VM frames so that they are participate in the rbp-chain
- This makes it possible to profile PHP programs running under HHVM using Linux's perf tool

Saved FP	Saved RIP	
Func	Bytecode offset	# args / flags
This / Cls	VarEnv / ExtraArgs / InvName	

100

Events: 4K cycles

```
1.02% hhvm
                hhvm
                              [.] HPHP::ObjectData::setProp(HPHP::Class*, HPHP::S
   - HPHP::ObjectData::setProp(HPHP::Class*, HPHP::StringData const*, HPHP::Typed
      - 73.64% HPHP::JIT::MInstrHelpers::setPropCO(HPHP::Class*, HPHP::TypedValue
        + 57.02% PHP::/www/wordpress/wp-includes/post.php::WP_Post::__construct
        + 16.74% PHP::/www/wordpress/wp-includes/post.php::WP_Post::__construct
        + 16.61% PHP::/www/wordpress/wp-includes/post.php::sanitize_post
     + 16.47% HPHP::ObjectData::o_setArray(HPHP::Array const&)
     + 3.98% PHP::/www/wordpress/wp-includes/post.php::WP_Post::__construct
   0.80% hhvm hhvm
                              [.] HPHP::f_in_array(HPHP::Variant const&, HPHP::Va
   - HPHP::f_in_array(HPHP::Variant const&, HPHP::Variant const&, bool)
     + 49.94% PHP::/www/wordpress/wp-includes/post.php::sanitize_post_field
     + 26.36% PHP::/www/wordpress/wp-includes/post.php::sanitize_post_field
     + 5.86% PHP::/www/wordpress/wp-includes/option.php::get_option
     + 3.90% PHP::/www/wordpress/wp-includes/post.php::sanitize_post
   0.73% hhvm hhvm
                              [.] HPHP::Unit::GetNamedEntity(HPHP::StringData con
   - HPHP::Unit::GetNamedEntity(HPHP::StringData const*, bool, HPHP::String*)
      - 31.10% HPHP::f_is_a(HPHP::Variant const&, HPHP::String const&, bool)
        + 73.09% PHP::/www/wordpress/wp-includes/post.php::get_post
        + 25.37% PHP::/www/wordpress/wp-includes/post.php::get_post
        + 1.17% PHP::/www/wordpress/wp-includes/post.php::get_post
     - 27.98% HPHP::Unit::loadFunc(HPHP::StringData const*)
         - 85.99% HPHP::JIT::fpushCufHelperString(HPHP::StringData*, HPHP::ActRec
           + 97.73% PHP::/www/wordpress/wp-includes/plugin.php::apply_filters
           + 0.77% PHP::/www/wordpress/wp-includes/plugin.php::apply_filters
Press '?' for help on key bindings
```

Exception handling / unwinding

- setjmp / longjmp can be useful when implementing exception handling for a VM
- HHVM's initial EH implementation used setjmp and longjmp to skip over the enterTCHelper trampoline and VM frames when an exception was thrown
 - This scheme was a bit clunky but it got
 HHVM's EH system up and running quickly

Exception handling / unwinding

Later, HHVM switched over to using g++'s
 __register_frame() function to integrate
 with the C++ runtime's exception unwinding
 system by registering a "personality routine"

```
_Unwind_Reason_Code

tc_unwind_personality(
   int version,
   _Unwind_Action actions,
   uint64_t exceptionClass,
   _Unwind_Exception* exceptionObj,
   _Unwind_Context* context
);
```

C++ / machine code interop

- On x64 and other popular platforms,
 C++ has a well-defined ABI for function calls
- Generating calls from machine code to C++ functions is light-weight and easy*

^{*} Provided that parameters are pointers or primitive types and virtual methods are not involved

C++ / machine code interop

- Making C++ call into machine code is also relatively simple
- C++ allows the programmer to take a void* that points to machine code and cast it to a function pointer and invoke it:

```
void* p = getMachineCodeAddress();
typedef int(*FuncPtr)(int);
FuncPtr fn = (FuncPtr)p;
int result = fn(123);
```

enterTCHelper example

enterTCHelper:

```
push %rbp
push %rcx
mov %rdi,%rbx
mov %r8,%r12
mov %rsi,%rbp
call *%rdx
pop %rbx
mov %rdi,0x0(%rbx)
mov %rsi,0x8(%rbx)
...
pop %rbp
ret
```

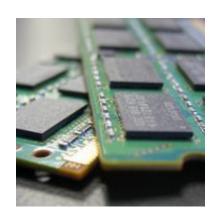
```
// Parameters: rdi, rsi, rdx, rcx, r8
extern "C" void
enterTCHelper(Cell* sp, ActRec* fp,
              TCA start, TReqInfo* info,
              void* tlBase);
inline void
enterTC(TCA start, TReqInfo& info) {
  asm volatile("" : : :
    "rbx", "r12", "r13", "r14", "r15");
  auto& regs = vmRegsUnsafe();
  enterTCHelper(regs.stkTop(), regs.fp,
                start, &info, tl base());
  asm volatile("" : : :
    "rbx", "r12", "r13", "r14", "r15");
```

Manual Memory Management

- C++ gives the programmer freedom to manually manage memory; this was essential for HHVM
- High compatibility with the php.net interpreter was important; given PHP's semantics, refcounting was less risky
 - Thus tracing GC was not an attractive option at the outset

Native allocation

 C++ provides an easy and light-weight means for the programmer to choose which implementation of malloc they want to use

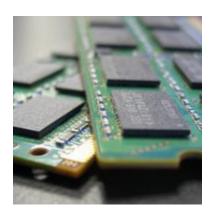


jemalloc

 HHVM uses jemalloc to handle calls to malloc APIs such as malloc(), free(), calloc(), etc.

jemalloc

 jemalloc is a world-class concurrent memory allocator used by the FreeBSD operating system, Firefox, and many other software projects



jemalloc

- Very efficient both for single- and multithreaded programs
 - Particularly good at minimizing fragmentation for long-running processes

jemalloc

- HHVM takes advantage of jemalloc-specific APIs that aren't part of the standard malloc interface
- Some features were added to jemalloc to specifically to help out HHVM
 - jemalloc's mallocx() and allocm() APIs were updated to support allocating memory with "low" addresses that can fit within 32 bits
 - Sophisicated heap profiling functionality was added to jemalloc to aid with investigating how to improve HHVM's performance

Low addresses

 Allocating memory with low addresses helps make HHVM's data structures smaller and helps with generating more more efficient machine code

```
struct LowClassPtr {
  int32_t m_raw;
  Class* get() {
    return reinterpret_cast<Class*>(m_raw);
  }
  void set(Class* c) {
    m_raw = reinterpret_cast<int32_t>(c);
  }
};
```

Low addresses

```
movl $0x7671a20, -0x20(%rbx)
[c7 43 e0 20 1a 67 07]
```

7 bytes total

```
movabs $0x7fffe8408000, %rax

mov %rax, -0x20(%rbx)

[48 b8 00 80 40 e8 ff 7f 00 00 48 89 43 e0]

14 bytes total
```

Memory Management for PHP

- The PHP language requires that engines provide some form of automatic memory management for programs written in PHP
- PHP models concurrency using separate requests
- Each request has its own distinct heap, and at the end of a request the entire heap dies
- A given request cannot directly access the heap of another request

HHVM's request allocator

- HHVM implements its own custom request allocator
 - Maintains separate isolated heap for each request; this means refcounting doesn't need to use atomic inc/dec
 - Optimizes reclamation at the end of the request
 - Avoids some of the overheads of malloc
 - Controlling the allocator implementation makes it easier to optimize how JIT'd code interacts with the allocator

Huge pages

- Manual memory management also made it possible for HHVM to take advantage of madvise()'s MADV_HUGEPAGE feature which takes advantage of larger page sizes supported by the TLB hardware
 - By using fewer iTLB entries we were able to significantly reduce iTLB misses, which gave us a nice boost for larger PHP codebases

Obstacles

- There were some features of C++ that posed challenges for us when building HHVM
- For most of these obstacles there
 was a way to hack around them,
 which is a testament to C++'s power
 and flexibility

Unions and non-POD types

- Before C++11, unions could only work with "plain old data" (POD) types
- For HHVM, we wanted to reuse parts of HPHPc's runtime that dealt with iterating over PHP arrays, but we encountered a problem where we needed to make a union of non-POD types

Unions and non-POD types

```
struct Iter {
 bool init(TypedValue* c) {
    new (&arr()) ArrayIter(c);
 bool minit(TypedValue* v) {
    new (&marr()) MArrayIter(v);
 void free() { arr().~ArrayIter(); }
 void mfree() { marr().~MArrayIter(); }
 ArrayIter& arr() { return *(ArrayIter*)m u; }
 MArrayIter& marr() { return *(MArrayIter*)m_u; }
 char m_u[MAX(sizeof(ArrayIter),
               sizeof(MArrayIter))];
} __attribute__ ((aligned(16)));
```

C++11's Unrestricted Unions

```
struct Iter {
  bool init(TypedValue* c) {
    new (&m u.arr) ArrayIter(c);
  bool minit(TypedValue* v) {
    new (&m_u.marr) ArrayIter(v);
  void free() { m_u.arr.~ArrayIter(); }
  void mfree() { m_u.marr.~MArrayIter(); }
  union Data {
    ArrayIter arr;
   MArrayIter marr;
  } m_u;
```

Unnecessary refcounting with smart pointers

 Before move constructors and rvalues were introduced in C++11, smart pointers would often do unnecessary refcounting:

```
class Variant {
   Variant(String s) { ... }
};
class String {
   ...
};
String foo() { ... }
Variant v(foo());
```

Avoiding unnecessary refcounting

 Prior to C++11, we did some awkward dances to avoid unnecessary refcounting:

```
class Variant {
  enum NoInc { noInc = 0 };
  Variant(StringData* s, NoInc) { .. }
};
class String {
  StringData* detach() {
    auto p = m_data;
    m_data = nullptr;
    return p;
  }
};
String foo() { .. }

Variant v(foo().detach(), Variant::noInc);
..
```

C++11's Move Constructors

 Move constructors made it a lot easier to avoid unnecessary refcounting with smart pointers:

```
class Variant {
   Variant(String&& s) { .. s.detach() .. }
};
class String {
   StringData* detach() {
     auto p = m_data;
     m_data = nullptr;
     return p;
   }
};
String foo() { .. }

Variant v(foo());
..
```

C++ instance methods

- Generating machine code that calls into a non-virtual C++ instance method was a bit difficult
 - Getting at the machine code address for the method not straight-forward
- For HHVM, we wanted to reuse some existing parts of HPHPc's runtime that used non-virtual C++ instance methods
- We found a way to make it work for g++

C++ instance methods

```
template <typename MethPtr>
void* getMethodPtr(MethPtr p) {
  union U { MethPtr meth; void* ptr; };
  return ((U*)&p)->ptr;
}
class C {
  public: int foo(int x) { .. }
};
void generateCallSite() {
 void* addr = getMethodPtr(&C::foo);
 printf("callq 0x%x\n", addr);
}
// Example output
callq 0x400570
```

C++ instance methods

```
template <typename MethPtr>
void* getMethodPtr(MethPtr p) {
  union U { MethPtr meth; void* ptr; };
  return ((U*)&p)->ptr;
}
class C {
  public: int foo(int x) { .. }
};
void test(C* c) {
  typedef int (*FuncPtr)(void*,int);
  void* addr = getMethodPtr(&C::foo);
  ((FuncPtr)addr)(c, 123);
```

C++ virtual methods

- Generating machine code that calls into a C++ virtual method was tricky
- We found a way to make it work for g++
- For HHVM, we wanted to reuse some existing parts of HPHPc's runtime that used C++ virtual methods

C++ virtual methods

```
template <typename MethPtr>
int getVTableOffset(MethPtr p) {
  union U { MethPtr meth; int64_t off; };
  return ((U^*)\&p)->off - 1;
class C {
  public: virtual int foo(int x) { .. }
};
void generateCallSite() {
  int off = getVTableOffset(&C::foo);
  printf("mov (%rdi), %rax\n");
  printf("callq *0x%x(%rax)\n", (int)off);
// Example output
mov (%rdi), %rax
callq *0x8(%rax)
```

C++ virtual methods

```
template <typename MethPtr>
int getVTableOffset(MethPtr p) {
  union U { MethPtr meth; int64_t off; };
  return ((U^*)\&p)->off - 1;
void* getVirtMethAddr(void* obj, int off) {
  return *(void**)(*(intptr t*)obj + off);
class C {
  public: virtual int foo(int x) { .. }
};
void test(C* c) {
  typedef int (*FuncPtr)(void*,int);
  int64 t off = getVTableOffset(&C::foo);
  ((FuncPtr)getVirtMethAddr(c, off))(c, 123);
```



Questions?

Website: hhvm.com

Facebook Page: facebook.com/hhvm

Github: github.com/facebook/hhvm