

#### Distributive Law

$$ab + ac = a(b + c)$$

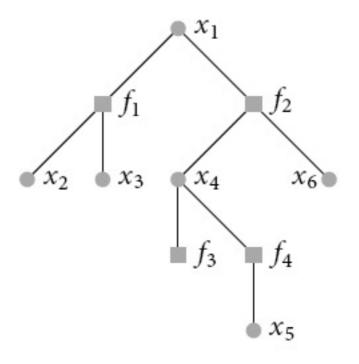
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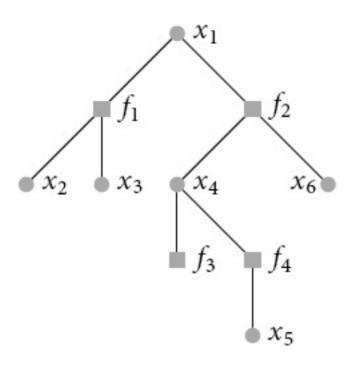
$$\sum_{i,j} a_i b_j \qquad (\sum_i a_i)(\sum_j b_j)$$

$$f(x_1, x_2, x_3, x_4, x_5, x_6) = f_1(x_1, x_2, x_3) f_2(x_1, x_4, x_6) f_3(x_4) f_4(x_4, x_5)$$

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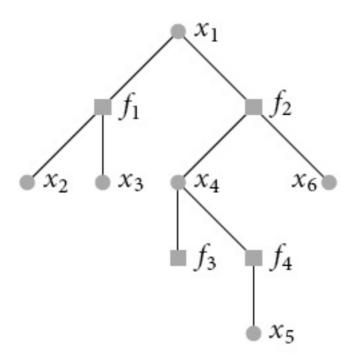


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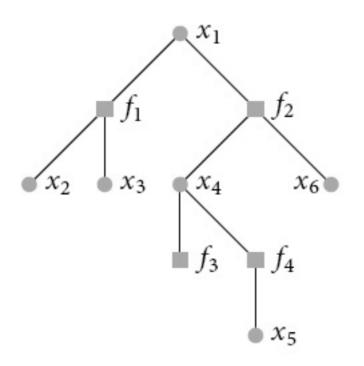
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Note: f(x<sub>1</sub>) is a function; therefore, it takes on a distinct value for each value of x<sub>1</sub>

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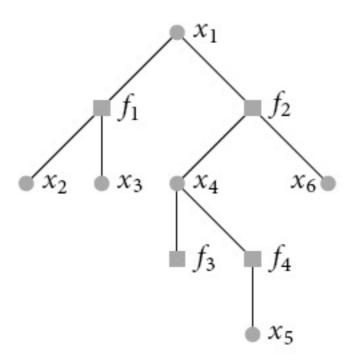


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 $|\mathcal{X}|$  alphabet

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 $\Theta(|\mathcal{X}|^6)$  brute force complexity

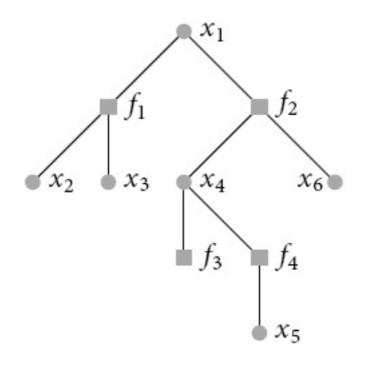
$$f(x_1) = \left[\sum_{x_2,x_3} f_1(x_1,x_2,x_3)\right] \left[\sum_{x_4} f_3(x_4) \left(\sum_{x_6} f_2(x_1,x_4,x_6)\right) \left(\sum_{x_5} f_4(x_4,x_5)\right)\right]$$

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 complexity



Does there exist a systematic way to find this low complexity scheme using the structure of the graph?



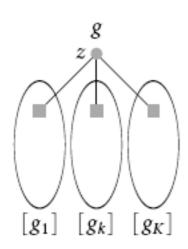
$$g(z) = \sum_{\sim z} g(z, \dots)$$

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$$g(z,\ldots)=\prod_{k=1}^K [g_k(z,\ldots)]$$

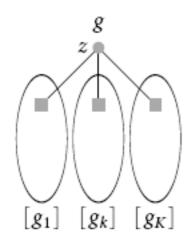
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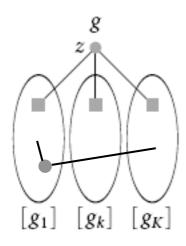
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Note: the individual functions  $g_k(z, ...)$  only share the variable z; all other variables are "private"

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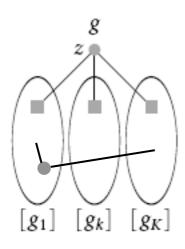


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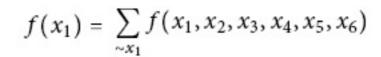
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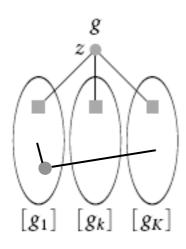
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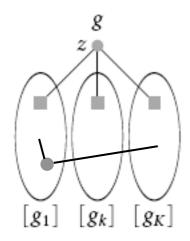
$$f(x_1,...) = [f_1(x_1,x_2,x_3)][f_2(x_1,x_4,x_6)f_3(x_4)f_4(x_4,x_5)]$$



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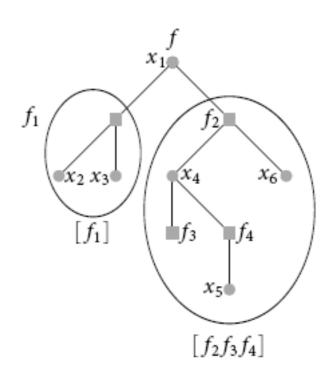
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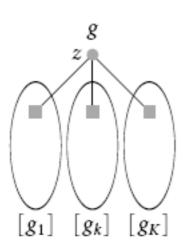
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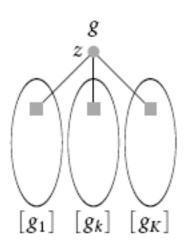




$$\sum_{z} g(z, \dots) = \sum_{z} \prod_{k=1}^{K} [g_k(z, \dots)] = \prod_{k=1}^{K} \left[ \sum_{z} g_k(z, \dots) \right]$$
marginal of product product of marginals



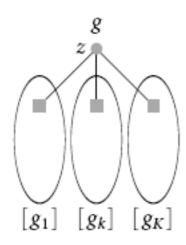
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marginal  $\sum_{\sim z} g(z, \dots)$  is the product of the individual marginals

recall that g(z) is a function, taking a distinct value for each value of z

$$\sum_{z} g(z, \dots) = \underbrace{\sum_{z} \prod_{k=1}^{K} [g_k(z, \dots)]}_{\text{marginal of product}} = \underbrace{\prod_{k=1}^{K} \left[\sum_{z} g_k(z, \dots)\right]}_{\text{product of marginals}}$$

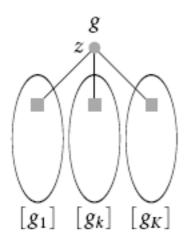


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instead of computing g(z) directly by brute force we can first compute each of the functions  $g_k(z)$ ; we then get g(z) by multiplying these functions  $g_k(z)$ 

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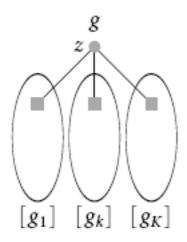
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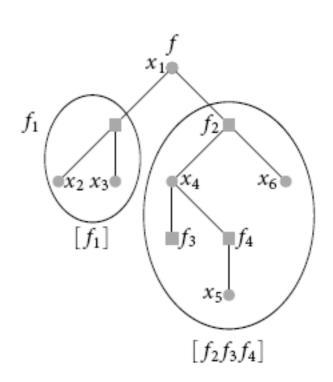
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marginal of product  $\sum_{z} g_k(z, \dots) = \sum_{k=1}^{K} \left[ \sum_{z} g_k(z, \dots) \right]$ 
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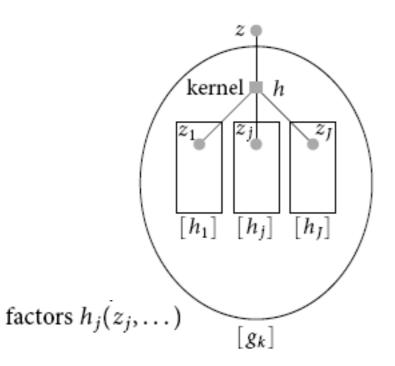
$$g_k(z,...) = \underbrace{h(z,z_1,...,z_J)}_{\text{kernel}} \prod_{j=1}^J \underbrace{\left[h_j(z_j,...)\right]}_{\text{factors}}$$

"kernel"  $h(z, z_1, \ldots, z_I)$ 

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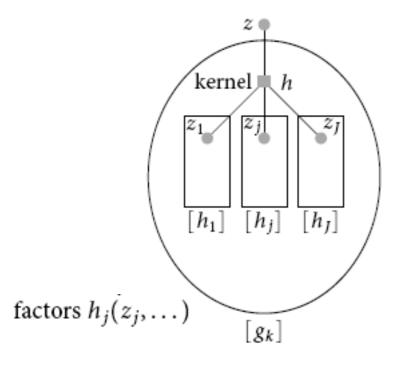
$$g_k(z,...)$$

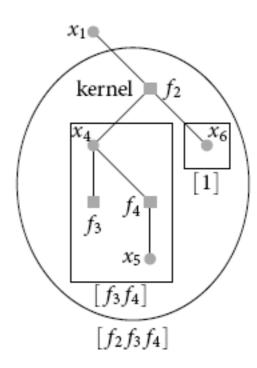


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"kernel"  $h(z, z_1, \ldots, z_J)$ 

$$g_k(z,...)$$





$$f_2(x_1, x_4, x_6) f_3(x_4) f_4(x_4, x_5) = \underbrace{f_2(x_1, x_4, x_6)}_{\text{kernel}} \underbrace{[f_3(x_4) f_4(x_4, x_5)]}_{x_4} \underbrace{[1]}_{x_6}$$

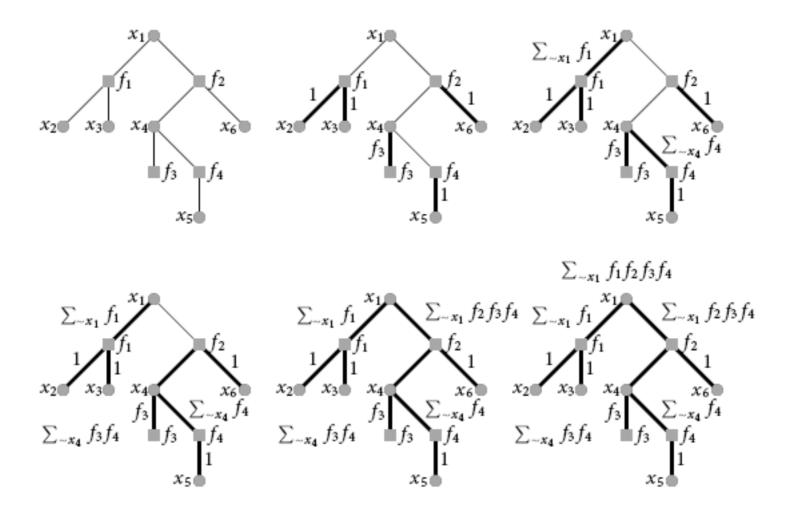
$$\sum_{z} g_k(z, \dots) = \sum_{z} h(z, z_1, \dots, z_J) \prod_{j=1}^{J} [h_j(z_j, \dots)]$$

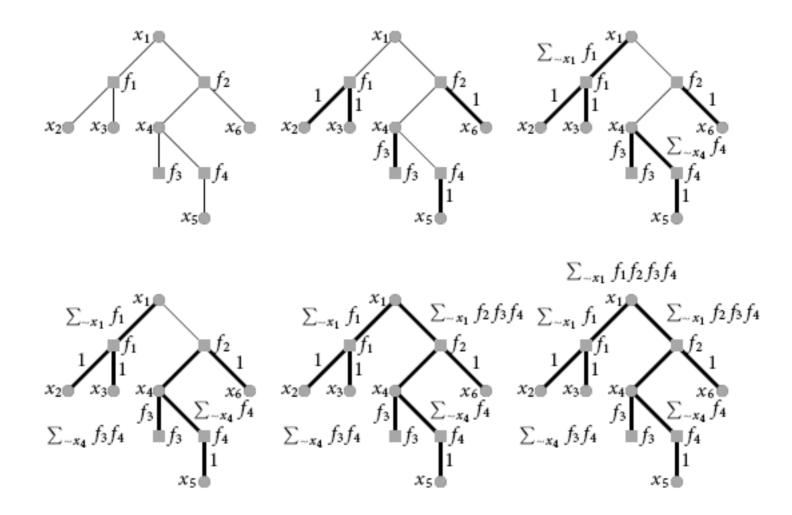
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product of marginals

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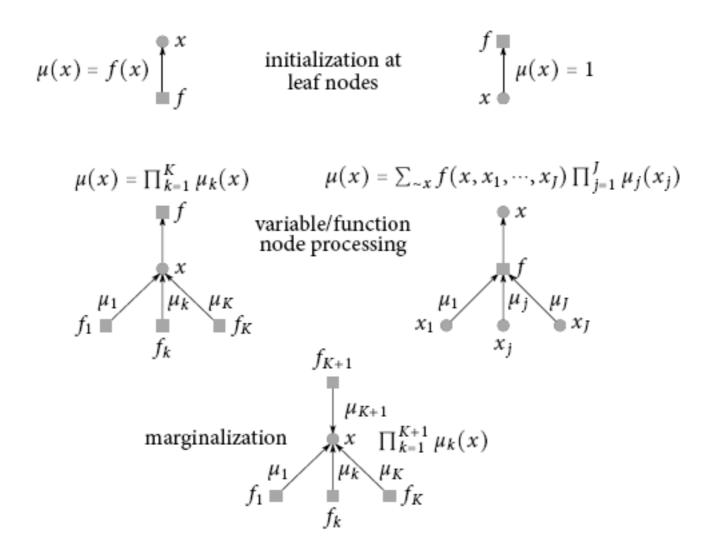
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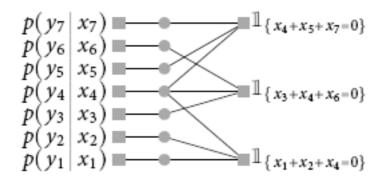
complexity proportional to highest degree

### Message Passing Rules



# Summary and Limitations

## Summary and Limitations



#### References

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