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star height

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The *star height* of a regular expression  $p$  over an alphabet  $\Sigma$ , denoted by  $\text{ht}(p)$ , is defined recursively as follows:

1.  $\text{ht}(\emptyset) = \text{ht}(a) = 0$  for any  $a \in \Sigma$ ;
2.  $\text{ht}(p \cup q) = \text{ht}(pq) = \max(\text{ht}(p), \text{ht}(q))$ ;
3.  $\text{ht}(p^*) = \text{ht}(p) + 1$ .

Let  $\Sigma = \{a, b, c\}$ . The following expressions have star height 0, 1, 2, 3:

$$(a \cup b)c \quad a^*b^* \quad (a^*b)^*c \quad ((ab^*a \cup c)^*b)^*$$

Since any regular expression  $p$  describes a regular language  $L(p)$ , we may extend the definition of star height to regular languages. However, since a regular language may be described by more than one regular expressions, we need to be a little careful:

The *star height* of a regular language  $L$  is the least integer  $i$  such that  $L$  may be described by a regular expression of star height  $i$ . In other words:

$$\text{ht}(L) = \min\{h(p) \mid L = L(p), p \in R(\Sigma)\},$$

where  $R(\Sigma)$  is the set of all regular expressions over  $\Sigma$ .

**Remarks.**

- Any regular language over a singleton alphabet has star height at most 1, which can be proved by the pumping lemma for regular languages.
- If the alphabet  $\Sigma$  consists of at least two letters, then for every positive integer  $n$ , there is a regular language whose star height is  $n$ . In fact, it can be shown that if  $|\Sigma| = 2$ , then for every  $n$ , there is a regular language  $L$  over  $\Sigma$  such that  $\text{ht}(L) = n$ .
- It was an open question whether an algorithm exists for determining  $\text{ht}(L)$  for an arbitrary regular language  $L$ . In 2005, the question was resolved by D. Kirsten in the positive, and the algorithm is that of a non-deterministic finite automaton.
- We may also define star height on generalized regular expressions. For a regular language  $L$ , define  $\text{ht}_g(L) = \min\{h(p) \mid L = L(p), p \in R_g(\Sigma)\}$ , where  $R_g(\Sigma)$  is the set of all generalized regular expressions. It is clear

that  $\text{ht}_g(L) \leq \text{ht}(L)$ . However, it is still an open question whether, for every integer  $n$ , there is a regular language  $L$  with  $\text{ht}_g(L) = n$ .

A star-free language has star height 0 with respect to representations by generalized regular expressions, but may be positive with respect to representations by regular expressions, for example  $L = \{ab\}^*$ .

## References

- [1] A. Salomaa, *Formal Languages*, Academic Press, New York (1973).
- [2] A. Salomaa, *Jewels of Formal Language Theory*, Computer Science Press (1981).