

Functional Blocks: Addition

- Binary addition used frequently
- Addition Development:
 - *Half-Adder* (HA), a 2-input bit-wise addition functional block,
 - *Full-Adder* (FA), a 3-input bit-wise addition functional block,
 - *Ripple Carry Adder*, an iterative array to perform binary addition, and
 - *Carry-Look-Ahead Adder* (CLA), a hierarchical structure to improve performance.

Functional Block: Half-Adder

- A 2-input, 1-bit width binary adder that performs the following computations:

X	0	0	1	1
+ Y	+ 0	+ 1	+ 0	+ 1
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C S	0 0	0 1	0 1	1 0

- A half adder adds two bits to produce a two-bit sum
- The sum is expressed as a sum bit, S and a carry bit, C
- The half adder can be specified as a truth table for S and C \Rightarrow

X	Y	C	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

Logic Simplification: Half-Adder

- The K-Map for S, C is:
- This is a pretty trivial map!
By inspection:

$$S = XY' + X'Y = X \oplus Y$$

$$S = (X+Y) (X'+Y')$$

Recall that $S' = X'Y' + XY$

and

$$C = XY$$

$$C = ((XY)')'$$

- These equations lead to several implementations.

S		Y	
		1	
X	1		

C		Y	
X		1	

Five Implementations: Half-Adder

- We can derive following sets of equations for a half-adder:

$$(a) S = XY' + X'Y$$

$$C = XY$$

$$(b) S = (X+Y)(X'+Y')$$

$$C = XY$$

$$(c) S = (C + X'Y')'$$

$$C = XY$$

$$(d) S = (X+Y) C'$$

$$C' = X'+Y'$$

$$(e) S = X \oplus Y$$

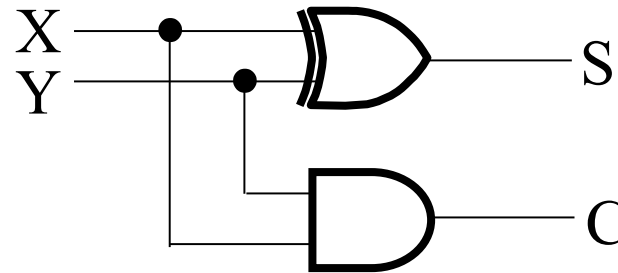
$$C = XY$$

- (a), (b), and (e) are SOP, POS, and XOR implementations for S.
- In (c), the C function is used as a term in the AND-NOR implementation of S, and in (d), the C' function is used in a POS term for S.

Implementations: Half-Adder

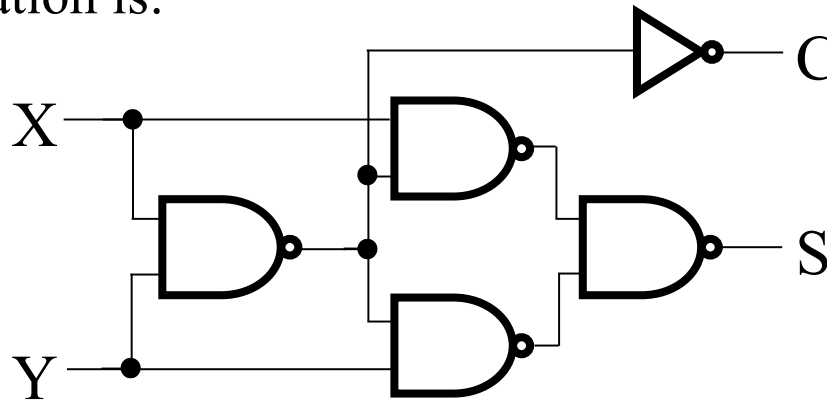
- The most common half adder implementation is:
(e)

$$S = X \oplus Y$$
$$C = XY$$



- A NAND only implementation is:

$$S = (X+Y) C'$$
$$C = ((XY)')'$$



Functional Block: Full-Adder

- A full adder is similar to a half adder, but includes a carry-in bit from lower stages. Like the half-adder, it computes a sum bit, S and a carry bit, C.

- For a carry-in (Z) of 0, it is the same as the half-adder:

Z	0	0	0	0
X	0	0	1	1
+ Y	+ 0	+ 1	+ 0	+ 1
C S	0 0	0 1	0 1	1 0

- For a carry- in (Z) of 1:

Z	1	1	1	1
X	0	0	1	1
+ Y	+ 0	+ 1	+ 0	+ 1
C S	0 1	1 0	1 0	1 1

Logic Optimization: Full-Adder

- Full-Adder Truth Table:

X	Y	Z	C	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

- Full-Adder K-Map:

K-Map for Sum (S):

			Y
		1	1
	0	1	3
			2
X	1		1
	4	5	7
			6
		Z	

K-Map for Carry (C):

			Y
		1	1
	0	1	3
			2
X		1	1
	4	5	7
			6
		Z	

Equations: Full-Adder

- From the K-Map, we get:

$$S = XY'Z' + X'YZ' + X'Y'Z + XYZ$$

$$C = XY + XZ + YZ$$

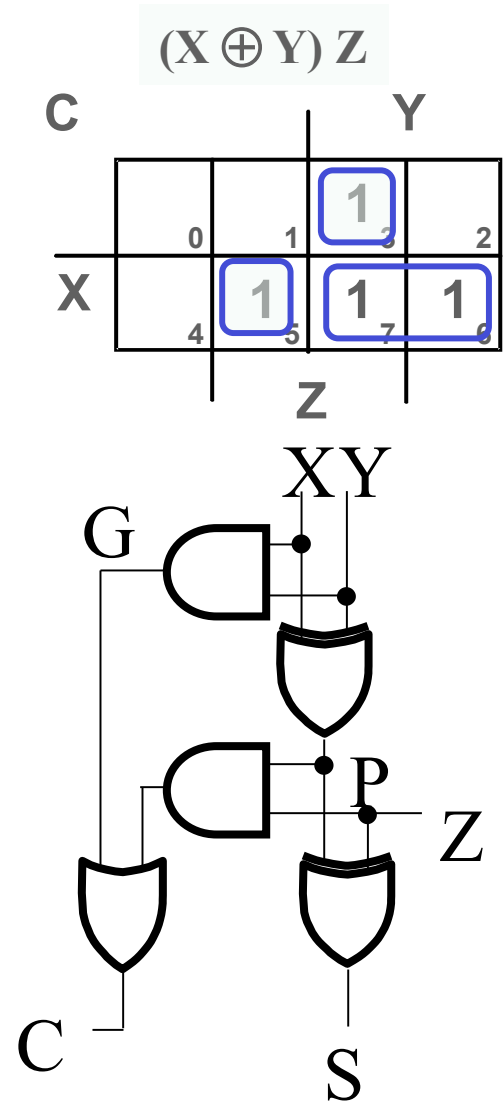
- The S function is the three-bit XOR function (Odd Function):

$$S = X \oplus Y \oplus Z$$

- The Carry bit **C** is 1 if both **X** and **Y** are 1 (the sum is 2), or if the sum ($X \oplus Y$) is 1 and a carry-in (**Z**) occurs. Thus C can be re-written as:

$$C = XY + (X \oplus Y) Z$$

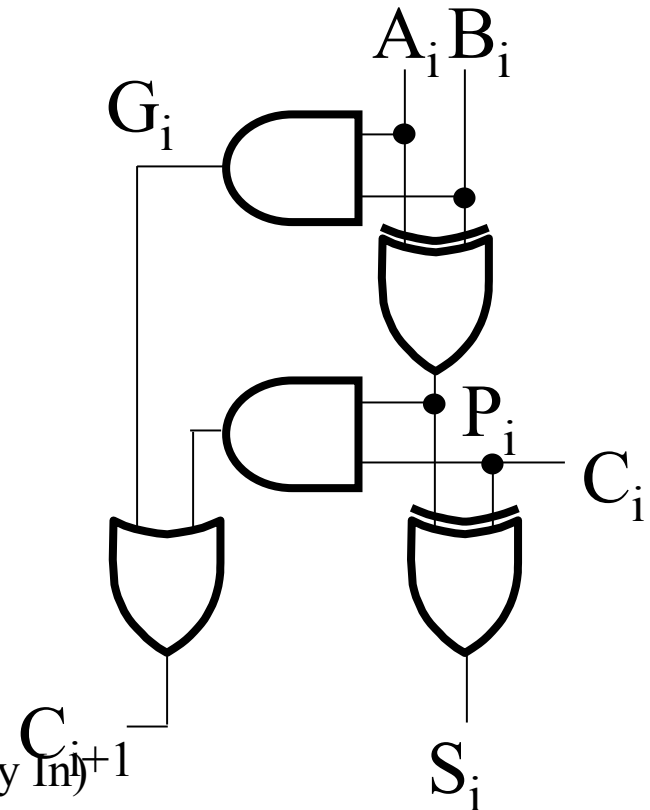
- The term $X \cdot Y$ is *carry generate* (G).
- The term $X \oplus Y$ is *carry propagate* (P).



Implementation: Full Adder

- Full Adder Schematic
- Here X, Y, and Z, and C (from the previous pages) are A, B, C_i and C_o , respectively. Also,
 G = generate and
 P = propagate.
- Note: This is really a combination of a 3-bit odd function (for S) and Carry logic (for C_o):

(G = Generate) OR (P = Propagate AND C_i = Carry In)
 $C_o = G + P \cdot C_i$



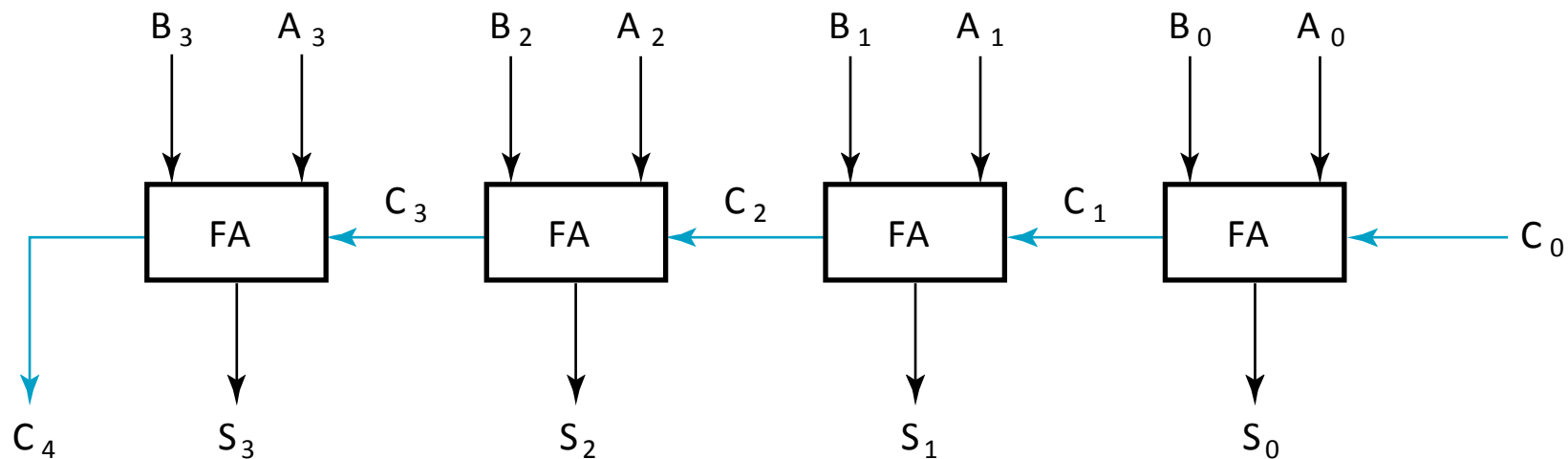
Binary Adders

- To add multiple operands, we “bundle” logical signals together into vectors and use functional blocks that operate on the vectors
- Example: 4-bit ripple carry adder: Adds input vectors $A(3:0)$ and $B(3:0)$ to get a sum vector $S(3:0)$
- Note: carry out of cell i becomes carry in of cell $i + 1$

Description	Subscript 3 2 1 0	Name
Carry In	0 1 1 0	C_i
Augend	1 0 1 1	A_i
Addend	<u>0 0 1 1</u>	B_i
Sum	1 1 1 0	S_i
Carry out	0 0 1 1	C_{i+1}

4-bit Ripple-Carry Binary Adder

- A 4-bit Ripple Carry Adder made from four 1-bit Full Adders:



Binary Adder/Subtractor

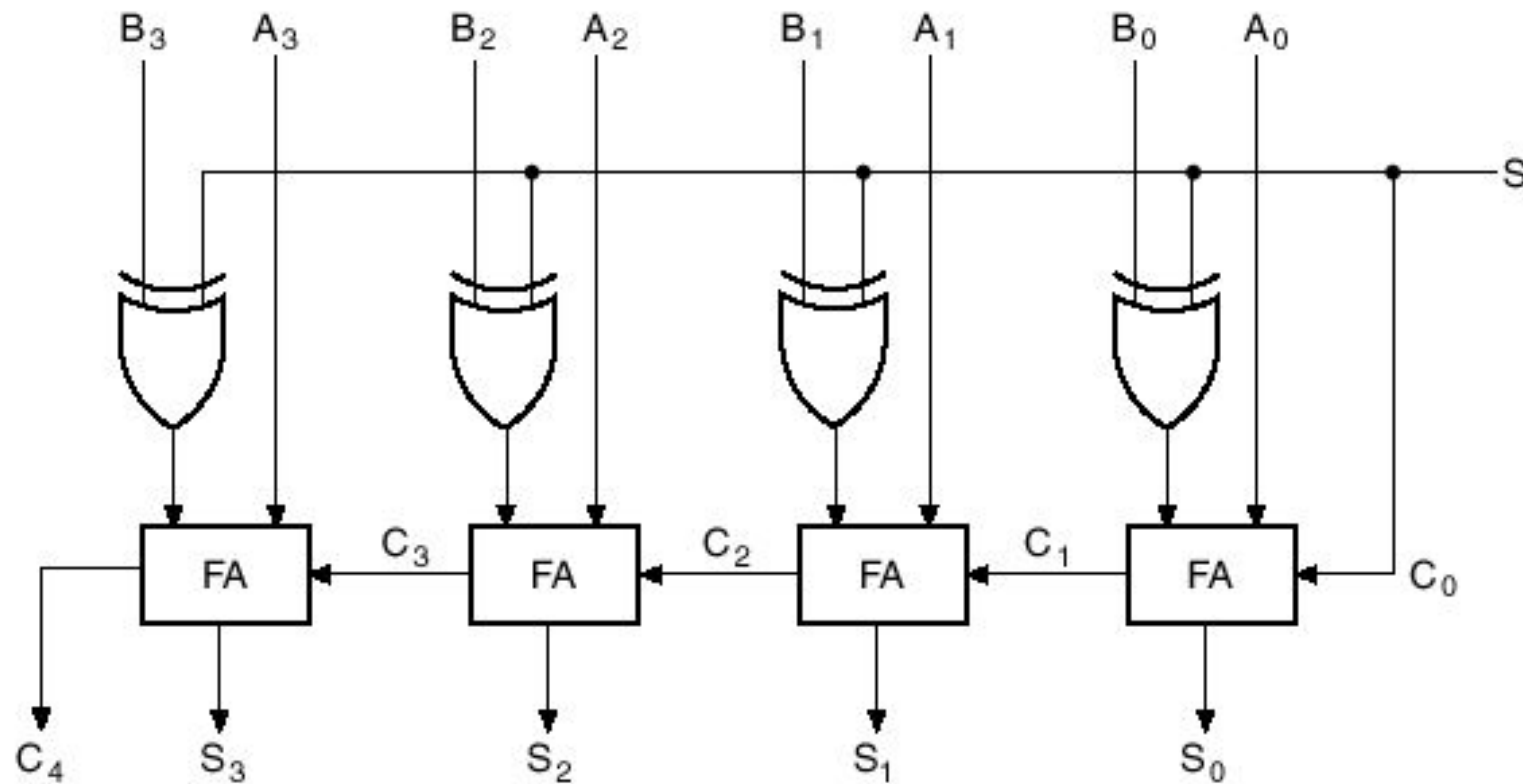


Fig. 3-31 Adder-Subtractor Circuit