

【1】Uboot 启动流程：

分析版本：uboot-2013

入口查看 u-boot.lds

通过链接脚本可知入口为 `_start`，位于

```
OUTPUT_FORMAT("elf32-littlearm", "elf32-littlearm", "elf32-littlearm")
OUTPUT_ARCH(arm)
ENTRY(_start)
SECTIONS
{
    . = 0x00000000;
    . = ALIGN(4);
    .text :
    {
        __image_copy_start = .;
        arch/arm/cpu/armv7/start.o (.text*)
        *(.text*)
    }
}
```

`arch/arm/cpu/armv7/start.o`。

第一阶段开始：

进入 `arch/arm/cpu/armv7/start.S`

```
.globl _start
_start: b reset
ldr pc, _undefined_instruction
ldr pc, _software_interrupt
ldr pc, _prefetch_abort
ldr pc, _data_abort
ldr pc, _not_used
ldr pc, _irq
ldr pc, _fiq
```

异常向量表设置

```
reset:
bl save_boot_params
/*
 * set the cpu to SVC32 mode
 */
mrs r0, cpsr
bic r0, r0, #0x1f
orr r0, r0, #0xd3
msr cpsr, r0
```

设置cpu处于svc工作模式

```
/* Set vector address in CP15 VBAR register */
ldr r0, =_start
mcr p15, 0, r0, c12, c0, 0 @Set VBAR
#endif
```

协处理器p15的c12寄存器来重新定位

`Bl cpu_init_cp15`（使分支预测无效，数据）

```
ENTRY(cpu_init_cp15)
```

```
/*
```

```
 * Invalidate L1 I/D
```

```
*/
```

```
mov r0, #0 @ set up for MCR
```

```
mcr p15, 0, r0, c8, c7, 0 @ invalidate TLBs
```

```
mcr p15, 0, r0, c7, c5, 0 @ invalidate icache
```

```
mcr p15, 0, r0, c7, c5, 6 @ invalidate BP array
```

```
mcr p15, 0, r0, c7, c10, 4 @ DSB
```

```
mcr p15, 0, r0, c7, c5, 4 @ ISB
```

关闭数据预取功能

DSB:多核cpu对数据处理指令

流水线清空指令

```
*/
```

```
6 mcr p15, 0, r0, c1, c0, 0
```

```
7 bic r0, r0, #0x00002000 @ clear bits 13 (--V-)
```

```
8 bic r0, r0, #0x00000007 @ clear bits 2:0 (-CAM)
```

```
9 orr r0, r0, #0x00000002 @ set bit 1 (--A-) Align
```

```
0 orr r0, r0, #0x00000800 @ set bit 11 (Z---) BTB
```

```
1 #ifdef CONFIG_SYS_ICACHE_OFF
```

```
2 bic r0, r0, #0x00001000 @ clear bit 12 (I) I-cache
```

```
3 #else
```

```
4 orr r0, r0, #0x00001000 @ set bit 12 (I) I-cache
```

```
5 #endif
```

```
6 mcr p15, 0, r0, c1, c0, 0
```

```
7 mov pc, lr @ back to my caller
```

```
8 ENDPROC(cpu_init_cp15)
```

关闭MMU，使能I-cache

#### NOTE:

分支预测：在流水线里，会将后面的代码优先加载到处理器中，由于是循环，会使后面加载的代码无效，故出现了分支预测技术。(统计跳的次数来选择装载循环的代码还是下面的代码)。

#### Bl cpu\_init\_crit

```
ENTRY(cpu_init_crit)
```

```
/*
```

```
 * Jump to board specific initialization...
```

```
 * The Mask ROM will have already initialized
```

```
 * basic memory. Go here to bump up clock rate and handle
```

```
 * wake up conditions.
```

```
*/
```

```
b lowlevel_init @ go setup pll,mux,memory
```

```
ENDPROC(cpu_init_crit)
```

跳到 Low\_level\_init,位于 board/samsung/fs4412/lowlevel\_init.S

```
ldr r0, =0x1002040c
```

```
ldr r1, =0x00
```

```
str r1, [r0]
```

关闭看门狗

```
ldr r0, =0xffffffff /* r0 <- Mask Bits*/  
bic r1, pc, r0 /* pc <- current addr of code */  
/* r1 <- unmasked bits of pc */  
ldr r2, _TEXT_BASE /* r2 <- original base addr in ram */  
bic r2, r2, r0 /* r2 <- unmasked bits of r2*/  
cmp r1, r2 /* compare r1, r2 */  
beq 1f /* r0 == r1 then skip sdram init */
```

```
/* init system clock */  
bl system_clock_init
```

```
/* Memory initialize */  
bl mem_ctrl_asm_init
```

```
1:
```

```
/* for UART */
```

```
bl uart_asm_init
```

```
/*bl tzpc_init*/
```

```
pop {pc}
```

返回start S

比较当前pc指针与TEXT\_BASE的高8位是否一样来判断，当前代码是否在内存中

对系统时钟初始化

对内存初始化

对串口初始化

第一阶段结束，总结如下：

- 1 前面总结过的部分，初始化异常向量表，设置 svc 模式
- 2 配置 cp15，初始化 mmu cache tlb
- 3 板级初始化，clk,memory,uart 初始化

第二阶段开始：

Bl \_main, 跳转到 arch/arm/lib/crt0.S

初始 c 运行环境

```
1 #if defined(CONFIG_NAND_SPL)
2 /* deprecated, use instead CONFIG_SPL_BUILD */
3 ldr sp, =(CONFIG_SYS_INIT_SP_ADDR)
4 #elif defined(CONFIG_SPL_BUILD) && defined(CONFIG_SPL_STACK)
5 ldr sp, =(CONFIG_SPL_STACK)
6 #else
7 ldr sp, =(CONFIG_SYS_INIT_SP_ADDR)
8 #endif
9 bic sp, sp, #7 /* 8-byte alignment for ABI compliance */
10 sub sp, #GD_SIZE /* allocate one GD above SP */
11 bic sp, sp, #7 /* 8-byte alignment for ABI compliance */
12 mov r8, sp /* GD is above SP */
13 mov r0, #0
14 bl board_init_f
```

初始化sp,为分支c语言做准备

8位对齐

保留128B放GD结构体，存放全局信息，GD的地址存放在r8中

跳转到 arch/arm/lib/board.c

```
void board_init_f(ulong bootflag)
{
    bd_t *bd;
    init_fnc_t **init_fnc_ptr;
    gd_t *gd;
    ulong addr, addr_sp;
#ifdef CONFIG_PRAM
    ulong reg;
#endif
    void *new_fdt = NULL;
    size_t fdt_size = 0;

    memset((void *)gd, 0, sizeof(gd_t));

    for (init_fnc_ptr = init_sequence; *init_fnc_ptr; ++init_fnc_ptr) {
        if ((*init_fnc_ptr)() != 0) {
            hang();
        }
    }

    gd->fdt_blob = (void *)(_end_ofs + _TEXT_BASE);
    /* Allow the early environment to override the fdt address */
    gd->fdt_blob = (void *)getenv_ulong("fdtcontroladdr", 16,
        (uintptr_t)gd->fdt_blob);
```

对全局信息GD结构体进行填充，mon\_len通过链接脚本可以知道放的是uboot代码大小

循环执行init\_fnc\_t数组中的函数，作硬件初始化

结构如下：

```

init_fnc_t *init_sequence[] = {
    arch_cpu_init,      /* basic arch cpu dependent setup */
    mark_bootstage,
#ifdef CONFIG_OF_CONTROL
    fdtdec_check_fdt,
#endif
    #if defined(CONFIG_BOARD_EARLY_INIT_F)
        board_early_init_f,
    #endif
    timer_init,        /* initialize timer */
#ifdef CONFIG_BOARD_POSTCLK_INIT
    board_postclk_init,
#endif
#ifdef CONFIG_FSL_ESDHC
    get_clocks,
#endif
    env_init,          /* initialize environment */
    init_baudrate,      /* initialize baudrate settings */
    serial_init,        /* serial communications setup */
    console_init_f,     /* stage 1 init of console */
    display_banner,     /* say that we are here */
    #if defined(CONFIG_DISPLAY_CPUINFO)
        print_cpuinfo,  /* display cpu info (and speed) */
    #endif
    #if defined(CONFIG_DISPLAY_BOARDINFO)
        checkboard,     /* display board info */
    #endif
    #if defined(CONFIG_HARD_I2C) || defined(CONFIG_SOFT_I2C)
        init_func_i2c,
    #endif
    dram_init,          /* configure available RAM banks */
    NULL,
};

```

初始化各硬件

Dram\_init 初始化成功之后，剩余代码将会对 sdram 空间进行规划。

```
addr = CONFIG_SYS_SDRAM_BASE + gd->ram_size;
```

addr 的值由 CONFIG\_SYS\_SDRAM\_BASE 加上 ram\_size。也就是到了可用 sdram 的顶端。

```

#if !(defined(CONFIG_SYS_ICACHE_OFF) && defined(CONFIG_SYS_DCACHE_OFF))
/* reserve TLB table */
gd->tlb_size = 4096 * 4;
addr -= gd->tlb_size;

/* round down to next 64 kB limit */
addr &= ~(0x10000 - 1);

gd->tlb_addr = addr;
debug("TLB table from %08lx to %08lx\n", addr, addr + gd->tlb_size);
#endif

/* round down to next 4 kB limit */
addr &= ~(4096 - 1);
debug("Top of RAM usable for U-Boot at: %08lx\n", addr);

```

如果icache与dcache是打开的，就留出64K的空间作为tlb空间，最后addr就是tlb的地址，4K对齐

继续对 gd 结构体填充

```
memcpy(id, (void *)gd, sizeof(gd_t));
```

填充完成将信息拷贝到内存指定位置。

继续回到 main

```
ldr sp, [r8, #GD_START_ADDR_SP] /* r8 = gd->start_addr_sp */
bic sp, sp, #7 /* 8-byte alignment for ABI compliance */
ldr r8, [r8, #GD_BD] /* r8 = gd->bd */
sub r8, r8, #GD_SIZE /* new GD is below bd */
```

将r8指向新的gd地址

```
adr lr, here
ldr r0, [r8, #GD_RELOC_OFF] /* lr = gd->start_addr_sp */
add lr, lr, r0
ldr r0, [r8, #GD_START_ADDR_SP] /* r0 = gd->start_addr_sp */
mov r1, r8 /* r1 = gd */
ldr r2, [r8, #GD_RELOCADDR] /* r2 = gd->relocaddr */
b relocate_code
pre:
```

代码重定位

对 lr 的操作为了让返回时，返回的是重定位的 here 处。

代码自搬移，防止与内核冲突，代码位于 `arch/arm/cpu/armv7/start.S`

```
ENTRY(relocate_code)
    mov r4, r0 /* save addr_sp */
    mov r5, r1 /* save addr of gd */
    mov r6, r2 /* save addr of destination */

    adr r0, _start /* ldr r0, =_start*/
    cmp r0, r6
    moveq r9, #0 /* no relocation. relocation offset(r9) = 0 */
    beq relocate_done /* skip relocation */
    mov r1, r6 /* r1 <- scratch for copy_loop */
    ldr r3, _image_copy_end_ofs
    add r2, r0, r3 /* r2 <- source end address */

    copy_loop:
        ldmia r0!, {r9-r10} /* copy from source address [r0] */
        stmia r1!, {r9-r10} /* copy to target address [r1] */
        cmp r0, r2 /* until source end address [r2] */
        blo copy_loop
```

循环将代码搬移到指定高地址

这里只是将链接脚本中 `_image_copy_end` 到 `_start` 中的代码，其它段还没有操作。

在这里我们有疑惑就是将代码重定位到高地址，那运行的地址不就和链接地址不一样了，那运行可能不正常？这个疑惑就是 `.rel.dyn` 帮我们解决了，主要还是编译器帮我们做的工作，在链接中有如下：【参考：<http://blog.csdn.net/skyflying2012/article/details/37660265>】

```
ifndef CONFIG_NAND_SPL
LDFLAGS_u-boot += -pie
endif
```

重

定位到高地址之后，再次回到 main

here:

```
/* Set up final (full) environment */ 关icache,保证数据从sdram中更新,更新异常向量表,因为代码被重定位了
bl c_runtime_cpu_setup /* we still call old routine here */

ldr r0, =__bss_start /* this is auto-relocated! */
ldr r1, =__bss_end__ /* this is auto-relocated! */

mov r2, #0x00000000 /* prepare zero to clear BSS */ ← 清BSs

clbss_l: cmp r0, r1 /* while not at end of BSS */
strlo r2, [r0] /* clear 32-bit BSS word */
addlo r0, r0, #4 /* move to next */
blo clbss_l
```

```
/* call board_init_r(gd_t *id, ulong dest_addr) */
mov r0, r8 /* gd_t */
ldr r1, [r8, #GD_RELOCADDR] /* dest_addr */
/* call board_init_r */
ldr pc, =board_init_r /* this is auto-relocated! */
```

R0=gd R1=RELOCADDR  
调用board\_init\_r

调用 board\_init\_r 主要是对外设的初始化。

```
for (;;) {
    main_loop();
}
```

进入超循环

Main\_loop 函数主要功能是处理环境变量，解析命令

install\_auto\_complete(); //安装自动补全的函数，分析如下。

getenv(bootcmd)

bootdelay(自启动)

```
if (bootdelay != -1 && s && !abortboot(bootdelay)) {
```

如果延时大于等于零，并且没有在延时过程中接收到按键，则引导内核。abortboot 函数的分析见下面

```
rc = run_command(lastcommand, flag); 处理用户输入命令
```

uboot 启动流程分析如下：

第一阶段：

设置 cpu 工作模式为 SVC 模式

关闭中断，mmu,cache

关看门狗

初始化内存，串口

设置栈

代码自搬移

清 bss

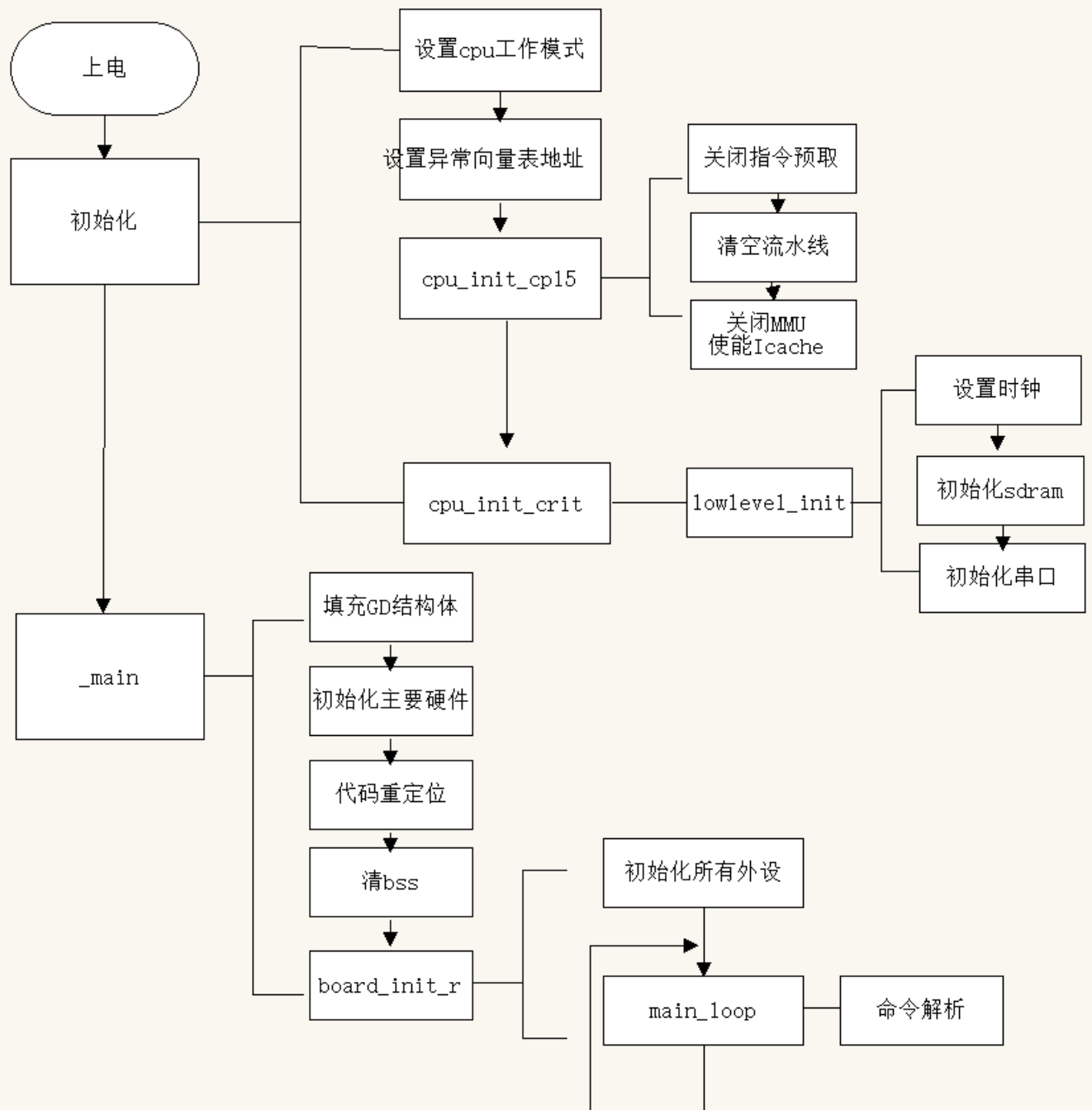
跳 c

## 第二阶段

初始化外设，进入超循环

## 超循环处理用户命令

针对 uboot2013 启动流程图如下:



下面是最精简的自己实现 uboot 的代码形式【代码见附件】

```

.
├── arch
│   └── arm
│       └── arm_cortexa8
│           ├── map.lds
│           ├── start.o
│           └── start.S
├── board
│   └── samsung
│       ├── low_levelinit.o
│       ├── low_levelinit.S
│       ├── mem_setup.o
│       ├── mem_setup.S
│       ├── nand.c
│       ├── nand.o
│       └── test.c
├── common
│   ├── do_go.c
│   ├── do_go.o
│   ├── main.c
│   └── main.o
├── driver
│   ├── string.c
│   ├── string.o
│   ├── uart.c
│   ├── uart.h
│   └── uart.o
├── include
│   ├── s5pc100.h
│   └── setup.h
├── lib_arm
│   ├── board.c
│   └── board.o
├── Makefile
├── mini_uboot.bin
├── mini_uboot.dis
└── mini_uboot.elf

```

针对a8开发板自己编写uboot的代码结构

【2】命令解析过程 run\_command:

```

int run_command(const char *cmd, int flag)
{
#ifdef CONFIG_SYS_HUSH_PARSER
    /*
     * builtin_run_command can return 0 or 1 for success, so clean up
     * its result.
     */
    if (builtin_run_command(cmd, flag) == -1)
        return 1;

    return 0;
#else
    return parse_string_outer(cmd,
        FLAG_PARSE_SEMICOLON | FLAG_EXIT_FROM_LOOP);
#endif
}

```



进入 builtin\_run\_command,分析 cmd\_process

```
if (cmd_process(flag, argc, argv, &repeatable))
```

```
cmdtp = find_cmd(argv[0]);
if (cmdtp == NULL) {
    printf("Unknown command '%s' - try 'help'\n", argv[0]); 查找命令
    return 1;
}
```

```
/* If OK so far, then do the command */
if (!rc) {
    rc = cmd_call(cmdtp, flag, argc, argv); 真正执行命令
    *repeatable &= cmdtp->repeatable;
}
```

【3】Uboot 添加自定义命令：uboot 中的命令使用 U\_BOOT\_CMD 这个宏声明来注册进系统，链接脚本会把所有的 cmd\_tbl\_t 结构体放在相邻的地方。

```
U_BOOT_CMD(
    help,    CONFIG_SYS_MAXARGS, 1, do_help, 这个宏将信息添加到
    "print command description/usage",      .u_boot_list.cmd段
    "\n"
    " - print brief description of all commands\n"
    "help command ... \n"
    " - print detailed usage of 'command'"
);
```

uboot 命令结构体如下：

```
struct cmd_tbl_s {
    char *name; /* Command Name */
    int maxargs; /* maximum number of arguments */
    int repeatable; /* autorepeat allowed? */
    /* Implementation function */
    int (*cmd)(struct cmd_tbl_s *, int, int, char * const []);
    char *usage; /* Usage message (short) */
#ifdef CONFIG_SYS_LONGHELP
    char *help; /* Help message (long) */
#endif
#ifdef CONFIG_AUTO_COMPLETE
    /* do auto completion on the arguments */
    int (*complete)(int argc, char * const argv[], char last_char, int maxv, char *cmdv[]);
#endif
};
```

通过链接脚本查看如下段

```
_u_boot_list_cmd_start = .;
*(SORT(.u_boot_list.cmd.*));
_u_boot_list_cmd_end = .;
```

```

#define U_BOOT_CMD_MKENT_COMPLETE(_name, _maxargs, _rep, _cmd, \
    _usage, _help, _comp) \
    { #_name, _maxargs, _rep, _cmd, _usage, \
      _CMD_HELP(_help) _CMD_COMPLETE(_comp) }

#define U_BOOT_CMD_MKENT(_name, _maxargs, _rep, _cmd, _usage, _help) U_BOOT_CMD_MKENT_COMPLETE(_name, _maxargs, _rep, \
#define U_BOOT_CMD_MKENT_COMPLETE(_name, _maxargs, _rep, _cmd, _usage, _help, _comp) \
    ll_entry_declare(cmd_tbl_t, _name, cmd, cmd) = \
        U_BOOT_CMD_MKENT_COMPLETE(_name, _maxargs, _rep, _cmd, \
            _usage, _help, _comp);

#define U_BOOT_CMD(_name, _maxargs, _rep, _cmd, _usage, _help) \
    U_BOOT_CMD_MKENT_COMPLETE(_name, _maxargs, _rep, _cmd, _usage, _help, NULL)

```

```

#define ll_entry_declare(_type, _name, _section_u, _section_d) \
    _type _u_boot_list_##_section_u##_##_name __attribute__(( \
        unused, aligned(4), \
        section(".u_boot_list."#_section_d"."#_name)))

```

指定相关代码指定段

【4】 bootm 分析：

```

U_BOOT_CMD(
    bootm, CONFIG_SYS_MAXARGS, 1, do_bootm,
    "boot application image from memory", bootm_help_text
)

```

位于函数 do\_bootm

```

if (bootm_start(cmdtp, flag, argc, argv))
    return 1;

```

执行bootm

位于函数 bootm\_start

```

/* get kernel image header, start address and length */
os_hdr = boot_get_kernel(cmdtp, flag, argc, argv,
    &images, &images.os.image_start, &images.os.image_len);
if (images.os.image_len == 0) {
    puts("ERROR: can't get kernel image!\n");
    return 1;
}

/* get image parameters */
switch (genimg_get_format(os_hdr)) {
case IMAGE_FORMAT_LEGACY:
    images.os.type = image_get_type(os_hdr);
    images.os.comp = image_get_comp(os_hdr);
    images.os.os = image_get_os(os_hdr);
}

```

填充内核相关信息

回到 do\_bootm

```

static boot_os_fn *boot_os[] = {
#ifdef CONFIG_BOOTM_LINUX
    [IH_OS_LINUX] = do_bootm_linux,
#endif
#ifdef CONFIG_BOOTM_NETBSD
    [IH_OS_NETBSD] = do_bootm_netbsd,
#endif
#ifdef CONFIG_LYNXKDI
    [IH_OS_LYNXOS] = do_bootm_lynxkdi,
#endif
#ifdef CONFIG_BOOTM RTEMS
    [IH_OS RTEMS] = do_bootm_rtems,
#endif
#ifdef CONFIG_BOOTM_OSE
    [IH_OS_OSE] = do_bootm_ose,
#endif
};

```

boot\_os 函数指针数组，从中可以知道，上面的执行是do\_bootm\_linux

位于 do\_bootm\_linux

```

boot_jump_linux(images);

```

位于 boot\_jump\_linux

```

/* Subcommand: GO */
static void boot_jump_linux(bootm_headers_t *images)
{
    unsigned long machid = gd->bd->bi_arch_number;
    char *s;
    void (*kernel_entry)(int zero, int arch, uint params);
    unsigned long r2;

    kernel_entry = (void (*)(int, int, uint))images->ep;

    s = getenv("machid");
    if (s) {
        strict_strtoul(s, 16, &machid);
        printf("Using machid 0x%lx from environment\n", machid);
    }

    debug("## Transferring control to Linux (at address %08lx)" \
        "... \n", (ulong) kernel_entry);
}

#ifdef CONFIG_BOOTM_OSE
bootstage_mark(BOOTSTAGE_ID_RUN_OS);
announce_and_cleanup();
#endif

```

【5】u-boot 编译流程分析说下：（也可以从 build.sh 分析）  
Make fs4412\_config

```
%_config:: unconfig
@$(MKCONFIG) -A $(@:_config=)
```

替换目标名fs4412\_config 为 fs

```
MKCONFIG := $(SRCTREE)/mkconfig
```

分析 mkconfig 脚本

解析 boards.cfg fs4412 相关数据

```
fs4412 arm armv7 fs4412 samsung exynos
```

作链接

```
ln -s ../arch/${arch}/include/asm asm
ln -s ${LNPREFIX}arch-${cpu} asm/arch
ln -s ${LNPREFIX}proc-armv asm/proc
```

针对平台作了一系列链接

导出信息到 config.mk

```
( echo "ARCH = ${arch}"
  if [ ! -z "$spl_cpu" ] ; then
    echo 'ifeq ($(CONFIG_SPL_BUILD),y)'
    echo "CPU = ${spl_cpu}"
    echo "else"
    echo "CPU = ${cpu}"
    echo "endif"
  else
    echo "CPU = ${cpu}"
  fi
  echo "BOARD = ${board}"

  [ "${vendor}" ] && echo "VENDOR = ${vendor}"
  [ "${soc}" ] && echo "SOC = ${soc}"
  exit 0 ) > config.mk
```

Include/config.mk 内容如下:

```
1 ARCH = arm
2 CPU = armv7
3 BOARD = fs4412
4 VENDOR = samsung
5 SOC = exynos
```

导出平台数据到 config.h

```
echo "#define CONFIG_SYS_ARCH \"${arch}\"" >> config.h
echo "#define CONFIG_SYS_CPU \"${cpu}\"" >> config.h
echo "#define CONFIG_SYS_BOARD \"${board}\"" >> config.h

[ "${vendor}" ] && echo "#define CONFIG_SYS_VENDOR \"${vendor}\"" >> config.h
```

Include/config.h 导出结果如下：

```
/* Automatically generated - do not edit */
#define CONFIG_SYS_ARCH "arm"
#define CONFIG_SYS_CPU "armv7"
#define CONFIG_SYS_BOARD "fs4412"
#define CONFIG_SYS_VENDOR "samsung"
#define CONFIG_SYS_SOC "exynos"
#define CONFIG_BOARDDIR board/samsung/fs4412
#include <config_cmd_defaults.h>
#include <config_defaults.h>
#include <configs/fs4412.h>
#include <asm/config.h>
```

Make -jxx 执行过程如下：

找第一个目标 all:

```
all: $(ALL-y) $(SUBDIR_EXAMPLES)

$(ALL-y) += $(obj)u-boot.srec $(obj)u-boot.bin $(obj)System.map
```

```
$(obj)u-boot.bin: $(obj)u-boot
$(OBJCOPY) ${OBJCFLAGS} -O binary $< $@
$(BOARD_SIZE_CHECK)
@+make -C sdfuse_q/
@./sdfuse_q/add_padding
```

上面代码是对 u-boot 进行格式转换，变成二进制 bin 格式之后，再加一些校验与 4412 开如平台加密信息。

依赖 u-boot: (573 行左右)

```
$(obj)u-boot: depend \
$(SUBDIR_TOOLS) $(OBJS) $(LIBBOARD) $(LIBS) $(LDSCRIPT) $(obj)u-boot.lds
$(GEN_UBOOT)
ifeq ($(CONFIG_KALLSYMS),y)
smap=`$(call SYSTEM_MAP,$(obj)u-boot) | \
awk '$$2 ~ /[tTwW]/ {printf $$1 $$3 "\\000"}'` ; \ 编译相关文件
$(CC) $(CFLAGS) -DSYSTEM_MAP="\"${smap}\"" \
-c common/system_map.c -o $(obj)common/system_map.o
```

## 【6】扩展题目练习

给学生留一个课后扩展练习题，如何通过现有 u-boot 将代码加载并烧录到外存，如何调整当前 emmc 分区大小，应该看哪块代码？