

# 问求算法设计与实践08

哈希

# Hash Functions

- 将对象从原空间映射到一个'小空间'中(维度低, 值域范围小)
- 计算快
- 冲突少 (函数设计, 散列表设计)

# Hash Functions

- 哈希对象
  - 数字
  - 字符串
  - 树, 图
  - 文件
  - ...

# Hash Functions

- $F(x) = x \bmod P$ ;
- $F(x) = x^{10} \bmod P$ ;
- ...

# 问题来了..

- 维护一个数据集
- 每次可能插入一个新的数字
- 每次可能查询数字 $x$ 是否在集合内

# 排序+二分

- 插入 $O(N\log N)$
- 查询 $O(\log N)$

# 平衡二叉树

- 插入 $O(\log N)$
- 查询 $O(\log N)$

# Hash Table!

- 插入 $O(1)$
- 查询 $O(1)$



# 字符串哈希

## Sample Input

```
[expelliarmus] the disarming charm
[rictusempra] send a jet of silver light to hit the enemy
[tarantallegra] control the movement of one's legs
[serpensortia] shoot a snake out of the end of one's wand
[lumos] light the wand
[obliviate] the memory charm
[expecto patronum] send a Patronus to the dementors
[accio] the summoning charm
@END@
4
[lumos]
the summoning charm
[arha]
take me to the sky
```

## Sample Output

```
light the wand
accio
what?
what?
```

# 字符串哈希算法

- BKDRHash, APHash, DJBHash, JSHash, RSHash, SDBMHash, PJWHash, ELFHash...
- <https://www.byvoid.com/blog/string-hash-compare/>

# BKDRHash

- 对于字符串S='acb'
- $H(S) = 'a' * 31^2 + 'c' * 31 + 'b'$

# BKDRHash

- 递推计算
- for  $i=0:\text{len}-1$ 
  - $H = H * 31 + S[i];$

# BKDRHash

- 如果有两个串的hash值相同，我们'几乎'可以说他们就是同一个串！
- 实际做题中，合适的选取seed，基本可以全对
- 实际应用中，可能还是会有一些反例

# Rabin-Karp Algorithm

- 给定字符串A, B
- 问B是否是A的子串
- 暴力  $O(\text{length}(A) * \text{length}(B))$
- KMP  $O(\text{length}(A) + \text{length}(B))$
- RK 平均 $O(\text{length}(A) + \text{length}(B))$

# 预处理

- 给定一个字符串A
- 给定多个查询  $(B, i, j)$ , 表示B是否与 $A_i \dots A_j$ 为相同的字符串
- 怎么在 $O(1)$ 时间内求出 $H(A_i \dots A_j)$ ?

# 前綴和trick

- $\text{Pre}[i] = H(A_1 \dots A_i)$
- $\text{Pre}[i] = \text{Pre}[i-1] * 31 + A_i$
- $H(A_i \dots A_j) = ?$



# 前綴和trick

- $\text{Pre}[i] = H(A_1 \dots A_i)$
- $\text{Pre}[i] = \text{Pre}[i-1] * 31 + A_i$
- $H(A_i \dots A_j) = \text{Pre}[j] - \text{Pre}[i-1] * 31^{(j-i+1)}$

# 溢出int?

- 用unsigned longlong
- 天然的模系统
- 加减乘法都满足模的性质

# Problem B

- 给定一个字符串
- 找出至少出现M次的最长子串

# Problem B

- 用前缀和预处理好哈希值
- 二分长度len
- 找出所有长度为len的哈希值 ( $O(N)$ 个)
- 将哈希值排序，看是否有至少M个相同的哈希值(或者用一个哈希表来判断)
- 复杂度 $O(N\log\log N)$ 或 $O(N\log N)$

# BNUOJ 34490

## Justice String

- 给定两个串A, B
- 问A中是否有一个子串, 与B长度相等, 且最多只有2个字母不一样

# BNUOJ 34490

## Justice String

- 完全一样
  - KMP, RK
- 一个字母不一样
  - 在A中枚举起点
  - 二分一个最远的距离使得匹配
  - 看后半段是否匹配

# BNUOJ 34490

## Justice String

- 两个个字母不一样
  - 在A中枚举起点  $O(N)$ 
    - 二分一个最远的距离使得匹配  $O(\log N)$
    - 以匹配点后一个点为起点
    - 二分一个最远的距离使得匹配  $O(\log N)$
    - 看最后一段是否匹配  $O(1)$
- $O(N \log N)$

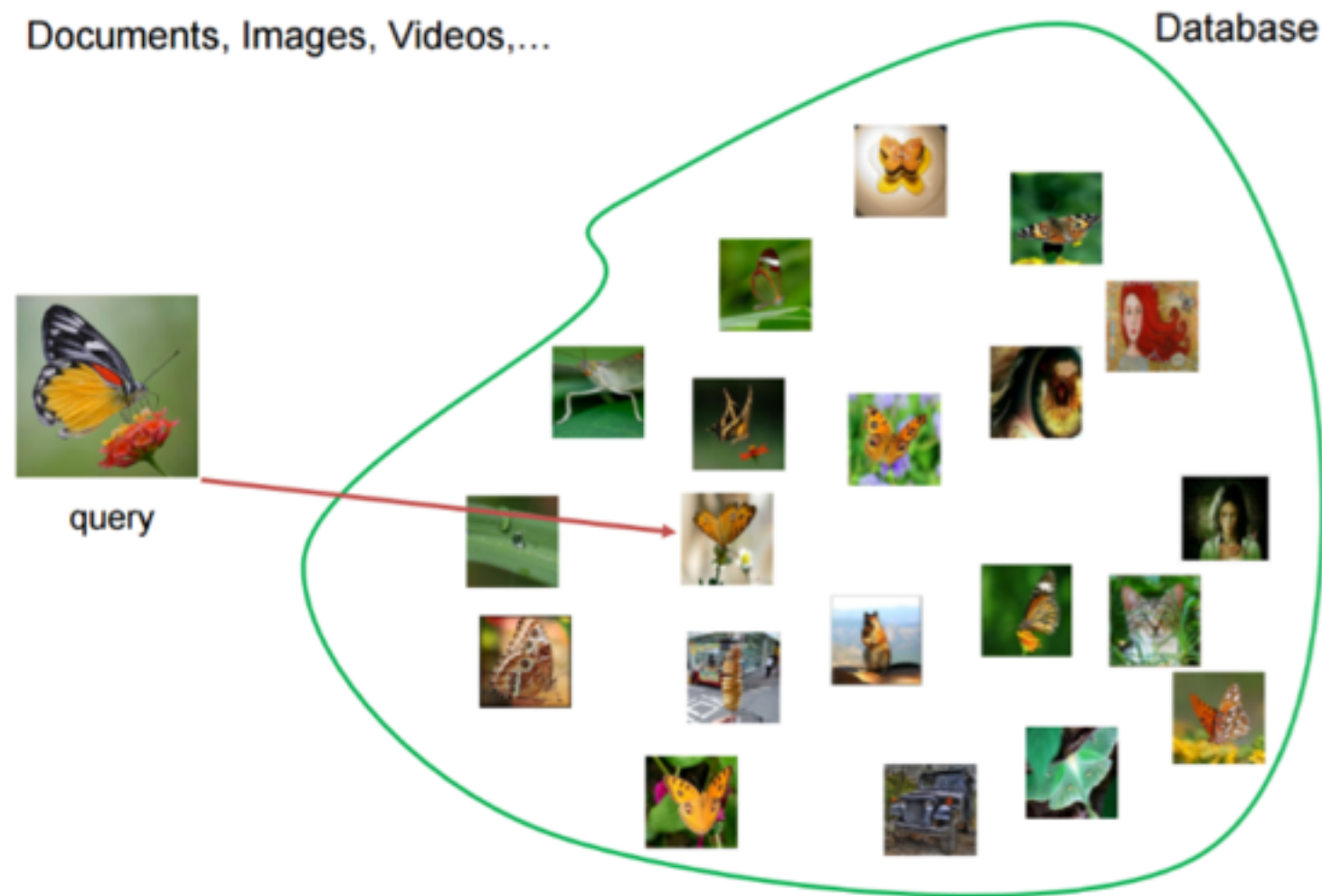
# Similarity-Preserving Hashing

- 当我们不仅仅考虑如何避免冲突
- 由于哈希函数一般是映射到一个海明空间
- 我们希望相似的对象，拥有‘相似’的哈希值(海明距离更近)



# Nearest Neighbor (NN) Search

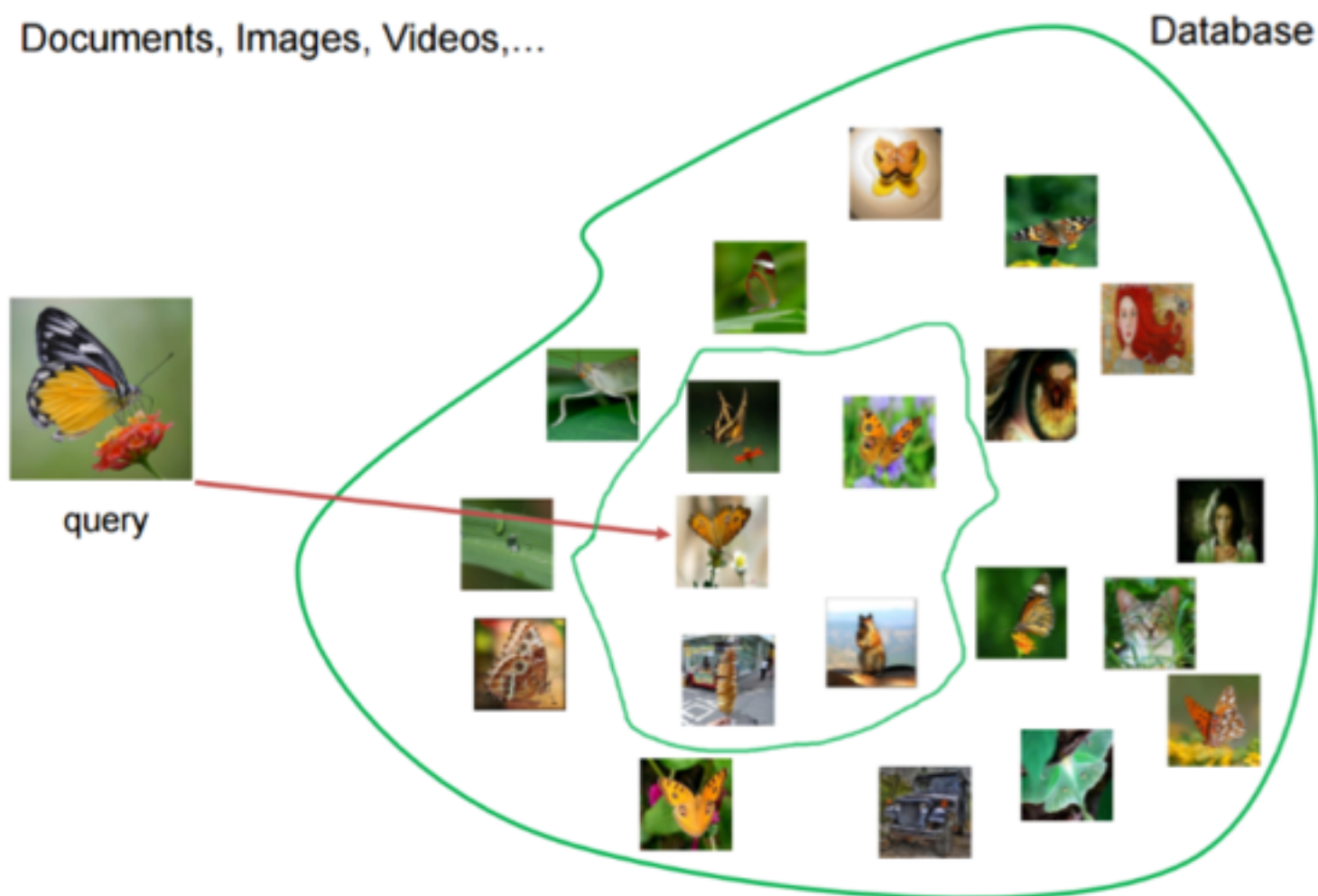
Given a database  $X = \{x_1, x_2, \dots\}$   $x_i \in \mathbb{R}^D$  and a query  $q \in \mathbb{R}^D$ .  
Return a point  $y$  which  $y = \arg \min_{y \in X} \text{dist}(y, q)$ .



There are lots of methods proposed to fast NN search (i.e K-D tree), but they fail in the **large scale or high-dimensional** cases.

# Approximate Nearest Neighbor (ANN) Search

A point  $z$  is the **approximate nearest neighbor** to  $q$ , if  $dist(z, q) \leq (1 + \epsilon)dist(y, q)$



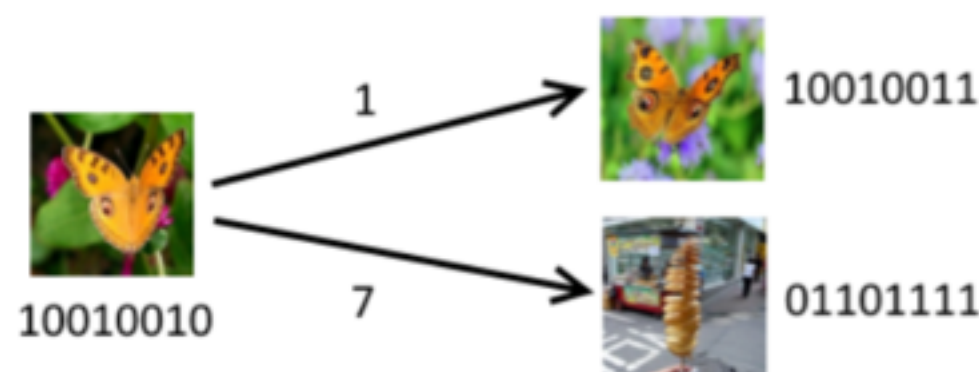
The most famous method of ANN is Locality Sensitive Hashing (LSH), which is based on hashing.

# Hashing Technique

Hash function map a object to a limited space

$$\text{i.e } f : \mathbb{R} \rightarrow \{0, 1, \dots, 255\}$$

Unlike traditional hashing methods which try to reduce colliding, hashing methods of ANN try to **preserve distance** in origin space.



Use low-dimensional hamming distance to approximate euclidian distance in origin space.

pros

- constant or sub-linear retrieve time
- low storage cost

cons

- hard to optimize a discrete problem

# Supervised Version

For real complex application images, videos...

- Object may have semantic information
- Two images are semantically similar cannot fully depend on the distance in feature (GIST, SIFT...) space
- There are tags or other supervised information could be utilized

In the setting of Supervised version, we are given the ground-truth of similarity between any two training samples.

# Locality-Sensitive Hashing

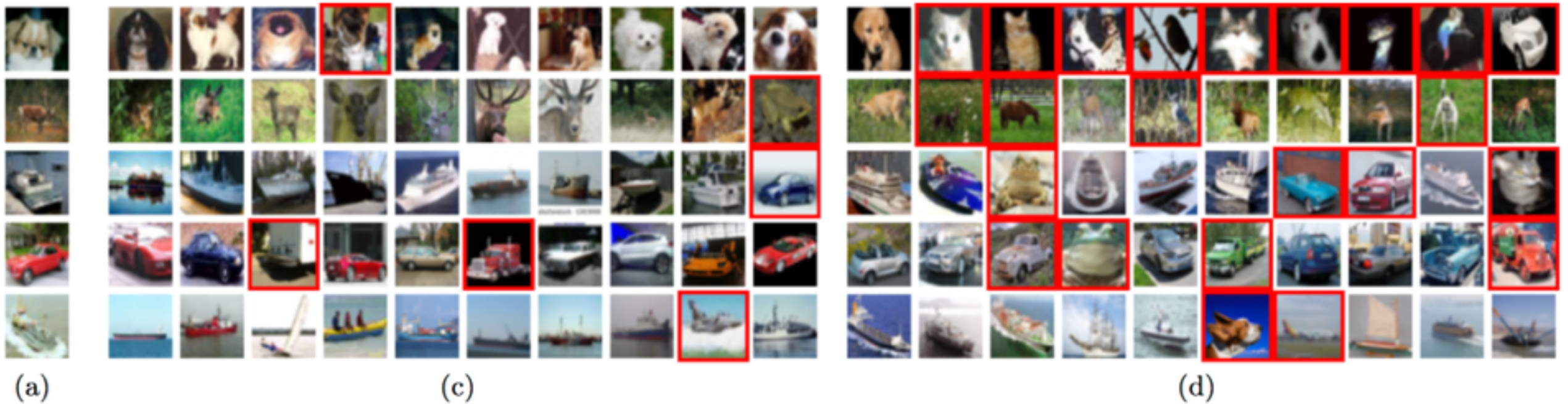
- Bit Sampling
- Random Projection
- ...



# Learning to Hash

- Generating Hashing Functions from data, rather than being independent of data
- Recent research indicates that Learning to Hash can achieve the same accuracy with much smaller codes length!
- Gong, Yunchao, and Svetlana Lazebnik. "Iterative quantization: A procrustean approach to learning binary codes." *Computer Vision and Pattern Recognition (CVPR), 2011 IEEE Conference on*. IEEE, 2011.

# LFH(Zhang et al. 2014)



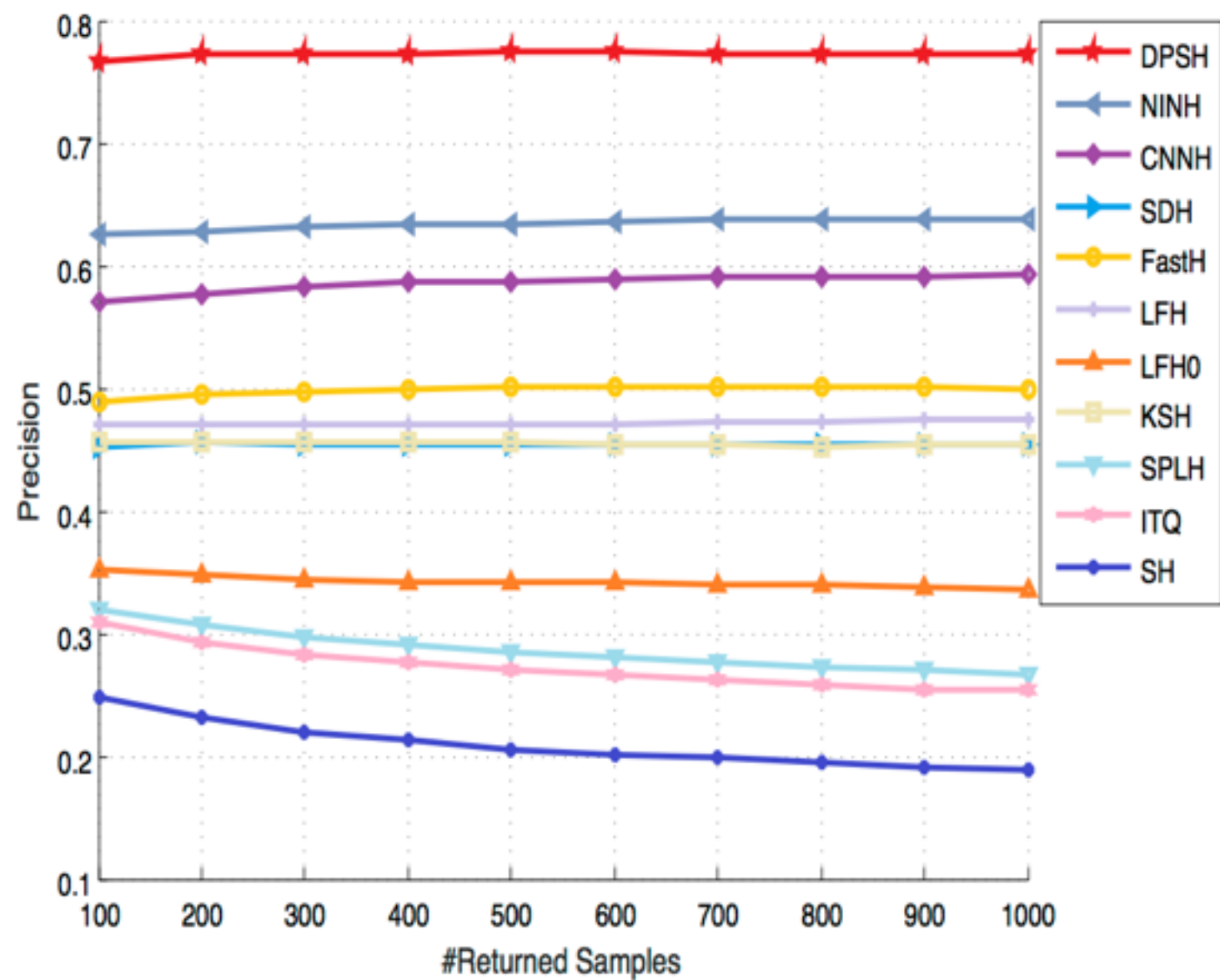
# COSDISH(Kang et al. 2016)

- MAP

Method	CIFAR-10 (60K)			
	8-bits	16-bits	32-bits	64-bits
COSDISH	<b>0.4986</b>	<b>0.5768</b>	<b>0.6191</b>	<b>0.6371</b>
SDH	0.2642	0.3994	0.4145	0.4346
LFH	0.2908	0.4098	0.5446	0.6182
TSH	0.2365	0.3080	0.3455	0.3663
KSH	0.2334	0.2662	0.2923	0.3128
SPLH	0.1588	0.1635	0.1701	0.1730
COSDISH_BT	<b>0.5856</b>	<b>0.6681</b>	<b>0.7079</b>	<b>0.7346</b>
FastH	0.4230	0.5216	0.5970	0.6446



# DPSH(Li et al. 2015)



- <http://arxiv.org/abs/1511.03855>

# DPSH(Li et al. 2015)



Figure 3. Retrieval results (top 15 returned images) for ten query images from CIFAR-10 using Hamming ranking on 48-bits hash code. Red rectangles indicate mistakes.