#### **Structured Systems Analysis**

Three popular graphical specification methods of '70s

- DeMarco
- Gane and Sarsen
- Yourdon

All equivalent

All equally good

Many U.S. corporations use them for commercial products

Gane and Sarsen used for object-oriented design

# **Structured Systems Analysis Case Study**

Sally's Software Store buys software from various suppliers and sells it to the public. Popular software packages are kept in stock, but the rest must be ordered as required. Institutions and corporations are given credit facilities, as are some members of the public. Sally's Software Store is doing well, with a monthly turnover of 300 packages at an average retail cost of \$250 each. Despite her business success, Sally has been advised to computerize. Should she?

Better question

What sections?

Still better

How? Batch, or online? In-house or out-service?

# Case Study (contd)

- Fundamental issue
  - What is Sally's objective in computerizing her business?
- Because she sells software?
  - o She needs an in-house system with sound and light effects
- Because she uses her business to launder "hot" money?
  - o She needs a product that keeps five different sets of books, and has no audit trail
- Assume: Computerization "in order to make more money"
  - o Cost/benefit analysis for each section of business

#### **Case Study (contd)**

- The danger of many standard approaches
  - o First produce the solution, then find out what the problem is!
- Gane and Sarsen's method
  - o Nine-step method
  - Stepwise refinement is used in many steps

## Case Study (contd)

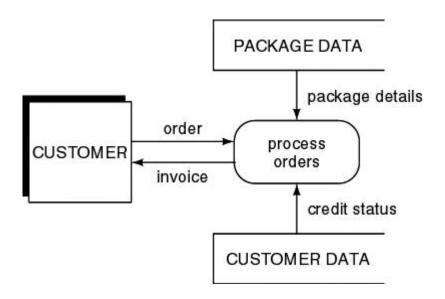
Data flow diagram (DFD) shows logical data flow

"what happens, not how it happens"

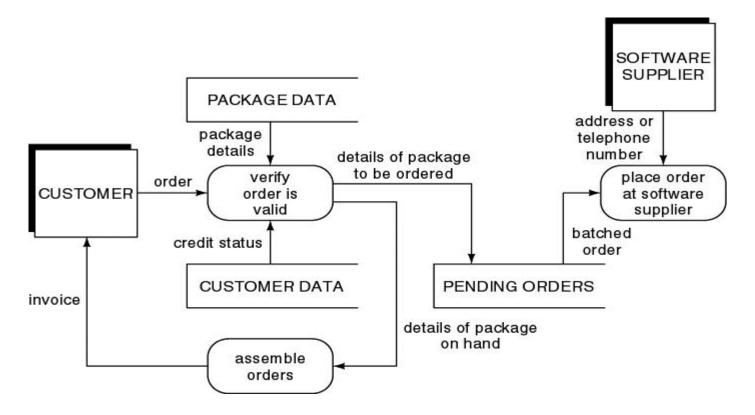
#### Step 1. Draw the DFD

First refinement

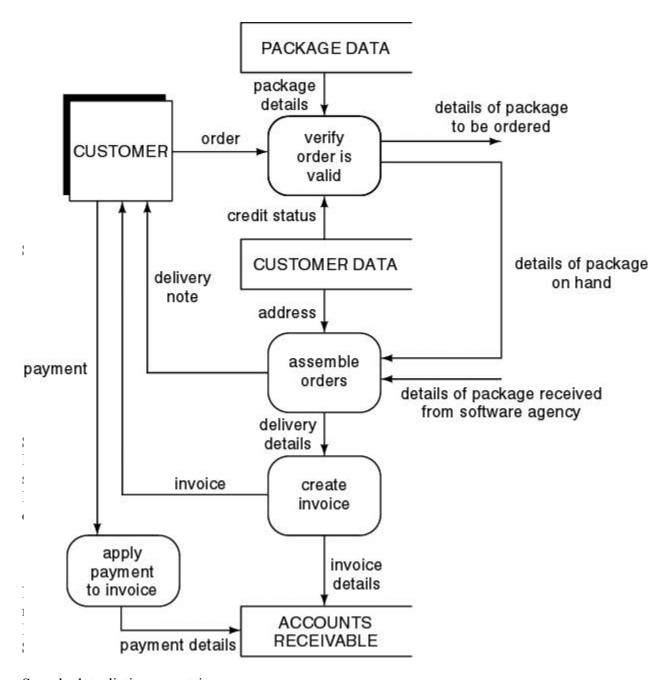
Infinite number of possible interpretations



Step 1 (contd)
Second refinement
pending orders scanned daily



Step 1 (contd)
Portion of third refinement



Sample data dictionary entries

## **Step 4. Refine Logic of Processes**

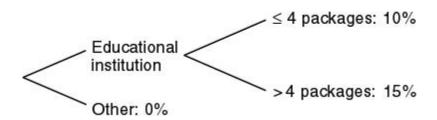
Determine what happens with in each process.

Have process give educational discount

- Sally must explain discount for educational institutions
- 10% on up to 4 packages, 15% on 5 or more

Translate into decision tree

## give educational discount

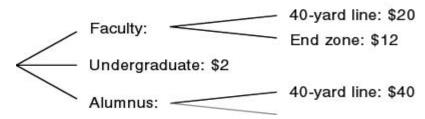


## Step 4 (contd)

Advantage of decision tree

Missing items are quickly apparent

#### football seats



Can also use decision tables

CASE tools for automatic translation

#### **Step 5. Define Data Stores**

Define exact contents of each store and representation (format)

- COBOL: specify to pic level
- Ada: specify digits or delta

Specify where immediate access is required

• Data immediate access diagram (DIAD)

#### **Step 6. Define Physical Resources**

For each file, specify

- File name
- Organization (sequential, indexed, etc.)
- Storage medium
- Blocking factor(decisions can be made bcz of this)
- Records (to field level)

If a database management system(DBMS) is to be used, then the relevant information for each table is specified here

# **Step 7. Determine Input/Output Specifications**

Specify input forms, input screens, printed output

### **Step 8. Determine Sizing**

Computing Numerical data for Step 9 to determine hardware requirements

Volume of input (daily or hourly)

Size, frequency of each printed report and its deadlines

Size, number of of each type passing between CPU and mass storage

Size of each file

# **Step 9. Hardware Requirements**

- Mass storage for back-up
- Input needs
- Output devices
- Is existing hardware adequate?
- If not, recommend buy/lease

After approval by client: Specification document is handed to design team and software process continues

#### Petri Nets

- A major difficulty with specifying real-time systems is coping with timing
  - o This difficulty itself manifest in different ways
  - Synchronization problems
  - Race conditions
  - o Deadlock

Often a consequence of poor specifications

#### Petri Nets (contd)

- o Petri nets
  - o Powerful technique for specifying systems with potential timing problems
- o A Petri net consists of four parts:
  - Set of places P
  - o Set of transitions T
  - o Input function I
  - Output function O

### Petri Nets (contd)

#### Petri Nets (contd)

- $\circ$  More formally, a Petri net is a 4-tuple C = (P, T, I, O)
- o  $P = \{p_1, p_2, ..., p_n\}$  is a finite set of *places*,  $n \ge 0$
- $\begin{array}{ll} \circ & T = \{t_1,\,t_2,\ldots,t_m\} \text{ is a finite set of } \textit{transitions},\, m \geq 0, \text{ with } P \text{ and } T \text{ disjoint} \\ \circ & I:T \ {\rm \rlap{\@ }} P^{\rm \rlap{\@ }} P^{\rm \rlap{\@ }} \text{ is } \textit{input } \text{function, mapping } \text{from transitions to bags of places} \\ \end{array}$

- o (A bag is a generalization of sets which allows for multiple instances of element in bag, as in example above)
- o Marking of a Petri net is an assignment of tokens to that Petri net

#### Petri Nets (contd)



Four tokens, one in p<sub>1</sub>, two in p<sub>2</sub>, none in p<sub>3</sub>, and one in p<sub>4</sub>

Represented by vector (1,2,0,1)

Transition is enabled if each of its input places has as many tokens in it as there arcs from the place to that transition

# Petri Nets (contd)

Transition  $t_1$  is enabled (ready to fire)

If  $t_1$  fires, one token is removed from  $p_2$  and one from  $p_4$ , and one new token is placed in  $p_1$ 

Important: Number of tokens is not conserved

Transition t<sub>2</sub> is also enabled

## Petri Nets (contd)

Petri nets are indeterminate

Suppose t<sub>1</sub> fires

Resulting marking is (2,1,0,0)

## Petri Nets (contd)

Now only t<sub>2</sub> is enabled

It fires

Marking is now (2,0,2,0)

# Petri Nets (contd)

More formally, a marking M of a Petri net places P to the non-negative integers N

C = (P, T, I, O) is a function from the set of

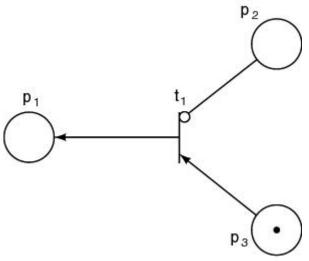
 $M: P \otimes N$ 

A marked Petri net is then 5-tuple (P, T, I, O, M)

#### Petri Nets (contd)

Inhibitor arcs

Inhibitor arc is marked by small circle, not arrowhead



Transition  $t_1$  is enabled

In general, transition is enabled if at least one token on each (normal) input arc, and no tokens on any inhibitor input arcs

Elevator Problem: Petri Net

- Product is to be installed to control n elevators in a building with m floors
- Each floor represented by place  $F_f$ ,  $1 \le f \le m$
- Elevator represented by token

Token in F<sub>f</sub> denotes that an elevator is at floor F<sub>f</sub>

Elevator Problem: Petri Net (contd)

First constraint

1. Each elevator has a set of m buttons, one for each floor. These illuminate when pressed and cause the elevator to visit the corresponding floor. The illumination is canceled when the corresponding floor is visited by an elevator

Elevator button for floor f is represented by place  $EB_f$ ,  $1 \le f \le m$ Token in  $EB_f$  denotes that the elevator button for floor f is illuminated

Elevator Problem: Petri Net (contd)

A button must be illuminated the first time button is pressed and subsequent button presses must be ignored .this is specified in below fig

If button EB<sub>f</sub> is not illuminated, no token in place and transition EB<sub>f</sub> pressed is enabled Transition fires, new token is placed in EB<sub>f</sub>

Now, no matter how many times button is pressed, transition EB<sub>f</sub> pressed cannot be enabled Elevator Problem: Petri Net (contd)

When elevator reaches floor g, token is in place  $F_g$ , transition Elevator in action is enabled, and then fires

Tokens in EB<sub>f</sub> and F<sub>g</sub> removed

This turns off light in button EB<sub>f</sub>

New token appears in F<sub>f</sub>

This brings elevator from floor g to floor f

Elevator Problem: Petri Net (contd)

Motion from floor g to floor f cannot take place instantaneously

Timed Petri nets

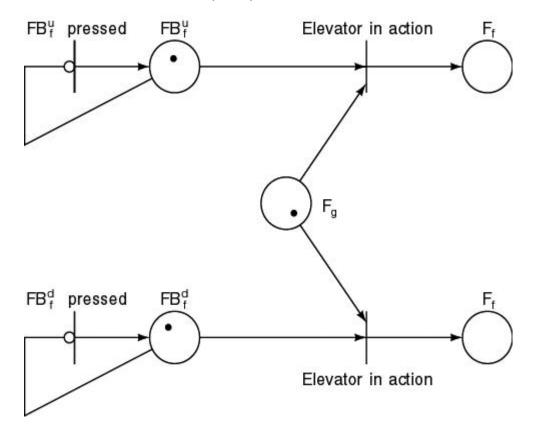
Elevator Problem: Petri Net (contd)

Second constraint

2. Each floor, except the first and the top floor, has 2 buttons, one to request an upelevator, one to request a down-elevator. These buttons illuminate when pressed. The illumination is canceled when the elevator visits the floor, and then moves in desired direction

Floor buttons represented by places FB<sup>u</sup><sub>f</sub> and FB<sup>d</sup><sub>f</sub>

Elevator Problem: Petri Net (contd)



Elevator Problem: Petri Net (contd)

The situation when an elevator reaches floor f from floor g with one or both buttons illuminated If both buttons are illuminated, only one is turned off

(A more complex model is needed to ensure that the correct light is turned off)

Elevator Problem: Petri Net (contd)

Third constraint

C <sub>3</sub> . If an elevator has no requests, it remains at its current floor with its doors closed If no requests, no Elevator in action transition is enabled
The Data Dictionary  Building a Data Dictionary
Data Dictionary Notation
Data Dictionary Example