prime detection

miller-rabin

```
ll qmul(ll a, ll b, ll mod)//快速乘
   11 c=(long double)a/mod*b;
   11 res=(unsigned long long)a*b-(unsigned long long)c*mod;
   return (res+mod)%mod;
}
ll qpow(ll a, ll n, ll mod)//快速幂
   11 res=1;
   while(n)
       if(n&1) res=qmul(res,a,mod);
       a=qmul(a,a,mod);
       n>>=1;
   return res;
}
bool millerrabin(ll n)//Miller Rabin Test
   if(n<3||n%2==0) return n==2;//特判
   11 u=n-1, t=0;
   while(u\%2==0) u/=2,++t;
   ll ud[]={2,325,9375,28178,450775,9780504,1795265022};
   for(11 a:ud)
       11 v=qpow(a,u,n);
       if(v=1||v=n-1||v=0) continue;
        for(int j=1;j<=t;j++)</pre>
        {
           v=qmul(v,v,n);
           if(v==n-1&&j!=t){v=1;break;}//出现一个n-1,后面都是1,直接跳出
           if(v==1) return 0;//这里代表前面没有出现n-1这个解,二次检验失败
       if(v!=1) return 0;//Fermat检验
   }
   return 1;
}
```

1e18 fac decompose

```
using 11 = long long;
11 max_factor, n;
std::map<11, int> MP;
11 gcd(11 a, 11 b) {
    while (b) {
       11 t = a \% b;
        a = b;
        b = t;
    }
    return a;
}
11 qpow(11 x, 11 p, 11 mod) {
    11 \text{ ans} = 1;
    x\%=mod;
    while (p) {
        if (p \& 1) ans = ans * x % mod;
        x = x * x % mod;
        p >>= 1;
    return ans;
}
bool Miller_Rabin(ll p) { // 判断素数
    if (p < 2) return false;
    if (p == 2 \mid \mid p == 3) return true;
    11 d = p - 1, r = 0;
    while (!(d & 1)) ++r, d >>= 1; // 将d处理为奇数
    for (11 k = 0; k < 10; ++k) {
        11 a = rand() \% (p - 2) + 2;
        11 x = qpow(a, d, p);
        if (x == 1 \mid \mid x == p - 1) continue;
        for (int i = 0; i < r - 1; ++i) {
            x = (_int128) x * x % p;
            if (x == p - 1) break;
        if (x != p - 1) return false;
    return true;
}
11 Pollard_Rho(11 x) {
    11 s = 0, t = 0;
    11 c = (11) rand() \% (x - 1) + 1;
    int step = 0, goal = 1;
    11 \text{ val} = 1;
    for (goal = 1;; goal *= 2, s = t, val = 1) { // 倍增优化
        for (step = 1; step <= goal; ++step) {</pre>
```

```
t = ((\underline{\ }int128) \ t * t + c) % x;
            val = (_int128) val * abs(t - s) % x;
            if ((step % 127) == 0) {
                11 d = \gcd(val, x);
                if (d > 1) return d;
            }
        }
        11 d = \gcd(val, x);
        if (d > 1) return d;
   }
}
void fac(11 x) {
    if (x < 2) return;
    if (Miller_Rabin(x)) {
                                        // 如果x为质数
        max_factor = std::max(max_factor, x); // 更新答案
        MP[x]++;
        fac(n /= x);
        return;
    }
    11 p = x;
    while (p \ge x) p = Pollard_Rho(x); // 使用该算法
    while ((x \% p) == 0) x /= p;
    //fac(x);
    fac(p); // 继续向下分解x和p
}
11 cul(ll x, ll y, ll k, ll mod) {
    11 \text{ ans} = 1;
    ans = qpow(x\mbox{mod}, (y+1), mod);
    11 ans2 = qpow(ans, k, mod);
    ans = (ans2 - 1 + mod)\%mod;
    int inv = (qpow(x\%mod, k, mod) - 1 + mod) \% mod;
    inv = qpow(inv, mod - 2, mod);
    return ans * inv % mod;
}
int main()
{
    fac(....)
}
```

segtree

维护区间最值为例:

```
struct T {
```

```
int 1,r,mid;
    int lazy,sum;
}tree[N<<2];</pre>
void build(int rt,int l,int r)
{
    tree[rt].lazy=0;
    tree[rt].l=l;
    tree[rt].r=r;
    if(1==r)
        tree[rt].sum=a[1];//依据初始化类型具体而定,置0可以看为空树,但是1,r均初始化。
        return ;
    }
    int mid=tree[rt].mid=l+r>>1;
    build(rt<<1,1,mid);</pre>
    build(rt << 1 | 1, mid+1, r);
}
void push_down(int rt)
    if(tree[rt].lazy)
        tree[rt<<1].sum+=tree[rt].lazy;</pre>
        tree[rt<<1].lazy+=tree[rt].lazy;</pre>
        tree[rt<<1|1].sum+=tree[rt].lazy;</pre>
        tree[rt<<1|1].lazy+=tree[rt].lazy;</pre>
        tree[rt].lazy=0;
    }
}
void push_up(int rt){
    tree[rt].sum = max(tree[rt<<1].sum,tree[rt<<1|1].sum);//看情况而定
}
void update(int rt,int 1,int r,int v)
{
    if(tree[rt].r<1||tree[rt].l>r) return ;
    if(tree[rt].l>=l&&tree[rt].r<=r)</pre>
    {
        tree[rt].sum+=v;
        tree[rt].lazy+=v;
        return ;
    }
    push_down(rt);
    if(tree[rt].mid>=1) update(rt<<1,1,r,v);</pre>
    if(tree[rt].mid<r) update(rt<1|1,1,r,v);
    push_up(rt);
}
int query(int rt,int 1,int r)
    if(tree[rt].r<1||tree[rt].l>r) return 0;
    if(tree[rt].l>=l&&tree[rt].r<=r) return tree[rt].sum;</pre>
    push_down(rt);
    if(tree[rt].mid<1) return query(rt<<1|1,1,r);
    if(tree[rt].mid>=r) return query(rt<<1,1,r);</pre>
```

```
return max(query(rt<<1|1,1,r), query(rt<<1,1,r));
}</pre>
```

BIT (Fenwick)

template from jiangly 注意 add 当下, sum (x) = x-1的sum 有超时风险, 因为vector

```
template <class T>
struct Fenwick {
    int n;
    vector<T> a;
    Fenwick(int n = 0) {
        this->n = n;
        a.assign(n, T());
    }
    void add(int x, T v) {
        for (int i = x + 1; i \le n; i += i \& -i) {
            a[i - 1] += v;
        }
    }
    T sum(int x) {
        auto ans = T();
        for (int i = x; i > 0; i -= i \& -i) {
            ans += a[i - 1];
        return ans;
    }
    T rangeSum(int 1, int r) {
        return sum(r) - sum(1);
    }
};
```

exgcd

求解: ax+by = gcd(a,b)

```
11 exgcd(11 a,11 b,11 &x,11 &y)
{
    if(!b)
    {
        x=1;y=0;
        return a;
    }
}
```

DSU

```
struct DSU {
   vector<int> f, siz;
    DSU() {}
    DSU(int n) {
       init(n);
    }
   void init(int n) {
       f.resize(n);
       iota(f.begin(), f.end(), 0);
        siz.assign(n, 1);
   }
    int find(int x) {
       while (x != f[x]) {
           x = f[x] = f[f[x]];
        }
       return x;
    }
    bool same(int x, int y) {
       return find(x) == find(y);
   }
    bool merge(int x, int y) {
       x = find(x);
       y = find(y);
       if (x == y) {
           return false;
       }
        siz[x] += siz[y];
       f[y] = x;
       return true;
    }
   int size(int x) {
```

```
return siz[find(x)];
};
```

tarjan

使用tarjan算法求强连通分量。

有向图 (对于无向图还需要加一个int fa(若需要求环))

```
vector<int>adj[N];
int dfn[N],low[N],stk[N],top,dn,vis[N];
int nsz,nl[N];
void tarjan(int u)
    vis[u] = 1;
    dfn[u] = low[u] = ++dn;
    stk[++top] = u;
    for(auto v:adj[u]){
        if(!dfn[v]){
            tarjan(v);
            low[u] = min(low[u],low[v]);
        else if(vis[v]){
            low[u] = min(low[u],dfn[v]);
        }
   }
    if(low[u]==dfn[u]){//强连通分量
        nl[u] = ++nsz;
       while(stk[top]!=u)nl[stk[top]] = nl[u],vis[stk[top]] = 0,top--;
       vis[u] = 0, top--;
}
```