```
3.
       a) addition
       b) simplification
       c) modus ponens
       d) modus tollens
       e) hypothetical syllogism
4.
       a) simplification
       b) disjunctive syllogism
       c) hypothetical syllogism
       d) addition
       e) hypothetical syllogism
13.
       a)
S(x): "x is a student at the class"
J(x): "x knows how to write programs in JAVA"
H(x): "get a high-paying job"
S(Doug)
J(Doug)
\forall x (J(x) \rightarrow H(x))
conclusion :\exists x (S(x) \land J(x))
       1. \forall x (J(x) \rightarrow H(x))
       2. J(Doug)
                                          using modus ponens
       3. H(Doug)
       4. S(Doug) \( \Lambda \) H(Doug)
                                          conjunction 4 and given
       5. \exists x (S(x) \land H(x))
                                          using existential generalization
       b)
C(x): "x in the class"
W(x): "enjoys whale watching"
P(x): "cares about ocean pollution"
\exists x (C(x) \land W(x))
\forall x (W(x) \rightarrow P(x))
conclusion : \exists x(C(x) \land P(x))
   1. \exists x (C(x) \land W(x))
   2. C(a) \( \text{W(a)} \)
                                    existential instantiation
   3. C(a)
                                           simplification
   4. \forall x (W(x) \rightarrow P(x))
   5. W(a) \rightarrow P(a)
                                    existential instantiation
   6. W(a)
   7. P(a)
                                           simplification
   8. C(x) \(\Lambda\) W(a)
                                   conjunction from 3 and 7
   9. \exists x(C(x) \land P(x))
                                   exitential generalization
```

```
c)
S(x): "x student in the class"
C(x): "x owns a personal computer"
W(x): "x use a word processing program"
\forall x (S(x) \rightarrow C(x))
A^{\times} (C(x) \rightarrow M(x))
conclusion : S(zake) → W(zake)
       1. \forall x (S(x) \rightarrow C(x))
       2. S(zake) \rightarrow C(zake)
                                         using universal instantiation
       3. S(zake)
                                         given
       4. C(zake)
                                         modus ponens
       5. \forall x (C(x) \rightarrow W(x))
       6. C(zake) \rightarrow W(zake)
                                         using universal instantiation
       7. C(zake)
                                         modus ponens
       8. W(zake)
                                         from 4
       9. S(zake) \rightarrow W(zake)
                                         hypothetical syllogism
       d)
J(x): "x lives in New Jersey"
O(x): "x lives within 50 miles of the ocean"
N(x): "x have seen the ocean"
\forall x (J(x) \rightarrow O(x))
\exists x(J(x) \land -N(x))
conclusion : \exists x(0(x) \land -N(x))
       1. \forall x (J(x) \rightarrow O(x))
       2. J(a) \rightarrow O(a)
                                         using universal instantiation
       3. J(a)
       4. 0(a)
                                                modus ponens
       5. \exists x(J(x) \land -N(x))
```

using exitential instantiation

simplification

exitential generalization

using conjunction

13.

a)

S(x): "x is a student in the class"
R(x): "x owns a red convertible"
T(x): "x has gotten a speeding ticket" $S(x) \land R(x)$ $\forall x (R(x) \rightarrow T(x))$ R(Linda)

6. J(a) \(\Lambda\) -N(a)

8. 0(a) Λ -N(a)

9. $\exists x (O(x) \land -N(x))$

7. -N(a)

```
conclusion : \exists x(s(x) \land T(x))
   1. S(Linda) \( \Lambda \) R(Linda)
   2. S(Linda)
                                     simplification
   3. \forall x (R(x) \rightarrow T(x))
   4. R(Linda) → T(Linda)
                                     universal instantiation
   5. R(Linda)
                                     given
   6. T(Linda)
                                     modus ponens
   7. T(Linda) \( \Linda \)
                                     from 2 using conjunction
   8. \exists x (T(x) \land S(x))
                                     exitential generalization
       b)
R(x): "x is a roommate"
S(x): "x is a student"
D(x): "x has taken a course in discrete mathematics"
A(x): "x can take a course in algorithms"
\forall x (S(x) \land D(x))
\forall x (D(x) \rightarrow A(x))
R(x) \rightarrow S(x)
conclusion : \forall x (R(x) \rightarrow A(x))
   1. R(Aaron) \rightarrow S(Aaron)
   2. R(Aaron)
    S(Aaron)
                                     modus ponens
   4. \forall x (S(x) \land D(x))
   S(Aaron) Λ D(Aaron)
                                     universal instantiation
   6. D(Aaron)
                                     simplification
   7. \forall x (D(x) \rightarrow A(x))
   8. D(Aaron) \rightarrow A(Aaron)
                                     universal instantiation
   9. R(Aaron) \rightarrow A(Aaron)
                                     from 1 and 8 using hypothetical syllogism
   10. \forall x (R(x) \rightarrow A(x))
                                     universal generalization
       c)
M(x): "x is a movies produced by John Sayles"
W(x): "x is a wonderful movie"
C(x): "x is a movie about coal miners"
Ax(W(x) \rightarrow M(x))
\exists x (M(x) \land C(x))
conclusion : \exists x(W(x) \land C(x))
   1. \forall x (M(x) \rightarrow W(x))
   2. M(c) \rightarrow W(c)
                                     universal instantiation
   3. \exists x (M(x) \land C(x))
   4. M(c) \( \Lambda \) C(c)
                                     existential instantiation
   5. M(c)
                                     simplification
   6. M(c) \rightarrow W(c)
```

```
7. W(c) modus ponens
```

8. $W(c) \land C(c)$ conjunction from 7 and 4 by simplification

9. $\exists x(W(x) \land C(x))$ existential generalization

d)

S(x): "x in the class"

F(x): "x has been to France"

L(x): "x had visited the Louvre"

 $\exists x(S(x) \land F(x))$ $\forall x(F(x) \rightarrow L(x))$

conclusion : $\exists x(S(x) \land L(x))$

- 1. $(S(x) \wedge F(x))$
- 2. S(a) \(\Lambda \) F(a) existential instantiation
- 3. S(a) simplification
- 4. $\forall x(F(x) \rightarrow L(x))$
- 5. $F(a) \rightarrow L(a)$ universal instantiation 6. F(a) from 2 by simplification
- 7. L(a) modus ponens
- 8. $S(a) \wedge L(a)$ conjunction 7 and 3
- 9. $\exists x (S(x) \land L(x))$ existential generalization

19.

- a) NOT a valid argument p \rightarrow q doesn't mean that I can say q \rightarrow p (misusing of the converse)
 - b) valid argument using modus tollens
- c) NOT a valid argument $p \rightarrow q$ doesn't mean that I can say $-p \rightarrow -q$ (misusing of the inverse)

20.

- a) not a valid argument (misusing of converse)
- b) Valid argument (universal instaitation)

23.

error in step 5, we are assuming that the same c can be true for both P(x) and Q(x) which maybe NOT! We can rewrite the argument as follows:

5. Q(a) Existential instantiation

6. P(c) \(\Lambda \) Q(a) Conjunction

7. $\exists x \exists y (P(x) \land Q(y))$ Existential generalization

8. S(a) \(\Lambda \) R(a)

9. $\forall x(R(x) \land S(x))$

in step 2 c should be refering to an arbitrary value from the domain not a specific value.

error in step 3, we can simplify only from statement having Λ , same for step 5. another error in step 7, using conjunction is using the AND operator as follows: $\forall x \ P(x) \ \mathbf{v} \ \forall x \ O(x)$.

```
26.
\forall x (P(x) \rightarrow Q(x))
\forall x (Q(x) \rightarrow R(x))
conclusion : \forall x (P(x) \rightarrow R(x))
   1. \forall x (P(x) \rightarrow Q(x))
    2. P(a) \to 0(a)
                                        universal instantiation (a is an arbitrary value)
    3. \forall x (Q(x) \rightarrow R(x))
    4. Q(a) \rightarrow R(a)
                                        universal instantiation (a is an arbitrary value)
    5. P(a) \rightarrow R(a)
                                        Hypothetical syllogism from 2 and 4
   6. \forall x (P(x) \rightarrow R(x))
                                        universal generalization
27.
if \forall x (P(x) \rightarrow (Q(x) \land S(x))) and \forall x (P(x) \land R(x)) then \forall x (R(x) \land S(x))
    1. \forall x (P(x) \rightarrow (Q(x) \land S(x)))
    2. P(a) \rightarrow (Q(a) \land S(a))
                                                universal instanitation (a is arbitary value)
    3. \forall x(P(x) \land R(x))
   4. P(a) \( \Lambda \) R(a)
                                                universal instanitation (a is arbitary value)
    5. P(a)
                                                simplification
    6. Q(a) \( \Lambda \) S(a)
                                                modus ponens from 5 and 2
    7. S(a)
                                                simplification
```

conjunction and simplification from 4

universal generalization