

上 海 交 通 大 学 试 卷 (A 卷)

(2020 至 2021 学 年 第 1 学 期)

班级号 _____ 学号 _____ 姓名 _____

课程名称 _____ 计算机系统基础 (汇编) _____ 成绩 _____

Problem 1

[1] [2]

[3] [4]

[5] [6]

Problem 2

[1] [2]

[3] [4]

[5] [6]

[7] [8]

[9] [10]

Problem 3

1. [1] [2]

[3] [4]

2. FP=

sign=

exp=

frac=

Problem 4

1. [1] [2]

[3] [4]

[5] [6]

[7] [8]

我承诺，我将严
格遵守考试纪律。

题号	1	2	3	4	5	6			
得分									
批阅人(流水阅 卷教师签名处)									

2.

3.

Problem 5

1.

2. [1]
- [2]
- [3]
- [4]
- [5]
- [6]
- [7]

3.

str	output
"62"	0x2
"26"	
"84"	
"4791"	

4.

Problem 6

1. [1] [2] [3]

2.

3 [4] [5]

[6] [7]

[8] [9]

4 4.a

[10]

[11]

4.b

Problem 1 (18 points)

1. Consider the following C program

```
int a = 0x62;  
int b = a << 2;  
unsigned int c = a - 0x63;  
short d = !a | 0;  
int e = a & 9;
```

Assume the program will run on an **8-bit** machine and use two's complement arithmetic for signed integers. A 'short' integer is encoded in **4 bits**, while a normal 'int' is encoded in **8 bits**. Please fill in the blanks below. (3*6=18')

Expression	Binary Representation
a	0110 0010
b	[1]
c	[2]
d	[3]
e	[4]
a ^ a	[5]
~a	[6]

Problem 2 (10 points)

Suppose a **64-bit little-endian** machine has the following memory and register status.

Memory status

Address	Low	High
0x4000	0xaf 0xbe 0x00 0x00 0x00 0x00 0x00 0x00	
0x4008	0x56 0x34 0x12 0x00 0x00 0x00 0x00 0x00	
0x4010	0xfd 0xff 0xff 0xff 0xff 0xff 0xff 0xff	
0x4018	0x12 0x23 0x34 0x45 0x56 0x67 0x78 0x89	
0x4020	0xf0 0xff 0xff 0xff 0x00 0x00 0x00 0x00	
0x4028	0xf0 0xe0 0xd0 0xc0 0xb0 0xa0 0x90 0x80	

Register status

Register	Hex Value
%rax	0x00000000 00002020
%rbx	0x8fffffff 000000ab
%rcx	0xffffffff ffffffff
%rdx	0x00000000 00004030
%rsi	0x00000000 00000004
%rsp	0x00000000 00004028

The following instructions are executed **sequentially**.

	Operation
1	subq \$2020, %rax
2	movq %rax, (%rdx,%rcx,8)
3	imulq %rcx, %rsi
4	sarq \$0x4, %rbx
5	pushq %rcx

After executing five instructions above, please fill in the blanks below. For 'Hex Value', write in **8-byte hex value**. For example, the value on the address from 0x4000 to 0x4007 are 0xaf 0xbe 0x00 0x00 0x00 0x00 0x00 0x00, the hex value should be 0xbeaf. (1' *10=10')

Address	Hex Value
0x4000 ~ 0x4007	0xbeaf
0x4008 ~ 0x400f	[1]
0x4010 ~ 0x4017	[2]
0x4018 ~ 0x401f	[3]
0x4020 ~ 0x4027	[4]
0x4028 ~ 0x402f	[5]

Register	Hex Value
%rax	[6]
%rbx	[7]
%rcx	[8]
%rdx	0x4030
%rsi	[9]
%rsp	[10]

Problem 3 (12 points)

The following figure is a 16-bit floating point representation based on the IEEE floating point format. Assume we use the IEEE round-to-even mode to do the approximation.

sign (1bit)	exp (8bits)	frac (7bits)
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- Fill the blanks with proper values. ($2^4=16$)
 - Normalized: $(-1)^{\text{sign}} \times (1.\text{fraction}) \times 2^{\text{exp}-\text{bias}}$, where bias= [1] ;
 - Infinite ($-\infty$) (in hexadecimal form): [2] ;
 - Largest Negative Normalized Value (in hexadecimal form): [3] ;
 - Smallest Positive Denormalized Value (in hexadecimal form): [4] ;
- Consider the number $(20.77)_{10}$. Please convert it into the floating point format (hexadecimal) we designed above. You need also provide sign, exp, and frac part separately in hexadecimal form. (4')

Problem 4 (20 points)

Please answer the following questions according to the definition of **heterogeneous data structures in x86-64**. (**NOTE**: the size of data types is shown in Figure 3.1 in ICS book.)

```
union u {  
    union u *ptr;  
    struct s {  
        char data[2];  
        int *loc;  
        short x;  
    } s;  
} u;
```

1. Fill in the following blocks. (2'*8=16')

Representation	Value
<code>sizeof(u)</code>	[1]
<code>sizeof(u.s)</code>	[2]
<code>sizeof(u.ptr)</code>	[3]

Please represent address with **Hex**

<code>&u</code>	<code>0x550000001000</code>
<code>&(u.ptr)</code>	[4]
<code>&(u.s)</code>	[5]
<code>&(u.s.data)</code>	[6]
<code>&(u.s.loc)</code>	[7]
<code>&(u.s.x)</code>	[8]

2. How many bytes are **WASTED** in **struct s**? Explain your solution. (2')
3. Rearrange the above fields in **struct s** to conserve the most space in the memory. How many bytes are **WASTED** in **rearranged struct s**? Explain your solution. (2')

Problem 5 (25 points)

One of TA wrote a simple C program and the assembly code is provided. Suppose both of them are executed on a **64-bit little-endian** machine. Please read the code and answer the following questions.

```
// ASCII (0~9) : 0x30~0x39
int aaa(char *str){
    int result = str[0];
    int* ptr = str;
    switch(__[1]){
        case '2':
            result = __[2];
            break;
        case '7':
            result = (*ptr) >> 1;
        case '5':
            result = __[3];
            break;
        case __[4]:
            result = 9;
        default:
            result = result * 2 - 1;
    }
    return result;
}

int main(){
    char *str = "62";
    printf("0x%x\n", aaa(str));
    return 0;
}
```

<pre> 1: .section .text 2: <aaa>: 3: pushq %rbp 4: movq %rsp, %rbp 5: movsbl (%rdi), %eax 6: movl %eax, -4(%rbp) //result 7: movsbl 1(%rdi), %eax 8: subl [immediate], %eax 9: cmpl \$6, %eax 10: ja .L3 11: movq [5], %rax 12: jmp *%rax 13: .L4: 14: cmpl \$0x35, -4(%rbp) 15: jg .L10 16: movl -4(%rbp), %eax 17: jmp .L11 18: .L10: 19: movl \$2, %eax 20: .L11: 21: movl %eax, -4(%rbp) 22: jmp .L2 23: .L7: 24: movl (%rdi), %eax 25: sarl %eax 26: movl %eax, -4(%rbp) </pre>	<pre> 27: .L8: 28: movl -4(%rbp), %eax 29: notl %eax 30: movl %eax, -4(%rbp) 31: jmp .L2 32: .L9: 33: movl \$9, -4(%rbp) 34: [6] 35: .L3: 36: movl -4(%rbp), %eax 37: addl %eax, %eax 38: subl \$1, %eax 39: movl %eax, -4(%rbp) 40: .L2: 41: movl -4(%rbp), %eax 42: popq %rbp 43: ret 44: 45: .section .rodata 46: .align 8 47: .L5: 48: .quad .L4 49: .quad .L3 50: .quad .L9 51: .quad .L8 52: .quad .L3 53: .quad [7] </pre>
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- Given that the minimum case number is '2', what immediate number should be filled in Line 8 of the assembly code?(2')
- Please fill in rest of the blanks within C code and assembly code. **If you think nothing is required to write, please write NONE.** (2'*7=14')
- If we change the value of 'str' in 'main', what is the output of this program?(1*3=3')

str	output
"62"	0x2
"26"	[1]
"84"	[2]
"4791"	[3]

- Please explain the advantage and limitation of using "Jump Table" to implement 'switch' statement, and provide a simple code which is not suitable to be translated into a "Jump Table" (6').

Problem 6 (15 points)

One of the TAs wrote a **strange function bar**. Answer the questions below.

<pre> void bar(long a, long b, long c, long d, long e, char f[], char g) { if (a <= 0) { printf("&g is: %p, g is: %d\n", &g, g); return; } bar(a-1, b-2, 0, 0, 0, 0, g); bar(b-1, a-1, 0, 0, 0, 0, g + 1); } </pre>	<pre> int main() { bar(2, 3, 0, 0, 0, 0, 0); } </pre>
<pre> <bar>: 1 pushq %r14 2 pushq %rbx 3 subq \$8, %rsp 4 movb 32(%rsp), %al 5 testq %rdi, %rdi 6 jle .L1 7 movq %rsi, %r14 8 movq %rdi, %rbx 9 subq \$1, %rdi 10 subq \$2, %rsi 11 movb %al, (%rsp) 12 callq bar 13 subq \$1, %r14 14 subq \$1, %rbx 15 movb 32(%rsp), %al 16 addb \$1, %al 17 movb %al, (%rsp) 18 movq %r14, %rdi 19 movq %rbx, %rsi 20 callq bar </pre>	<pre> 21 addq \$8, %rsp 22 popq %rbx 23 popq %r14 24 retq .L1: 25 movsbl %al, %edx 26 leaq 32(%rsp), %rsi 27 movl \$.L.str, %edi 28 xorl %eax, %eax 29 callq printf // code for return <main>: 34 subq \$8, %rsp 35 movl \$0, (%rsp) 36 movl \$2, %edi 37 movl \$3, %esi 38 callq bar // code for return </pre>

NOTE: For **ALL** the following questions, assume **BEFORE** the execution of instruction at **line 34 "subq \$8, %rsp"**, the value of the register **%rsp** is **0x7fffffffde68**.

1. Here are **3 options** to describe the purpose of instructions:

A. Passing arguments to functions

B. Save callee-saved registers

C. Save caller-saved registers

Choose an option for each of the instructions below: ($2 \times 3 = 6$)

Line 1 "pushq %r14": [1]

Line 7 "movq %rsi, %r14": [2]

Line 17 "movq %al, (%rsp)": [3]

2. Can we remove **line 15** "movb 32(%rsp), %al"? Why? (3')

3. The program outputs multiple lines, each of which contains the address of `g("&g is %p:")` and the value of `g("g is %d: ")`. Fill the table to show **the output of the program**. ($0.5 \times 6 = 3$)

"&g is: %p"	"g is: %d"
[4]	0
0x7fffffffde20	[5]
[6]	1
0x7fffffffde00	[7]
[8]	[9]

4. Jack observes that before **execution of instruction at line 20 "callq bar"**, the **current stack frame is NOT needed any more**. Then he figures out that the **"callq bar"** at **line 20** can be removed and **"retq"** at **line 24** can be replaced by a simple **"jmp"** instruction while the output value for **"g is:%d"** will remain **UNCHANGED**. (3')

a. The table below details how he modifies the assembly (Note: **nop** is an instruction that does nothing). Fill in the table.

Line before modification	Line after modification
24: retq	24: jmp bar
20: callq bar	20: nop
[10]	[11]

b. Show the output of the program after the modification of assembly.