### **COMPILER DESIGN**

### **SOLUTIONS**

$$S \rightarrow aSAb \mid bSBc \qquad \Rightarrow First(S) = \{a, b\}$$

$$A \rightarrow +AB \mid \epsilon \qquad \Rightarrow First(A) = \{+, \epsilon\}$$

$$B \rightarrow *BC \mid \epsilon \qquad \Rightarrow First(B) = \{*, \epsilon\}$$

$$C \rightarrow aC \mid d \qquad \Rightarrow First(C) = \{a, d\}$$

**1.** What is in the Follow(S)?

(a) 
$$\{a, b, c, +, \$\}$$

(b) 
$$\{a, c, +, *, \$\}$$

(c) 
$$\{b, c, +, *, \$\}$$

(d) 
$$\{a, b, d, *, \$\}$$

**Solution:** Option (c)

**Explanation:** 

$$S \rightarrow aSAb \mid bSBc$$

$$A \rightarrow +AB \mid \epsilon$$

$$B \rightarrow *BC \mid \epsilon$$

$$C \rightarrow aC \mid d$$

 $Follow(S) = \{First(A), First(B), b, c, \$\}$ 

**2.** What is in the Follow(B)?

(a) 
$$\{a, b, c, d, *\}$$

(b) 
$$\{a, b, d, \epsilon, \$\}$$

(c) 
$$\{a, c, d, *, \$\}$$

(d) 
$$\{c, d, b, +, *\}$$

**Solution:** Option (a)

**Explanation:** 

$$Follow(B) = \{C, Follow(A), First(C)\}$$

$$= \{c, b, *, a, d\}$$

- **3.** Choose the False statement.
- (a) No left recursive/ ambiguous grammar can be LL(1)
- (b) The class of grammars that can be parsed using LR methods is proper subset of the class of grammar that can be parsed by LL method
- (c) LR parsing is non-backtracking method
- (d) LR parsing can describe more languages than LL parsing

**Solution:** Option (b)

### **Explanation:**

FALSE, as  $LL(1) \subseteq LR(k)$ 

**4.** Consider the following SDT.

$$A \rightarrow BC *(I) B.i = f(A.i)$$
(II)  $B.i = f(A.S)$ 
(III)  $A.S = f(B.s)$ 

Which of the above is violating L – attributed definition?

(a) I only

(b) II only

(c) I, II

(d) I, II, III

**Solution:** Option (b)

## **Explanation:**

It does not follow L-attribute definition.

5

$$X \rightarrow YZ$$

$$Y \rightarrow Y + Z \{ print ('+'); \}$$
  
 $T \{ Y.val = T.val \}$ 

$$Z \rightarrow *Y \{ print (`*'); \} Z$$

$$T \{ Z.val = T.val \}$$

$$\varepsilon$$

 $T \rightarrow num \{print(num.val);\}$ 

For 2+3\*2, the above translation scheme prints

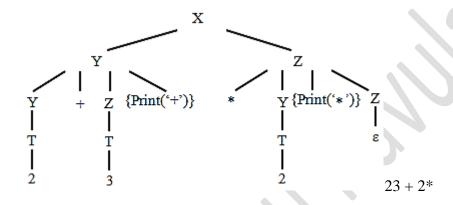
(a) 
$$2+3*2$$

(b) 
$$23+2*$$

(c) 
$$232*+$$

**Solution:** Option (b)

### **Explanation:**



**6.** Consider the following expression

$$x = a*b - c*d + e$$

For generating target code how many register will be required apart from accumulator A?

(a) 1

(b) 2

(c) 3

(d) 5

**Solution:** Option (a)

## **Explanation:**

$$x = a * b - c * d + e$$

MOV A, a

MUL b

MOV R<sub>1</sub>, A

MOV A, C

MUL d

SUB R<sub>1</sub>, A

ADD R<sub>1</sub>, e

So, One Register required.

## **7.** Consider the following two grammars

$$G_1{:}\; A \to A1 \mid 0A1 \mid 01$$

$$G_2: A \rightarrow 0A \mid 1$$

Which of the following is True regarding above grammars?

(a)  $L_1$  is LR(k)

(b)  $L_2$  is LR(k)

(c) Both  $L_1$  and  $L_2$  is LR(k)

(d) None is LR(k)

**Solution:** Option (b)

### **Explanation:**

$$G_1: A \rightarrow A1 \mid 0A1 \mid 01 --- Ambiguous grammar$$

$$G_2$$
:  $A \rightarrow 0A \mid 1$  --- Regular grammar

Ambiguous grammar is not LR(k)

Above Regular grammar is LR(k)

## **8.** Consider the following grammar.

$$S \to aB \mid aAb$$

$$A \to bAb \mid a$$

$$B \to aB \mid \epsilon$$

How many back tracks are required to generate the string aab from the above grammar?

(a) 1

(b) 2

(c) 3

(d) 4

**Solution:** Option (b)

# **Explanation:**

 $S \Rightarrow aB$ 

 $S \Rightarrow aAb$ 

S ⇒ aAb

⇒ aaB

⇒ abAbb

⇒ aab

⇒ aa

⇒ ababb

Backtrack

Backtrack

So, 2 backtracking is required.