

Introduction to Computer Systems

2017 Fall Middle Examination

Name_____ Student No._____ Score_____

Problem 1: (14 points)

[1]	[2]	[3]
[4]	[5]	[6]
[7]		

Problem 2: (12 points)

[1]	[2]
[3]	[4]
[5]	[6]
[7]	[8]
[9]	[10]
[11]	[12]

Problem 3: (18 points)

1. [1]	[2]
[3]	[4]
[5]	[6]
[7]	[8]
[9]	[10]

2.

3.

4.

Problem 4: (11 points)

1. [1]	[2]	[3]
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2.

3. [1] [2]

4.

Problem 5: (22 points)

1	[1]	[2]
	[3]	[4]
	[5]	[6]
	[7]	[8]
	[9]	[10]
	[11]	[12]

2

3

Problem 6: (23 points)

1	[1]	
2	[1]	[2]
3	[1]	[2]
	[3]	[4]
	[5]	[6]
	[7]	
	[8]	[9]
	[10]	[11]
	[12]	[13]
	[14]	[15]
4	[16]	

Problem 1: (14 points)

1. Consider the following C program

```
int a = 0x18;
unsigned short ua = a;
int b = ua >> 1;
short c = (b && 0) + 1;
unsigned int d = (~(unsigned int)a) ^ 0xe5;
int e = 0xf4 & 0xe3;
```

Assume the program will run on an **8-bit machine** and use **two's complement arithmetic** for signed integers. A 'short' integer is encoded in **4 bits**, while a normal 'int' is encoded in **8 bits**. Please fill in the blanks below. ($2^7=128$)

Expression	Binary Representation
a	0001 1000
ua	[1]
b	[2]
c	[3]
d	[4]
e (!0)	[5]
(e + 0x20) + (0x11 >> 2)	[6]
(ua >> d) + (b << d)	[7]

Problem 2: (12points)

Suppose a **64-bit little endian** machine has the following memory and register status. (NOTE: **Instructions are independent**). ($1^7=128$)

Memory status

Address	Low	High
0x4000	0x00 0x00 0x00 0x00 0x35 0x00 0x00 0x00	
0x4008	0x00 0x00 0x00 0x00 0x00 0x00 0x00 0x71	
0x4010	0x00 0x00 0x00 0x00 0x80 0xff 0x04 0x08	
0x4018	0x00 0x00 0x00 0x00 0xde 0xad 0xbe 0xef	
0x4020	0xde 0xad 0xbe 0xef 0xab 0xcd 0xdc 0xba	

Register status

Register	Hex Value
%rax	0x00000000 00004000
%rcx	0xffffffff ffffffff
%rdx	0xaabbccdd ababcdcd
%rbx	0x00000000 00000002
%rsp	0x00000000 00004010

Please fill in the blanks below. For 'Value', write in **8-byte hex value**. If the instruction does not change any register or memory, fill the corresponding two blanks with '--'. If the instruction changes **multiple** destinations, write all of them in blanks and make sure the destinations and updated values are listed in the **same** order.

Operation	Destination	Hex Value
subq %rcx, %rax	[1]	[2]
movl \$17, (%rax, %rbx, 8)	[3]	[4]
sarq \$32, %rdx	[5]	[6]
cmpq \$0x71, 8(%eax)	[7]	[8]
leaq 8(%rax, %rcx, 8), %rdx	[9]	[10]
popq %rcx	[11]	[12]

Problem 3: (18points)

Please answer the following questions according to the definition of heterogeneous data structures. (NOTE that the size of data types in x86-64 is shown in the Figure 3.1 in ICS book.)

```
struct node_t {
    char type;
    union data_t {
        struct {
            long lsn;
            char loaded;
            struct node_t **child_cache;
            short keys[7];
            long children[8];
        } intern;
```

```

    struct {
        struct node_t * (*split) (int);
        short keys[7];
        char values[32];
    } leaf;
} data;
char status;
} node;

union data_t *data = &(node.data);

```

This declaration illustrates that structures can be embedded within unions.

1. Fill in the following blocks. (please represent address with **Hex**) (10')

Representation	x86-64
sizeof (node)	[1]
sizeof (node.data)	[2]
sizeof (node.data.leaf)	[3]
sizeof (node.data.intern)	[4]
node	0x601060
&(node.data)	[5]
&(data->intern.loaded)	[6]
&(data->intern.children)	[7]
&(data->leaf.keys)	[8]
&(data->leaf.values)	[9]
&(node.status)	[10]

2. How many bytes are **WASTED** in **struct intern** under x86-64? Explain your solution. (3')
3. If you can **rearrange** the declarations in the **struct intern**, how many bytes of memory can you **SAVE** in **struct intern** compared to the **original declaration** under x86-64? Explain your solution (3')
4. Alice thought the struct leaf is too small, so he changed "**struct { ... } leaf;**" into "**struct { ... } leaf[2];**" based on the **UNOPTIMIZED** data structure, what's the size of **node** now? (2')

Problem 4: (11 points)

The following figure shows part of codes compiled on **an x86-64 machine**. Please answer the following question according to the code.

<pre> int test[5][7]; int main(void) { int a[3][4]; int b[4][2]; int sum = 0; int col = __[1]__; int row = __[2]__; for (int i = 0; i < 4; i++) sum += a[row][i] * b[i][col]; printf("%d\n", sum); return 0; } </pre>	<pre> 6 main: 1 pushq %rbp 2 movq %rsp, %rbp 3 subq \$0x60, %rsp 4 movl \$0x0, -0x4(%rbp) 5 movl \$0x0, -0x8(%rbp) 6 movl \$0x1, -0x10(%rbp) 7 movl \$0x2, -0xc(%rbp) 8 jmp .L2 9 .L3: 10 movl -0x8(%rbp), %eax 11 cltq 12 movl -0x10(%rbp), %edx 13 movslq %edx, %rdx 14 salq \$0x2, %rdx 15 addq %rdx, %rax 16 movl -0x40(%rbp,%rax,4), %edx 17 movl -0xc(%rbp), %eax 18 cltq 19 movl -0x8(%rbp), %ecx 20 movslq %ecx, %rcx 21 addq %rcx, %rcx 22 addq %rcx, %rax 23 movl -0x60(%rbp,%rax,4), %eax 24 imull %edx, %eax 25 addl %eax, -0x4(%rbp) 26 addl \$0x1, -0x8(%rbp) 27 .L2: 28 cmpl \$0x3, -0x8(%rbp) 29 jle .L3 30 movl -0x4(%rbp), %eax 31 movl %eax, %esi 32 movl \$0x400634, %edi 33 movl \$0x0, %eax 34 call printf 35 movl \$0x0, %eax 36 leave 37 ret </pre>
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1. Suppose the address of `test` is `0x601060`, what's the value of following expressions? (NOTE: `n` is a variable and please represent address with **Hex**) (3')

- 1) `&(test[1][3]):` [1] ;
- 2) `&(test[3][n]):` [2] ;
- 3) `&(test[n][2]):` [3] ;

2. What's the address of array `a` and array `b`? (NOTE: please represent with expressions based on `%rbp`) (2')
3. What's the value of `row` and `col`? Please complete the initialization code in C program (2')
4. Please show the value of the underlined `%rax` in line 16 and 23 when `i=2` and explain how the value is calculated using `col`, `row` and `i`. (4')

Problem 5: (22points)

<pre>// ASCII(0~9):0x30~0x39 int aaa(char *str, int len){ int i = 0, result = len; for(i; i < len; i++){ switch(str[i]){ case '1': result = <u>[1]</u>; break; case <u>[2]</u>: result = str[i] >> 3; case <u>[3]</u>: result++; break; case '5': result = <u>[4]</u>; break; case '7': result = 9; default: result = <u>[5]</u>; } } return result; } int main(){ char *str = "54749110"; int pos = aaa(str, 8); printf("pos:%d\n", pos); long l = 0x123456789abcdef; unsigned char *bytel=(char *)&l; printf("0x%.2x%.2x\n", bytel[pos-1], bytel[pos]); return 0; }</pre>	<div>①</div>	<pre>.section .rodata .align 8 .L5: .quad .L4 .quad .L3 .quad .L6 .quad .L7 .quad .L8 .quad .L3 .quad .L9 .text <aaa>: pushq %rbp movq %rsp, %rbp movq %rdi, -24(%rbp) movl <u>[6]</u>, -28(%rbp) movl \$0, -4(%rbp) movl -28(%rbp), %eax movl %eax, -8(%rbp) jmp .L2 .L13: movl -4(%rbp), %eax movslq %eax, %rdx movq -24(%rbp), %rax addq %rdx, %rax movzbl <u>[7]</u>(%rax), %eax movsbl %al, %eax subl \$49, %eax cmpl \$6, %eax ja .L3 movq <u>[8]</u>, %rax jmp *%rax</pre> <div>②</div>
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<pre> .L4: cmp1 \$8, -8(%rbp) jg .L10 movl -8(%rbp), %eax jmp .L11 .L10: movl \$2, %eax .L11: movl %eax, -8(%rbp) jmp .L12 .L7: movl -4(%rbp), %eax movslq %eax, %rdx movq [9], %rax addq %rdx, %rax movzbl (%rax), %eax sarb \$3, %al movsbl %al, %eax movl %eax, -8(%rbp) .L6: addl \$1, -8(%rbp) jmp .L12 </pre>	③	<pre> .L8: movl -4(%rbp), %eax addl %eax, -8(%rbp) jmp .L12 .L9: movl \$9, -8(%rbp) [10] .L3: movl -8(%rbp), %eax addl %eax, %eax subl \$1, %eax movl %eax, -8(%rbp) .L12: addl \$1, -4(%rbp) .L2: movl -4(%rbp), %eax cmp1 -28(%rbp), %eax [11] movl [12], %eax popq %rbp ret </pre>	⑤
	④		⑥

Suppose the C and assembly code are executed on a 64-bit little-endian machine. Read the code and answer the following questions.

- Please fill in the blanks within C and assembly code. (1.5' * 12)
NOTE: no more than one instruction/statement per blank. If you think nothing is required to write, please write NONE.
- What is the output of the main function? (2')
- Suppose the machine is big-endian, what will be the output of the main function? (2')

Problem 6: (23points)

One of TAs of ICS wrote a stupid program. The following C code and assembly code are executed on a **64-bit little endian machine**.

<pre> void f(long a, long b, long c, long d, long e, long f, int g, int h) { h = g; printf("%p\n", ____ [16] ____); } </pre>	<pre> int main(void) { f([8], [9], [10], [11], [12], [13], [14], [15]); return 0; } </pre>
<pre> 400526 <f>: 400526: 55 pushq %rbp 400527: 48 89 e5 movq %rsp, %rbp 40052a: 8b 45 10 movl 0x10(%rbp), %eax 40052d: 89 45 14 movl %eax, 0x14(%rbp) // Assembly code for calling printf 40054d: c9 leaveq 40054e: c3 retq 40054f <main>: 40054f: 55 pushq %rbp 400550: 48 89 e5 movq %rsp, %rbp // Assembly code for preparing arguments 400577: e8 aa ff ff ff callq 400526 <f> 40057c: 48 83 c4 10 ____ [1] ____ 400580: b8 00 00 00 00 movl \$0, %eax 400585: c9 leaveq 400586: c3 retq </pre>	

1. Fill in the blank in the **assembly code**. (2').

40057c: 48 83 c4 10 [1]

2. There is a bug in assembly code of **f**. Please give the address of the instruction and fix it. You don't need to give the binary code of fixed instruction. (2'*2=4')

 [1] : [2]

3. Assume **BEFORE** the execution of instruction at 400577 (**callq <f>**), the memory and register states are as follows:

register	value
%rax	0x00000000 00000000
%rbx	0x00000000 0000000a
%rcx	0x00000000 00000007
%rdx	0x00000000 00000002
%rsi	0x00000000 00000008
%rdi	0x00000000 00000005
%rbp	0x00000000 1000df40
%rsp	0x00000000 1000df30
%r8	0x00000000 00000001
%r9	0x00000000 00000003
%r10	0x00000000 00000004

memory	value
0x1000df40	0x00000000 00400590
0x1000df38	0x00000000 00000004
0x1000df30	0x00000000 00000006

Please fill in the following table that show the state **AFTER** the execution of instruction at **40052a (movl 0x10(%rbp), %eax)**: (1'*7=7')

register	value
%rax	[1]
%rbp	[2]
%rsp	[3]

memory	value
0x1000df38	[4]
0x1000df30	[5]
0x1000df28	[6]
0x1000df20	[7]

Please fill in the C code ([8]~[15]) of **main** function: (1'*8=8')

4. Please fill in the blank in C code ([16]) of **f** function, so that the **printf** will always **output the return address of f**. (2')