Problem 1: Memory Allocation (14 points)

There are some **unallocated memory** and **heap**, as shown below. Heap is organized as a sequence of **contiguous** allocated and free blocks. **Allocated** blocks are shaded, and **free** blocks are blank (each block represents 1 word = 4 bytes). **Headers** and **footers** are labeled with the number of bytes and allocated bit. The allocator maintains **double-word** alignment. You are given the execution sequence of memory allocation operations (i.e., malloc() or free()) from 1 to 6.

| | ↓ P1 | | | | | | | |
|-------------|-------------|-------------|------|------|------|------|--|------|
| 16/1 | | | 16/1 | 32/0 | | | | |
| | | ↓ P2 | | | | | | |
| | 32/0 | 16/1 | | | 16/1 | 16/0 | | 16/0 |
| ↓ P3 | | | | | | | | |
| 16/1 | | | 16/1 | 24/0 | | | | 24/0 |

- 1. P4 = malloc(4)
- 2. free (P1)
- 3. P5 = malloc(9)
- 4. P6 = malloc(5)
- 5. free (P3)
- 6. P7 = malloc(2)

Please answer the questions below. Assume that **immediate coalescing** strategy and **splitting free blocks** are employed.

- 1. Assume **first-fit** algorithm is used to find free blocks. Please draw the **final** status of memory and mark with block size in headers and footers after the operation sequence is executed (5').
- 2. Assume **best-fit** algorithm is used to find free blocks. Please draw the **final** status of memory and mark with block size in headers and footers after the operation sequence is executed (5').
- 3. Please calculate the total bytes of **internal fragments** (Note: **DO NOT** consider P1, P2 and P3 for internal fragments) (4').

Problem 2: Lock (13 points)

```
2. __node_t* next;
3.
    tid t tid;
4. } node t;
5.
6. typedef struct lock t {
7.
        node_t* head;
8. } lock t;
9.
10. void lock init(lock t* lock) {
        lock->head = NULL;
11.
12. }
13. /* each caller should alloc and hold a node by itself */
14. void lock(lock t* lock, node t* node) {
15.
        if (lock == NULL || node == NULL) {
16.
           /* output error message... */
17.
           exit(-1);
18.
       }
19.
      node->next = NULL;
20.
       node->tid = gettid();
21.
       node_t* old = test_and_set(lock->head, node);
22.
       if (old == NULL) return;
23.
       old->next = node;
    park();
24.
25. }
26.
27. void unlock(lock t* lock, node t* node) {
28.
        if (lock == NULL || node == NULL) {
29.
           /* output error message... */
30.
           exit(-1);
31.
        }
32.
       if (node->next == NULL) {
33.
           if (compare and swap([1], [2], [3])) {
34.
              return;
35.
          }
36.
          while(node->next == NULL)
37.
              continue;
38.
       }
39.
        unpark(node->next->tid);
40.}
```

Above code is an implementation of a kind of lock called MCS. Each lock_t structure _ A _ 卷 总 _ 14 _ 页 第 _ 2 _ 页

represents a lock. Each thread who wants a lock will maintain a node_t structure. Nodes are connected through their arriving order. The head of lock will always point to the newest node. Test_and_set and compare_and_swap operations are atomic. DO NOT consider the CPU may reorder the instructions.

- compare_and_swap(a, b, c) will set a's value to c and return 1 if and only if a's old value equals b, otherwise return 0 and do nothing. Fill the blanks in the code. Hint: consider the situation that another thread acquires the same lock between line 32 and 33. Your code is used to keep the lock running correctly under this situation. (3')
- 2. Why do we need line **36** and **37**? (2')
- 3. There are two problems in this implementation. Figure them out and try to fix them **if possible**. **Hint:** Consider about **correctness** and **security**. (4')
- 4. Analyze the lock about its **fairness** and **performance**. (4')