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CS61C Summer 2012 Midterm

Your Name:			SID:	
Your TA (Circle): Raphael	Paul	Brandon	Sung Roa	
Name of person to your LEFT:				
Name of person to your RIGH ⁻	Γ:			

This exam is worth 110 points and will count for 20% of your course grade.

The exam contains 9 questions on 13 numbered pages, including the cover page. Put all answers on these pages; don't hand in stray pieces of paper.

Question 0: You will receive 1 point for properly filling out this page as well your login on every page of the exam.

Question	Points (Minutes)	Score
0	1 (0)	
1	15 (26)	
2	10 (14)	
3	10 (18)	
4	9 (16)	
5	10 (18)	
6	17 (26)	
7	12 (22)	
8	26 (40)	
Total	110 (180)	

All the work is my own. I had no prior knowledge of the exam contents nor will I share the contents with others in CS61C who have not taken it yet.

Signature:	
Signature:	

Question 1: Potpourri – Hard to spell, nice to smell... (15 points, 26 minutes)

a) True/	Fal	lse:
-----------------	-----	------

- T F We prefer two's complement over the unsigned representation because two's complement can represent more values.
- T F The assembler uses symbol tables to resolve absolute addresses.
- T F A program will always execute faster in a RISC architecture than a CISC architecture.
- T F The greater the number of memory accesses in a program, the greater the AMAT.
- T F Pseudo-instructions do not always use \$at.
- T F Since \$s0 is a "saved register," it does not need to be saved before any function calls.
- b) Fill in the function below, which returns a new copy of the argument (struct definition not shown):

return new;

}

c) Here is what is currently on the heap:

_	_	_		_	_	_	
Λ .	Λ	Λ			_ n	_ n	i
~	_						i

The order of allocation/frees: A allocated, B allocated, C allocated, B freed, D allocated

What allocation strategy was used?

d) What is the average CPI of the program described in the table to the right? Which is better: halving memory access cycles or arithmetic cycles and why?

Instruction Category	Cycles	Frequency
Memory Access	10	0.1
Arithmetic	2	0.4
Branch	3	0.2
Comparison	1	0.3

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- e) I'm getting the message "cannot execute binary file". In one or two sentences, explain what the problem is and how to fix it.
- f) In our 32-bit single-precision floating point representation, we decide to convert one significand bit to an exponent bit. How many **denormalized numbers** do we have relative to before? (Circle one)

More Fewer

Rounded to the nearest power of 2, how many denorm numbers are there in our new format? (Answer in IEC format)

Question 2: Flippin' Fo' Fun (10 points, 14 minutes)

Assume that the most significant bit (MSB) of x is a 0. We store the result of flipping x's bits into y. Interpreted in the following number representations, how large is the <u>magnitude</u> of y relative to the <u>magnitude</u> of x? Circle ONE choice per row.

Unsigned	y < x	y = x	y > x	Can't Tell
One's Complement	y < x	y = x	y > x	Can't Tell
Two's Complement	y < x	y = x	y > x	Can't Tell
Sign and Magnitude	y < x	y = x	y > x	Can't Tell
Biased Notation (e.g. FP exponent)	y < x	y = x	y > x	Can't Tell

Question 3: Doctor Who?!? (10 Points, 18 Minutes)

The Daleks are invading the Earth again, and we need the help of the Doctor! Find the errors in this code and fix them so that the code correctly prints "The 10th Doctor and the Blue Police Box". There is exactly one coding error for each function and function call pair and can be fixed by changing 5 or 6 lines total. Fill in the corrections in the blanks on the opposite page.

```
1
     void whichDoctor(int* input) {
2
        input = 10;
3
     void doctorChanger(char** input1, char** input2) {
4
5
         char* temp = *input1;
6
         *input1 = *input2;
7
         *input2 = temp;
8
     }
9
     char* policeBoxGiver(char* input) {
10
         *input = "The Master";
11
         return "Police Box";
12
13
     char* colorMaker(void) {
14
         char* color = malloc(sizeof(char) * 4);
15
         color[0] = 'B';
16
         color[1] = '1';
17
         color[3] = 'u';
18
        color[2] = 'e';
19
        color[4] = 0;
20
        return color;
21
     }
22
     char* colorFixer(char* input) {
23
         char temp = *(input+2);
24
         *(input+2) = *(input+1);
25
         *(input+1) = temp;
26
27
     int main(void) {
28
         int * ith = malloc(sizeof(int));
29
         whichDoctor(ith);
30
         char* doctor = "Master";
31
         char* master = "Doctor";
32
         char* details = "and the";
33
         char* color = colorMaker();
34
         colorFixer(color);
35
         char* box = "David Tennant";
36
         doctorChanger(doctor, master);
37
        policeBoxGiver(box);
38
        printf("The %dth %s %s %s %s", *ith, doctor, details, color, box);
39
      }
```

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Line #	Corrected Code
	

Question 4: Let Me Float This Idea By You (9 Points, 16 Minutes)

For a very simple household appliance like a thermostat, a more minimalistic microprocessor is desired to reduce power consumption and hardware costs. We have selected a **16-bit** microprocessor that does not have a floating-point unit, so there is no native support for floating point operations (no float/double). However, we'd still like to represent decimals for our temperature reading so we're going to implement floating point operations in software (in C).

a) Define a new variable type called fp:

We have decided to use a representation with a **5-bit exponent field** while following all of the representation conventions from the MIPS 32-bit floating point numbers **except denorms**.

Fill in the following functions. Not all blanks need to be used. <u>You can call these functions and assume proper behavior regardless of your implementation</u>. Assume our hardware implements the C operator ">>" as shift right arithmetic.

b)

```
/* returns -num */
fp negateFP(fp num) {
    return _____;
}
```

c)

```
/* returns the signed value of the exponent */
int getExp(fp num) {
    return
    ;
}
```

d)

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Question 5: Who Says Less is Better? (10 points, 18 Minutes)

We're going to take a page out of the ARM book and design a new instruction set architecture with just **16 32-bit registers**. This means that we only need 4-bit register fields in our instructions.

a) How many extra bits do we have now for other fields in the following formats?

R: _____ J: ____

b) For R-format instructions, would you give the extra bits to opcode, shamt, or funct? ______ Explain your choice in a sentence or two (no credit without explanation):

For I-format instructions, we naturally give the extra bits to the immediate field, resulting in the following format:

[opcode (6) | rs (4) | rt (4) | immediate (18)]

c) What fraction of our address space can we now reach with a branch instruction?

d) Assume our PC currently contains the address 0x08000000.

What is the LOWEST address (in hex) we can reach with a branch?

e) Write out the Verilog pseudocode (as in the OPERATION column on the MIPS Green Sheet) for beq. Make sure you specify what BranchAddr is.

Question 6: Cache in While You Can (17 points, 26 Minutes)

Со	nsider a single 4KiB cach	e with 512B blocks and a writ	e-back policy. Assume a 32-bit address space.
a)	If the cache were direct	t-mapped,	
	# of rows?	# of offset bits?	
b)	If the cache were 4-way	y set associative,	
	# of tag bits?	# of index bits?	# of bits per cache slot?
Со	nsider an array of the fol	llowing location structs:	
	<pre>pedef struct { // some under int visited; int danger; location;</pre>	defined number of othe	r struct members
10	cation locs[NUM_LO	CS];	
soı			we've visited. Assume this gets executed held in a register, and the size of the array is
fo	r(int i = 0; i < N if(locs[i].vis	_	
c)	What's the fewest poss	sible number of bytes written	to main memory?
d)	What's the greatest po	ssible number of bytes writte	n to main memory?
No	w consider if we store th	ne visited and danger info	rmation in individual arrays instead:
	t visited[NUM_LOCS t danger[NUM_LOCS]		
e)	This way, the cache car	n exploit better	for the above task.
	We can expect a	(higher or	lower) miss rate
	because of the change	in the number of	(type of cache miss) misses.

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Consider the following code with NUM LOCS $> 2^10$.

```
for(int i = 0; i < NUM_LOCS; i++)
    if(visited[i] && danger[i] > 5) count++;
```

Two memory accesses are made per iteration: one into visited, the other into danger. Assume that the cache has no valid blocks initially. You are told that in the worst case, the cache has a miss rate of 100%. Consider each of the following possible changes to the cache individually.

f)	Mark each as E, if it eliminates the chances of this worst-case scenario miss rate, R if it reduces the
	chances, or N if it's not helpful.

•	More sets, same I	block size, same asso	ociativity	

- Double associativity, half block size, same total cache size

- Everything stays the same but use a write-through policy instead

Question 7: Can't Make Copies Fast Enough (12 points, 22 Minutes)

We are revisiting our friend the Fast String Copy from lecture! Recall that the function prototype in C is as follows:

```
char *strcpy(char *dst, char *src);
```

Consider the following MIPS implementation of this function:

```
jal strcpy # begin function call
...
strcpy:
   addi $v0,$a0,0
loop: lb $t0,0($a1)
   sb $t0,0($a0)
   addiu $a0,$a0,1
   addiu $a1,$a1,1
   beq $t0,$zero,exit
   j loop
exit: jr $ra
```

Suppose we are running code on a machine with the following cache parameters:

- Unified L1\$ with a hit time of 2 cycles and a hit rate of 95%
- Miss Penalty to main memory of 200 cycles
- Base CPI of 1.5 (in the absence of cache misses)
- a) Calculate our machine's AMAT:
- b) What is the CPI of a single call to strcpy with src = "" (the function call includes the jal)?
- c) We decide to add a L2\$ to reduce our AMAT to 6. Our L2\$ has a hit time of 20 cycles. What's the worst Local Hit Rate that will still meet our AMAT goal?

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d)	In addition to speeding up our architecture, we want to speed up our code, so we decide to
	eliminate the return value (presumably the caller retains a copy of the destination pointer). In this
	case, the strcpy function above can be rewritten in just 6 instructions. Write out this
	implementation in the blanks below, introducing any necessary labels (don't worry about any label
	name clashes with strcpy).

strcpy2:		

e) If we call strcpy and strcpy2 on the same src string of length n+1=N (N includes '\0'), what is the ratio of instructions executed in strcpy2 (including the jal)? Leave your answer in terms of N.

f) Is the ratio in part (e) the same as the relative performance between these two functions? In a sentence or two, explain why or why not.

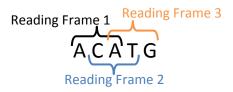
Question 8: Putting the Science in Computer Science (26 points, 40 minutes)

DNA can be called the "alphabet of life." From a *very* simplified view, DNA within a cell produces amino acids, which in turn produce proteins, which are the building blocks for most of your body. Here we'd like to write some code for examining a strand of DNA.

a) DNA is made up of nucleotides, which we write shorthand as A, C, G, and T. DNA is in base 4
(quaternary)! Fill in the table below, using the DNA nucleotide symbols in alphabetical order
(A < C < G < T).

Decimal	DNA
	CAT
50	

An amino acid is encoded by three nucleotides. Because DNA is found in long strands, the following 5 nucleotides can be read 3 different ways:



The sequence **ATG** (as seen in the 3rd reading frame) signals the beginning of a protein ("start codon").

b) Fill in the blanks on the opposite page for the **recursive** function find_start in MIPS that <u>returns</u> the position of the first start codon found in the given strand of DNA. Assume each nucleotide is stored as a char in memory. *Blanks do not necessarily need to be filled*. Maximum points awarded for using the *fewest* amount of registers and memory.

[Answer the following AFTER looking at the code]

t and a second s				
Ass	sume we call find_star	t from main with char dr	na[] = "GCATGC";.	
c)	c) How many total frames are created on the Stack (not including main)?			
d)) What is the maximum depth of the Stack (in # of frames, not including main)?			
e)) What will the line j ret look like once this file is run through the assembler?			
f)	Where will the label ret show up? (Circle one)			
	Symbol Table	Relocation Table	Both	Neither

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```
/* pos: search position from start of strand */
               int find start(char *dna, int pos);
                                                         this is a recur. function don't
                                                         forget dumby
find start:
     addiu $sp,$sp,-4 # PROLOGUE
     sw $ra, 0($sp) #don't need to set up par. since dna stays the same for jal call
     jal strlen # call strlen(dna); Assume strlen doesn't
                            # change $a0 or $a1
    slti $t0, $v0, 3 # make sure we don't read past the end of
                            # the array
           $t0,$0,chk # 'chk' for check if start codon
     beq
     addi $v0,$0,-1  # return -1 (start codon not found)
                       # 'ret' for return
     j
chk: Text
          $t1,0($t0)
     lb
                       why do you add 65?
     addi $t2,$0,65
     bne $t1,$t2,rec # 'rec' for recurse
           $t1,1($t0)
     lb
     addi $t2,$0,84
     bne $t1,$t2,rec
           $t1,2($t0)
     lb
     addi $t2,$0, 71
           $t1,$t2,rec
      add $v0, $t2, $0
                      # return current position
           ret
rec: lb $t1, 4($a1)
                   # recurse at next position
     addi $t1, $t1, 1
     Text
           find start
     jal
     Text
     Text
                      # EPILOGUE
ret:
     sw $ra, 0($sp)
     addiu $sp,$sp,4
           $ra
     jr
```

C function prototype: /* dna: start address of DNA strand */

BACK OF EXAM

(Any work on this page will not be graded)