# Chapter 27: Interlude - Thread API

### Overview

- Main Idea: POSIX threads (pthreads) provide interfaces for thread creation, synchronization (locks/condition variables), and management.
- **Key Terms**: pthread\_create, pthread\_join, mutex, condition variable, critical section.

#### 1. Thread Creation

```
pthread_create Signature
```

#### **Example: Passing Arguments**

```
typedef struct { int a; int b; } myarg_t;
void *mythread(void *arg) {
    myarg_t *args = (myarg_t *) arg;
    printf("%d %d\n", args->a, args->b);
    return NULL;
}
int main() {
    pthread_t p;
    myarg_t args = { 10, 20 };
    pthread_create(&p, NULL, mythread, &args); // Pass struct as argument    pthread_join(p, NULL);
}
```

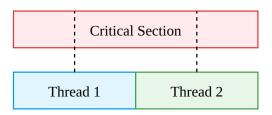


(Main thread creates worker thread with arguments.)

## 2. Thread Completion

pthread\_mutex\_unlock(&lock);

```
pthread_join Signature
int pthread_join(pthread_t thread, void **retval); // Waits for thread to finish
Example: Returning Values
typedef struct { int x; int y; } myret_t;
void *mythread(void *arg) {
   myret_t *rvals = malloc(sizeof(myret_t));
   rvals->x = 1;
   rvals -> y = 2;
   return (void *) rvals; // Return heap-allocated struct
}
int main() {
   pthread_t p;
   myret t *rvals;
   pthread_create(&p, NULL, mythread, NULL);
   pthread_join(p, (void **) &rvals);
   printf("Returned %d %d\n", rvals->x, rvals->y);
   free(rvals);
}
Warning: Never return stack-allocated data from a thread!
// BAD CODE: Returning stack address
myret_t oops;
return (void *) &oops; // Crash-prone!
3. Locks (Mutexes)
Initialization
pthread_mutex_t lock = PTHREAD_MUTEX_INITIALIZER; // Static
pthread_mutex_init(&lock, NULL);
                                                   // Dynamic
Usage
pthread_mutex_lock(&lock);
x = x + 1; // Critical section
```



<sup>\*(</sup>Threads serialize access to critical sections.)\_

#### **Common Pitfalls**

1. Missing Initialization:

```
pthread_mutex_t lock; // Uninitialized → undefined behavior!
pthread_mutex_lock(&lock);
```

2. Ignoring Error Codes:

Always check return values or use wrappers:

```
void Pthread_mutex_lock(pthread_mutex_t *mutex) {
   int rc = pthread_mutex_lock(mutex);
   assert(rc == 0);
}
```

#### 4. Condition Variables

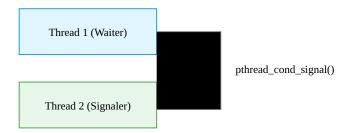
#### Initialization

```
pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
```

#### Waiting and Signaling

```
// Thread 1: Waits for condition
pthread_mutex_lock(&lock);
while (ready == 0) {
    pthread_cond_wait(&cond, &lock); // Releases lock while waiting
}
pthread_mutex_unlock(&lock);

// Thread 2: Signals condition
pthread_mutex_lock(&lock);
ready = 1;
pthread_cond_signal(&cond);
pthread_mutex_unlock(&lock);
```



(Signaling thread wakes up waiting thread.)

#### **Key Points:**

- Always use a while loop (not if) for conditions due to spurious wakeups.
- Hold the mutex when modifying shared state (e.g., ready).

#### 5. Thread API Guidelines

Do	Don't
Initialize locks/condition variables.	Use uninitialized synchronization primitives.
Check return codes.	Ignore error handling.
Pass heap-allocated data to threads.	Return stack addresses from threads.
Use condition variables for signaling.	<pre>Spin-wait on flags (while (flag == 0);).</pre>

# Homework Insights

#### 1. Race Conditions:

valgrind --tool=helgrind main-race # Detects data races

#### 2. Deadlocks:

Circular lock acquisition (e.g., T1 holds A, waits for B; T2 holds B, waits for A).

#### 3. Efficiency:

Spin-waiting (while (done == 0);) wastes CPU; condition variables are preferred.